Overview

■ 1. Finite-State Morphological Parsing — An Introduction

Lexicon, morphotactics and orthographic rules

■ 2. The Lexicon and Morphotactics

■ 3. Morphological parsing with Finite-State Transducers

Two-level morphology

Finite-State Transducer

■ 4. Orthographic Rules and Finite-State Transducers

Spelling rules

Tell me and I forget. Teach me and I remember. Involve me and I learn.
-- Benjamin Franklin

Finite-State Morphological Parsing

Parsing English morphology

 The second column contains the stem of each word + assorted morphological features

Features specify additional information about stem

+N means word is a noun; +SG means singular

Input	Morphological parsed output
cats	cat +N +PL
cat	cat +N +SG
cities	city +N +PL
geese	goose +N +PL
goose	(goose +N +SG) or (goose +V)
gooses	goose +V +3SG
merging	merge +V +PRES-PART
caught	(caught +V +PAST-PART) or (catch +V +PAST)

Finite-State Morphological Parsing

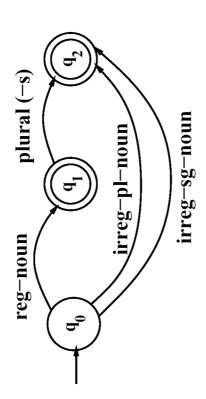
We need at least the following to build a morphological parser:

- Lexicon: the list of stems and affixes, together with basic information about them (Noun stem or Verb stem, etc.)
- which classes of morphemes can follow other classes of morphemes. E.g., the rule that English plural morpheme follows the noun rather • Morphotactics: the model of morpheme ordering that explains than preceding it.
- changes that occur in a word, usually when two morphemes combine Orthographic rules: these spelling rules are used to model the (e.g., the $y \rightarrow ie$ spelling rule changes city + -s to cities).



A lexicon is a repository for words.

- The simplest one would consist of an explicit list of every word of the language. Inconvenient or impossible!
- Computational lexicons are usually structured with
- a list of each of the stems and
- Affixes of the language together with a representation of morphotactics telling us how they can fit together.
- The most common way of modelling morphotactics is the finite-state automaton.



An FSA for English nominal inflection

Reg-noun	Irreg-pl-noun	Irreg-sg-noun	plural
fox	geese	goose	S-
cat	sheep	sheep	
table	mice	mouse	
book			

An FSA for English verbal inflection

Reg-verb-stem	Irreg-verb-stem	Irreg-past-verb	past	Past-part	Past-part Pres-part	3sg
walk	cut	caught	pə-	pə-	-ing	s-
fry	speak	ate				
talk	sing	eaten				
impeach		sang				
		spoken				



English derivational morphology is more complex than inflectional morphology

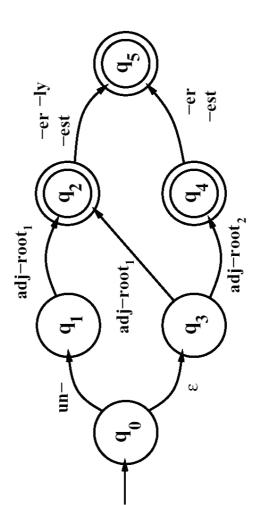
- So automata of modeling English derivation tends to be quite complex
- Some even based on CFG
- [Antworth, 1990] A small part of morphotactics of English adjectives
- Adjectives can have an optional prefix (un-)
- An obligatory root (big, cool, etc)
- − And an optional suffix (-er, -est, or -ly)

An FSA for a fragment of English adjective Morphology #1 big, bigger, biggest cool, cooler, coolest, coolly redder, reddest clear, clearer, clearest, clearly, unclear, unclearly happy, happier, happiest, happily unhappier, unhappiest, unhappily real, unreal, really

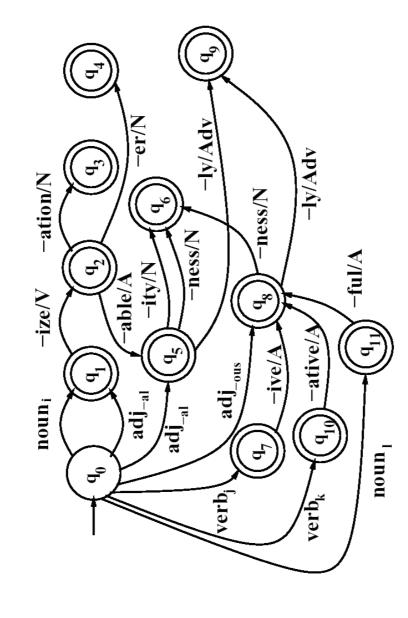
The FSA#1 recognizes all the listed adjectives, and ungrammatical forms like unbig, unfast, redly, smally and realest.

Thus #1 is revised to become #2.

The complexity is expected from English derivation.



