



Parsing with Context-Free Grammars

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Overview

- Background – Parsing
- Parsing as Search
 - Top-down Parsing
 - Bottom-up Parsing
- A Basic Top-Down Parser
 - Adding bottom-up filtering
- Problems: Basic Top-Down Parser
 - Left-recursion
 - Ambiguity
 - Repeated Parsing subtrees
- Dynamic Programming

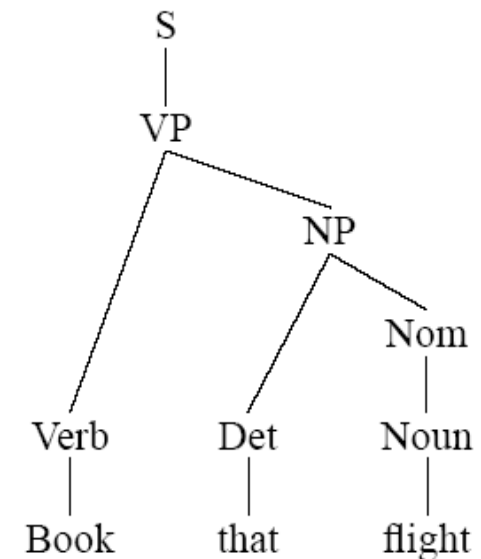
Background

- ♦ Syntactic parsing
 - The task of recognizing a sentence and assigning a syntactic structure to it
- ♦ Since CFGs are a declarative formalism – do not specify how the parse tree for a given sentence should be computed
- ♦ Parse trees are useful in applications such as :
 - Grammar checking
 - Semantic analysis
 - Machine translation
 - Question answering
 - Information extraction

Parsing as Search

- ♦ The parser can be viewed as searching through the space of all possible parse trees to find the correct parse tree for the sentence.
- ♦ ***How can we use the grammar to assign the parse tree?***

$S \rightarrow NP VP$	$Det \rightarrow that \mid this \mid a$
$S \rightarrow Aux NP VP$	$Noun \rightarrow book \mid flight \mid meal \mid money$
$S \rightarrow VP$	$Verb \rightarrow book \mid include \mid prefer$
$NP \rightarrow Det Nominal$	$Aux \rightarrow does$
$Nominal \rightarrow Noun$	$Prep \rightarrow from \mid to \mid on$
$Nominal \rightarrow Noun Nominal$	$Proper-Noun \rightarrow Houston \mid TWA$
$NP \rightarrow Proper-Noun$	$Nominal \rightarrow Nominal PP$
$VP \rightarrow Verb$	
$VP \rightarrow Verb NP$	



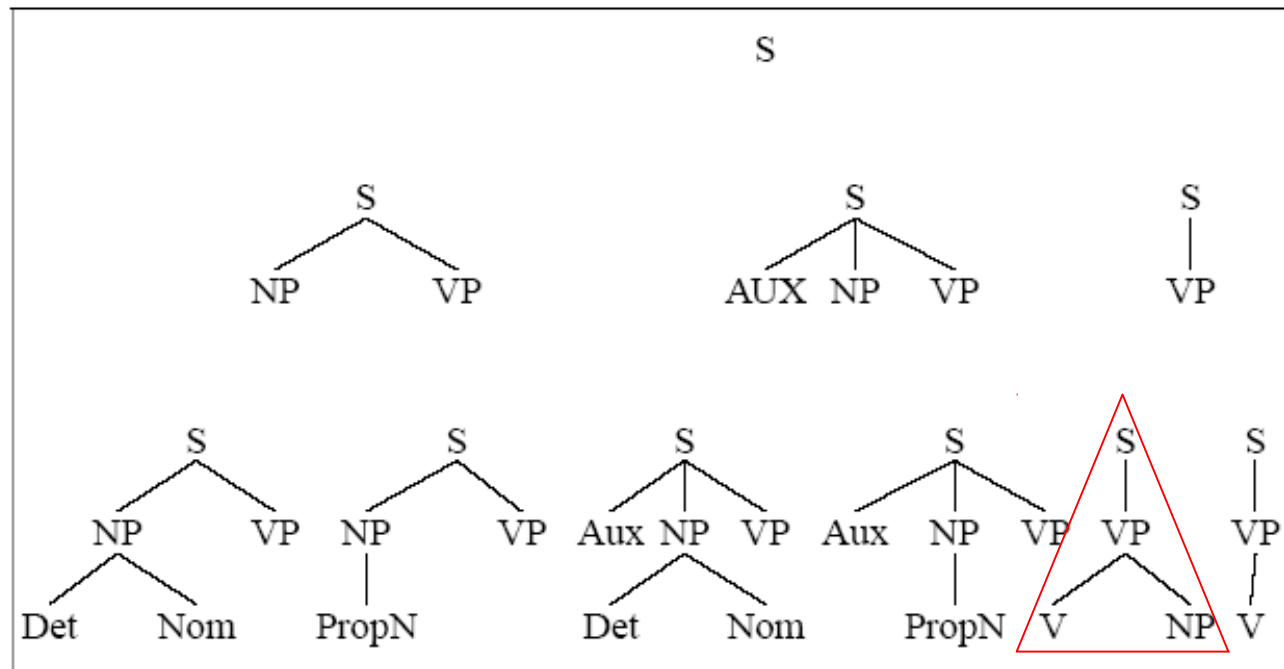
Parsing as Search

- ♦ The goal of a parsing search is to find all trees whose root is the start symbol S , which cover exactly the words in input.
- ♦ Two kinds of constraints:
 - Data – input words as leaves (*book, that, flight*)
 - Grammar – must have one root, which must be the start symbol S
- ♦ Leads to the two search strategies:
 - *Top-down* or goal-directed search
 - *Bottom-up* or data-directed search

