

Array Manipulation Patterns

1) Product of Array except self.

Given an array `nums`, return an array `answer` where `answer[i] = product of all the elements of nums except nums[i]`

required: $O(N)$ solution

Note: do not use division operation.

Approach #1: (Brute Force Approach):

```
int n = nums.length;
```

```
int[] answer = new int[n];
```

```
for (int i = 0; i < n; i++) {
```

```
    int ans = 1;
```

```
    for (int j = 0; j < i; j++) {
```

```
        ans = ans * nums[j];
```

```
    }
```

```
    for (int k = i + 1; k < n; k++) {
```

```
        ans = ans * nums[k];
```

```
    }
```

```
    answer[i] = ans;
```

```
}
```

```
return answer;
```

T.C: $O(N^2)$

S.C: $O(1)$

Approach #2:- (using division operators)

find # of zeroes if # of zeroes > 1 then all elements of answer array is 0

else

find product of non zero elements.

```
int n = nums.length;
```

```
int[] answer = new int[n];
```

```
int zeroes = 0;
```

```
int prod = 1;
```

```
for (int i = 0; i < n; i++) {
```

```
    if (nums[i] == 0) {
```

```
        zeroes++;
```

```
    } else { prod *= nums[i];
```

```
    }
```

```
if (zeroes > 1) {
```

```
    return answer;
```

```
}
```

```
for (int i = 0; i < n; i++) {
```

```
    if (zeroes == 1) {
```

```
        for (int i = 0; i < n; i++) {
```

```
            if (nums[i] == 0) {
```

```
                answer[i] = prod;
```

```
            }
```

```
        }
```

```
    }
```

```
else {
```

```
    for (int i = 0; i < n; i++) {
```

```
        answer[i] = prod / nums[i];
```

```
    }
```

```
}
```

```
return answer;
```

T.C: $O(N)$

S.C: $O(1)$

Approach #3: (using prefix and postfix arrays)

Find postfix and prefix product arrays.

Say if $\text{nums} = [1, 2, 3, 4]$

then

prefix = $[1, 1, 2, 6]$

postfix = $[24, 12, 4, 1]$

then final output = $[24, 12, 8, 6]$

```
int n = nums.length;
```

```
int[] answer = new int[n];
```

```
int[] prefix = new int[n];
```

```
for
```

```
int[] postfix = new int[n];
```

```
prefix[0] = 1;
```

```
for (int i = 1; i < n; i++) {
```

```
    prefix[i] = prefix[i-1] * nums[i-1];
```

```
}
```

```
postfix[n-1] = 1;
```

```
for (int i = n-1; i > 0; i--) {
```

```
    postfix[i-1] = postfix[i] * nums[i];
```

```
}
```

```
for (int i = 0; i < n; i++) {
```

```
    answer[i] = prefix[i] * postfix[i];
```

```
}
```

```
return answer;
```

TC: $O(N)$

SC: $O(N)$

Approach #4:- (Prefix array with O(1) space)

```
int n = nums.length;
```

```
int[] answer = new int[n];
```

```
answer[0] = 1;
```

```
for (int i = 1; i < n; i++) {
```

```
    answer[i] = answer[i-1] * nums[i-1];
```

```
}
```

```
int lastZeroIndex = -1;
```

```
int postProd = 1;
```

```
for (int i = n-1; i >= 0; i--) {
```

```
    answer[i] = answer[i+1] * postProd;
```

```
    postProd *= nums[i];
```

```
}
```

```
return answer;
```

T.C: $O(N)$

S.C: $O(1)$

2. Move zeroes

Given an array `nums`, you have to move all zeroes to end restoring the relative order of non-zero elements.

example:-

`nums = [0, 1, 0, 3, 12]`

`o/p :- [1, 3, 12, 0, 0]`

Approach #1:- (Brute Force)

Count zeroes

Store non zero elements in new array then store them back to

original array.

```
int zeroes = 0;
```

```
int n = nums.length;
```

```
int[] ans = new int[n];
```

```
for (int i = 0; i < n; i++) {  
    if (nums[i] == 0) {  
        zeroes++;  
    }  
}
```

not needed.

```
int k = 0;
```

```
for (int i = 0; i < n; i++) {
```

```
    if (nums[i] != 0) {
```

```
        ans[k++] = nums[i];
```

```
    }
```

```
while (k < n) {
```

```
    ans[k++] = 0;
```

```
for (int i = 0; i < n; i++) {
```

```
    nums[i] = ans[i];
```

```
}
```

```
return ans;
```

T.C: $O(N)$

S.C: $O(N)$

Approach #2: (Using 2 pointers) [In place modify]

```
int n = nums.length;
```

```
int lastZeroAtIndex = 0;
```

```
for (int i = 0; i < n; i++) {
```

```
    if (nums[i] != 0) {
```

```
        nums[lastZeroAtIndex++] = nums[i];
```

```
    }
```

```
}
```

```
while (lastZeroAtIndex < n)
```

```
    nums[lastZeroAtIndex++] = 0;
```

```
}
```

```
return nums;
```

T.C: $O(N)$

S.C: $O(1)$

Summary

How to avoid extra space in array manipulation problems?

- 1) Use two pointers to in place modify array.
- 2) Instead of making copy, overwrite existing array.
- 3) encode information into array itself.
- 4) Reuse input array as output
- 5) Sometimes, you can encode two values into one array slot, e.g. Bit manipulation, modulo 3