An FSA for a fragment of English adjective Morphology #2



An FSA for another fragment of English derivational morphology

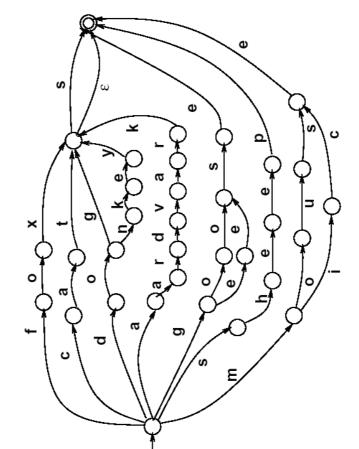
FSA can be used to solve the

problem of morphological

recognition:

- Determining whether an input string of letters makes up a *legitimate*English word or not
- The resulting FSA can then be defined as the level of the individual

letter.



An FSA for English nominal inflection

→ Two-level morphology

→ Finite-State Transducer [FST]



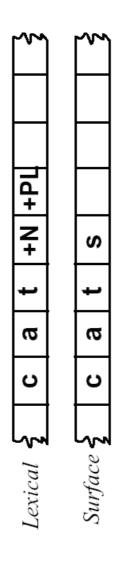
Given the input, for example, cats, we would like to produce:

- Two-level morphology, by Koskenniemi (1983)
- Represent a word as a correspondence between a lexical level
- represents a concatenation of morphemes making up a word
- and the **surface level**
- Represents the actual spelling of the final word.



Morphological parsing is implemented by building mapping rules:

- that <u>maps letter sequences</u> like *cats* on the surface level
- into morpheme and features sequence like cat +N +PL on the lexical level.





The automaton we use for performing the mapping between these two levels is the **finite-state transducer** or **FST**.

- A transducer maps between one set of symbols and another;
- An FST does this via a finite automaton.
- Thus an FST can be seen as a two-tape automaton which recognizes or generates pairs of strings.
- The FST has a more general function than an FSA:
- An FSA defines a formal language
- An FST defines a relation between sets of strings.
- Another view of an FST:
- A machine reads one string and generates another.



C B A

FST as recognizer:

a transducer that takes a pair of strings as input and output accept if the string-pair is in the string-pair language, and a reject if it is not.

FST as generator:

a machine that outputs pairs of strings of the language. Thus the output is a yes or no, and a pair of output strings.

FST as translator:

A machine that reads a string and outputs another string.

FST as set relater:

A machine that computes relation between sets.

D C B A

- A formal definition of FST (based on the **Mealy machine** extension to a simple FSA):
- Q: a finite set of N states q_0, q_1, \ldots, q_N
- composed of an input-output pair i:o; one symbol I from an input Σ: a finite alphabet of complex symbols. Each complex symbol is $\Sigma \subseteq I / O$. *I* and *O* may each also include the epsilon symbol ε . alphabet *I*, and one symbol *o* from an output alphabet *O*, thus
- q_0 : the start state
- F: the set of final states, $F \subseteq Q$
- new state $q' \in Q$. δ is thus a relation from $Q \times \Sigma$ to Q.

 B.E. VII Sem Natural Language Processing Given a state $q \in Q$ and complex symbol $i:o \in \Sigma$, $\delta(q, i:o)$ returns a • $\delta(q,i:o)$: the transition function or transition matrix between states.



D C B A

- FSAs are isomorphic to regular languages, FSTs are isomorphic to regular relations.
- Regular relations are sets of pairs of strings, a natural extension of the regular language, which are sets of strings.
- FSTs are closed under union, but generally they are not closed under difference, complementation, and intersection.
- Two useful closure properties of FSTs:
- **Inversion:** If T maps from I to O, then the inverse of T, T^{-1} maps
- **Composition:** If T_1 is a transducer from I_1 to O_1 and T_2 a transducer from I_2 to O_2 , then T_1 o T_2 maps from I_1 to O_2



Inversion is useful because it makes it easy to convert a FST-as-parser into an FST-as-generator. Composition is useful because it allows us to take two transducers than run in series and replace them with one complex transducer. $T_1 \circ T_2(S) = T_2(T_1(S))$

View an FST as having two tapes:

- *upper* or *lexical tape*: characters from the left side of the *a:b* pairs

lower or *surface tape*: characters from the right side of the *a:b* pairs

Each symbol *a:b* expresses how the symbol *a* from one tape is mapped to the symbol *b* on the another tape

reg-noun-stem $\begin{pmatrix} \mathbf{q}_1 \\ \mathbf{q}_2 \end{pmatrix}$ +N: ϵ $\begin{pmatrix} \mathbf{q}_4 \\ +\mathrm{SG}:\# \\ \mathbf{q}_2 \end{pmatrix}$ +SG:# irreg-pl-noun-form $\begin{pmatrix} \mathbf{q}_2 \\ +\mathrm{N}: \epsilon \\ \mathbf{q}_3 \end{pmatrix}$ +N: ϵ $\begin{pmatrix} \mathbf{q}_4 \\ +\mathrm{C}:\# \\ \mathbf{q}_4 \end{pmatrix}$ +PL:#

