convert into byte. not include sivialitation.

class user implements devializable. premate steing asomame. primate string passessied:

-> private duancient Etting passervered.

e) Reflection breaks the OOPs protection (frameworks)

& INHEITANCE.

a) Slingle inherentence class z extends x, y { } // X class c extends B//~

b) Constructor chaining

cy Method according us hideing (static methods are hidden

d'y final method council be overviolen

e) Private Member are not unherited

by Object class is Auperclass of All classes.

37 Shadow voviables (super, shis)

parent to child. A) IS A Relationship buildyega Inheritance.

if Upcasting & downcasting.

Arrival a = new Deg(); alog of = new Dog(); Dog d = (Dog) a; //downcahing Animal a = d; // upcashing

> Animal a = new Animal(); dog d = (Dog) a; // Ever Keunstur

Generics?-* Java seemoves all generic type inforemation at suntime is cally Type Eraswee. Meaning: - at compile time, gueric help enforces stype safety. al suntime, the type becomes naw (list) - no you can't overland based on generic types leg: - public word schow (List < schwing > dist) {} public void schow (List< Struig> void) {} x Compile nohy? Bublic woid show (List dist) // after type exame Inlidewed! Eg: List< Pertends Number> list = new ? extends T AlwayList (Integer > (); list add (10); X can't add anything dist. add (null); // only neell is allowed, Number n = leat.get(0); // Integer i = hist get (0); // (casted) 9 Duper And T eg: List < ? super Anteger > leat = new Asercy List < N comberX) lest. add (10), // ~ list add (new Integer (5)); // ~ Object obj = lest-get(0); // allowed ~ Integer val = list get (0); // Compeler doesn't allow. > Because the compiler doesn't know if it's a Mumber, Object etc. - il's 200 generic 20 secret! Note.

? extends T => I have a box of some subtype of T: I can slook, but I can't add anything to it.

9 Super T > I have a box that excepts T on anything about wit, I can vafely put T into us, but not read sepecific types from it

```
Hash Code & Equals.
```

) . Eg: abbeing S1 = "FB"; // S1 hashcode = 2236 Steering S2 = "Ea"; // S2. hash (code() = 2236. As. o. pth (si. harh Code () = = S2. hash Code ()); // drue flash Map < String, Integer > map = new flash Map <>(); map. pul (SI, 100); map. pul (\$2, 200); Ø.o.ptfn(map); // Both key exect

hashCode () helps find the beeckers. equals () check conside the bucket for actual key identify.

2) . How hashMap recoolues Collesions. ind hash = key. hash code(); ient involve = hash % capacity; _ Then, in that under (bucket): if empty - add entry. if not emply - check existing key via equals.

if key abready exists - conse ourweit value.

else - chain the new mode (linked list / frece).

Eg: HashMap < key, value>

L Bucket [5] → [key = wals] -> [key 2 = val2]

of key1. hash Code () == key2. hash Code () but ! key1. equals (key2), both are

- 3 TrueMap uses compare To () for Souting, Not hash Coole().
- (4) Key Overweisting un HashNap: equals() + hashlade() Same Hash Code, different equals => Stored separentely. Sam hash Code, same equals > Value is replaced 1/4 you overvide equals (), you must overvide hasscock ()else unexpected behavious un flashMap. // Equals () behaveour.

Collection.

-> Acceay List remove (cint under) VS Accray List remove (object) Ec: Asercaylist< Integer > liet = new Aerray List<>1); clist add (1); list, adol(2); leaf.add(3); elist. seemove (2); -> ist stakes as the cle at index 2. 080, it remove 3.

lied remove (Integur value Of (2)); // il removes the cle ?

> TreeMap & Treesed > add nell therow NPE Haah Set & flashMap > use this if neels are needed.

→ of two objects have dame hashcode() -> check (quals() If equals () redwons deue, vet's a deeplicate le won't bead Set Duplicates not allowed.

> Natural Ordering + Amplement Default order + eles. comparti (Comparable) > Comparable < T> in Trucket, Truckly (eles) i compare To

Custom Ordering > You define when > compare (eles, eles). (Comparator <T>) creating the collection

TreeMap & Trece Bet require sorting. comparatol) (comparal) = 0 means desplicate, Even if 206 are not equals ()

foreach loop & use stereator. -> Derect list remove () concurrent Modification Exception

Eg: for (steering s: lead) { if (s. equals 'h ')) {
 list. vermore (s);

ghow & dhewers mismatch geine compete fine cereou for checked exception day -> catch -> finally -> redword throw * Only checked exception needs thousand like Joenephion, Rule: If finally contains a refuer, it were - earlier rection or show gets ignored. (cotch A B c) // Compile ever Aug & step 1 Throw X -> looks for catch block of type X/paren Step 2 No match ? -> skip to finally, then crash. show X Match found ? -> catch runs -> finally runs -> catch (Y) { estep 3 *If finally has a shower or refuser, it ouverdes inhateuer mas ein sier. finally { * throws clause must match the actual expressive expressive expressive uts pavent. void m2() throws duperclass void m3() throws farent I show new subclass () show new child () void m() shown sibling (IOEnception) Compile time.