Top-down Parser

- ♦ Trying to build from the **root node** S down **to the leaves**
- ♦ Start by start symbol S.
- ♦ Find the grammar rules with S on the left-hand side and expand the constituents in new trees
- ♦ At each level or **ply**, use the right-hand sides of the rules to provide new sets of expectations
- ♦ Tree grows until reaches the part-of-speech categories at the bottom
- ♦ Trees whose leaves fail to match all the words in the input can be rejected

Top-down Parser



$S \rightarrow NP VP$

$S \rightarrow Aux NP VP$

$S \rightarrow VP$

$NP \rightarrow Det Nominal$

$Nominal \rightarrow Noun$

$Nominal \rightarrow Noun Nominal$

$NP \rightarrow Proper-Noun$

$VP \rightarrow Verb$

$VP \rightarrow Verb NP$

$Det \rightarrow that \mid this \mid a$

$Noun \rightarrow book \mid flight \mid meal \mid money$

$Verb \rightarrow book \mid include \mid prefer$

$Aux \rightarrow does$

$Prep \rightarrow from \mid to \mid on$

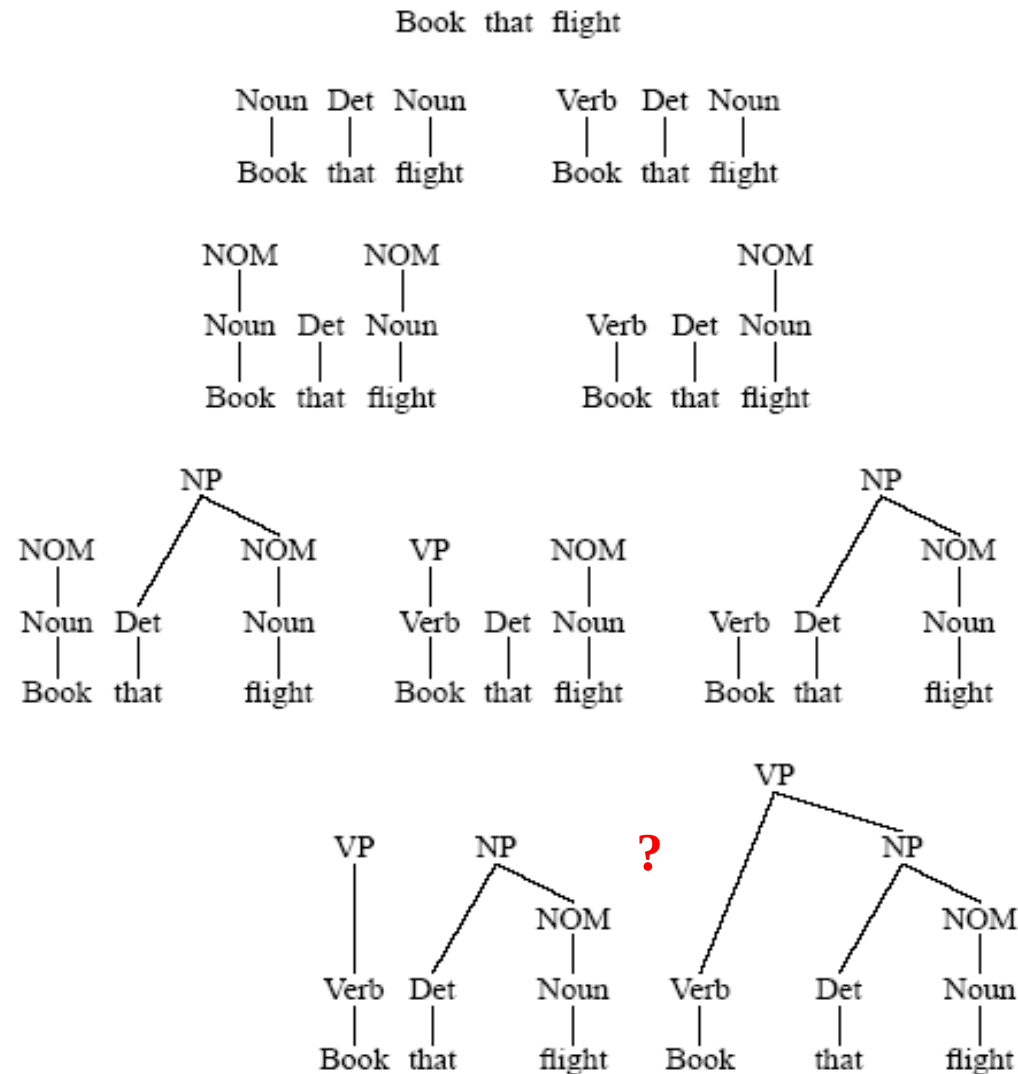
$Proper-Noun \rightarrow Houston \mid TWA$

$Nominal \rightarrow Nominal PP$

Bottom-up Parsing

- ♦ Starts with the words of the input and tries to build trees from the words up, by applying rules from the grammar
- ♦ Initially *for each input word* build partial trees with the part-of-speech
- ♦ For each ply, find the right-hand side of the rule that match the sequence of non-terminals
- ♦ Parse is success, if parser succeeds in building a tree rooted in the start symbol S

