

Peer-to-Peer Middleware

**Reference: George Coulouris, Jean Dollimore and Tim Kindberg,
“Distributed Systems Concepts and Design”, Fifth Edition, Pearson
Education, 2012**



Peer-to-Peer Middleware

- Peer To Peer Middleware
 - To provide **mechanism** to **access data resources** anywhere in **network**
 - **Functional Requirements :**
 - Simplify construction of services across many hosts in wide network
 - Add and remove resources at will
 - Add and remove new hosts at will
 - Interface to application programmers should be simple and independent of types of distributed resources



Peer-to-Peer Middleware

□ Peer To Peer Middleware (contd)

■ **Non-Functional Requirements :**

- Global Scalability
- Load Balancing
- Optimization for local interactions between neighboring peers
- Accommodation to highly dynamic host availability
- Security of data in an environment simplify construction of services across many hosts in wide network
- Anonymity, deniability and resistance to censorship



Peer-to-Peer Middleware

- Peer To Peer Middleware (contd)
 - Global scalability, dynamic host availability and load sharing and balancing across large numbers of computers pose major design challenges.
 - Design of **Middleware** layer
 - Knowledge of **locations** of **objects** must be **distributed** throughout **network**
 - Use of **replication** to achieve this



Routing Overlays

- Routing Overlays
 - Responsible for **locating nodes** and **objects**
 - Implements a **routing** mechanism in the application layer
 - Separate from any other routing mechanisms such as IP routing
 - Ensures that any **node** can **access** any **object** by routing each **request** through a **sequence** of **nodes**
 - Exploits **knowledge** at each **node** to **locate** the **destination**



Routing Overlays

□ GUIDs

- ‘pure’ names or opaque identifiers
 - Reveal nothing about the locations of the objects
 - Building blocks for routing overlays
- Computed from all or part of the state of the object using a function that deliver a value that is very likely to be unique. Uniqueness is then checked against all other GUIDs
- Not human readable



Routing Overlays

- Tasks of a routing overlay
 - **Routing Request to Objects:** Client **submits** a **request** including the **object GUID**, routing overlay **routes** the request to a **node** at which a **replica** of the **object** resides
 - **Insertion of Objects:** A node introduces a **new object** by **computing** its **GUID** and announces it to the routing overlay
 - **Deletion of Objects:** Clients can **remove** an **object**
 - **Node addition and removal:** Nodes may **join** and **leave** the service



Routing Overlays

- Types of Routing Overlays
 - DHT – Distributed Hash Tables
 - DHT – GUIDs are stored based on the hash value
 - (128 bit hash value using SHA-1 algorithm)
 - DOLR – Distributed Object Location and Routing
 - DOLR is a layer over the DHT that maps GUIDs and address of nodes at which replicas of objects are located. put() and get() APIs
 - DOLR – GUIDs host address is notified using the Publish() operation



P2P Vs Distributed Processing

Peer to Peer

- Millions of nodes cooperate to achieve a common goal
- Distribution of resources are usually explicit but location not known
- WAN based
- Home rather than enterprise based resource servers
- No overall management – insecure resources
- Intermittent connectivity – probabilistic access
- Application level protocols

No fundamental difference

Distributed Processing

- Smaller numbers of nodes cooperate
- May provide a single virtual machine concept with transparent distribution
- Mostly LAN based
- Within a single or a few enterprises
- Managed system – resources can be more trusted
- Tries to provide deterministic access to resources
- Middleware protocols supporting application level interaction



Summary

- ❑ Napster – immutable data, unsophisticated routing
- ❑ Current – mutable data, routing overlays, sophisticated algorithms
- ❑ Internet or company intranet support
- ❑ Distributed Computing (SETI)



Summary

□ Benefits of Peer-to-Peer Systems

- Ability to **exploit unused resources** (storage, processing) in the host computers
- **Scalability** to support large numbers of clients and hosts with **load balancing** of network links and host computer resources
- **Self-organizing** properties of the middleware platforms **reduces costs**



Summary

- **Weaknesses of Peer-to-Peer Systems**
 - Costly for the **storage** of **mutable data** compared to **trusted, centralized** service
 - **Can not** yet **guarantee anonymity** to hosts



Thank You