Visualization tool to compare HTTP/1.1, SPDY, QUIC

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Code Repository URL:

https://github.com/pushkargarg/compare-protocols

Visualization Tool URL:

http://allv28.all.cs.stonybrook.edu/shsinghal/FCN/

1. Introduction

In today's internet, the tolerance threshold for how long a user will wait for a slow loading website continues to decrease. Users demand near real time responses which is not possible with HTTP/1.1. One of the drawbacks of HTTP/1.1 is opening of too many TCP connections to achieve concurrency, as, a large number of connections may lead to network congestion and hence to an even worse performance. Second disadvantage is that, the server cannot initiate a transfer and hence response time in case of embedded objects is significantly high. SPDY rectified above shortcomings of HTTP/1.1 by introducing server push, request prioritization and combining of multiple small http requests and response into one packet. It uses multiplexing over single TCP connection and hence avoids congestion. Server push allows server to send data before explicit request from client improving latency in case of embedded objects. It also provides header compression resulting in fewer bytes that need to be transferred. One issue that arises with multiplexing in SPDY is that since it is done over single TCP connection that needs to provide ordered delivery, if one packet is lost all subsequent packets in all streams will have to wait for its retransmission (Head of Line Blocking). QUIC has been developed over UDP and retransmission and congestion control

modules for it have been developed at application layer. Since UDP does not enforce in-order delivery, QUIC provides multiplexing with no head of line blocking arising in case of packet loss. Also, in case a connection was already established between client and server in recent past, QUIC uses Connection ID (CID) to identify connections instead of IP addresses which might change on change with network interface, resulting into 0-RTT re-connection cost. The ACK frames in QUIC also contain the delay between the time packet was received and its acknowledgement was sent resulting into estimation of precise RTT. But, the 0-RTT reconnection cost is applicable only if same server is hit otherwise new connection would need to be established.

2. Problem Description

The goal of this project is to create a tool to visualize and differentiate between the flow of packets for HTTP, SPDY and QUIC and to investigate the effectiveness of SPDY and QUIC over HTTP by comparing their performances on general web traffic and heavy latency websites such as YouTube. The tool will allow comparison of any two or all three protocols at a time. The contribution of this project is twofold: i) We provide a tool to visualize differences in handshake and flow of the three protocols. ii) We asses and provide a tool to visualize the performances of the three protocols in terms of average RTT per object and Handshake cost.

3. Measurement Environment

We used a regular laptop with Chrome and Firefox browsers and Wireshark to capture network packets as pcap files. We, then developed packet parsers for all the three protocols in order to parse the pcap files so as to

obtain the flow information and page load time and handshake cost values for comparison. Packets were captured for all three protocols on general web traffic as well as streaming data traffic. One limitation we faced here is that since QUIC has been deployed only on Google and YouTube there wasn't much option available to capture OUIC packets. The websites for which data was captured were: google.com (flow while loading home page), gmail.com (flow while loading sign in page), google translate (flow while translating a sample German text), google image search (flow when icon size logos are searched), youTube.com (flow during full streaming of a 50 seconds long video), spdy.centminmod.com/flags.html, facebook.com (flow while loading home page of a signed in user), linkedin.com (flow while loading home page of a signed in user), twitter.com(flow while loading sign in page), en.wikipedia.com(flow while loading home page).

4. Solution Methodology

The module has been developed using JAVA language and JnetPcap library. For visualization tool, D3 has been used. Pcap packets were captured through Wireshark. Pcap files are then given as input to OUIC Parser and SPDY/HTTP 1.1 Parser which output the CSV files containing information such as source and destination IP address. port number. sequence acknowledgement numbers, TLS information (in case of SPDY) and server configurations, client hello and server rejection (in case of OUIC) and performance matrix (page load time). The visualization tool uses these csv files as input in order to display data. It has a drop down menu from which a user can select whether he wants to compare protocols based on flow (handshake and reconnection) or based on page load time.

Figure 1 shows a snapshot of the visualization tool. On upper left hand corner, a drop-down menu provides users with options such as to compare protocols based on Handshake flow, Reconnection and Page load Time. Another drop down menu provides them to compare data for different websites, such as, google.com, gmail.com, facebook.com etc.

