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Examination Roll No. - 18312 9 11011

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Title of Paper - Computer Language, design

and Engineering

3.5

First & hollow table for Crammar:

 $E \rightarrow TE'$ $E' \rightarrow +TE' \mid E$ $T \rightarrow FT'$ $T' \rightarrow *FT' \mid E$ $F \rightarrow ial(E)$

first & Follow cuts:

	First	Follow
E → TE'	٤id, (3	{\$,)}
E'->+TE'/E	名+, 色子	を参り3
T→FT'	{id, C}	を +, \$,) 子
T' -> *FT'/E	{*, 2}	着+,ま,つま
F -> id/(E)	{id, €}	至*,+,\$,)3

The LL(1) Parsing Table is;

E	id E→TE'	+	×	(E→TE')	\$
E'		E'->+TE'			E'>E	Ε'→ε
	T→FT'			T→FT'		
T'		T'>E	T'->*FT		T'->E	T'->E
	F->id			F->(E)		

As you can see that all the null productives are put under the feelen set of that symbol and all the remaining productions are his under the first of that symbol.

Every grammar is not frasible for LL(1) Parsing table. It may be possible that one cell may contain more than one production.