Graph with BFS and DFS

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2019年10月25日 下午 2:28
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Common Methods:DFS, BFS, Union Fine, Topologic Sorting

核心是想办法把一个问题看成graph问题

★ 图的DFS遍历,用hashmap存每个节点的neighbo

★ 图的BFS遍历,用hashmap存每个节点的 neighbors,queue去遍历

```
class Solution {
public:
    Node* cloneGraph(Node* node) {
        unordered_map<Node*, Node*> m;
        return helper(node, m);
    }
    Node* helper(Node* node, unordered_map<Node*, Node*>& m) {
        if (!node) return NULL;
        if (m.count(node)) return m[node];
        Node* clone = new Node(node -> val);
        m[node] = clone;
        for (auto& neighbor: node->neighbors)
        {
            clone -> neighbors.push_back(helper(neighbor,m));
        }
        return clone;
    }
};
```

```
class Solution {
public:
    Node* cloneGraph (Node* node) {
        if (!node) return NULL;
        unordered map<Node*, Node*> m;
        queue<Node*> q{{node}};
        Node *clone = new Node(node->val);
        m[node] = clone;
        while (!q.empty()) {
            Node *t = q.front(); q.pop();
            for (Node *neighbor : t->neighbors) {
                if (!m.count(neighbor)) {
                    m[neighbor] = new Node(neighbor->val);
                    q.push(neighbor);
                m[t]->neighbors.push back(m[neighbor]);
            }
        }
        return clone;
    }
```

Leetcode <mark>200 islands or 695</mark> 以及 547
Friend Circles 一模一样的题目,联通分量,DFS
只不过双重for循环,一个for循环放在dfs里面
visited 一维数组表示friends circles

18 行不是多余的,当recursive调用的时候很有用 <mark>用visited,</mark>表示这个人的所有朋友都访问过,所以不用再访问;

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- ★ 能返回int值就返回int值的比较快,不用传入值进去;
- ★ 上下进行比较 void dfs(max...) or int dfs()

```
Solution {
                 int maxAreaOfIsland(vector<vector<int>>& grid) {
   if (grid.empty()) return 0;
                        int max = 0;
int m = grid.size();
int n = grid[0].size();
for(int i = 0; i < m; ++i)</pre>
 10
                                   r(int j = 0; j < n; ++j)
                                      int cur = 0;
if(grid[i][j]) dfs(grid,cur,i,j,m,n);
max = cur>max?cur:max;
 13
14
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18
                         return max;
 19
 20
21 v
                 void dfs(vector<vector<int>>& grid,int& cur, int r, int c, int m, int n)
                        if(r < 0 || c < 0 || r >= m || c >=n) return;
if(!grid[r][c]) return;
grid[r][c] = 0;
curit
 24
25
                       dfs(grid,cur,r,c-1,m,n);
dfs(grid,cur,r,c+1,m,n);
dfs(grid,cur,r,c+1,m,n);
 28
 1 -
                 int maxAreaOfIsland(vector<vector<int>>& grid) {
                        int h = grid.size();
if (h == 0) return 0;
 4
                        int w = grid[0].size();
                        int ans = 0:
                               (int i = 0; i < h; ++i)
for (int j = 0; j < w; ++j)
    ans = max(ans, area(grid, j, i, w, h));</pre>
10
               }
te:
14
                int area(vector<vector<int>>& grid, int x, int y, int w, int h) { if (x < 0 \mid | y < 0 \mid | x >= w \mid | y >= h \mid | grid[y][x] == 0) rei
15 🔻
16
18
                        grid[y][x] = 0;
                        return area(grid, x - 1, y, w, h)
+ area(grid, x + 1, y, w, h)
+ area(grid, x, y - 1, w, h)
+ area(grid, x, y + 1, w, h)
                                  + 1:
         };
                                 PC
                                                                                                             PC
                   Union(2, 4), Find(2) = 3
                                                                                          Union(5, 2) = Union(Find(5), Find(2))
          Ref: <a href="https://www.cs.princeton.edu/~rs/AlgsDS07/01UnionFind.pdf">https://www.cs.princeton.edu/~rs/AlgsDS07/01UnionFind.pdf</a>
                                                                                                                        PC = path compression
```

□★ 遇到图想着用 adjacency lists 来表示 map<int, vector<int>> graph --> graph.at(边 1);

Union Find!! 查找cycles

https://zxi.mytechroad.com/blog/datastructure/sp1-union-find-set/

```
class UnionFindSet {
public:
    UnionFindSet(int n) {
    ranks_ = vector<int>(n + 1, 0);
    parents_ = vector<int>(n + 1, 0);

    for (int i = 0; i < parents_.size(); ++i)
        parents_[i] = i;
}

// Merge sets that contains u and v.
// Return true if merged, false if u and v are already in one set.
bool Union(int u, int v) {</pre>
```