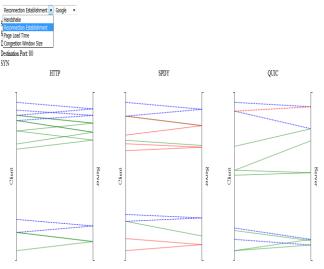


Figure 1: Visualization Tool

5. Evaluation And Results

5.1 Handshake and Reconnection in OUIC:

The initial handshake in QUIC protocol requires 1 RTT. When a client tries to connect with a server for the first time, it sends a Client Hello message which has a SNI (Server Name Indication) tag but not a STK (Source address token) or SCID (Server Config ID) [Figure 1.1] which the server requires in order to verify client's authentication. The server then responds with a connection rejection message sending along its configurations SCFG (Server Config) [Figure 1.2] that contains its configuration ID, information about authentication encryption algorithm, etc. and is valid in general for a few days. The client then re-sends a Client Hello with valid configurations and an ACK frame acknowledging receipt of configurations [Figure 1.3]. Now, Once the connection has been established, in case of no network activity, it automatically shuts-down after ICSL (Idle Connection State Lifetime) time that is decided during Client Hello. By-default its value is 30 sec. Since the server configuration has already been sent, an attempt at connection after passing of ICSL time does not require another handshake hence resulting in 0 RTT re-connection time [Figure 1.5].

This can be seen from a captured packet flow as seen in Wireshark:

Figure 1.1 - Initial connection attempt by Client without any valid STK or SCID

No.	Destination	Source	Protocol	Length	Info				Time													
Г	78 216.58.217.164	10.0.2.15	QUIC	1392	Client H	ello, CI	0: 120946	3957:	8.912559													
	79 10.0.2.15	216.58.217.164	QUIC	1392	Rejectio	n, CID:	120946395	7207:	8.926858													
Т	80 10.0.2.15	216.58.217.164	QUIC	1392	Payload	(Encrypte	ed), CID:	120	8.926890													
	81 216.58.217.164	10.0.2.15	QUIC	82	Payload	(Encrypte	ed), CID:	120	8.931121													
	82 216.58.217.164	10.0.2.15	QUIC	1392	Client H	ello, CI	0: 120946	3957:	8.932043													
	83 10.0.2.15	216.58.217.164	QUIC	1392	Payload	(Encrypte	ed), Seq:	3	8.946872													
4																						
▶ Frame 79: 1392 bytes on wire (11136 bits), 1392 bytes captured (11136 bits) ▶ Ethernet II, Src: RealtekU_12:35:02 (52:54:00:12:35:02), Dst: CadmusCo_82:ce:f5 (08:00:2) ▶ Internet Protocol Version 4, Src: 216.58.217.164, Dst: 10.0.2.15 ▶ User Datagram Protocol, Src Port: 443 (443), Dst Port: 60432 (60432) ■ QUIC (Quick UDP Internet Connections)																						
														D Public Flags: 0x0c								
														CID: 1209463957207	1071594							
														Sequence: 1								
													Message Authentication Hash: 5ee948ca069e34d40f8dd152									
▷ Private Flags: 0x00																						
		ame Type) Strea	m ID:1,	Type:	REJ (Rej	ection)																
	▶ Frame Type: STR	EAM (Special Fr	ame Typ	e) (0x	80)																	
	Stream ID: 1																					
	Tag: REJ (Rejec	tion)																				
	Tag Number: 7																					
Padding: 0000																						
	Tag/value: STK	(Source Address	Token)	(1=60))																	
	Tag/value: SNO	(Server nonce)	(1=56)																			
	Tag/value: PROF																					
	Tag/value: SCFG																					
	Tag/value: RREJ	(Bassans for s	erver c	ending'	\ (1-4) •	Code 12																
				ciluzing,) (1-4).	Code 12																
	Tag/value: CSCT			enuing,) (1-4).	coue 12																

