

FOURIER ANALYSIS (751799001, 701866001, 114-1) - HOMEWORK 7

Return by November 12, 2025 (Wednesday) 23:59

Total marks: 50

Special requirement. All homework must be prepared by using L^AT_EX.

Exercise 1 (10+10 points). For a fixed $x \in \mathbb{R}$, we see that $t \mapsto e^{ix \sin t}$ is a 2π -period function, therefore one sees that its Fourier series on $(0, 2\pi)$ is well-defined, and converges uniformly on $[0, 2\pi]$. Write

$$e^{ix \sin t} = \sum_{n \in \mathbb{Z}} J_n(x) e^{int} \quad \text{for all } t \in \mathbb{R}.$$

The function J_n is called the *Bessel function of the first kind* of order $n \in \mathbb{Z}$.

- (a) Compute $\sum_{n \in \mathbb{Z}} |J_n(x)|^2$ for all $x \in \mathbb{R}$.
- (b) For any $x \neq 0$, express $x^2 J_n''(x)$ in terms of $J_n'(x)$ and $J_n(x)$, where $J_n'(x) = \frac{d}{dx} J_n(x)$ and $J_n''(x) = \frac{d}{dx} J_n'(x)$.

Exercise 2 (10+10+10 points). We now consider the normalized Fourier transform as

$$\hat{f}(\xi) = \int_{\mathbb{R}^n} f(x) e^{-i2\pi x \cdot \xi} dx \quad \text{for all } f \in \mathcal{S}(\mathbb{R}^n),$$

which is often used in signal processing (see, e.g., Python's Numpy package). In this case, the inverse Fourier transform is given by

$$\check{f}(\xi) = \int_{\mathbb{R}^n} f(x) e^{i2\pi x \cdot \xi} dx \quad \text{for all } f \in \mathcal{S}(\mathbb{R}^n).$$

Now we fix $n = 1$.

- (a) Show the Poisson summation formula

$$\sum_{k \in \mathbb{Z}} \hat{f}(k) = \sum_{n \in \mathbb{Z}} f(n) \quad \text{for all } f \in \mathcal{S}(\mathbb{R}).$$

[Hint. Compute the Fourier series $\Phi(x) = \sum_{n \in \mathbb{Z}} f(n+x)$ and find that value $\Phi(0)$]

- (b) Suppose that $g \in \mathcal{S}(\mathbb{R})$. Show that

$$\sum_{k \in \mathbb{Z}} \hat{g} \left(\xi - \frac{k}{T} \right) = T \sum_{n \in \mathbb{Z}} g(nT) e^{-i2\pi n T \xi} \quad \text{for all } \xi \in \mathbb{R}.$$

In particular, if $g \in \mathcal{S}(\mathbb{R})$ and $\text{supp}(\hat{g}) \subset [-\frac{1}{T}, \frac{1}{T}]$,

$$\hat{g}(\xi) = T \sum_{n \in \mathbb{Z}} g(nT) e^{-i2\pi n T \xi} \quad \text{for all } \xi \in \left[-\frac{1}{T}, \frac{1}{T} \right].$$

This implies that g can be reconstructed based on partial knowledge of g , namely $\{g(nT)\}_{n \in \mathbb{Z}}$. **[Hint.** Choose $f(x) = g(Tx) e^{-i2\pi T x \xi}$ in (a)]

(c) Show that if $g \in \mathcal{S}(\mathbb{R})$ satisfies $\text{supp}(\hat{g}) \subset [-B, B]$, then

$$g(x) = \sum_{k \in \mathbb{Z}} g\left(\frac{k}{2B}\right) \text{sinc}(2Bx - k) \quad \text{for all } x \in \mathbb{R},$$

where sinc is the normalized sinc function given by

$$\text{sinc}(x) = \begin{cases} \frac{\sin(\pi x)}{\pi x} & \text{if } x \neq 0, \\ 1 & \text{if } x = 0. \end{cases}$$

This result is called the *Shannon interpolation theorem*.

Remark. In mathematics, audio signals are modeled by a continuous function g . However, in practice, it is not possible to record the full audio signal g , instead, the signals are typically recorded at a *sample rate* of f_s , i.e., f_s samples per second, with each sample represented as a floating-point number¹ between -1 and 1 . If we want to capture the signal contains no frequency components above B Hz (called the *bandwidth*), i.e. $\text{supp}(\hat{f}) \subset [-B, B]$, in view of Exercise 2(b)–(c) with $T = 1/B$, we see that the audio signal g can be reconstructed based on the discrete point $\{g(\frac{k}{2B})\}_{k \in \mathbb{Z}}$, which means that a nearly perfect reconstruction is guaranteed provided that $B < f_s/2$. For example, the sample rate $f_s = 16$ kHz captures frequencies up to 8 kHz, which is sufficient for speech recording (<https://arxiv.org/abs/2501.01650>). Human speech frequency is close to the piano's middle C (≈ 262 Hz) based on the modern A440 pitch standard². In the standard piano with 88 key, the first key is tuned to be 27.5 Hz and the last key is tuned to be 4186 Hz. The standard rate $f_s = 44.1$ kHz captures frequencies up to 22050 Hz, which covers the human hearing range of approximately 20 Hz to 20 kHz.

¹For storage efficiency, the audio is usually stored in 16-bit integer format (16-bit PCM) by multiplying each sample by 32767 and rounding to the nearest integer. There are exactly $2^{16} = 65,536$ integers ranging from -32767 to 32767 , which explains the term “16-bit”.

²https://en.wikipedia.org/wiki/Piano_key_frequencies