#### CS 186 Discussion 10

Recovery

# Logistics

Homework 5 due 11/24

HKN Course Surveys 11/30

• Midterm 3 on 12/02

### Recovery

- Logging enables atomicity and durability
- Works tightly with buffer pool and lock manager
- Strict 2PL

- Steal okay to evict dirty pages
  - Watch out for:
  - Log Info:
- No-Force can delay disk writes
  - Watch out for:
  - Log Info:

- Steal okay to evict dirty pages
  - Watch out for: Atomicity (xact doesn't commit)
  - Log Info: UNDO
- No-Force can delay disk writes
  - Watch out for:
  - Log Info:

- Steal okay to evict dirty pages
  - Watch out for: Atomicity (xact doesn't commit)
  - Log Info: UNDO
- No-Force can delay disk writes
  - Watch out for: Durability (commit lost to crash)
  - Log Info: REDO

### Homework 5 Demo

- Steal okay to evict dirty pages
- No-Force can delay disk writes
- So when are:
  - Updates persisted to disk?
  - Update records written?

- Steal okay to evict dirty pages
- No-Force can delay disk writes
- So when are:
  - Updates persisted to disk? on eviction (or earlier)
  - Update records written?
    - In the homework: before commit or eviction
    - In ARIES: whenever you update

# ARIES Logging

- Log Sequence Numbers (increasing "timestamps")
- pageLSN each page, most recent update LSN
- Record Types:
  - Update, Commit, Abort
  - Checkpoint
  - CLR (an undo)
  - End (done committing/aborting)

# ARIES Logging

- Transaction Table
  - Tracks <u>active</u> xacts
  - lastLSN most recent LSN written by xact
- Dirty Page Table
  - Tracks dirty pages in buffer
  - recLSN first record that dirtied page
- flushedLSN max LSN flushed so far

#### ARIES - Commit

- 1. Write commit record
- 2. Flush log to disk
- 3. Write end record

#### ARIES - Abort

- 1. Write abort record
- 2. Rollback fix buffer pool page with before image
  - Write CLR records per undo, update pageLSN
  - Redo CLRs if you crash
- 3. Write end record

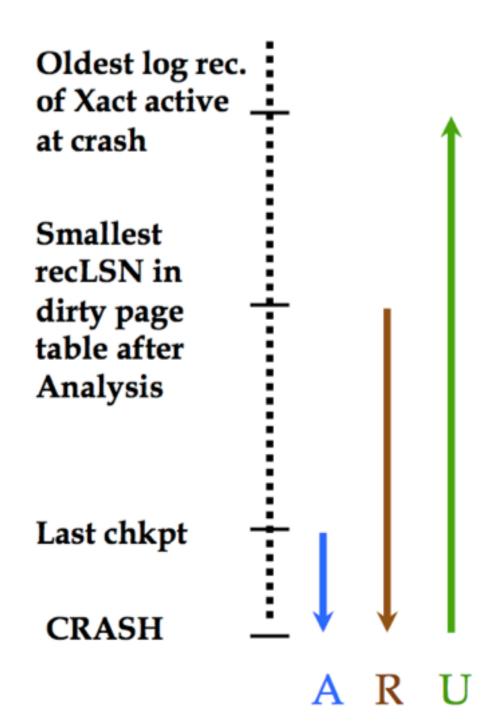
### ARIES - Checkpoint

- Faster recovery! Store most recent checkpoint
- Checkpoint record stores current xact table + DPT
- No need to force dirty pages to disk (recLSN)

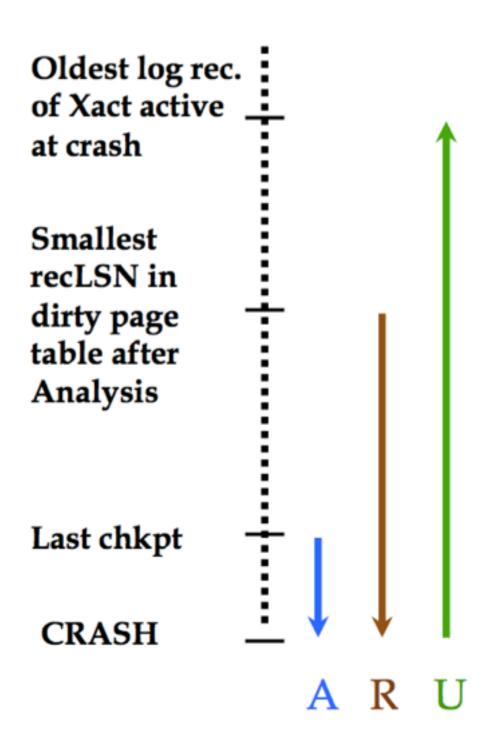
If we crash:

What do we need to restore?