Figure 1.2 - Server Rejection Containing Valid STK and SCFG

```
        Destination
        Source
        Protocol
        Length Info
        Time
        Time

        82 216.58.217.164
        10.02.15
        QUIC
        1392 Citient Hello, CID: 1209463957207118.932643

        83 10.00.2.15
        216.58.217.164 QUIC
        1392 Payload (Encrypted), Seq: 3
        8.946872

      IIC (Quick UDP Internet Connectation,
Public Flags: 0x0c
CID: 1209405972671071594
Sequence: 3
Nessage Authentication Hash: 63cdcd9fa48893f4d22dc918
> Private Flags: 0x00

**TREAM (Special Frame Type) Stream ID:1, Type: CHLO (Client Hello)

P Frame Type: STREAM (Special Frame Type) (0xa4)
4 QUIC (Quick UDP Internet Connections)
             Stream ID: 1
Offset Length: 1300
Data Length: 1204
Tag: CHLO (Client Hello)
Tag Number: 27
Padding: 6000
Tag/value: PAD (Padding) (1=262)
Tag/value: PAD (Fadding) (1=262)
Tag/value: SNI (Server Name Indication) (1=14): www.google.com
Tag/value: SNI (Server nonce) (1=56)
Tag/value: SNO (Server nonce) (1=56)
Tag/value: SNO (Server nonce) (1=56)
              Tag/value: SMO (Server nonce) (1-56)
Tag/value: YER (Version) (1-4) (028)
Tag/value: CFR (Version) (1-4) (028)
Tag/value: NONC (Client Nonce) (1-32)
Tag/value: MSPC (Max streams per connection) (1-4): 100
Tag/value: AEAD (Authenticated encryption algorithms) (1-4), AES-OCM with a 12-byte tag a Tag/value: AEAD (Authenticated encryption algorithms) (1-4), AES-OCM with a 12-byte tag a Tag/value: SID (Server config ID) (1-32): Chrome/50.0.2661.75 Linux x86_64
Tag/value: SID (Server config ID) (1-16)
Tag offset end: 500
Tag longeth 161
                     Tag value: 274b2f800e1069e197a6691b6b876103
Server Config ID: 274b2f800e1069e197a6691b7a6691b6
                                                                                                       Protocol Length Info
                   718 216.58.217.164 10.0.2.15
5401 216.58.217.164 10.0.2.15
                                                                                                      QUIC
                                                                                                                           80 Payload (Encrypted), CID: 120946395 12.833033
1392 Client Hello, CID: 1311394979007324 77.612913
                   5402 10.0.2.15
                                                                   216.58.217.164 QUIC
                                                                                                                           1392 Payload (Encrypted), Seq: 1
                                                                                                                                                                                                                                 77.629258
           ■ QUIC (Quick UDP Internet Connections)
                 ▶ Public Flags: 0x0d
CID: 13113949790073247891
                       Version: 0030
                        Sequence: 1
                       Message Authentication Hash: ca8abfa43b88a15cadd54ca4
                 ▶ Private Flags: 0x01

■ STREAM (Special Frame Type) Stream ID:1, Type: CHLO (Client Hello)

                        ▶ Frame Type: STREAM (Special Frame Type) (0xa0)
Stream ID: 1
                             Data Length: 1024
Tag: CHLO (Client Hello)
                              Tag Number: 26
                             Padding: 0000
Tag/value: PAD (Padding) (1=328)
                              Tag/value: SNI (Server Name Indication) (l=14): www.google.com
                             Tag/value: STK (Source Address Token) (1=58)
                              Tag/value: VER (Version) (1=4) Q030
Tag/value: CCS (Common Certificate Sets) (1=16)
                          rag/value: CCS (Commind Cet Italet Sets) / Tag/value: NSPC (Max streams per connection) (1=4): 100
Tag/value: NSPC (Max streams per connection) (1=4): 100
Tag/value: AED (Authenticated encryption algorithms) (1=4), AES-6CM with a 12-byte tag and Tag/value: UAID (Client's User Agent ID) (1=32): Chrome/50.0.2661.75 Linux x86.64

■ Tag/value: SCID (Server config ID) (1=16)
                                     Tag Type: SCID (Server config ID)
Tag offset end: 508
                                     [Tag length: 16]
Tag/value: 274b2f800e1069e197a6691b6b876103
```

Figure 1.3 - New Client Hello with valid Server Config ID Figure 1.4 New Connection Attempt after Expiration of ICSL

