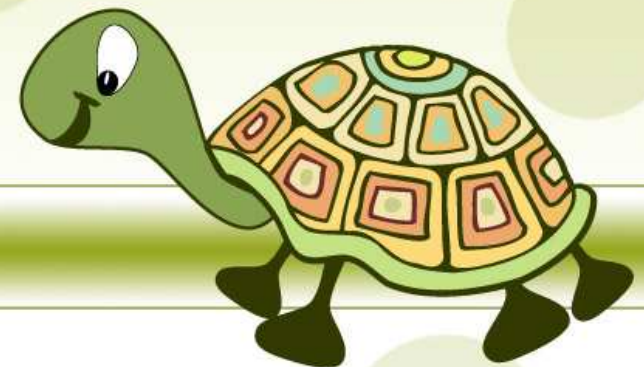


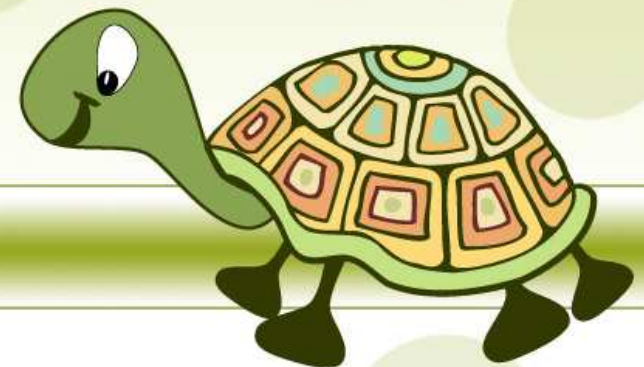
Resource Management

Reference: Pradeep K Sinha,
"Distributed Operating Systems: Concepts
and Design", Prentice Hall of India, 2007



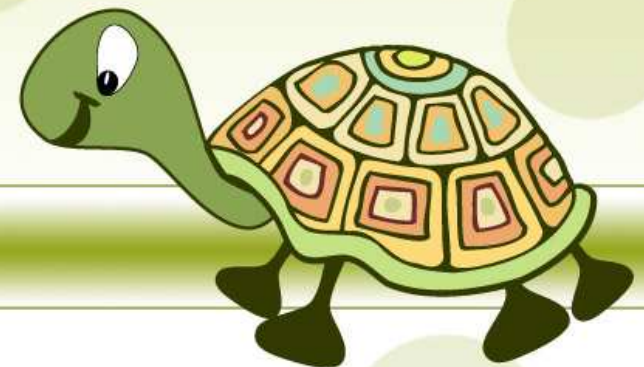
Overview

- Resource Management Techniques
- Desirable features of Scheduling Algorithms
- Task Assignment



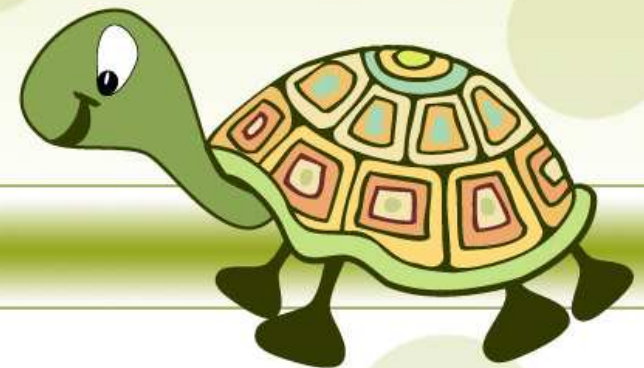
Introduction

- A **resource** can be **logical**, such as a shared **file** or **physical** such as **CPU**.
- The set of **available resources** in a **distributed system** acts like a single **virtual system**.