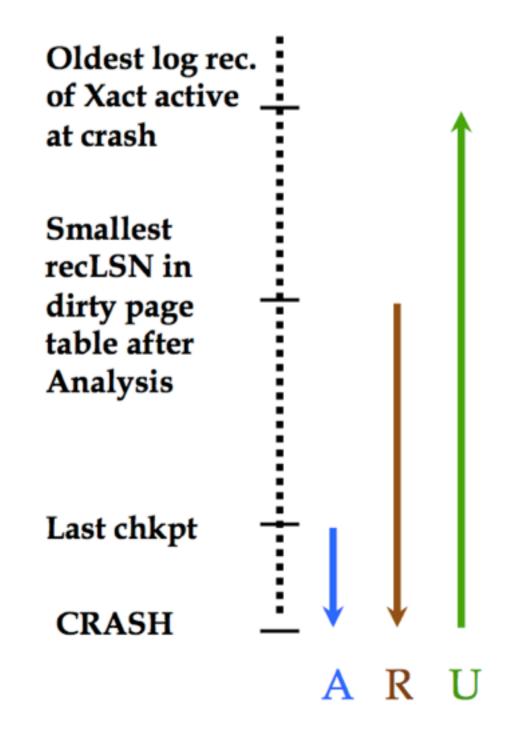
What do we need to undo?



- If we crash:
  - What do we need to restore?
    - Anything that committed
  - What do we need to undo?
    - Anything that aborted
    - or... didn't finish



- Analysis / REDO / UNDO
  - Need to know which xacts started/committed/aborted after checkpoint
  - xact table = loser xacts
    - Remove xact if END
    - Add xact if anything else
  - Add new pages to DPT for updates (+ recLSN)



LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	70	Running
T2	60	Running
T3	30	Running
T4	50	Running

PageID	recLSN
P5	50
P1	40

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
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Transaction	lastLSN	Status
T1	90	Running
T2	60	Running
T3	30	Running
T4	50	Running

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	90	Running
T2	110	Running
T3	30	Running
T4	50	Running

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	90	Running
T2	110	Committed
Т3	30	Running
T4	50	Running

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
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Transaction	lastLSN	Status
T1	90	Running
T2	110	Running
T3	30	Running
T4	130	Running

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
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180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	90	Running
T3	30	Running
T4	130	Running

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
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110	update: T2 writes P3
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130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	90	Running
Т3	30	Running
T4	130	Aborting

PageID	recLSN
P5	50
P1	40
P3	90

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

Transaction	lastLSN	Status
T1	90	Running
T5	160	Running
T3	30	Running
T4	130	Aborting

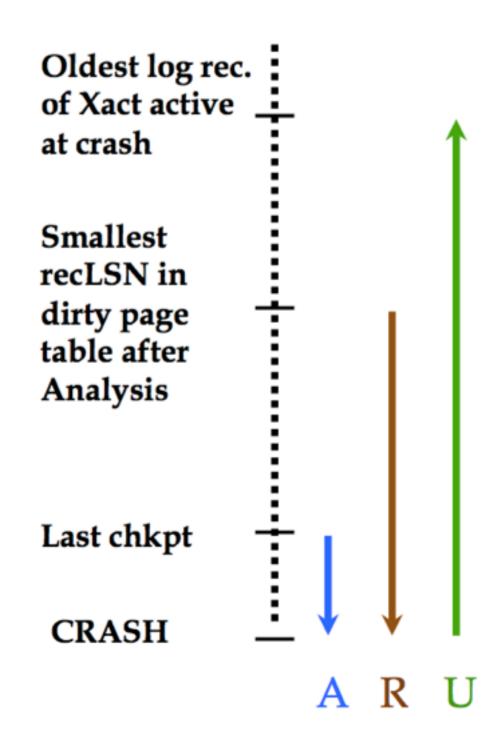
PageID	recLSN
P5	50
P1	40
P3	90
P2	160

LSN	Record
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
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160	update: T5 writes P2
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Transaction	lastLSN	Status
T1	90	Running
T5	160	Running
T3	30	Running
T4	180	Aborting

PageID	recLSN
P5	50
P1	40
P3	90
P2	160

- Analysis / REDO / UNDO
  - Redo <u>everything</u> from min recLSN, unless:
    - page not in DPT
    - recLSN of page > LSN
    - pageLSN ≥ LSN
  - Redo with after-image
  - Update pageLSNs