Server Config ID: 274b2f800e1069e197a6691b6b876103

```
Length Info
       Destination
                              Source
                                                     Protocol
                                                                                                                                                                                              Time
   718 216.58.217.164
                              10.0.2.15
                                                                 80 Payload (Encrypted), CID: 12094639572071071594, Seq: 154
                                                                                                                                                                                              12.833033
                                                     QUIC
  5401 216.58.217.164
                              10.0.2.15
                                                     OUIC
                                                               1392 Client Hello, CID: 13113949790073247891, Seq: 1
                                                                                                                                                                                              77.612913
  5402 10.0.2.15
                              216.58.217.164
                                                    QUIC
                                                               1392 Payload (Encrypted), Seq: 1
                                                                                                                                                                                              77.629258
Frame 5402: 1392 bytes on wire (11136 bits), 1392 bytes captured (11136 bits)
Ethernet II, Src: RealtekU 12:35:02 (52:54:00:12:35:02), Dst: CadmusCo 82:ce:f5 (08:00:27:82:ce:f5)
 Internet Protocol Version 4, Src: 216.58.217.164, Dst: 10.0.2.15
User Datagram Protocol, Src Port: 443 (443), Dst Port: 56863 (56863) QUIC (Quick UDP Internet Connections)
  ▷ Public Flags: 0x00
    Sequence: 1
    Payload: aee9038e873fb3d7a2b0151f72201a7ab0bef050cdf8cdec...
```

Figure 1.5 Server Responds without Rejection Hence, O RTT Reconnection

Same observation has been displayed in our visualization tool as below:

Handshake

Source IP: 10.0.2.15

Source Port: 60432

Destination IP: 216.58.217.164

Destination Port: 443

CHLO

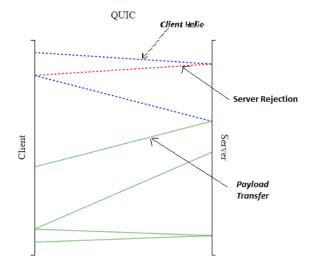


Figure 3.1: QUIC Handshake

Reconnection Establishment ▼ Google ▼

Source IP: 10.0.2.15
Source Port 60432
Destination IP: 216.58.217.164
Destination Port: 443
CHLO

QUIC

Reconnection
Attempt after expiration of ICSL, no server rejection hence 0 RTT

Figure 3.2: QUIC Reconnection

The blue dotted lines represent a connection attempt from client(CHLO) and red dotted line represents rejection by server (REJ). Green lines show payload transfer. On hovering over each line, the packet information such as source/destination IP, source/destination port number and whether it's a CHLO or server REJ is shown in the upper left corner of the display. Here, the first three flow lines show the 1 RTT initial handshake and during reconnection attempt there's no server rejection and hence 0 RTT.

5.2 Handshake and Reconnection in SPDY:

The packets captured were encrypted and marked as TLSv1.2 on Wireshark [Figure 2.1], but still the flow of packets could be observed manually. The initial handshake in SPDY requires 3 RTT. 1 RTT for initial TCP handshake and 2 more RTT for SSL handshakes. In case of reconnection, SSL handshake only needs to exchange keys between the client and the server and does not require server authentication thus reconnection establishment in SPDY requires 1 RTT for TCP handshake and 1 RTT for SSL handshake in needed. [Figure 2.2]

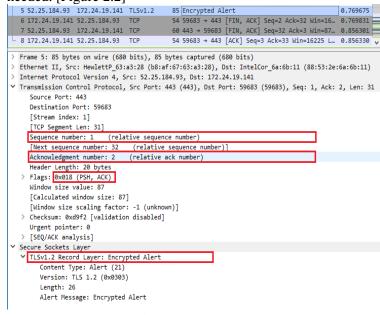


Figure 2.1 – Encrypted SPDY packet showing push acknowledgment

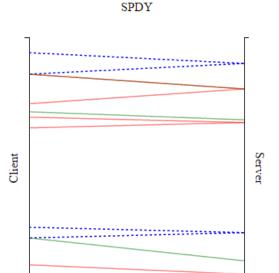


Figure 2.2 SPDY Handshake and Reconnection Flow as seen in Visualization Tool

In Figure 2.2, the blue dotted lines represent a TCP handshake. In the initial connection as well as the reconnection there is only one single TCP handshake. The handshake starts with the client sending a SYN and receiving a SYN ACK from the server. Following this a green line representing the ACK sent by the client in reply to the SYN ACK. This packet might be sent in parallel with the request for SSL server authentication represented by the brown line in the initial handshake and is followed by another handshake for SSL key exchange, represented by the red lines. As there is no need for server authentication in reconnection establishment the latter connection establishment only has one SSL handshake for key exchange.