Introduction

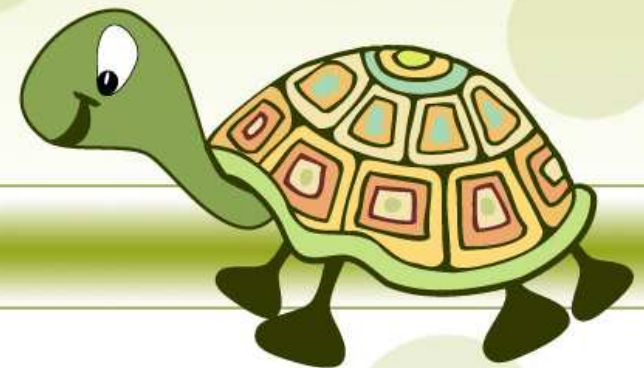
- Resource manager:
 - Controls the **assignment** of **resources** to processes.
 - Routes the **processes** to **suitable nodes** of the system in such a manner that **resource usage**, **response time**, **network congestion**, and **scheduling overhead** are optimized.



Resource Management Techniques

Task assignment approach:

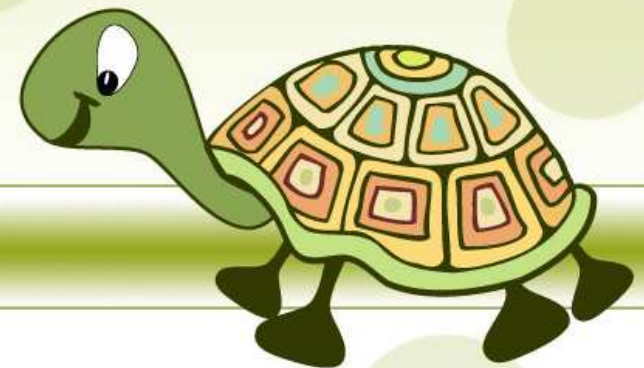
- Each **process submitted** by a **user** for processing is viewed as a **collection** of related **tasks**.
- **Tasks** are **scheduled** to suitable **nodes** to improve performance.



Resource Management Techniques

Load-balancing approach:

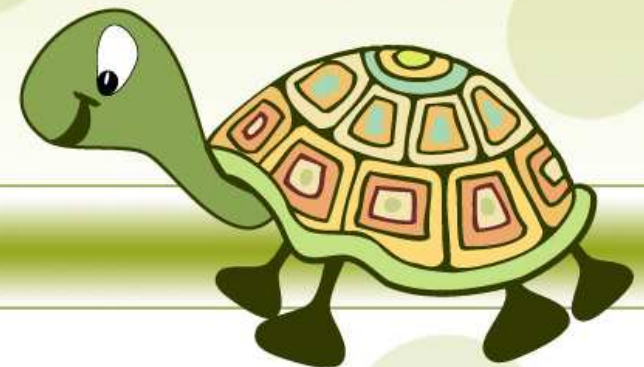
- All the **processes** submitted by the **users** are **distributed** among the **nodes** of the system.
- **Equalizes** the **workload** among the **nodes**.



Resource Management Techniques

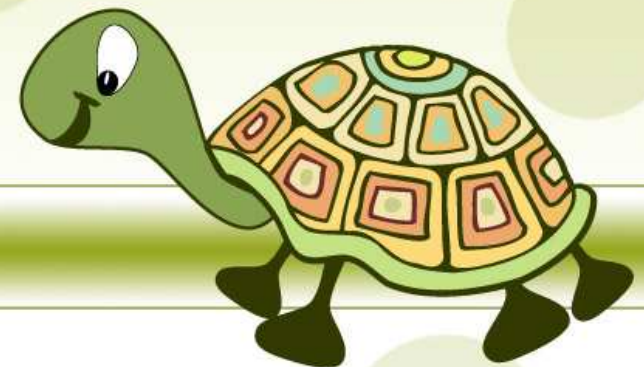
Load-sharing approach:

- Attempts to conserve the ability of the system, assuring that **no node** is **idle** while **processes wait** for being **processed**.