E shown new Exception ();

Polymorphism 41 allowers one reference to point to objects of diff fypes and invoke the sight behaviour at suntime. Mothed Querloading of True Bolymorphism. Static resolve at suntime compile time by Method Ownioling = True Polymorphism. Overeding own time ia appasting ey Method Overviding Rules. Child class ka modifier bada hona iy Access Modiefir la chaige parent se 1). Retwen type -> can be considered and. my Exception Rule > Subclass method can therow novewer Or stame checked exceptions. d.y Polymouphism with Interfaces. Anderforence desperence - Bouwerful abstraction e. Polymouphions & Aurays. Animald] animals = new Dog [3]; // consulant average allow. animals [0] = new Dog (); // L animals [3] = new Cal(); // X Runtim devaysfore f. Private/ Static / Final methods are not our voiden gy unstance of :- used to check whether an object is an. wistomes of specific class, subclass on interfac return Ex boolean Sype Null safety: null unslance of sulype -> X false (show

ABSTRACTION.

- It can be achieved by:
- Abolian classes.
- Interfaces
- 1) Abstract:
- ay You can define <u>constructors</u> inside alesteact classes. constructors runs when subclass vis unstantialed.
- by Abstract class can have both abstract & non abstract methods
- of Abertract class can have instance voviables & static variables.
- dy You can declare static & final methods (BUT not alestrae (abouted static word prient () & //x not allowed.
- e) Can implement unterfaces.
- of Can ordered concrete class.
- g) Can be partially implemented.
- by Can't be final or private.
 - / final abstract clase x {} //x
 - / private abstract class 477 // x

ANTERFACE All vaerables are public estertic final - constants. > can't make them. puniate, Buteled/nonfinal. by All methods are public abstract (by default) of you cannot create object of an interface. d) From Jana &. inderface can have alfaelt methods. default word greet () { Static methods Sopthi(" flello"); Calling. Jone. super. goeset (); et Methods cannot be primate ou predected. J'y Supports multiple Amheridance gy an underfaces can ordered multiple unterfaces. underface A {} uinterface 18 {} undreface C exclands A, B { 3 // valid. * class extend only 1 class but jundonface can extends many einterfaces hy Interfaces can't have constructors. if It is used to represent capability, not extension # MARKERS THTERFACES.

=> Empty interfaces, no methods.

=> SERIALIZABLE, CLONEABLE, REMOTE

Abolialtion. Encapsulation. Hides implementation (logic) Hides data (variables) design concern. Becurity concorn Achieved using alestract Achieved using perinate, public, class, underface getter/ deters focus: Whail alward be focus: heho can access my expressed ? . douba 9 Final Classes & Final Methods. final class - cannot be unhavisted. final method - cannot be overviolen. Final usuiable - constant # Where to use HAS-A (Association, Aggregation, Composition) Association. Use: - Two classes were related but can work independently. Eg: Doctor - Palient, Student - College. class Doctor & / // word ein loose relationships patient assigned Patient; / where no ownership is required Aggregation. One object "has-a" another, but shey can live exparally. Eg: Depardment has Teachers, but Teacher ean change dept. class Department {

Dist (Teacher) leachers;

Of the child object. Composition. One object completely own another. Eg: Library has Books, Can has Engine, fluman has fleart class Engine {} ans Car {

Sed when the Barent

class is responsible

for child's lifety cl. class Car f

050LJD Principles

- Each class should do one thing & do it well.

 Each class should do one thing & do it well.

 otherwise, it skecomes harden to maintain & test.
- 2. Sport Closed Preinciple (OCP)

 Structure should be open for extension, but closed for modification.

2. You ashould be able to add new features without changing orde

- 3. Liskou Substitution Principle (LSP).
 Objects of a Duporclass schould be replacable with objects of its subclass without breaking the app.
- idea: A child class must behave in a way that does not georgewise / break the expectations from the parent class
- 4. Interface Seguation Principle (ISP)
 Clients should not be forced to depend on unforface they do not use.
- idea. Big indoefaces virdo semaller ones.
- 5 Dependency Invasion Principle (DIP). Depend on abstraction, not on convicte classes:
- idea tigh level modules should not depend on low-level medules. Both should depend on abotractions (interfaces).

Mars adding)

((+2) how & 16).