A lexical transducer for English nominal inflection T_{num}

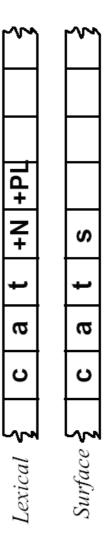
Reg-noun	Irreg-pl-noun	Irreg-sg-noun
fox	g o:e o:e s e	goose
cat	sheep	sheep
gop	m o:i u:ɛ s:c e	mouse
aardvark		

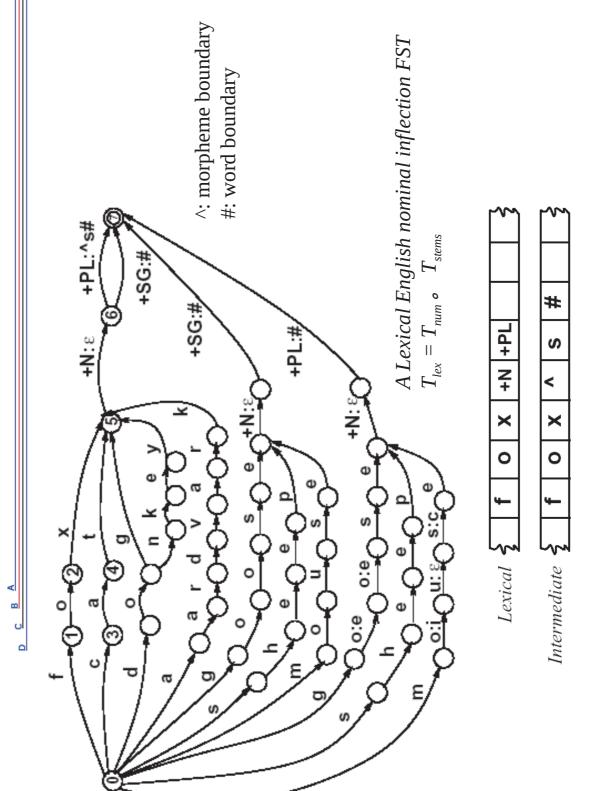
D C B A

- The transducer will map:
- plural nouns = stem + morphological marker +PL
- singular nouns = stem + morphological marker +SG
- Thus a surface cats will map to cat +N +PL

```
-c:c a:a t:t +N:epsilon +PL:^s#
```

-c maps to itself, as do a and t, while morphological feature +N maps to nothing, and the feature +PL maps to ^s −The symbol ^ indicates *morpheme boundary*, while symbol # indicates a word boundary





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Orthographic Rules and FSTs

Spelling rules (or orthographic rules)

Name	Description of Rule	Example
Consonant doubling	1-letter consonant doubled before -ing/-ed	beg/begging
E deletion	Silent e dropped before -ing and -ed	make/making
E insertion	e added after -s, -z, - x , - ch , -s h , before -s	watch/watches
Y replacement	-y changes to -ie before -s, -i before -ed	try/tries
K insertion	Verb ending with vowel + -c add -k	panic/panicked

Spelling changes can be thought as taking as input a simple concatenation of morphemes and producing as output a slightly-modified concatenation of

morphemes.

 Lexical
 \$\frac{1}{2} | \frac{1}{2} | \f

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Orthographic Rules and FSTs

"insert an e on the surface tape just when the lexical tape has a morpheme ending in x (s or z) and the next morphemes is -s"

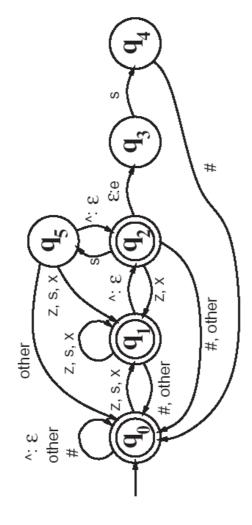
$$\begin{array}{c|c} x \\ x \\ \hline z \\ \end{array} \land \begin{array}{c} x \\ \\ z \\ \end{array}$$

"rewrite *a* and *b* when it occurs between *c* and *d*"

[Chomsky and Halle (1968)]

$$a \rightarrow b / c _d$$

Orthographic Rules and FSTs



The transducer for the E-insertion rule

other	0	0	0	ı	,	0
#	0	0	0	1	0	1
ə:3	,	,	3	ı	,	,
3.0	0	2	0	ı	ı	2
Z:Z	1	П	1	ı	,	ı
x:x	1	_	I	ı	ı	_
S * S	I	_	Ŋ	4		_
State \ Input	<i>q</i> 0:	q_1 :	<i>q</i> 2:	q_3	q_4	<i>q</i> 5

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References

Speech and Language Processing, [Chapter 3] Daniel Jurafsky and James H. Martin



