

1 Performance of Models

I ran my tests on the SEASnet GNU / Linux with Java version 1.8.0_51. I tried gathering memory information and cpu information by running `/proc/meminfo` and `/proc/cpuinfo` respectively but I did not have those permissions on the servers.

1.1 Synchronized

```
java UnsafeMemory Synchronized 4 1000000 6 5 6 3 0 3 - 1599.20 ns / transition
java UnsafeMemory Synchronized 8 1000000 6 5 6 3 0 3 - 2930.15 ns / transition
java UnsafeMemory Synchronized 16 1000000 6 5 6 3 0 3 - 4930.45 ns / transition
java UnsafeMemory Synchronized 8 1000000 12 5 6 3 0 3 - 2976.33 ns / transition
java UnsafeMemory Synchronized 8 1000 6 5 6 3 0 3 - 33684.4 ns / transition
```

The provided `SynchronizedState` class is a 100% reliable and DRF method because it utilizes Java's `synchronized` keyword which prevents thread interference and allows for atomicity when writing to variables. However, by ensuring this, it is also a very slow method.

1.2 Null

```
java UnsafeMemory Null 4 1000000 6 5 6 3 0 3 - 432.974 ns / transition
java UnsafeMemory Null 8 1000000 6 5 6 3 0 3 - 1934.32 ns / transition
java UnsafeMemory Null 16 1000000 6 5 6 3 0 3 - 4237.49 ns / transition
java UnsafeMemory Null 8 1000000 12 5 6 3 0 3 - 2060.19 ns / transition
java UnsafeMemory Null 8 1000 6 5 6 3 0 3 - 29488.6 ns / transition
```

The `NullState` class is reliable in the sense that it passes all the cases. As expected, it outperforms `Synchronized` in every test case since it just returns dummy values.

1.3 Unsynchronized

```
java UnsafeMemory Unsynchronized 1 1000000 6 5 6 0 3 - 51.0469 ns / transition
java UnsafeMemory Unsynchronized 2 1000000 6 5 6 0 3 - freezes
java UnsafeMemory Unsynchronized 2 100000 6 5 6 0 3 - 1013.94 ns / transition
java UnsafeMemory Unsynchronized 4 1000 6 5 6 0 3 - 14679.0 ns / transition
java UnsafeMemory Unsynchronized 8 1000 6 5 6 0 3 - 32599.8 ns / transition
```

The `UnsynchronizedState` class is not DRF, since it has the same code as `SynchronizedState` but does not use the `synchronized` keyword. It also is not reliable and runs into frequent sum mismatches, especially as you parallelize and increase the amount of threads, like I did in the last two test cases.

1.4 GetNSet

```
java UnsafeMemory GetNSet 2 1000 6 5 6 3 0 3 9208.09 ns / transition  
java UnsafeMemory GetNSet 4 1000 6 5 6 3 0 3 - 19755.2 ns / transition  
java UnsafeMemory GetNSet 8 1000 6 5 6 3 0 3 - 39110.5 ns / transition  
java UnsafeMemory GetNSet 8 10000 12 5 6 3 0 3 - 10305.3 ns / transition
```

I implemented this class using Java's AtomicIntegerArray. Although the get and set methods of the class are volatile / atomic by themselves, they are used together in the code and therefore can cause race conditions. As a result, it runs into frequent sum mismatches as well. Comparing the 3rd test case to the last tests for Unsynchronized and Synchronized, it seems that although GetNSet is between Unsynchronized / Synchronized in implementation, it is slower than both.

1.5 BetterSafe

```
java UnsafeMemory BetterSafe 4 1000000 6 5 6 3 0 3 - 570.405 ns / transition  
java UnsafeMemory BetterSafe 8 1000000 6 5 6 3 0 3 - 1208.86 ns / transition  
java UnsafeMemory BetterSafe 16 1000000 6 5 6 3 0 3 - 2462.92 ns / transition  
java UnsafeMemory BetterSafe 2 1000000 8 5 6 3 0 3 - 740.888 ns / transition  
java UnsafeMemory BetterSafe 8 1000 6 5 6 3 0 3 - 41923.0 ns / transition  
java UnsafeMemory BetterSafe 8 100 12 5 6 3 0 3 - 295414 ns / transition
```

This method is DRF and also reliable. Instead of using the synchronized keyword, I used Java's ReentrantLock to ensure thread safety. After some reason, it seems that ReentrantLock often performs better than using synchronized because it has extended capabilities, like lock polling, interruptible lock waits and also a configurable fairness policy. This allows more multiple condition variables per monitor while synchronized only have one wait / notify queue.

1.6 BetterSorry

```
java UnsafeMemory BetterSorry 2 1000000 8 5 6 3 0 3 - 413.319 ns / transition  
java UnsafeMemory BetterSorry 8 1000 12 5 6 3 0 3 - 27289.3 ns / transition  
java UnsafeMemory BetterSorry 8 100 12 5 6 3 0 3 - 210935 ns / transition
```

This method uses an array of AtomicIntegers instead of ReentrantLocks so as expected, it performs better than BetterSafe since locks are slower to use. It is also more reliable than Unsynchronized because it uses AtomicInteger's atomic operations to increment and decrement.

The first test case would hang forever on Unsynchronized so this avoids that race condition. However this method still has issues, e.g. while you're getting the value at an

index, it can be changed in the period between the get and the increment / decrement. I recreated this by making my initial values all 0's like so: `java UnsafeMemory BetterSorry 4 1000000 8 0 0 0 0 0`.

2 Conclusion

I would suggest using either BetterSafe or BetterSorry for GDI's applications. The safer but slower route would be use BetterSafe but the spec mentions that GDI does not mind "a small amount of breakage" so BetterSorry would also work, and plus, it's faster.