Bottom-up Parsing



Top-down Vs Bottom-up

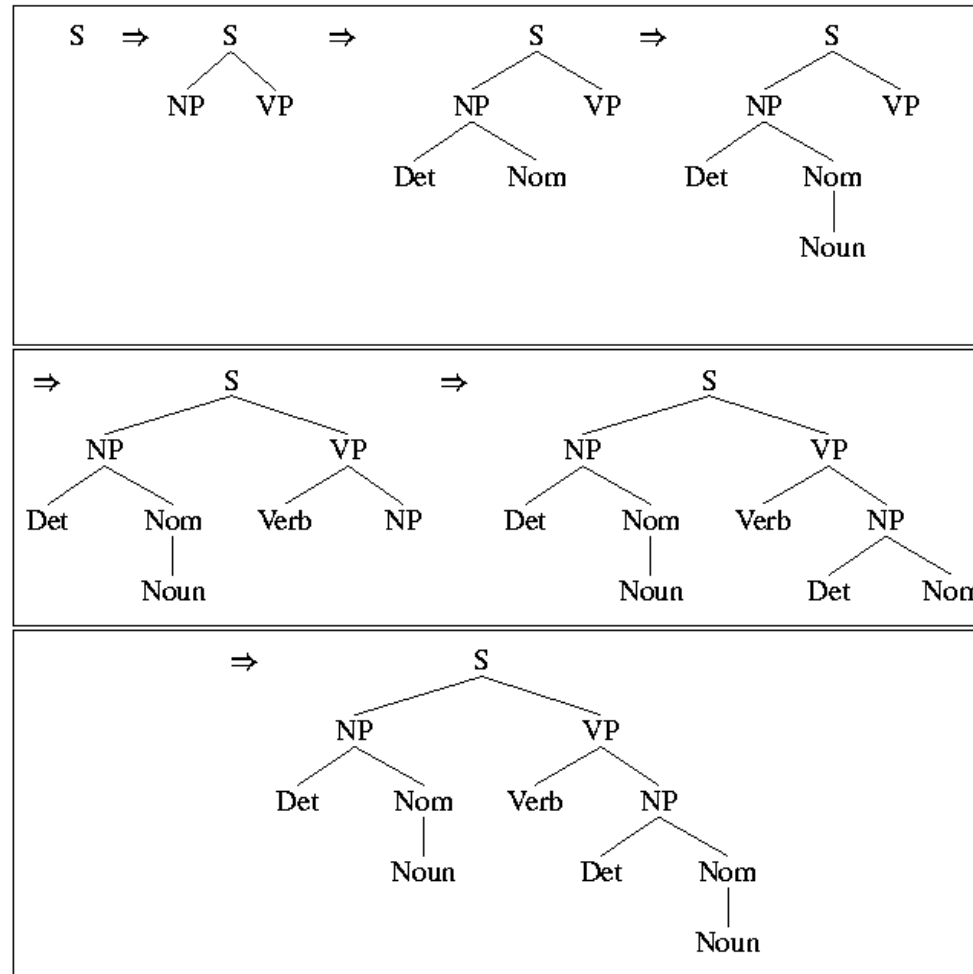
- ♦ ***Comparisons***

- ♦ The top-down strategy never wastes time exploring trees that cannot result in an *S*.
 - Spend considerable effort on *S* trees that are not consistent with the input.
 - Generate trees before ever examining the input.
- ♦ The bottom-up strategy never suggest trees that are *not at least locally grounded* in the actual input
 - Trees that have no hope to leading to an *S* are generated with wild abandon.

A Basic Top-Down Parser

- ♦ Solution: Incorporate features of both the top-down and bottom-up approaches
- ♦ In top-down depth-first approach:
 - Left-most unexpanded leaf node is expanded first
 - Grammar rules are applied according to their textual order

A Basic Top-Down Parser



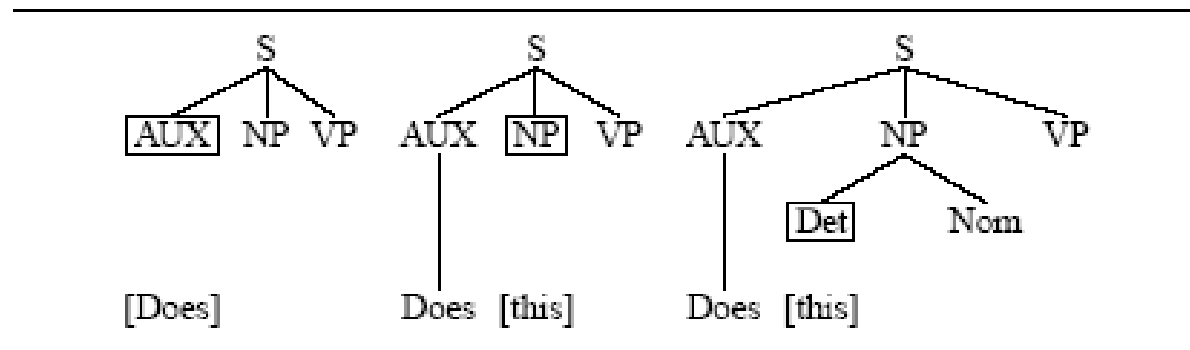
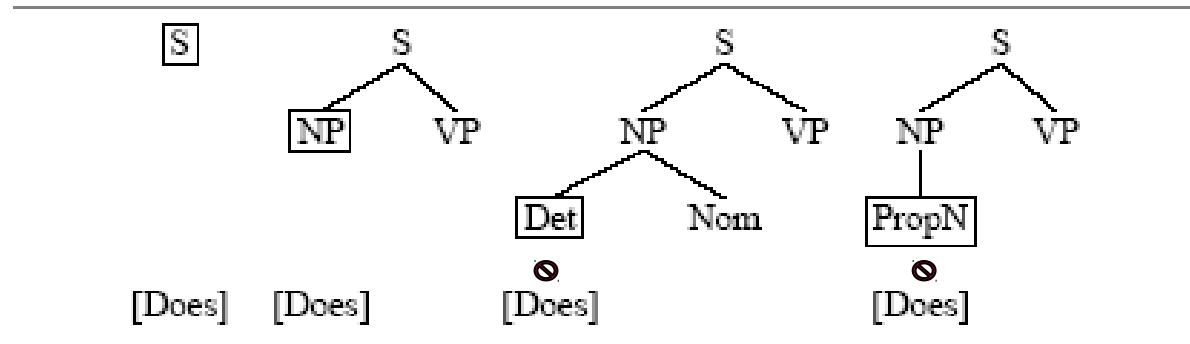
A top-down depth-first derivation

A Basic Top-Down Parser

- ♦ Top-down, depth-first, left-to-right approach: expand the left-most unexpanded node in the tree
 - The node currently being expanded is shown in a box
 - The current input word is bracketed

