

Assignment-2

Ques.

1) Write linear search code in a sorted array with minimum comparisons.

Ans -

function Is_Carr($a[]$, n, k)
(where n = number of elements and k is the element
to be searched)

{
flag = 0 // initial value of flag
for (i = 0; i < n; i++)
{

 if ($k == a[i]$)
 {
 flag = 1
 break
 }

 }

}
return flag // returns 1 if element found

3
return 0 // returns 0 if element not found

~~QUESTION~~

Q-2 Write pseudo code for iterative and recursive insertion sort. why is the cost added on line sorting. why?

Ans:-

a) Recursive approach:-

function insertion arr[n]

{

if ($n <= 1$)

{

return

}

insertion arr, n-1

last = arr[n-1]

$j = n - 2$

while ($j >= 0$ & $arr[i] > last$)

{

arr[j+1] = arr[j]

$j = j - 1$

arr[0] = last;

3

b) iterative approach:

function insertion (arr, τ , m)

{
for (i from 1 to $m-1$)

{
 key = arr[$i+1$] - to be inserted

$\tau = \tau - 1$

 while ($\tau \geq 0$ and arr[τ] $>$ key)

{
 arr[$\tau + 1$] = arr[τ];

$\tau = \tau - 1$

insert τ

arr[$\tau + 1$] = key;

3

3

new insertion point is called on the left, since in

the algorithm, the element ~~for~~ with ~~some~~ same
elements are to be inserted with the same
, value, in the rotated array too.

Algorithmic Tools:

1. insertion point →
2. Merge sort →
3. Heap sort →
4. quick sort →
5. Radix sort →
6. Count sort →
7. Bubble sort →
8. selection sort →

~~Time complexity = O(n log n)~~

1-3 Main concern is the sorting algorithm described in the class.

Algorithm	Time Complexity	Space Complexity
insertion sort	$O(n^2)$	$O(1)$
merge sort	$O(n \log n)$	$O(n)$
quick sort	$O(n^2)$	$O(1)$
bubble sort	$O(n^2)$	$O(1)$
selection sort	$O(n^2)$	$O(1)$
radix sort	$O(nk)$	$O(k)$
counting sort	$O(n+k)$	$O(n+k)$

→ Add for worst cases.

4 Discuss for all the sorting Algorithms,
whether they are In-place, stable or
on-line.

	In-place	stable	On-line
Sort	1	1	0
Bubble	1	0	0
Heap	1	1	1
Inversion	1	0	0
Duck	1	0	0
Count	0	1	0
Merge	0	1	0
Radix	0	1	0
Selection	1	0	0

- Write Pseudocode for Binary Search - compare
it's space and time complexity.

Ans. 5

Binary Search

iterative Approach:

function BS(arr[], n, key) /* condition →
array must be sorted */

{

 l = l;

 r = 0;

 while (l <= r)

{

 mid = (l + (r - 1)) / 2

 if (arr[mid] == key)

{

 return true;

}

 else if (arr[mid] > key)

{

 l = mid - 1;

}

 else

{

 r = mid + 1;

}

if loop ends if element not found
return false;

Recursive Approach:-

Function $BS(\text{arr}, l, r, \text{key})$

```
{ if (l <= r)
    {
        mid = (l + r - 1) / 2
        if (arr[mid] == key)
            return true;
        else if (arr[mid] < key)
            return Recursion BS(arr, mid+1,
                r, key)
        else
            return BS(arr, 0, mid-1, key)
    }
    return false;
}
```

Linear search (Iteration)

$$T.C = O(n)$$

$$S.C = O(1)$$

Binary Search (Iteration)

$$T.C = O(\log n)$$

$$S.C = O(1)$$

Linear Search (Recursive Ver)

$$T.C = O(n)$$

$$S.C = O(1)$$

No stack used

Binary Search (Recursive Tree)

$$T.C = O(\log n)$$

$$S.C = O(\log n)$$

Stack to Recursion calls

Q-6 Write down the recurrence relation for Binary Search.

Ans -

Pseudo code :-

$$T(n) \Rightarrow BS(a, i, j, n) = O(1)$$
$$\text{mid} = (i + j)/2 = O(1)$$

if ($a[\text{mid}] == x$)

return $a[\text{mid}]$

if ($a[\text{mid}] > x$)

return $BS(a, i, \text{mid}-1, n) \leftarrow T(n/2)$

^{ed RE}
for medium BS (a, mid, b, n) $\leftarrow T(n/2)$

Hence,

$$T(n) = T(n/2) + 1$$

Q-7 which is the best sorting for practical uses? explain

Ans - quicksort is often the preferred sorting algorithm for practical applications. It's one of the most efficient sorting algorithms and it's generally considered the fastest. Some real world-applications of sorting include: searching, element uniqueness, telephone dictionaries, ranks based on scores.

Q-13 Bubble sort scans whole array even when array is sorted. (or it be modified so that it doesn't scan the whole array once it is sorted.)

Ans - sol. Bubble sort can be optimized for better usage for partially or fully sorted arrays:-

function BubbleSort arr (arr[], n)

n = length

for (i=0; i<n; i++)

{

swapped = False;

for (j=0; j<n-i-1; j++)

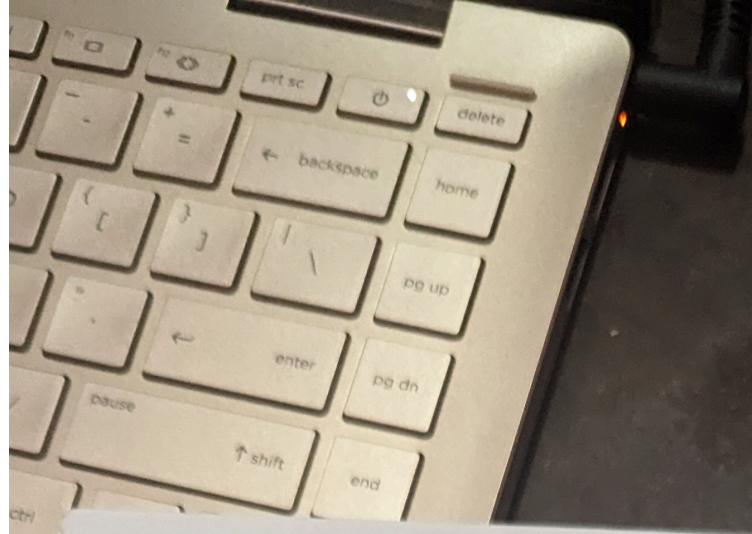
{ if (arr[j] > arr[j+1])

{

swap;

}

}
if (not-swapped)
break;



In this modification, the swatted flag has been used to define whether any swats were made during the inner loop.