

Load Sharing

Reference: Pradeep K Sinha,
"Distributed Operating Systems: Concepts and
Design", Prentice Hall of India, 2007



Load sharing approach

- For the **proper utilization** of the **resources** of a **distributed system** it is not required to **balance** the **load** on all the **nodes**.
- It is **necessary** and **sufficient** to prevent the **nodes** from being **idle** while some **other node** have **more** than **two processes**.
- This rectification is often called **dynamic load sharing** instead of **dynamic load balancing**.



Issues in designing load-sharing algorithms

- Load sharing algorithms do not attempt to balance the average workload on all the nodes of the system, rather they only attempt to ensure that no node is idle when a node is heavily loaded.
- The priority assignment policies and migration limiting policies are same as that for the load-balancing algorithms. Other policies are described here.



Load Estimation Policy

- It is **sufficient** to know whether a **node** is **busy** or **idle**.
- Methods for estimating load:
 - **Count** the **total number** of **processes** on a node.
 - Measure **CPU utilization**.



Process transfer policies

- All-or-nothing strategy:
 - Uses the **single threshold policy** with the **threshold value** of **all the node fixed** at **1** and some uses **2**.
 - **Drawback** : **Loss** of **available processing power** in the system.
 - **Solution** : use a **threshold** value of **2** instead of **1**.



Location policies

1. Sender-initiated policy
2. Receiver-initiated policy



Sender initiated location policy

- Heavily loaded nodes search for lightly loaded node to which work may be transferred.
- When load becomes more than the threshold value, it either broadcasts a message or randomly probes the other nodes one by one to find a lightly loaded node.



Sender initiated location policy

- In the **broadcasting method**, the **presence** or **absence** of a **suitable receiver node** is known as soon as the **sender node receives reply** messages from the **other nodes**.



Sender initiated location policy

- In the **random probing** method, the **probing continues until** either a **suitable node** is **found** or the **no. of probes** reaches a **static probe limit, L_p** .
- **Fixed limit** has better **scalability** than broadcast method.



Receiver-initiated location policy

- Lightly loaded node search for heavily loaded nodes from which work may be transferred.
- When a node's load falls below the threshold value either it broadcasts a message indicating its willingness to receive processes or randomly probes the other nodes one by one to find a heavily loaded node.



Receiver-initiated location policy

- In the **broadcast method**, a **suitable node** is **found** as soon as the **receiver node** receives **reply messages** from the other **nodes**.
- In the **random probing** method, the probing **continues** until either a **suitable node** is **found** or the **no. of probes** reaches a **static probe limit**, L_p .



Location Policies

- Both **sender-initiated** and **receiver-initiated policies** offer substantial **performance** advantages over the situation in which **no load sharing** is attempted.
- **Sender-initiated policies** are preferable at **light to moderate system loads**, while **receiver-initiated policies** are preferable at **high system loads**.



Location Policies

- If the cost of process transfer under receiver-initiated policies is significantly greater than under the sender-initiated policies due to the preemptive transfer of processes
- Sender-initiated policies provide uniformly better performance.



State information Exchange Policy

1. Broadcast when state changes.
2. Poll when state changes.



Thank You

