I/P → O/P

I/P → Array of length N ✓

TC ∝ f(size of input) ✓

↳ Amount of time taken by an algo. to run as function of size of input.

x → 10 machines/servers ✓

x → 3 machines ✓ ∝ cost ✓

Q→ check if given input N is prime?

↳ A +ve number with exactly 2 factors 1 & N.

1 → only 1 factor ✗

check if 'i' is a factor of N?

((N % i) == 0) ✓

Bruteforce

-1 → 1, -1

factors = 0;

for (i = 1; i <= N; i++) { 1–10

if (N % i == 0) factors++;

}

return (factors == 2);

N = 10 ✓

steps = 10 ✓

N = 10: a b → 1*10, 2*5 → {1, 2, 5, 10}

N = 20: 1*20, 2*10, 4*5 → {1, 2, 4, 5, 10, 20}

N = 13: 1*13 ⇒ prime

N = a*b, a <= b

↳ >1 ⇒ N is not prime ✓

b = N/a

a <= b

a <= N/a ⇒ a^2 <= N

⇒ a <= √N

N = 25: 1*25, 5*5

3√25 = 2...

✓ for (a = 2; a*a <= N; a++) {

if (N % a == 0)

return false;

}

return true;

N = 10

steps = √N = √10 = 3 ✓ ✓

$A_1 \rightarrow 2^N$ steps

$A_2 \rightarrow N^3$ steps ✓

1 step = 1 ms

Rate of growth of time wrt input size ✓

	2^N	N^3
N = 2	$2^2 = 4$ ✓	$2^3 = 8$
N = 5	$2^5 = 32$ ✓	$5^3 = 125$
N = 10	$2^{10} = 1024$	$10^3 = 1000$ ✓
N = 100	$2^{100} \approx 10^{30}$ very large	$100^3 = 10^6$ ✓

$TC = O(2^N)$

$N \rightarrow 2N^2 + N - 2$

$TC = O(N^2)$

$N \rightarrow 2 - 3N + \sqrt{N}$

$TC = O(N)$

$TC = O(N^3)$

$N \rightarrow N^{1.1} + 2 + 3N - \sqrt{N}$

$TC = O(N^{1.1})$

Quiz

1)
```
for(i = 0; i < N; i += 2) {
    print(i);
}
```

0, 2, 4, 6 ... N

$\frac{N}{2}$ steps $= \frac{1}{2} * N$

$TC = O(N)$

2)
```
for(i = 0; i < N; i++) {
    for(j = i+1; j < N; j++) {
        ans += i * j;
    }
}
```

i = 0	j → 1 — N
i = 1	j → 2 — N
i = 2	j → 3 — N
⋮	
N-1	j → N-1 — N

$N + (N-2) + (N-3) + \dots 1 = \leq N$

$= \frac{N*(N+1)}{2}$

$\frac{1}{2}N^2 + \frac{1}{2}N \leftarrow \frac{N^2}{2} + \frac{N}{2}$

$TC = O(N^2)$ ✓

3) `j = 0;`
`for (i = 0; i < N; i++) {`
`while (j < N && a[i] <= a[j]) { j++; }` ← N times
`sum += a[j];`
`}`

$i \to N$ times
$j \to N$ times
$N + N = 2N$

$j = 0 + 2 \dots N$ ✓

$TC = O(N)$ ✓✓

4) `i = N; a = 0;`
`while (i > 0) { a += i; i /= 2; }`

$1/2 = 0$ ✓

$N \to \frac{N}{2} \to \frac{N}{2^2} \to \frac{N}{2^3} \dots \frac{N}{2^k} = 1$

after k steps

$\Rightarrow N = 2^k \Rightarrow k = \log_2(N)$

$TC = O(\log(N))$ ✓

$O(\log_2(N))$

$\log_b(b^k) = k$ ✓

$\log_b(a) = \frac{\log(a)}{\log(b)}$

$\log_2 N = \frac{\log(N)}{\log(2)} = \frac{1}{\log(2)} * \log(N)$

4) `for (i = 1; i < N; i = i * 2) { print(i); }`

$1 \to 2 \to 2^2 \to 2^3 \to \dots \to 2^k = N$

after k steps

$k \approx \log_2(N)$

$TC = O(\log(N))$

5) `k = 0;`
`for (i = 0; k < N; i++) {`
`k = k + i;`
`}`

$i \to 0 \quad 1 \quad 2 \quad 3$

$k = 0 + 0 + (1) + (2) + (3)$

$(1 + 2 + 3 + \dots x) = N \Rightarrow \frac{x * (x+1)}{2} = N$

$\#steps = x$

$\Rightarrow x^2 + x = 2N$

$\Rightarrow x^2 + x - 2N = 0$

$ax^2 + bx + c = 0$

$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$

$x = \frac{-1 + \sqrt{1 - 4(1)(-2N)}}{2} = \frac{-1 + \sqrt{1 + 8N}}{2}$

$TC = O(\sqrt{N})$ ✓

Space Complexity → Extra space to convert given i/p to desired o/p.

