

## Category

1. A collection of **objects**.
2. A collection of **arrows**.
3. Operations assigning to each arrow  $f$ 
  - (a) An object  $\text{dom}(f)$  called the **domain** of  $f$ , and
  - (b) An object  $\text{cod}(f)$  called the **codomain** of  $f$  often expressed by  $f : \text{dom}(f) \rightarrow \text{cod}(f)$
4. An associative operator  $\circ$  assigning to each pair of arrows,  $f$  and  $g$ , where  $\text{dom}(f) = \text{cod}(g)$ , a **composite** arrow

$$f \circ g : \text{dom}(g) \rightarrow \text{cod}(f).$$

5. For each object  $A$ , an identity arrow,  $\text{id}_A : A \rightarrow A$  satisfying the law that for any arrow  $f : A \rightarrow B$  as

$$\text{id}_B \circ f = f = f \circ \text{id}_A.$$

## HASK

In Hask, **objects are Haskell types**, so the object part of a functor,  $F_0$ , is a function from types to types. Therefore, a type constructor is suitable for use as the object function of a functor. For example

$$F_0 A = [A]$$

is suitable. In this case

$$F_1 = \text{map}$$

completes the definition of a functor.

## Functors

A **functor**  $F$  from category  $\mathcal{C}$  to  $\mathcal{D}$  is a pair of functions.

1.  $F_0 : \mathcal{C}\text{-objects} \rightarrow \mathcal{D}\text{-objects}$
2.  $F_1 : \mathcal{C}\text{-arrows} \rightarrow \mathcal{D}\text{-arrows}$

such that

1. If  $f$  is an arrow in  $\mathcal{C}$  and  $f : A \rightarrow B$  then  $F_1(f) : F_0(A) \rightarrow F_0(B)$
2.  $F_1(\text{id}_a) = \text{id}_{F_0(A)}$ , and
3.  $F_1(f \circ g) = F_1(f) \circ F_1(g)$  whenever  $f \circ g$  defined.

## Functors - HASK

A typical way of to define a functor in haskell is to do

### Object Functor

```
data NAME type =  
    Base | Constructor1 type | Constructor2 type . . .
```

### Function Functor

```
instance Functor Name where  
    fmap f BASE = something  
    fmap f (Constructor x) = Constructor (f x)
```

## Small Category

Is a category where the collection of objects is a set and the collection of arrows is a set.

**Cat**

Is the category of small categories which has categories as objects all functors between these as arrows.

**Dual**

The dual of a category  $\mathcal{C}$  is identical to  $\mathcal{C}$  but all the arrows are reversed. This is sometimes denoted  $\mathcal{C}^{op}$ .