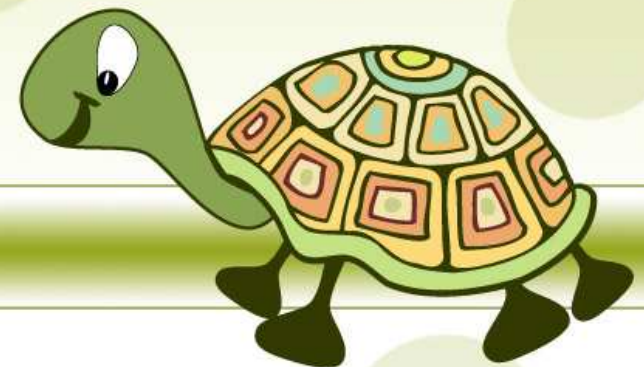
Desirable features of a good Scheduling Algorithms

- No a priori knowledge about the processes.
- Dynamic in nature.
- Quick decision-making capability.
- Balanced system performance.
- Stability.
- Scalability.
- Fault tolerance.
- Fairness of service.



Task assignment approach

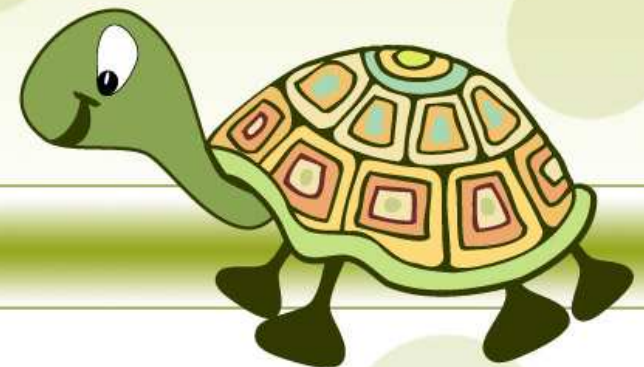
- A **process** is considered to be composed of **multiple task**.
- Goal is to find an **optimal assignment policy** for the **task** of an individual process.



Task assignment approach

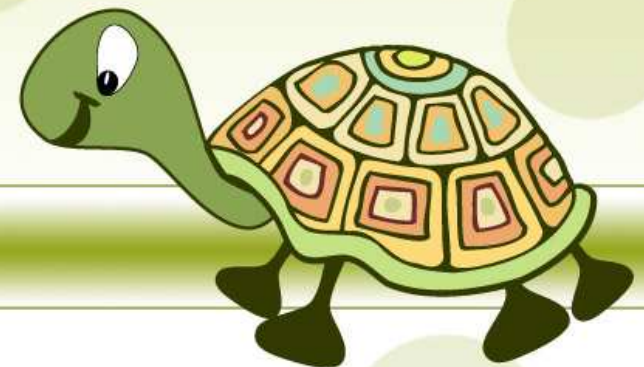
Assumptions:

1. A **process** has already been **split** into **pieces** called **tasks**.
2. Amount of **computation required** by each task and **speed** of each **processor** are known.
3. The **cost of processing** each task on every node of the system is known.
4. The **IPC costs** between every pair of task is known.



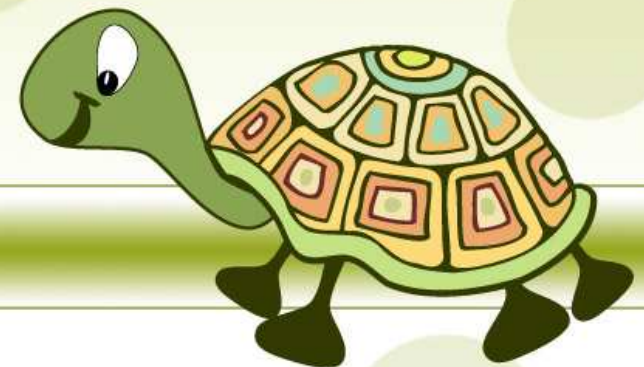
Task assignment approach

5. Other constraints, like **Resource requirements** of the **tasks** and the **available resources** at each node are also known.
6. **Reassignment** of the **tasks** is generally **not possible**.



