Data Science & AI

DATA STRUCTURES Through Python

STACK

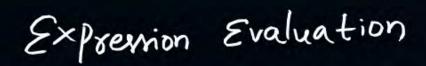
By- Satya sir

Lecture No.-05

Recap of Previous Lecture









Topics to be Covered











Stack Implementation

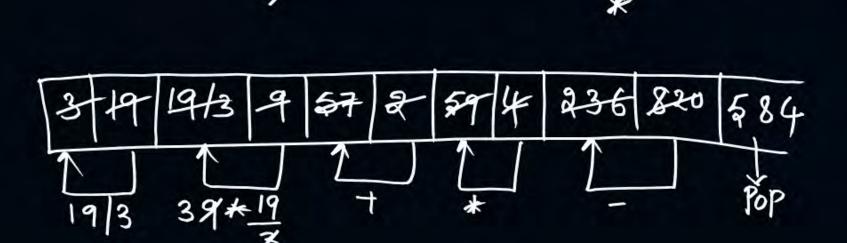
- Using Lists
- Using deque
- Using Lifoqueue
- Stack Using Queue (simple queue)



ExPression Evaluation

- Postfix Exp

- Refix EXP.



A 3 19/3 6 382 3647 H 396

POP

19 *62

Evaluate

/Postfix Expression: 34*7+3/6*2-47+*=396

Prefix expression: - 820 * 4 + 2 * 9 / 19 3 = 584





Stack (Last-In-first-out): Insertion (Rish), Deletion (Pop), Access (Peek)

overflow

Underflow

Stack can be Implemented Using Python in either of 3 ways:

- 1) Using Lists (arrays)
- 2) Using Lifoqueue clans
- 3) Using deque clan

Stack implementation in python Using Lists

Configuous Membra Drawback: Memory Wastarge: Internal [1234]

Creating a stack def create_stack(): stack = [] return stack

Push Time Complexity: ()(1) Pop Time Complexity: O(1) Peek Time Complexity: (O(1)

Creating an empty stack

→def check_empty(stack):

return len(stack) == 0 seturns Tour if Empty, False otherwise

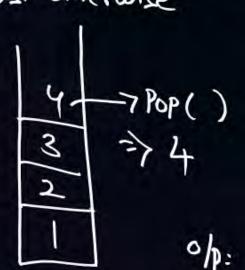
Adding items into the stack

→ def push(stack, item):

stack.append(item)

print("pushed item: " + item)

7 will Insert Element at the top of Stack.



In other lang: Arrays => State Memory

In Python: Arrays (lists) => Dynamic Memory

(No need of checking white

Removing an element from the stack def pop(stack):

if (check_empty(stack)): return "stack is empty" return stack.pop() -

stack = create stack()

But, In Pop() Or Peck(), Empty Condition need to be verified.

Driver Code

stack = create_stack() | list

push(stack, str(1))

push(stack, str(2))

push(stack, str(3))

push(stack, str(4))

print("popped item: " + pop(stack))

print("stack after popping an element: " + str(stack))

oh: Popped item: 4 Stack after Popping 1 2 3









> Not Suitable for multithreading applications

from collections import deque (Random memory No Internal allo Cation - No Internal print('\nElements poped from my_stack:') frommenlation rint(my_stack.pop()) # C print(my_stack.pop())# 6 my_stack = deque() # constructor $print(my_stack.pop()) \neq \omega$

append() function is used to push

element in the my_stack

my_stack.append('a')

my_stack.append('b')

my_stack.append('c')

print('Initial my_stack:') print(my_stack) # a b c Recursive Procen

TI, T2, T3, f(5) > f(4) 7, T4, T5 f(1) es(2) e f(3)

print('\nmy_stack after elements are poped:') print(my_stack) # Empty, Nothing is Printed

Count: Race condition data Inconsistency

T.C: Push Pop | Peek: O(1) P3

when thread of different





9, size () returns Number of Elementy Present

```
from queue import LifoQueue (Handle Race Condition issue, with Synchronization Tools)

print('\nElements poped from the my_stack')

print(my_stack.get()) # z

print(my_stack.get()) # y

print(my_stack.get()) # y

print(my_stack.get()) # z

print(my_stack.get()) # z

print(my_stack.get()) # z

print(my_stack.get()) # z

print("\nStack is Empty: ", my_stack.empty())

my_stack.put('x')

my_stack.put('y')

my_stack.put('z')
```

print("Stack is Full: ", my_stack.full()) # stack is Full: Folse print("Size of Stack: ", my_stack.qsize()) # 3

T.C: Posh Pop Peek(): O(1)





nplementation

- Using Lists: Contiguous memory allocation, Memory wastarge (or) Lens Utilization

- Using deque: Not suitable for multithreading applications, but fast in Revorming operations

- Using LifoQueue: Handle multithreading Effectively, but slow compared with degree Implementation.



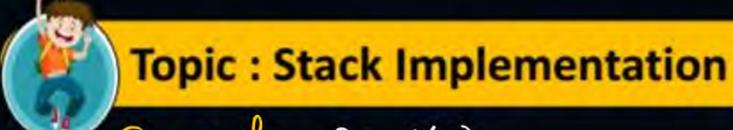


(LIFO)
Stack Implementation Using Simple Queue (FIFO)

If uses 2 queues for Implementing stack Nature

It can be Responsed in either of 2 ways:

- 1) By making Rush operation costly = Push: O(1)
- 2) By making Pop operation costly = Push: O(1)
 Ly Pop: O(n)



Let Q1, Q2 be 2 Quenes, S be Stack

Push(10), Rish(20), Rish(30), POP(), Push(40), POP(), POP(), POP()

Expected olp: 30,40,20,10

Pos	sh costly: Push: O(n)
i) de	equencall Elements from Q1) shift of
2) &	inquene them into Q2 Jai to E
3) 0	rqueue new Element into al
4) 5	hiff Elements back from 62 to Q1
Por	Poperation Pop. O(1)
deque	ne drom Q1. nd Print it.

Q1 (30 20 10)

shift from Rish (10) All to 622 Al [0]	Pop() dequeve from Q1 and Print Olp: 30 Q1 20 10 Q2	
している。 日 日 日 日 日 日 日 日 日 日	acqueue Osom KI and D.	

POP() dequene from 61 and Print 0/p: 20 Q1 10 92 POPC) Dequene from &1 and Print Actual of sequence 30,40,20,10 == Expected ofp.



2 mins Summary

Push(10), Rish (20), Rish(30), POP(), Rish(40), Pop(), Pop(), Pop(), Pop()

Expected 0/p: 30,40,20,10



Pop operation costly:

Push (): enquene Element înto 621

0(1)

Pop() Operation O(n)

- 1. Shift (n-1) Elements from Q1 to Q2
- 2. Dequeue last Element from Q1 and Print
- 3. Shift (n-1) Elements back to &1 from &2

Q1 [42 [Posh (10) Q1 10 Q2 Push (20) 61 [10 20] 62 Push (30) Q1 10/20/30) Q2 Pop() 1. Q1->Q2 Q1 (30) Q2 [10/20] 2. Degreve from Q1 and Print Q1p: 30 3- Q2->01 01 (10/10) Q7

	Posh (40)
	Q1 [10/20/40] &2 [
	Q1 40 (20 20)
3.	0 P:40 V Q 10120
	POPC) (8) 20 (52 [10)
g	
1.1	3. 01 10 62
	Actual 0/p: 30, 40,20,10
	11. 30, 40,20,10



THANK - YOU