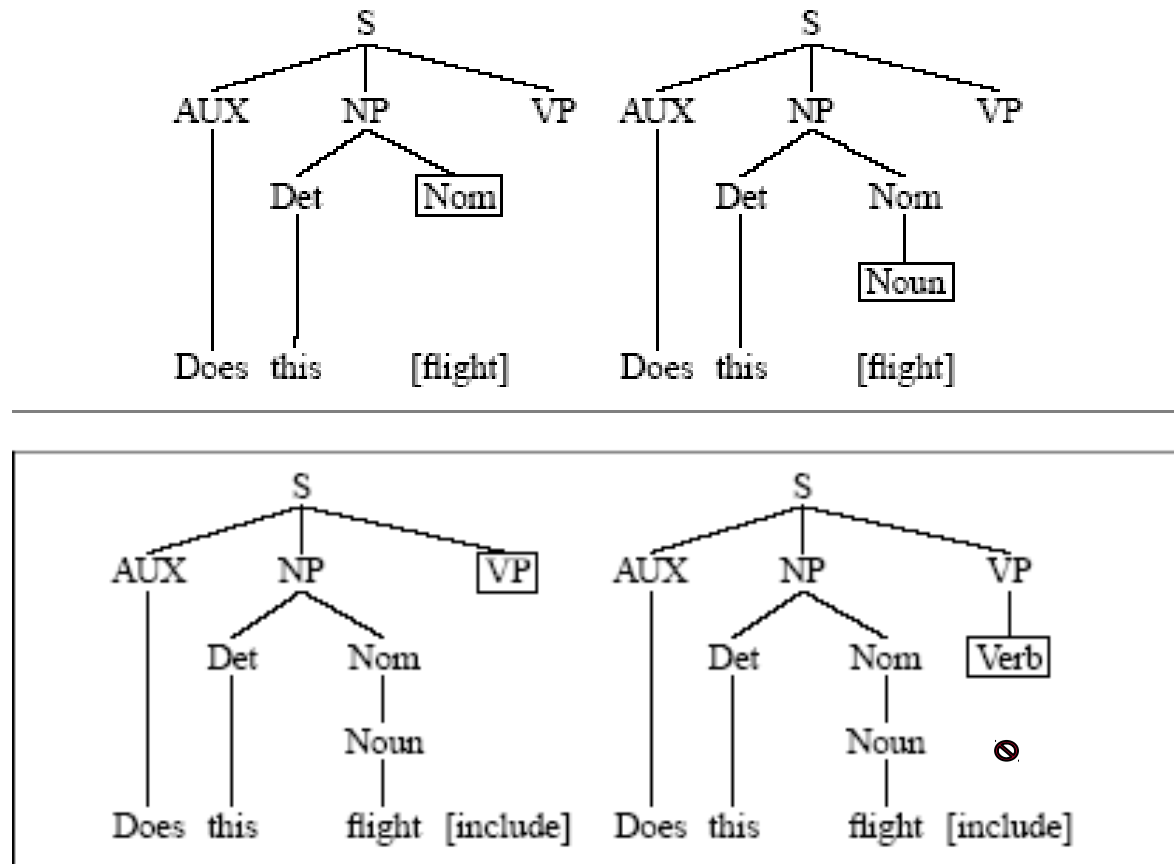
A Basic Top-Down Parser

Does this flight include a meal?



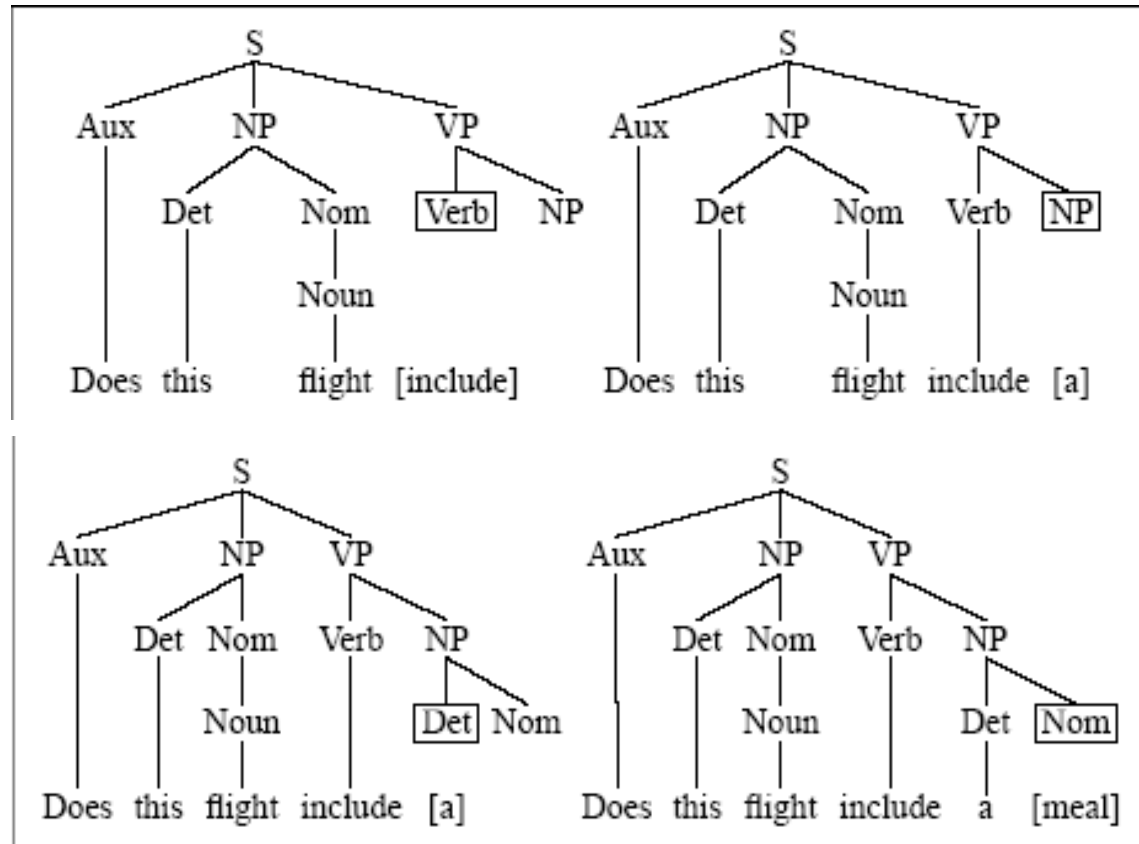
A top-down, depth-first, left-right derivation