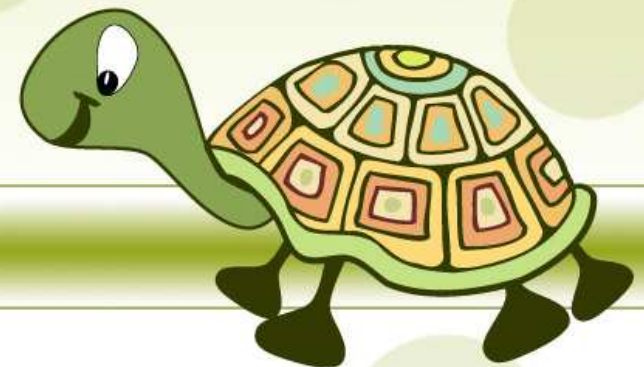
Task assignment approach

- Goals:
 - Minimization of IPC costs
 - Quick turnaround time for the complete process
 - A high degree of parallelism
 - Efficient utilization of system resources in general
- These goals often conflict with each other.



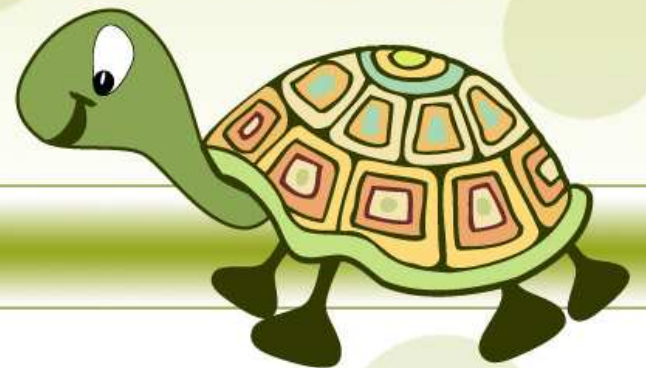
Task assignment approach

- Two task assignment parameters
 - Task execution cost &
 - Inter-task communication cost



Example

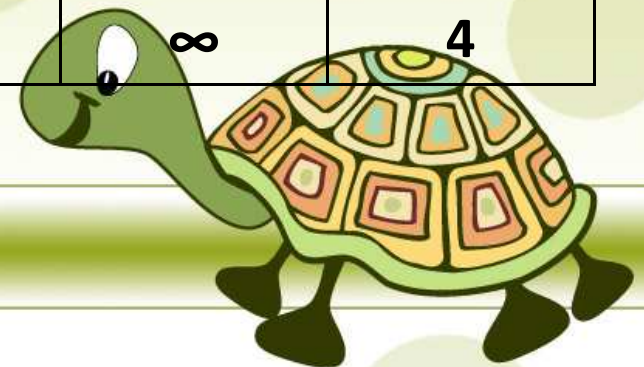
- Total tasks = 6
- Total nodes = 2



Example

Inter task communications cost						
	t 1	t 2	t 3	t 4	t 5	t 6
t 1	0	6	4	0	0	12
t 2	6	0	8	12	3	0
t 3	4	8	0	0	11	0
t 4	0	12	0	0	5	0
t 5	0	3	11	5	0	0
t 6	12	0	0	0	0	0

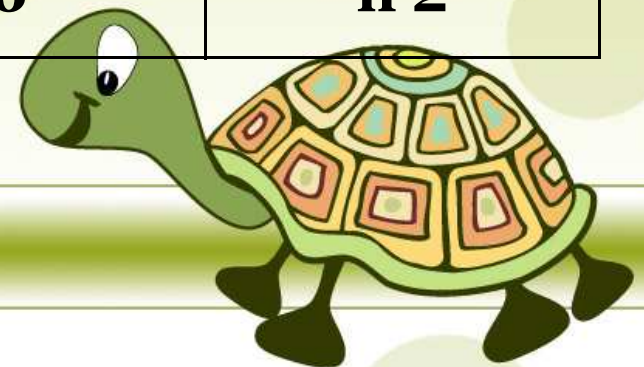
Execution costs		
Task	Nodes	
	n 1	n 2
t 1	5	10
t 2	2	∞
t 3	4	4
t 4	6	3
t 5	5	2
t 6	∞	4