#### Worksheet - REDO

LSN	Record
40	update: T4 writes P1
50	update: T4 writes P5
60	update: T2 writes P5
70	update: T1 writes P2
80	begin checkpoint
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

#### Skip redo if:

- page not in DPT
- recLSN of page > LSN
- pageLSN ≥ LSN

PageID	recLSN
P5	50
P1	40
P3	90
P2	160

#### Worksheet - REDO

LSN	Record
40	update: T4 writes P1
50	update: T4 writes P5
60	update: T2 writes P5
70	update: T1 writes P2
80	begin checkpoint
90	update: T1 writes P3
100	end checkpoint
110	update: T2 writes P3
120	T2 commit
130	update: T4 writes P1
140	T2 end
150	T4 abort
160	update: T5 writes P2
180	CLR: undo T4 LSN 130

#### Skip redo if:

- page not in DPT
- recLSN of page > LSN
- pageLSN ≥ LSN

PageID	recLSN
P5	50
P1	40
P3	90
P2	160

#### Worksheet - 2ab

Transaction	lastLSN	Status

PageID	recLSN

#### Worksheet - 2ab

Transaction	lastLSN	Status
T3	70	Running
T2	80	Aborting

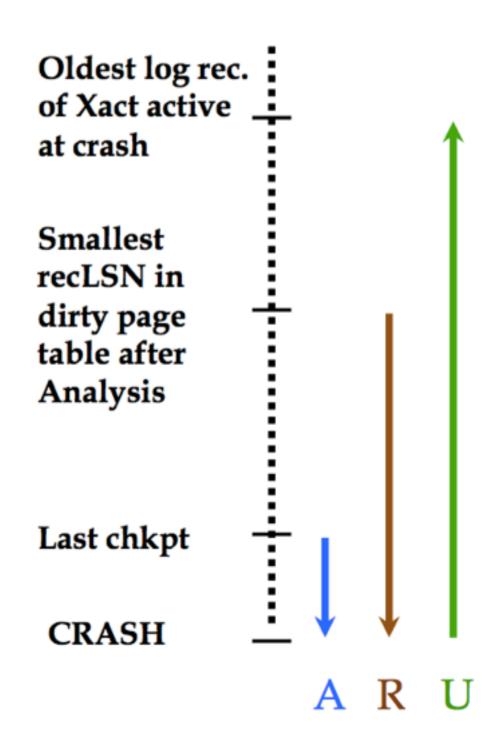
PageID	recLSN

#### Worksheet - 2ab

Transaction	lastLSN	Status
Т3	70	Running
T2	80	Aborting

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

- Analysis / REDO / UNDO
  - Undo <u>everything</u> in xact table -> abort these xacts
  - Performance optimization:
    - Repeat with decreasing LSN, until table empty
    - Undo updates and continue writing CLRs



Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

- Start from max lastLSN, work backwards
- ToUndo =  $\{70, 80\}$ 
  - prevLSN of 80?

Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

- Start from max lastLSN, work backwards
- ToUndo =  $\{50, 70\}$

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

- Start from max lastLSN, work backwards
- ToUndo =  $\{40, 50\}$

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

LSN	Record
100	CLR: T3 LSN = 70; undoNextLSN = 40

Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

- Start from max lastLSN, work backwards
- ToUndo =  $\{20, 40\}$

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

LSN	Record
100	CLR: T3 LSN = 70; undoNextLSN = 40
110	CLR: T2 LSN = 50; undoNextLSN = 20

Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

- Start from max lastLSN, work backwards
- ToUndo =  $\{20\}$

PageID	recLSN
P1	10
P2	70
P3	20
P4	40

LSN	Record	
100	CLR: T3 LSN = 70; undoNextLSN = 40	
110	CLR: T2 LSN = 50; undoNextLSN = 20	
120	CLR: T3 LSN = 40; undoNextLSN = null	
130	T3 end	

Xact	lastLSN	Status
Т3	70	Running
T2	80	Aborting

recLSN
10

70

20

40

P2

P3

P4

•	Start from	max	lastl	_SN,
	work bacl	kward	S	

LSN	Record
100	CLR: T3 LSN = 70; undoNextLSN = 40
110	CLR: T2 LSN = 50; undoNextLSN = 20
120	CLR: T3 LSN = 40; undoNextLSN = null
130	T3 end
140	CLR: T2 LSN = 20; undoNextLSN = null
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