A Basic Top-Down Parser



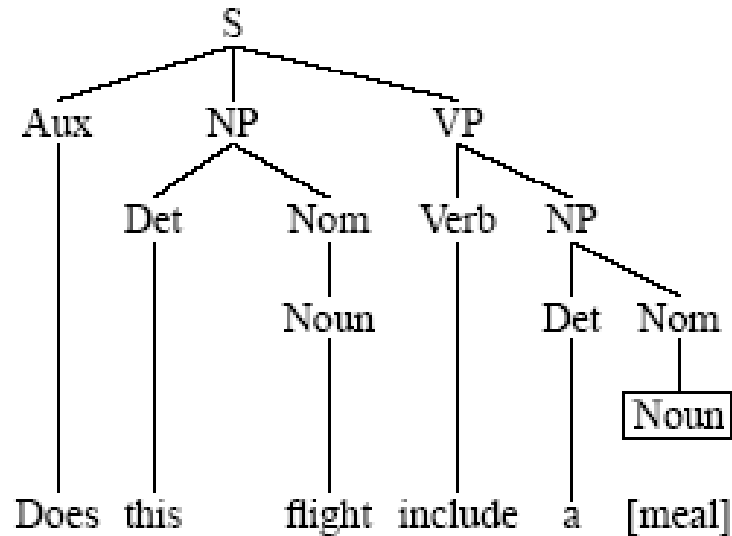
A top-down, depth-first, left-right derivation continued

A Basic Top-Down Parser



A top-down, depth-first, left-right derivation continued

A Basic Top-Down Parser



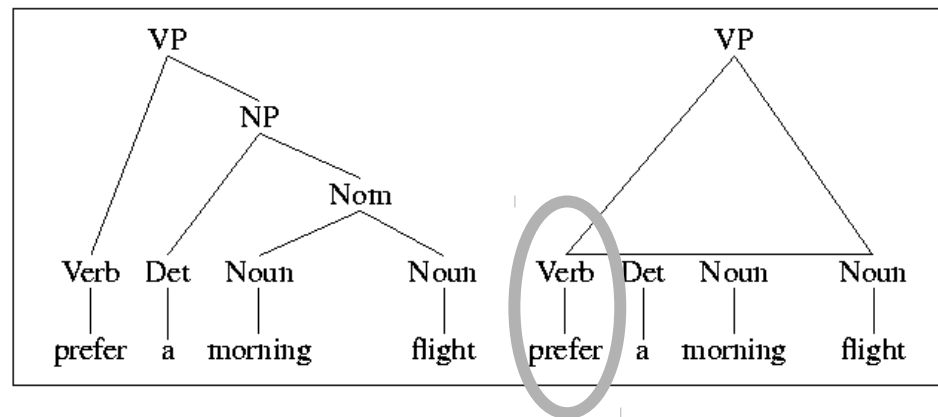
A top-down, depth-first, left-right derivation continued

A Basic Top-Down Parser

```
function TOP-DOWN-PARSE(input, grammar) returns a parse tree
  agenda  $\leftarrow$  (Initial S tree, Beginning of input)
  current-search-state  $\leftarrow$  POP(agenda)
  loop
    if SUCCESSFUL-PARSE?(current-search-state) then
      return TREE(current-search-state)
    else
      if CAT(NODE-TO-EXPAND(current-search-state)) is a POS then
        if CAT(node-to-expand)
           $\subset$ 
            POS(CURRENT-INPUT(current-search-state)) then
              PUSH(APPLY-LEXICAL-RULE(current-search-state), agenda)
            else
              return reject
          else
            PUSH(APPLY-RULES(current-search-state, grammar), agenda)
      if agenda is empty then
        return reject
      else
        current-search-state  $\leftarrow$  NEXT(agenda)
  end
```

A Basic Top-Down Parser

- ♦ ***Adding bottom-up filtering***
- ♦ The parser should not consider any grammar rule if the current input cannot serve as the *first word along the left edge of some derivation* from this rule
- ♦ The first word along the left edge of a derivation is called as the **left-corner** of the tree



A Basic Top-Down Parser

- ♦ ***Adding bottom-up filtering***

- ♦ Left-corner notion

- For non-terminals A and B , B is a left-corner of A if : $A \Rightarrow^* B\alpha$

- ♦ Three rules to expand S :

$S \rightarrow NP VP$

$S \rightarrow Aux NP VP$

$S \rightarrow VP$

Does this flight include a meal ?

- ♦ Using the left-corner notion, it is easy to see that only the $S \rightarrow Aux NP VP$ rule is a viable candidate
- ♦ Since the word *Does* can not serve as the left-corner of other two S -rules

A Basic Top-Down Parser

$S \rightarrow NP VP$	$Det \rightarrow that \mid this \mid a$
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$VP \rightarrow Verb$	
$VP \rightarrow Verb NP$	$Nominal \rightarrow Nominal PP$

Category	Left Corners
S	Det, Proper-Noun, Aux, Verb
NP	Det, Proper-Noun
Nominal	Noun
VP	Verb

Left-corner table for the above Grammar

Problems: Basic Top-Down Parser

- ◆ Problems with the top-down parser
 - Left-recursion
 - Ambiguity
 - Inefficient reparsing of subtrees
- ◆ Solution: Earley algorithm

Problem: Left-recursion

- ♦ Exploring infinite search space, when **left-recursive grammars** are used
- ♦ A grammar is left-recursive if it contains at least one non-terminal A , s.t. $A \xRightarrow{*} \alpha A \beta$, for some α and β and $\alpha \xRightarrow{*} \epsilon$.