Example

Serial assignment	
Task	Node
t 1	n 1
t 2	n 1
t 3	n 1
t 4	n 2
t 5	n 2
t 6	n 2

Optimal assignment	
Task	Node
t 1	n 1
t 2	n 1
t 3	n 1
t 4	n 1
t 5	n 1
t 6	n 2

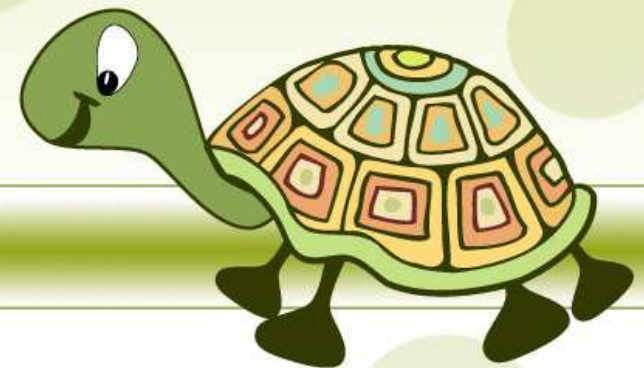


Example

$$\begin{aligned}\text{Serial assignment execution cost (x)} \\ &= x_{11} + x_{21} + x_{31} + x_{42} + x_{52} + x_{62} \\ &= 5 + 2 + 4 + 3 + 2 + 4 = 20\end{aligned}$$

$$\begin{aligned}\text{Serial assignment communication cost (c)} \\ &= c_{14} + c_{15} + c_{16} + c_{24} + c_{25} + c_{26} + c_{34} + c_{35} + c_{36} \\ &= 0 + 0 + 12 + 12 + 3 + 0 + 0 + 11 + 0 = 38\end{aligned}$$

$$\begin{aligned}\text{Total Serial assignment cost} \\ &= x + c = 20 + 38 = 58\end{aligned}$$



Example

Optimal assignment execution cost (x)

$$= x_{11} + x_{21} + x_{31} + x_{41} + x_{51} + x_{62}$$

$$= 5 + 2 + 4 + 6 + 5 + 4 = 26$$

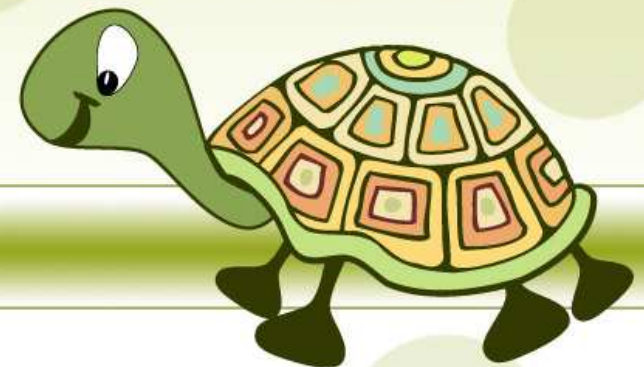
Optimal assignment communication cost (c)