$\propto f(\text{size of input})$

No extra space used ⇒ SC = O(1)

Q→ For a given input array, Find the product of freq. of every element in the array?

$|A| = N <= 10^5$

$1 <= A[i] = X <= 10^5$ → |F|

Eg→ A → [3 2 1 1 3 5 5 5] (indices 0 1 2 3 4 5 6 7)

A[i] → freq

3 → 2
2 → 1
1 → 2
5 → 3
4 → 0

Ans = 2 * 1 * 2 * 3 = 12

f(A, N, X)

∀ i F[i] = 0 → SC = O(X)

```
for(i=0; i<N; i++) { F[A[i]]++; }
ans = 1;
for(i=0; i<X; i++) {
    if(F[i] > 0)  ans = ans * F[i];
}
return ans;
```

Find freq. of all elements in A[].

F[] → output

SC = O(X) → incorrect

SC = O(1)

```
int a = 5;
```

4 byte = 4*8 = 32 bits

5 = 00---0---00000101

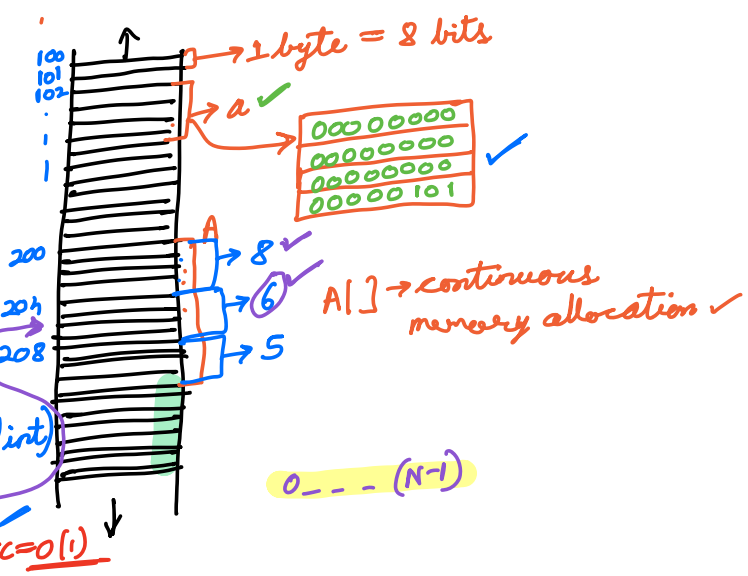
```
int A[] = {8, 6, 5};
```

4*3 = 12 bytes

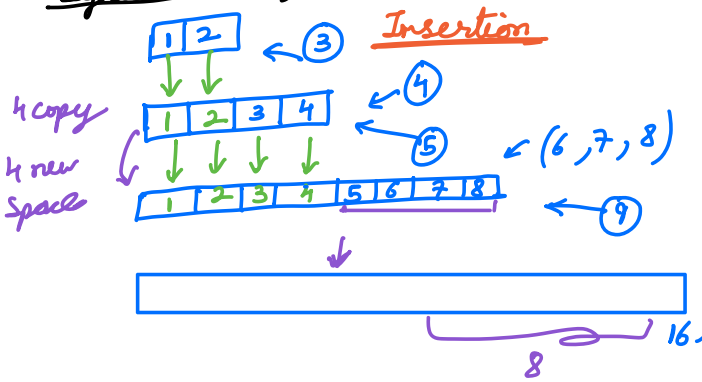
A[1] = base address + (1 * size of int)

= 200 + (1*4) = 204

TC = O(1)



Dynamic Array → ArrayList/Vector



TC = O(1) Access time ✓

↓

continuous memory ✓

X time to copy

⇒ TC = O(1) for X insertions after copy.

↳ TC ≈ O(2) ✓

↓

TC = O(1) ✓

✓ Inserting X elements is taking 2X time