

**TRANSLATION SCHEME FOR ASSIGNMENT STATEMENT**

The translation scheme is used to show how symbol table entries can be found. Terms used in the scheme is

lookup (id.name): checks for the entry name in the symbol table. If so it returns the pointer to that entry else it returns nil.

emit: used to emit three address statements to an output file

SNo	Productions	Semantic action
1	$S \rightarrow id=E$	$p = \text{lookup}(id.name)$ if $(p \neq \text{nil})$ then emit $(p=E.place)$ else error
2	$E \rightarrow E1+E2$	$E.place = \text{newtemp}$ emit $(E.place=E1.place+E2.place)$
3	$E \rightarrow E1 * E2$	$E.place = \text{newtemp}$ emit $(E.place=E1.place * E2.place)$
4	$E \rightarrow -E$	$E.place = \text{newtemp}$ emit $(E.place=E1.place = \text{uminus} E2.place)$
5	$E \rightarrow (E)$	$E.place = E1.place$
6	$E \rightarrow id$	$p = \text{lookup}(id.name)$ if $(p \neq \text{nil})$ then emit $(p=E.place)$ else error

**Fig. 1 Translation scheme for the assignment statement**

The above translation scheme is similar to the translation of three address code for the example  $a=b*-c+b*-c$ . During decalaration statement, variable names are entered into the symbol table. While executing the assignment statement, variable names are searched in the symbol table.

### **Type Conversions within Assignments**

There can be different types of variables or constants. So the compiler must either

1. reject mixed-type operations
2. or generate appropriate coercion (type conversion) instruction

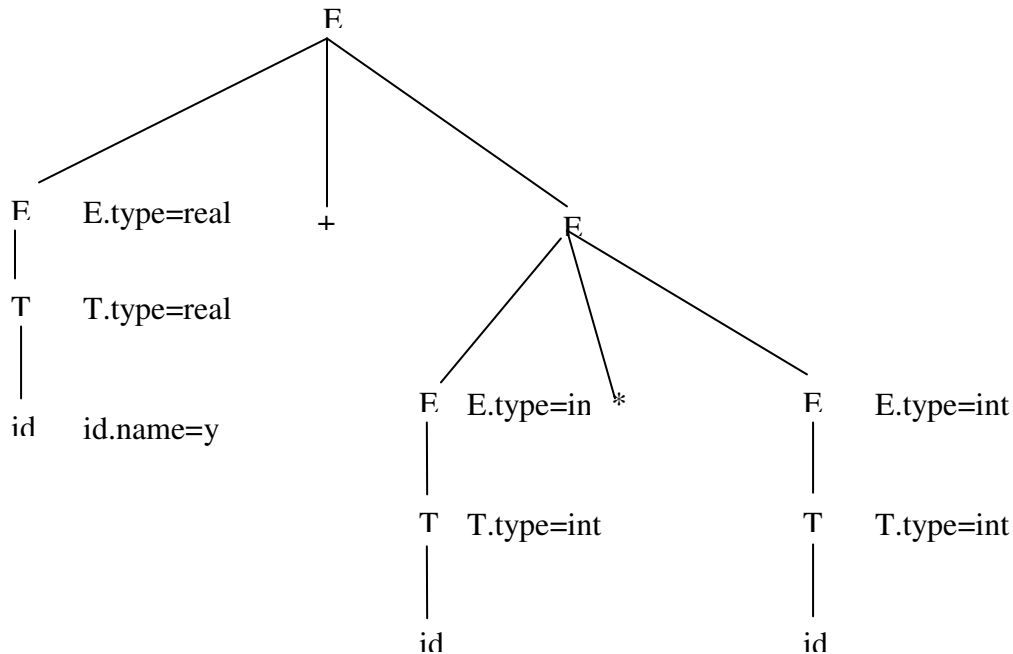
SNo	Productions	Semantic action
1	E->E+E	E.place=newtemp; if E1.type=integer and E2.type=integer then begin emit (E.place=E1.place+E2.place) E.Type=integer end else if E1.type=real and E2.type=real then begin emit (E.place=E1.place+E2.place) E.Type=integer end else if E1.type=integer and E2.type=real then begin u=newtemp emit (u=inttorealE2.place) emit (E.place=E1.place+E2.place) E.Type=real end else if E1.type=real and E2.type=integer then begin u=newtemp emit (u=inttoreal E1.place) emit (E.place=E1.place+E2.place) E.Type=real End
2	E->E1*E2	E.place=newtemp; if E1.type=integer and E2.type=integer then begin emit (E.place=E1.place*E2.place) E.Type=integer end else if E1.type=real and E2.type=real then begin emit (E.place=E1.place*E2.place) E.Type=integer end else if E1.type=integer and E2.type=real then begin u=newtemp emit (u=inttorealE2.place) emit (E.place=E1.place*E2.place) E.Type=real end else if E1.type=real and E2.type=integer then begin u=newtemp emit (u=inttoreal E1.place) emit (E.place=E1.place*E2.place) E.Type=real End
3	E->T	E.Type=T.Type

**Fig. 2 Translation scheme for the assignment statement including type conversion**

Example:

**$x = y + i * j$**

assume x and y are real variables, and i and j are integer variables



The output would be like this

**t1=i\*j  
t2=inttoreal(t1)  
t3=y\*t2  
x=t3**

First E is derived to E+E in which first E is derived to T and further derived to id. T.type will be real since the variable y is real. Second E is derived to E\*E in which first E is derived to T and further derived to id. T.type will be integer and the second E is further divide into id and T.type is integer since the variable j is integer. When the semantics at the appropriate places we'll get the TAC for the above assignment statement.