



11 February, 2021

Continuing previous example

$$\min z = \begin{cases} 31c_2 & \frac{c_1}{c_2} \geq \frac{5}{2} \quad (S_2 \leq S_1) \\ 10c_1 + 6c_2 & \text{if } \frac{3}{5} \leq \frac{c_1}{c_2} \leq \frac{5}{2} \quad (S_1 \leq S_2 \leq S_2) \\ 20c_1 & \text{if } \frac{c_1}{c_2} \leq \frac{3}{5} \quad (S_2 > S_2) \end{cases}$$

Now, Let the total requirement also changes (i.e. A Req. = 124 & B. Req. = 60)

$$3n_1 + 5n_2 \geq b_1$$

$$10n_1 + 4n_2 \geq b_2$$

$$3n_1 + 5n_2 \geq b_2$$

↳ Dual for the problem:

$$\max z = b_1 y_1 + b_2 y_2$$

$$\text{Subject to } 10y_1 + 3y_2 \leq 16$$

$$4y_1 + 5y_2 \leq 14$$

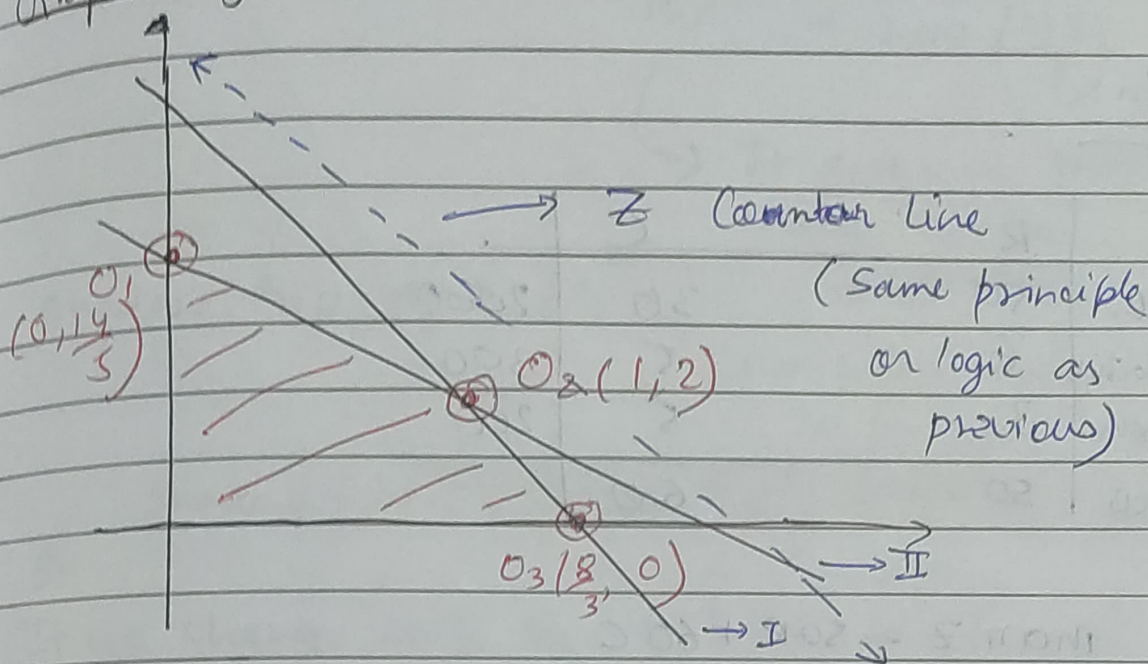
$$y_1, y_2 \geq 0$$



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Graphically



$$S_1 = -\frac{10}{3}, \quad S_2 = -\frac{4}{5}, \quad S_3 = -\frac{b_1}{b_2}$$

$$\text{If } S_2 \leq S_3 \Rightarrow O_1$$

$$S_1 \leq S_2 \leq S_3 \Rightarrow O_2$$

$$S_2 \leq S_1 \Rightarrow O_3$$

min value =  
(Dual)

$$\begin{cases} \frac{14}{5} b_2 & \text{if } \frac{b_1}{b_2} \leq \frac{4}{5} \\ b_1 + 2b_2 & \text{if } \frac{4}{5} < \frac{b_1}{b_2} \leq \frac{10}{3} \\ \frac{8}{3} b_1 & \text{if } \frac{b_1}{b_2} > \frac{10}{3} \end{cases}$$

→ min of primal problem



Shadow Price : Sol<sup>n</sup> of the dualEx:

	R	C	
Al.	50	30	2000
Machine	6	5	300
Labour	3	5	200
Price	50	60	

$$\text{max } Z = 50R + 60C$$

subject to

$$50R + 30C \leq 2000$$

$$6R + 5C \leq 300$$

$$3R + 5C \leq 200$$

$$R, C \geq 0$$

$$\hookrightarrow Z_{\text{max}} = 2750 \text{ at } (25, 25)$$

Dual:

$$\text{min } Z = 2000y_1 + 300y_2 + 200y_3$$

subject to

$$50y_1 + 6y_2 + 3y_3 \geq 50$$

$$30y_1 + 5y_2 + 5y_3 \geq 60$$

$$y_i \geq 0$$

$$\hookrightarrow Z_{\text{min}} \text{ at } = 2750$$

$$\text{at } \left( \frac{7}{16}, 0, \frac{75}{8} \right)$$

We can use Shadow price only to a point where our optimality is not gone.



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$$Z_{\max} = Z_{\min} = 2000 \left( \frac{7}{16} \right) + 300(0) + 200 \left( \frac{75}{8} \right)$$

It gives us the change in price difference.

Shadow Price of Dual

(No Information in primal)

$$\text{Shadow Price} = \left( \frac{7}{16}, 0, \frac{75}{8} \right) = Y^*$$

Sh

If we change 2000 to 2001 unit in primal problem, then for using one extra unit, the price difference will be given by  $\frac{7}{16}$ .

Dual sol<sup>n</sup>  $\rightarrow (y_1, y_2, y_3)$  is price &  $Y^*$  in the primal problem it is the shadow price.

We can't see these price in primal but

when we solved that in dual, we get the unit through which the change in  $Z_{fn}$  can be calculated.

$\therefore$  Shadow Price

$\frac{7}{16}$  tells us the difference in price in the original problem.

$Z_{\min} = Z_{\max} \rightarrow \frac{7}{16}$  is the price of resources in primal.

Resource limit or resources changed in primal  $\rightarrow$  Price diff.  $\frac{7}{16}$ .

11kg Labour  $\rightarrow \frac{75}{8}$ , Machine Time  $\rightarrow 0$





$$\text{Shadow Price} = \left( \frac{7}{16}, 0, \frac{75}{8} \right) = y^*$$

Doesn't Affect Any change ~~in~~

Why 0? Because in sol<sup>n</sup>, the slack variable is Basic (x<sub>4</sub>)

$$\begin{array}{rcl} 50R + 30C & \leq 2000 & \rightarrow = \\ 6R + 5C & \leq 300 & \rightarrow < \\ 3R + 5C & \leq 200 & \rightarrow = \end{array}$$

Because the slack is not in basic

Here the slack is in Basic sol<sup>n</sup> (~~Not~~)

∴ Strictly less than

It means that when manufacturing boats or products, our ~~laborer~~ Machine is underused. (using less than given value)

∴ Extra units is there

∴ If we change 300, no affect as we are using only around 275.

So, we can increase 300 and it doesn't affect the model.

Shadow Price when zero → Doesn't affect Model, can change without affecting model.



# Matrix Form of Simplex Algorithm

$$\min Z + Z_0 = \underbrace{(C \quad n)}_{\substack{\text{column} \\ \text{vector}}} \rightarrow \begin{matrix} \text{row} \\ \text{vector} \end{matrix} \quad \therefore \text{a number}$$

Subject to

$$Ax = b, \quad x \geq 0$$

$$\left[ \begin{array}{c|c} A & b \\ \hline C & Z_0 \end{array} \right]$$

$$\begin{array}{c|c} A & b \\ \hline C & Z_0 \end{array} \xrightarrow{\quad} \begin{array}{c|c} A^* & b^* \\ \hline C^* & Z_0^* \end{array}$$

(Initial)  (Final)

We find relations b/w initial & final tables, e.g. If we change something in  $C$ , then how that change will be reflected in  $C^*$ .

We need not do the full pivoting, we will just relations b/w initial & final.

Let for  $m$  constraints, there be  $n_{j1}, n_{j2}, \dots, n_{jm}$  Basic Variables.

Let  $A(j) = j^{\text{th}}$  col of  $A$   
 $A^{j_1} \dots A^{j_m} \rightarrow \text{Basic Variable Columns} = B$



$$\therefore B = \left[ A^{j_1} \mid A^{j_2} \mid \dots \mid A^{j_m} \right]$$

$\hookrightarrow$  Columns of Basic

$$C_B = [C_{j_1} \dots C_{j_m}] \rightarrow C \text{ values corresponding to Basic variables}$$

### Results

$$\left. \begin{array}{l} 1) A^* = B^{-1} A \\ 2) B^* = B^{-1} b \\ 3) C^* = C - C_B B^{-1} A = C - C_B A^* \\ 4) Z_o^* = Z_o - C_B B^{-1} b = Z_o - (C_B B^*) \end{array} \right\}$$