5.3 Handshake and Reconnection in HTTP/1.1:

HTTP/1.1 uses three-way handshake (SYN-SYN/ACK-ACK) which requires 1 RTT. Along with being persistent it also uses parallel connections as can be seen in below wire-shark capture [Figure 3.1]. The reconnection attempt in case of HTTP/1.1 after a connection has been terminated requires the same three-way handshake process as the initial connection.

No.	Source	Destination	Protocol	Leng	Info	Time
	1 172.24.21.152	31.13.80.36	TCP	66	51091 → 443 [SYN] Seq=0 Win=8192 Len= M	0.000000
	2 172.24.21.152	31.13.80.36	TCP	66	51092 → 443 [SYN] Seq=0 Win=8192 Len= M	0.013036
	3 31.13.80.36	172.24.21.152	TCP	66	443 → 51091 [SYN, ACK] Seq=0 Ack=1 Win=1	0.087314
	4 172.24.21.152	31.13.80.36	TCP	54	51091 -> 443 [ACK] Seq=1 Ack=1 Win=66240	0.087453
	5 31.13.80.36	172.24.21.152	TCP	66	443 → 51092 [SYN, ACK] Seq=0 Ack=1 Win=1	0.087692
	6 172.24.21.152	31.13.80.36	TCP	54	51092 -> 443 [ACK] Seq=1 Ack=1 Win=66240	0.087774

Figure 5.1 Parallel Connections and Three-Way handshake in HTTP/1.1

This information can be seen in the visualization tool as below [Figure 5.2 and Figure 5.3]. The dotted blue lines represent SYN and corresponding SYN-ACK. Subsequent ACK and data flow can be seen by green lines. The upper left corner shows same information as for other protocols except for showing whether its SYN, SYN-ACK or ACK. It can be seen from the connected green lines that in HTTP/1.1 no new request is sent over a connection until it receives an acknowledgement for last one.

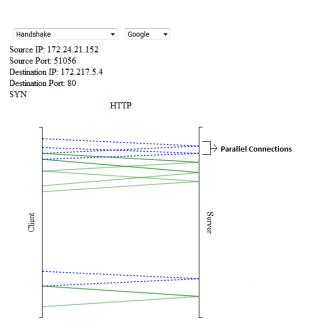


Figure 5.3 Handshake in HTTP/1.1



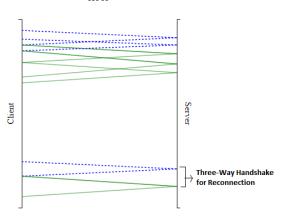


Figure 5.3 HTTP/1.1 Handshake and Parallel Connections

5.4 Page Load Time Comparison for Three Protocols:

Figure 6 shows the results as observed and displayed in the visualization tool:

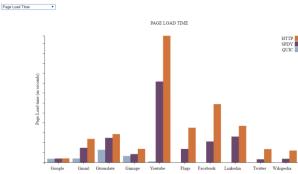


Figure 6 Page Load Time For Different Websites

It can be seen that QUIC performs better than SPDY and HTTP/1.1 wherever applicable especially in case of stream data (youTube.com). This can be because QUIC do not use packet sequence number while retransmitting it. This avoids retransmission timeouts by avoiding ambiguity about which packets have been received resulting in fewer re-buffers.

By looking at page load time for website *spdy.centminmod.com/flags.html* we can see that SPDY performs better than HTTP/1.1 when the page has large number of small objects as it combines multiple small http requests and response into one packet.

6. References

- QUIC Wire Layout Specifications https://docs.google.com/document/d/1W JvyZflAO2pq77yOLbp9NsGjC1CHetA XV8I0fQe-B_U/edit#heading=h.z2ju224lr24y
- HTTP Over UDP: An Experimental investigation of QUIC -http://c3lab.poliba.it/images/3/3b/QUIC_ sAC15.pdf
- Comparison of Web Transfer Protocols -<u>http://proprogressio.hu/wp-content/uploads/2016/01/MolnarSandor_2015.pdf</u>
- QUIC: A UDP-Based Secure and Reliable Transport for HTTP/2 -https://tools.ietf.org/html/draft-tsvwg-quic-protocol-00