$$= c_{16} + c_{26} + c_{36} + c_{46} + c_{56}$$

$$= 12 + 0 + 0 + 0 + 0 = 12$$

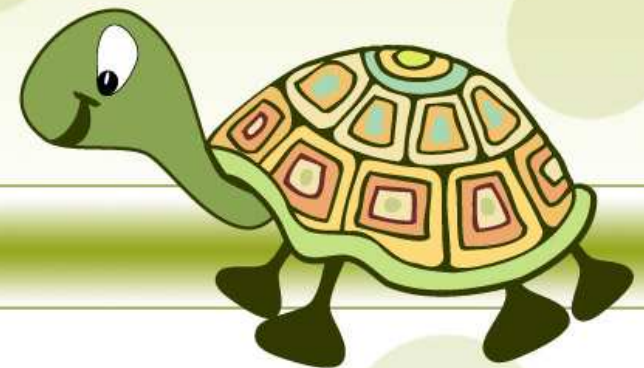
Total Optimal assignment cost

$$= x + c = 26 + 12 = 38$$



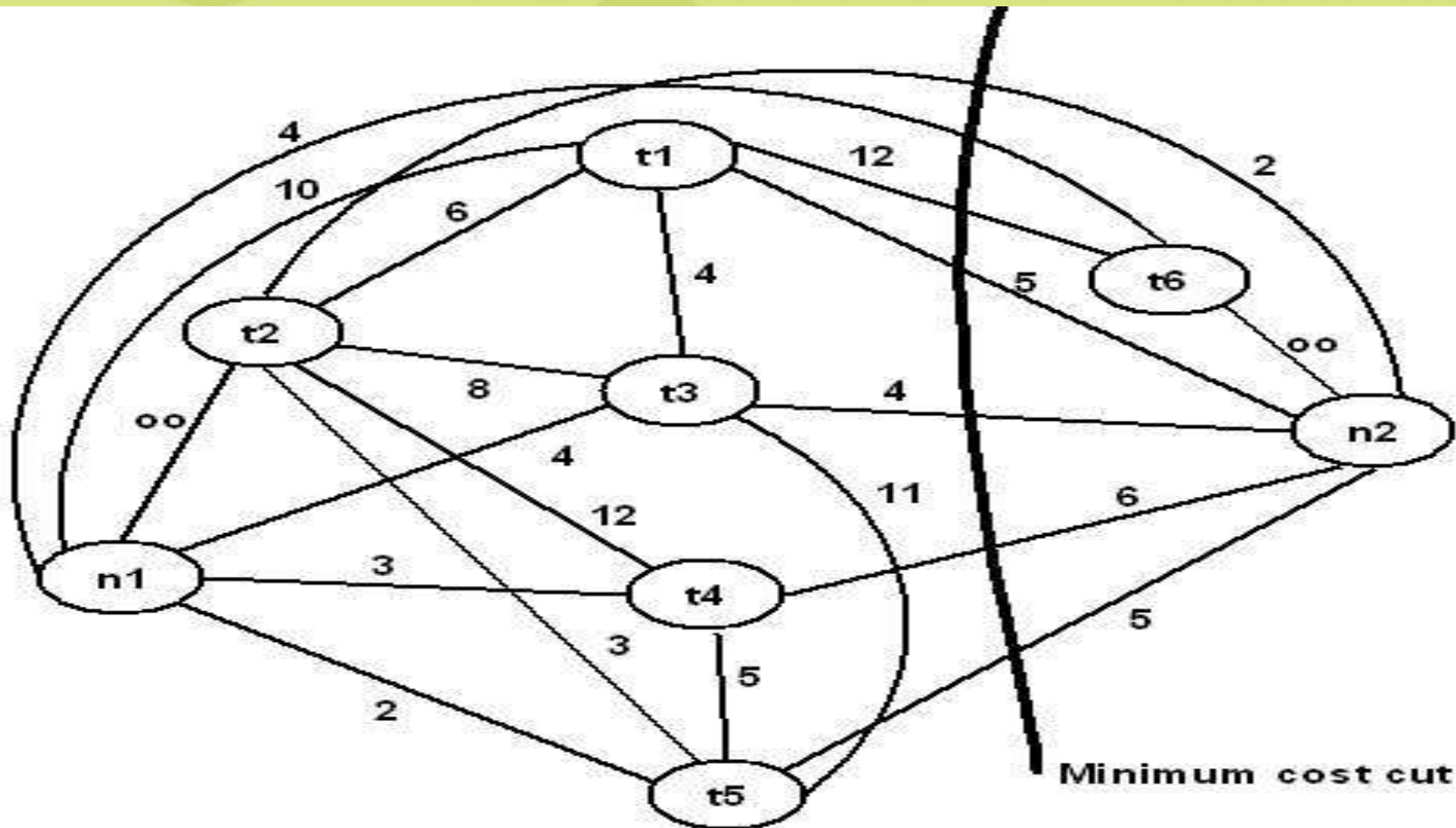
Finding an optimal assignment

- **Problem** : Finding an **assignment** of **tasks** to nodes that **minimizes total execution and communication costs**.
- An optimal assignment is found by creating a **static assignment graph**.