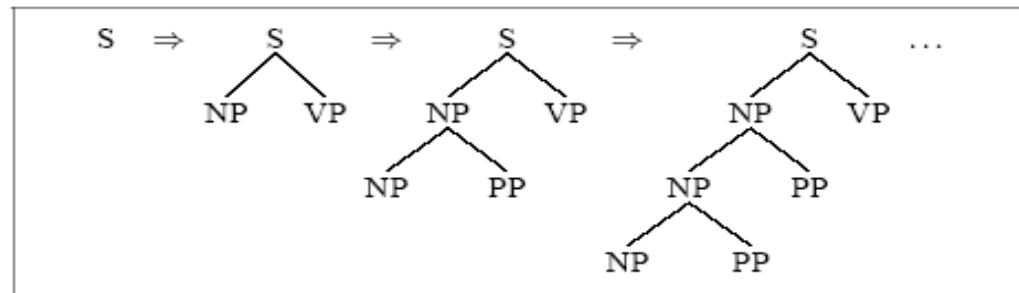
$NP \rightarrow NP PP$

$VP \rightarrow VP PP$

$S \rightarrow S \text{ and } S$

Left-recursive rules

$NP \rightarrow NP PP$



Problem: Left-recursion

- ♦ Two reasonable methods for dealing with left-recursion in a backtracking top-down parser:
 - Rewriting the grammar
 - Explicitly managing the depth of the search during parsing
- ♦ Rewrite each rule of left-recursion

$$A \rightarrow A\beta \mid \alpha \quad \Rightarrow \quad \begin{array}{l} A \rightarrow \alpha A' \\ A' \rightarrow \beta A' \mid \varepsilon \end{array}$$

left-recursive

Weakly equivalent non-left-recursive

- ♦ Rewriting may make semantic interpretation quite difficult

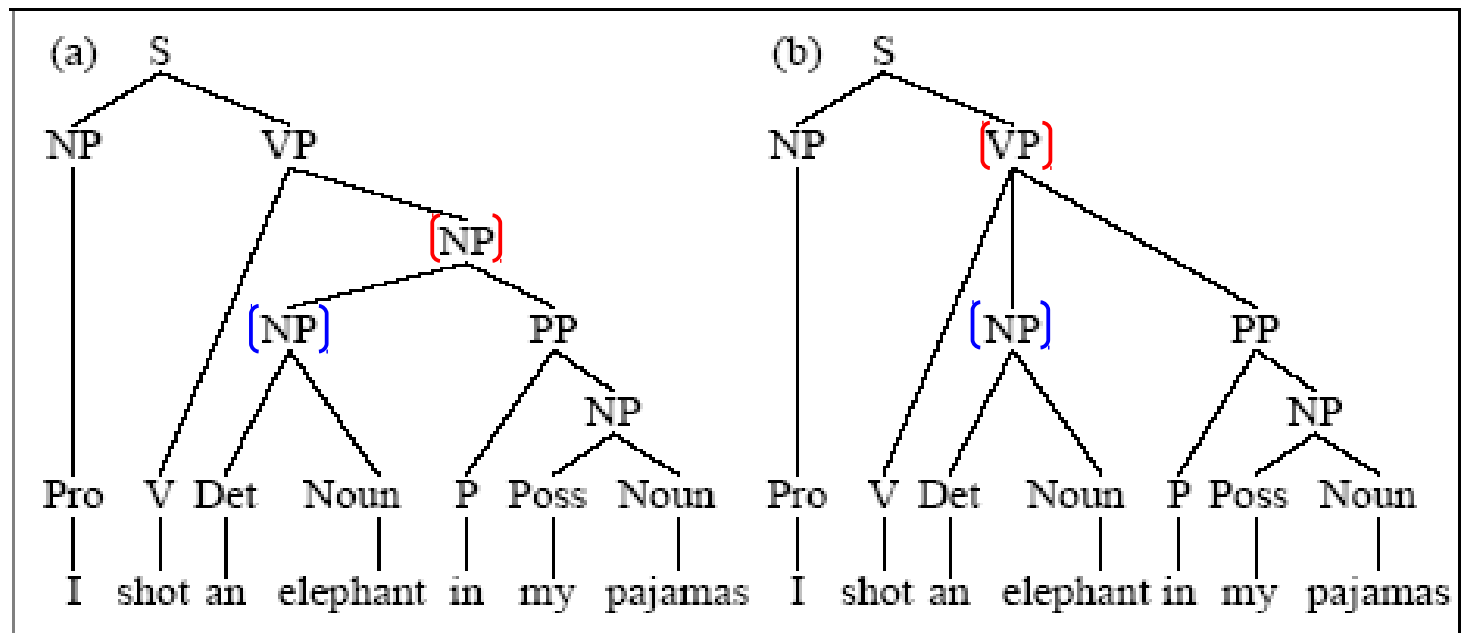
Problem: Ambiguity

- ◆ Lexical category ambiguity – word has more than one pos
- ◆ Disambiguation – choosing the correct pos for a word
- ◆ Structural ambiguity – grammar assigns *more than one possible parse* to a sentence
- ◆ Three kinds of structural ambiguity:
 - attachment ambiguity
 - coordination ambiguity
 - noun-phrase bracketing ambiguity

Problem: Ambiguity

- ◆ Attachment ambiguity

- A particular constituent can be attached to the parse tree at *more than one place*

