DISTRIBUTED FILE SYSTEMS

Reference: “Distributed Systems Concepts and Design”, 5th Edition,
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INTRODUCTION

- Sharing of stored information is the most important aspect of distributed resource sharing.
- Some of the problems are load balancing, reliability, availability and security.
- **Distributed File Systems** support sharing of information in the form of files and hardware resources.

INTRO(cntd)

- File systems provide a convenient programming interface to disk storage.
- The most effective in providing shared persistent storage for use in intranets.

INTRO(cntd)

	Sharing	Persistence	Distributed cache	*Consistency maintenance*	*Example*
Main memory	×	×	×	1	RAM
File system	×	✓	×	1	UNIX file system
Distributed file system	✓	✓	✓	✓	Sun NFS
Web	✓	✓	✓	×	Web server
Distributed shared memory	✓	×	✓	✓	Ivy (Ch. 18)
Remote objects (RMI/ORB)	✓	×	×	1	CORBA
Persistent object store	✓	✓	×	1	CORBA Persistent Object Service
Peer-to-peer storage system	✓	✓	✓	✓	OceanStore(Ch. 10)

Types of consistency between copies: 1 - strict one-copy consistency
√ - approximate consistency
X - no automatic consistency

Figure 1 Storage systems and their properties

INTRO(cntd)

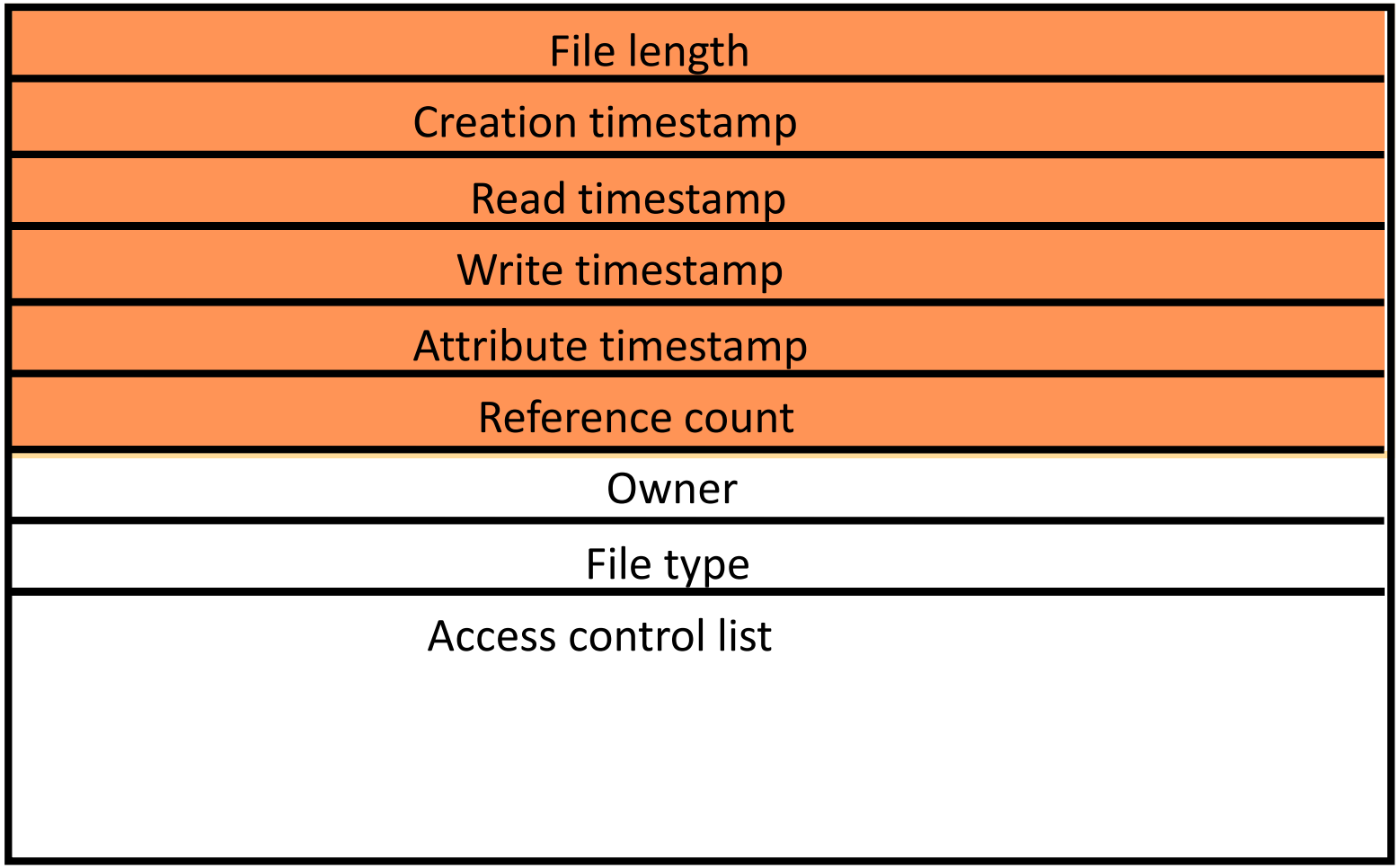
Directory module:	relates file names to file IDs
File module:	relates file IDs to particular files
Access control module:	checks permission for operation requested
File access module:	reads or writes file data or attributes
Block module:	accesses and allocates disk blocks
Device module:	disk I/O and buffering

Figure 2 File System Modules

CHARACTERISTICS OF FILE SYSTEMS

- File systems responsible for the organization, storage, retrieval, naming , sharing and protection of files.
- Programming interface that characterizes the file abstraction.
- Files contain both **data** and **attributes.**
- **Metadata –** extra information stored by a file system for management of files.

File attribute Record Structure



FILE SYSTEM OPERATIONS

Figure 4. UNIX file system operations

filedes = open(name, mode) *filedes = creat(name, mode)*	Opens an existing file with the given *name*. Creates a new file with the given *name*. Both operations deliver a file descriptor referencing the open file. The *mode* is *read*, *write* or both.
status = close(filedes)	Closes the open file *filedes*.
count = read(filedes, buffer, n) *count = write(filedes, buffer, n)*	Transfers *n* bytes from the file referenced by *filedes* to *buffer*. Transfers *n* bytes to the file referenced by *filedes* from buffer. Both operations deliver the number of bytes actually transferred and advance the read-write pointer.
pos = lseek(filedes, offset, whence)	Moves the read-write pointer to offset (relative or absolute, depending on *whence*).
status = unlink(name)	Removes the file *name* from the directory structure. If the file has no other names, it is deleted.
status = link(name1, name2)	Adds a new name (*name2*) for a file (*name1*).
status = stat(name, buffer)	Gets the file attributes for file *name* into *buffer*.

DISTRIBUTED FILE SYSTEM REQUIREMENTS

- Transparency

  Design of file service should support transparency requirements .

- Concurrent File updates

  Changes to a file should not interfere with operation of other clients accessing the same file .

DISTRIBUTED FILE SYSTEM REQUIREMENTS

- File replication

  A file may be represented by several copies of its contents at different locations

- Hardware and operating system heterogeneity

  Important aspect of openness

- Fault Tolerance

  File service must continue to operate in the face of client and server failures.

- Consistency

  One-copy semantics

DISTRIBUTED FILE SYSTEM REQUIREMENTS

- Security

  Need to authenticate client requests based on user identities and encryption of secret data.

- Efficiency

  Should offer facilities that are at least the same power and generality as found in conventional file systems.

THANK YOU