# Assignment 1

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October 4, 2006

# Assignment 2

## Question 1.2.4

This can be reduced in two ways:

# Question 1.3.2

Suppose f and g are strict then we have that  $f \perp = \perp$  and  $g \perp = \perp$ ,

$$\Rightarrow h\bot = f(g\bot) = f\bot = \bot$$
$$\Rightarrow h\bot = \bot$$
$$\Rightarrow \text{h is strict}$$

# Question 1.4.3

```
mylog :: Float -> Float -> Float
mylog x b = log(x)/log(b)
```

#### Question 1.4.7

```
uncurry :: (a->b->t) -> ((a,b)->t)
uncurry g (x,y) = g x y
```

#### Question 1.6.2

```
Defining swap(x, y) = (y, x) we verify:
```

```
flip(curry f) x y = (curry f) y x = f(y,x)
curry(f.swap) x y = (f.swap)(x,y) = f(swap(x,y))=f(y,x)
```

#### Question 2.4.1

## Question 2.4.3

```
age (d1,m1,y1) (d2,m2,y2)

| m1>m2 = y1-y2

| (d1>=d2) && (m1==m2) = y1-y2

| otherwise = (y1-y2)-1
```

#### Question 2.5.1

```
case2 (Left x) = 1 case2 (Right x) = 2
```

#### Question 2.5.2

```
=case(f,g).plus(h,k)
=case(f,g).case(Left.h,Right.k) (by defn.)
=case(case(f,g).Left.h,case(f,g).Right.l) (by prop. 3)
=case(f.h,g.k) (by prop 1. and prop 2)
```

#### Question 3.2.4

Let  $\Pi(p)$  be the proposition  $\Pi(p) \Leftrightarrow (m+n) + p = m + (n+p) \ \forall m, n \in \mathbb{N}$  where  $p \in \mathbb{N}$ .

Base:  $\Pi(0)$ 

Utilizing the definition x + 0 = x for p = 0 we have

LHS = 
$$(m + n) + 0 = (m + n) = m + n$$
  
RHS =  $m + (n + 0) = m + (n) = m + n$ 

So  $\Pi(0)$  is true.

**Assumption:** We will assume  $\Pi(p)$ , that is,  $(m+n)+p=m+(n+p) \ \forall m,n\in\mathbb{N}$ .

Induction step: (Show  $\Pi(P) \Rightarrow \Pi(suc(p))$ )

$$= (m+n) + suc(p)$$

$$= suc((m+n) + p) \text{ by defn. of } +$$

$$= suc(m+(n+p)) \text{ by assumption}$$

$$= m + suc(n+p) \text{ by defn. of } +$$

$$= m + (n + succ(p))$$

So we have that  $\Pi(p) \Rightarrow (m+n) + suc(p) = m + (n + succ(p)) \Rightarrow \Pi(suc(p))$ . So by induction we have that  $\Pi(p) \forall p \in \mathbb{N}$  as desired.