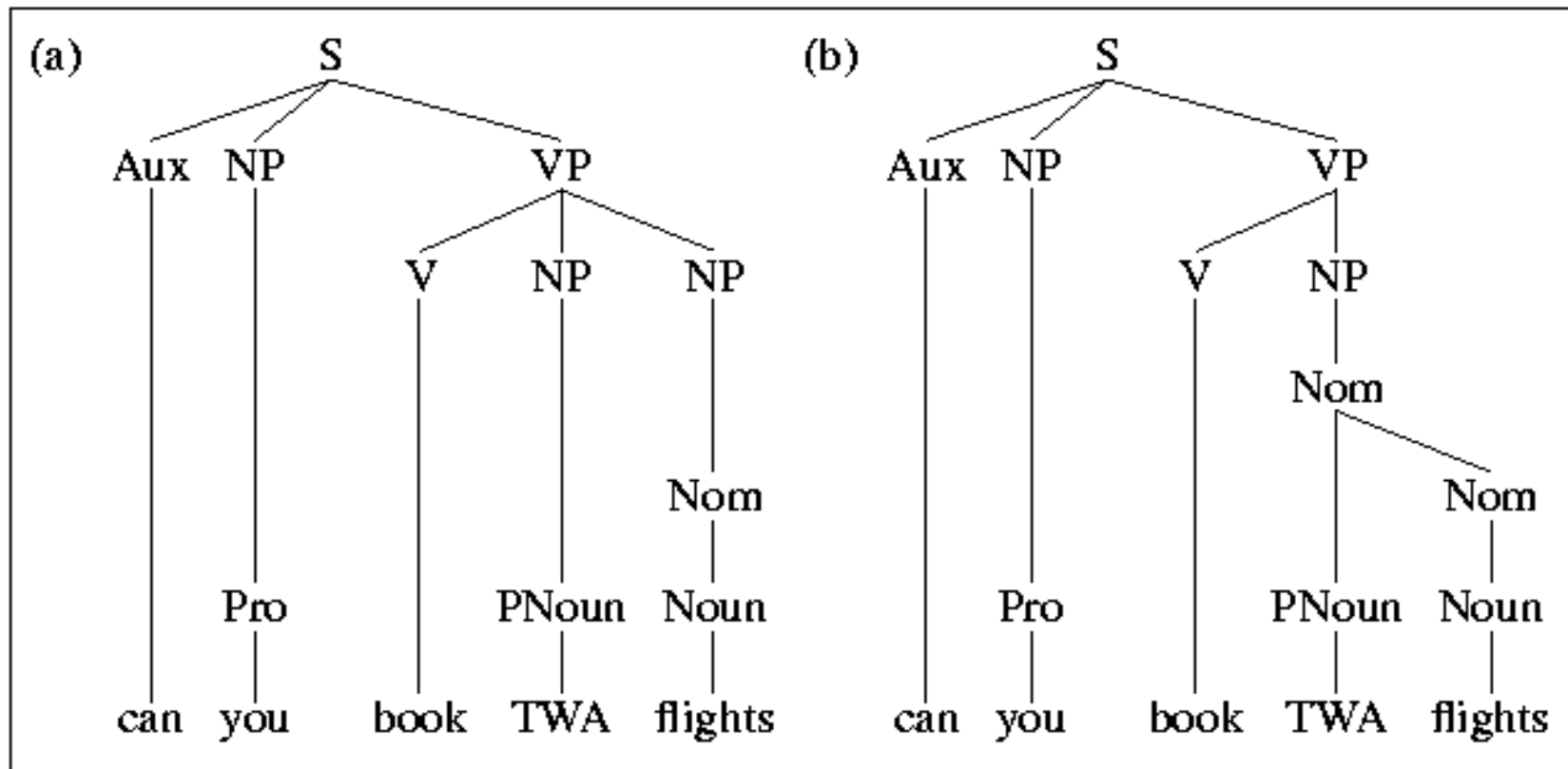


a) elephant is in the pajamas

b) Captain did the shooting in his pajamas

Problem: Ambiguity

- ♦ The sentence “Can you book TWA flights” is ambiguous:



a) Can you book flights on behalf of TWA

b) Can you book flights run by TWA

Problem: Ambiguity

- ♦ *Coordination ambiguity* – different sets of phrases can be conjoined by conjunction like *and*
old men and women
[old [men and women]] \Rightarrow *old men and old women*
[old men] and [women] \Rightarrow *only the men who are old*
- ♦ Parsing sentence thus requires disambiguation:
 - Choosing the correct parse from a multitude of possible parser
 - Requiring both statistical and semantic knowledge

Problem: Ambiguity

- ♦ Parsers which do not incorporate disambiguators may simply return all the possible parse trees for a given input.
- ♦ Potentially exponential number of parses that are possible for certain inputs
 - Show me the meal on Flight UA 386 from San Francisco to Denver.
 - The three PP's at the end of this sentence yield a total of 14 parse trees for this sentence.
 - Solution: use dynamic programming

