

[James Dolan]

⑥ 22/7/19

Given a monad  $T: C \rightarrow C$   
construct the  
algebras for the monad

left adjoint ..... t.b.c.

①

24/7/19

Example

theory of a cocategory  
object

syntactic category  
is category of  
small categories.

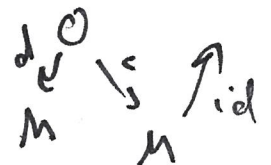
breaking the construction  
into stages:

primary operations:

domain, codomain, identity

secondary operations:

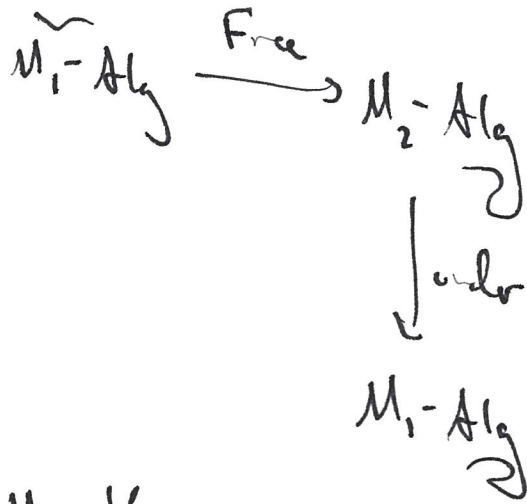
composition



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$M_2$ : Monad on  $\widetilde{M_1\text{-Alg}}$



Call this a monad tower.  
We get a chain of adjunctions:

$$\widetilde{C} \xrightleftharpoons[R_1]{L_1} M_1\text{-Alg} \xrightleftharpoons[R_2]{L_2} M_2\text{-Alg} \dots$$

Question:

compose adjunctions  
to get composite monad:

$$M_2 \circ M_1 := R_1 R_2 L_2 L_1$$

Call this  $N$ .

right  
adjoint  
comparison  
functor

$$\widetilde{M_2\text{-Alg}} \longrightarrow \widetilde{N\text{-Alg}}$$

when this is an  
equivalence of categories,  
we  
call these monads  
composable

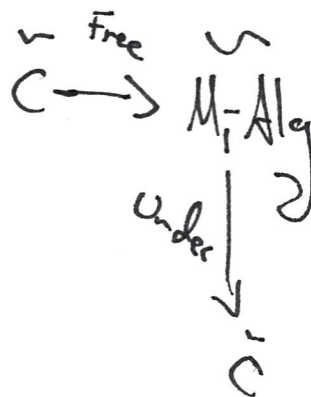
$M_1$ : Monad on a  
 $\widetilde{C}$ : locally pres category

Eilenberg Moore category  
of algebras

(Kleisli is subcategory)

is also locally  
presentable

ie. has a theory



the theory is: take the  
sketch for  $\widetilde{C}$  and append  
"co- $M_1$  algebra"

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Back to co-category example

Sketch



Object A  
"arrows"

primary ops:

$$A \xrightarrow{s} A$$

$$A \xrightarrow{t} A$$



with  
 $s^2 = s, t^2 = t,$   
 $st = s, ts = t$

The theory  
of a cograph.

Syn Category

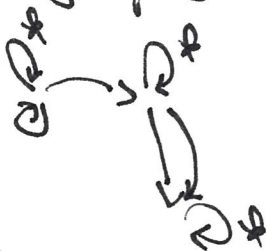


$$\tilde{C} = \tilde{\text{Set}}$$

The  $M_1$  monad

Algebras for  
this monad  
are actions of  
~~the~~ the monoid  
 $\{1, s, t\}$

ie. graphs



with vertices as  
degenerate special  
loop edges.

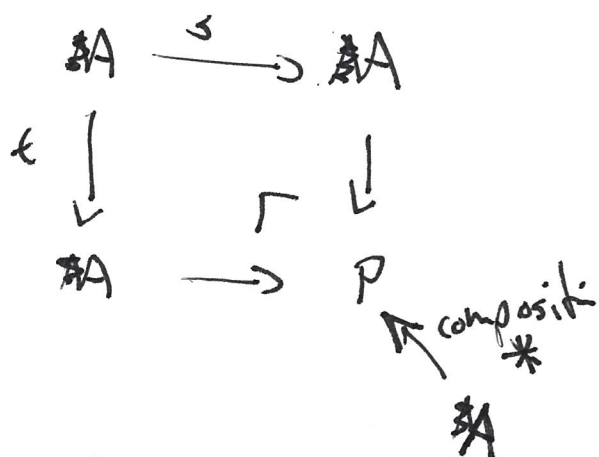
(simplicial sets  
with 0 & 1 in  
• objects)

Johns Hole

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Sketch



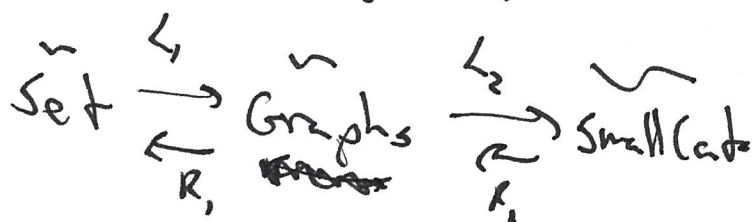
$\epsilon$  identity  
 $\delta$  assoc

we have a morphism  
 $\epsilon$  into a colimit

Synthetic category

The  $M_2$  monad on  
 the category of graphs  
 is the free category monad.