Finding an optimal assignment

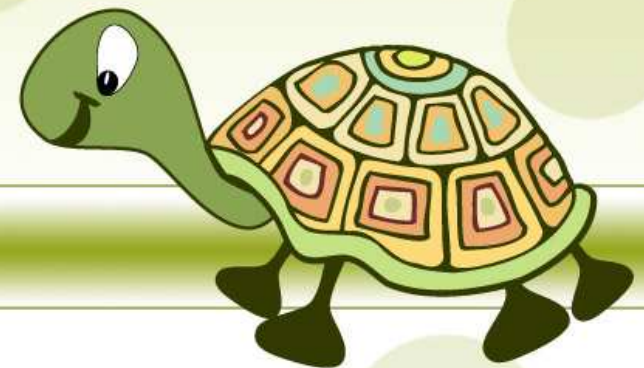
Assignment graph:



Assignment graph for the assignment problem with minimum cost cut

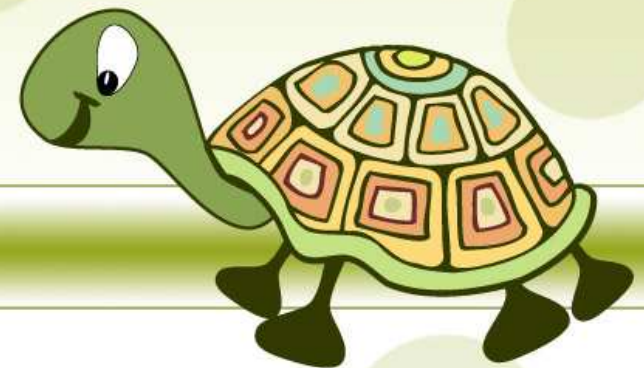
Finding an optimal assignment

- Node **n1** and node **n2** represent the two nodes (**processors**).
- Nodes **t1** through **t6** represent the **task** of the process.
- **Weight** of the edge represent **inter-task communication cost**.
- Weight on the edge joining a **task node** to **node n1** represent the **execution cost** of that **task** on **node n2** and vice versa.



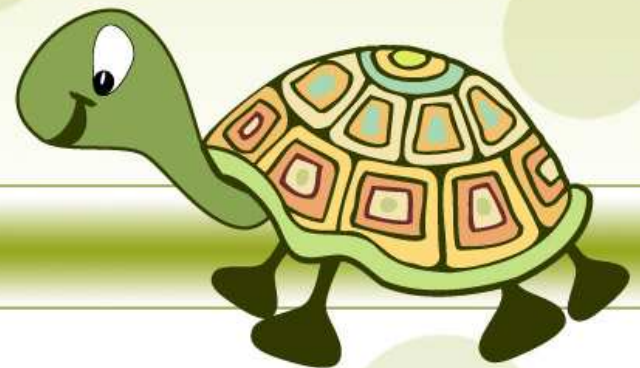
Finding an optimal assignment

A **cut-set** in this graph is defined to be a **set of edges** such that when these **edges** are **removed**, the nodes of the **graph** are **partitioned** into **two disjoint subset**.



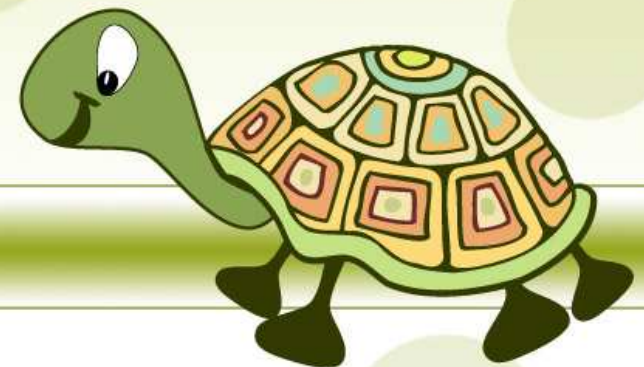
Finding an optimal assignment

- The weight of a **cut-set** is the **sum** of the **weight** of the **edge** in the **cut-set**.
 - represents the **cost** of the corresponding **task assignment**.
- An optimal assignment may be obtained by **finding a minimum-weight cut-set**.



Finding an optimal assignment

- The bold line indicates a **minimum weight cut-set** that corresponds to the **optimal assignment**.
- In a **two processor system**, an optimal assignment can be found in **polynomial time** by **utilizing max flow/ min cut algorithms**.



Thank You