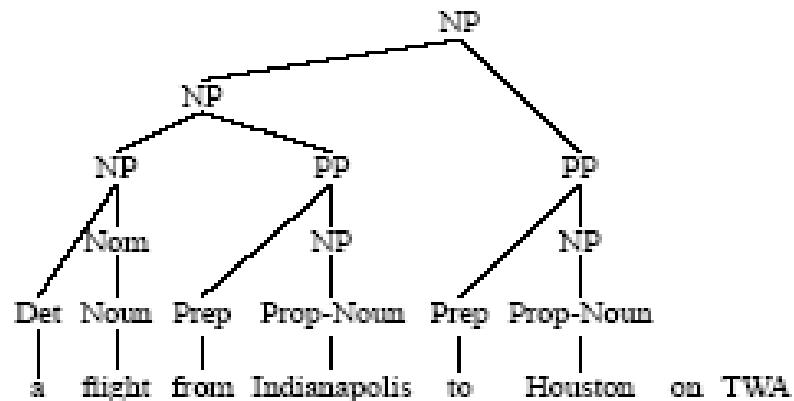
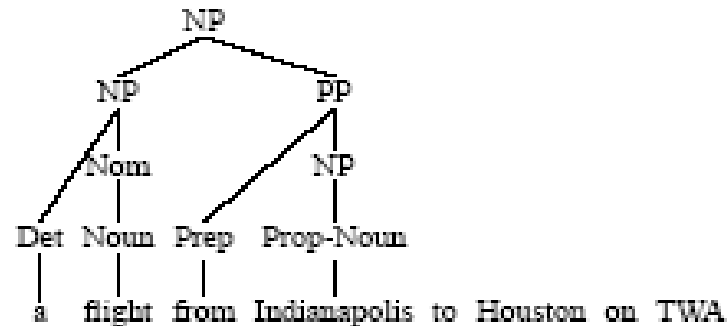
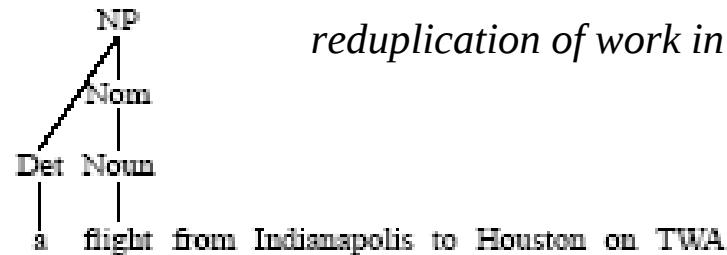
Problem: Repeated Parsing Subtrees

- ♦ The parser often builds valid trees for portions of the input, *then discards* them during backtracking, only *to find that it has to rebuild them again*.

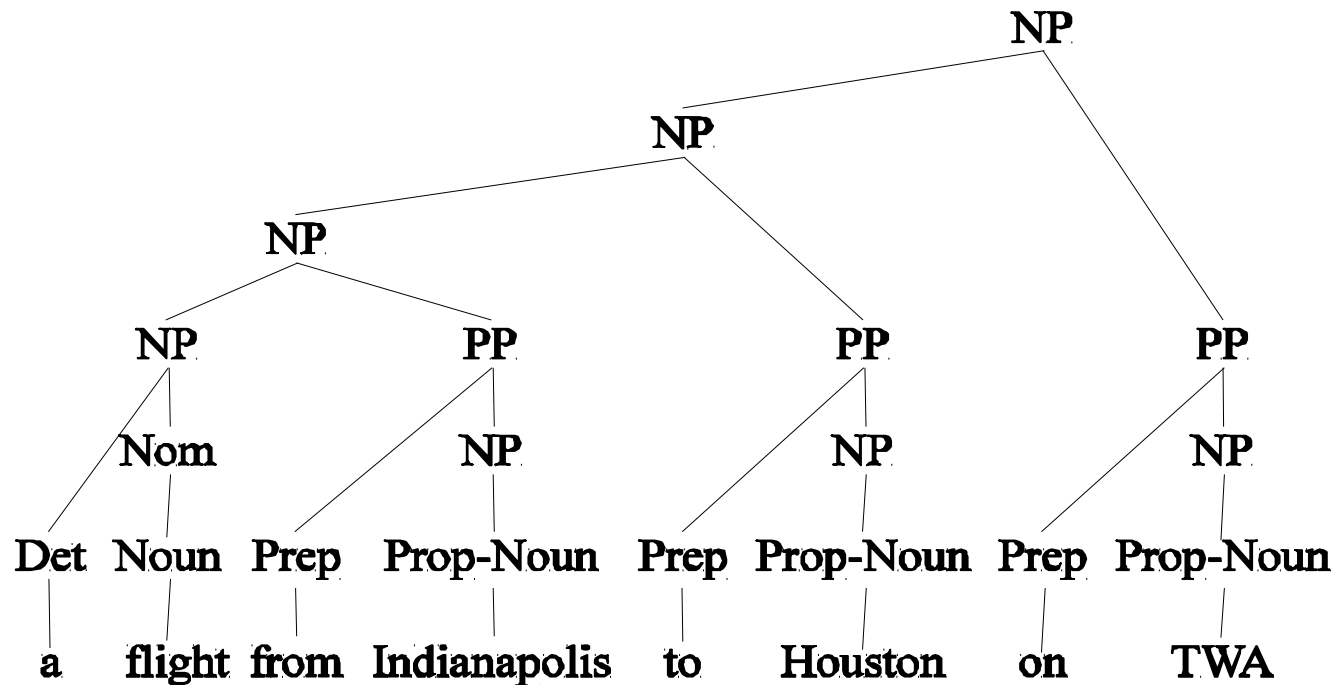
a flight	4
From Indianapolis	3
To Houston	2
On TWA	1
A flight from Indianapolis	3
A flight from Indianapolis to Houston	2
A flight from Indianapolis to Houston on TWA	1

Problem: Repeated Parsing Subtrees

reduplication of work in backtracking approach



Problem: Repeated Parsing Subtrees



reduplication of work in backtracking approach

Dynamic Programming

- ♦ Dynamic programming provides framework for solving the three kinds of problems afflicting top-down or bottom-up parsers
- ♦ Dynamic programming approaches systematically fill in tables of solutions to sub-problems
- ♦ When complete, the tables contain the solution to all the sub-problems needed to solve the problem
- ♦ Using tables to store the sub-trees for each constituents in the input solves reparsing and the ambiguity problem
- ♦ Also solves left-recursion problem

Dynamic Programming

- ♦ Three well-known dynamic parsers:
 - Cocke-Younger-Kasami (CYK) algorithm
 - Graham-Harrison-Ruzzo (GHR) algorithm
 - Earley algorithm

References

- ◆ Speech and Language Processing, *Jurafsky and H.Martin*
[Chapter 10. Parsing with Context-Free Grammars]

Thank You