The  $M_2$  algebras are the  
 small categories.



$L_2 L_1$ : free cat on free graph  
 on a set of arrows.

$$R_1 R_2 L_2 L_1 : X \mapsto 3 \cdot X$$

this is the same as

$$R_1 L_1 : X \mapsto 3 \cdot X$$

So the comparison functor

$$M_2\text{-Alg} \longrightarrow N\text{-Alg}$$

fails to be an equivalence  
 of categories!

Theorem you can't get this category  
 as the algebras of any monad  
 on Set.

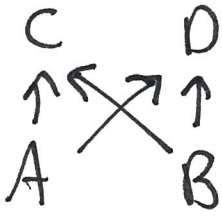
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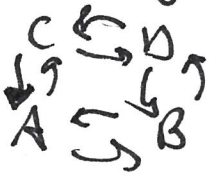
digression

What is a  
functor from a  
poset to a group?

eg 2



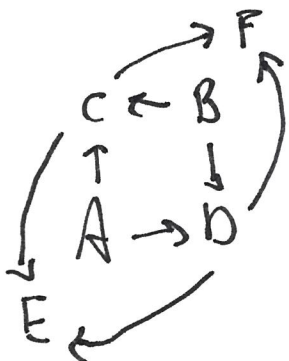
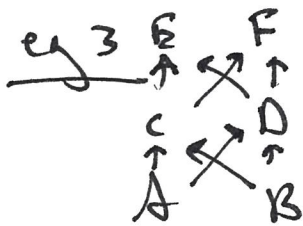
as a groupoid:



just retain  
end-morphisms of A:

$$A^D \cong \mathbb{Z}$$

the homotopy group  
of the circle.



this is the  
suspension of eg3.  
An octahedron (sphere).

Daniel Kan  
references  
Peter May  
references

Take

$$\text{Poset} \xrightarrow{\epsilon} \text{Small Cat}$$

Small Groupoid  
 $\uparrow L$  freely  
injected morphism

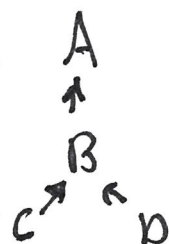
Consider adjoint pairs

$$\text{Small Gpd} \xrightleftharpoons[\epsilon]{R_G} \text{Sm Cat} \xrightleftharpoons[\epsilon]{L_G} \text{Sm Gpd}$$

$$\text{Poset} \xrightleftharpoons[\epsilon]{R_P} \text{Small Cat} \xrightleftharpoons[\epsilon]{L_P} \text{Poset}$$

Homotopy theory

eg. 1



$$F = \text{Fund} L_G R_P$$

Fundamental groupoid of a poset  
F preserves connectedness of a small category.

Geometric realization of a poset

the nerve of a poset is  
a simplicial set

defined as

For each composable  
n-chain of morphisms  
attach a n-simplex



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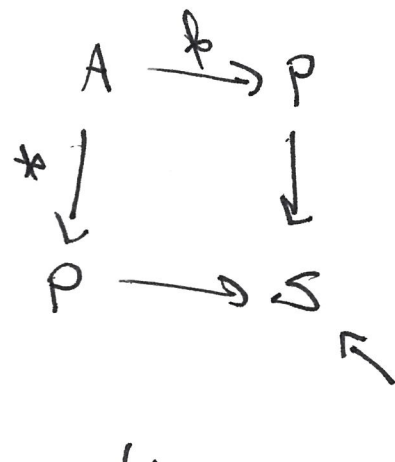
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Example sketch with a  
tertiary operation

building on a co-category  
sketch:

Silly example:

assigning an arrow for  
each commuting square



relies on composition and  
so is a tertiary operation.

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7/24/19

Given a monad:  $M: \tilde{C} \rightarrow \tilde{C}$

Kleisli: Free  $M$ -algebras

inclusion

Eilenberg-

Moore: all  $M$ -algebras

n.p.o.v.

The free  $M$ -algebra category  
is the 2-colimit  
of the monad.

The  $M$ -algebra ~~category~~  
is the 2-limit of the monad.

Kleisli category

the Kleisli  
freely introduces morphisms  
while preserving ~~equivalences~~  
colimits but ~~does~~ may not  
itself be co-complete, because  
of the new morphisms.

EM category

freely add colimits  
while preserving the  
old colimits.