Leetcode 684 & 685 redundant connection 1 and 2

Leetcode

Leetcode 990 Statifiablitity of equality equations

more concise Union Find

```
class Solution {
public:
  bool equationsPossible(vector<string>& equations) {
    iota(begin(parents_), end(parents_), 0);
    for (const auto& eq : equations)
        if (eq[1] = '=')
            parents_[find(eq[0])] = find(eq[3]);
        for (const auto& eq : equations)
        if (eq[1] == '!' && find(eq[0]) == find(eq[3]))
            return false;
        return true;
    }
    private:
    array<int, 128> parents_;
    int find(int x) {
        if (x != parents_[x])
            parents_[x] = find(parents_[x]);
        return parents_[x];
    }
};
```

iota 用顺序递增的值存储在指定范围内

0,, 126,127,128

```
int pu = Find(u);
   int pv = Find(v);
   if (pu == pv) return false;
   // Meger low rank tree into high rank tree
   if (ranks_[pv] > ranks_[pu])
     parents_[pu] = pv;
   else if (ranks_[pu] > ranks_[pv])
     parents_[pv] = pu;
   else {
     parents_[pv] = pu;
     ranks_[pv] += 1;
   return true;
 }
 // Get the root of u.
 int Find(int u) {
   // Compress the path during traversal
   if (u != parents_[u])
     parents_[u] = Find(parents_[u]);
   return parents_[u];
private:
 vector<int> parents_;
 vector<int> ranks_;
};
class Solution {
public:
 vector<int> findRedundantConnection(vector<vector<int>>& edges) {
   UnionFindSet s(edges.size());
   for(const auto& edge: edges)
     if (!s.Union(edge[0], edge[1]))
      return edge;
   return {};
 }
};
来自 < https://zxi.mytechroad.com/blog/tree/leetcode-684-redundant-connection/>
```

Templates

```
class UnionFindSet {
 ublic:
  UnionFindSet(int n) {
   ranks_ = vector<int>(n + 1, 0);
     parents_ = vector<int>(n + 1, 0);
     for (int i = 0; i < parents_.size(); ++i)</pre>
     parents_[i] = i;
  }
  // Merge sets that contains u and v.
 // Return true if merged, false if u and v are already in one set.
  bool Union(int u, int v) {
   int pu = Find(u);
    int pv = Find(v);
   if (pu == pv) return false;
   // Meger low rank tree into high rank tree
     if (ranks_[pv] < ranks_[pu])</pre>
     parents_[pv] = pu;
     else if (ranks_[pu] < ranks_[pv])</pre>
     parents_[pu] = pv;
```

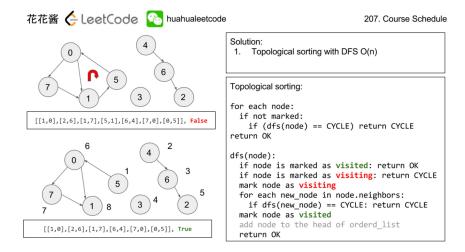
```
else {
    parents_[pv] = pu;
       ranks_[pu] += 1;
   return true;
  }
  // Get the root of u.
 int Find(int u) {
    // Compress the path during traversal
   if (u != parents_[u])
       parents_[u] = Find(parents_[u]);
   return parents_[u];
  }
private:
  vector<int> parents_;
 vector<int> ranks_;
};
```

来自 < https://zxi.mytechroad.com/blog/data-structure/sp1-union-find-set/>

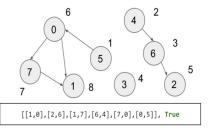
Leetcode course schedule 207

https://zxi.mytechroad.com/blog/graph/leetcode-207-course-schedule/

there are cycles, there is impossible





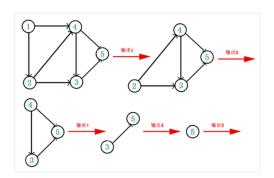


n	visiting	visited	ordered list
0	{0}	{}	{}
1	{0,1}	{}	{}
0	{0}	{1}	{1}
7	{0,7}	{1}	{1}
0	{0}	{1,7}	{7,1}
-	{}	{1,7,0}	{0,7,1}
2	{2}	{1,7,0}	{0,7,1}
-	{}	{1,7,0,2}	{2,0,7,1}
3	{3}	{1,7,0,2}	{2,0,7,1}
-	{}	{1,7,0,2,3}	{3,2,0,7,1}
4	{4}	{1,7,0,2,3}	{3,2,0,7,1}
6	{4,6}	{1,7,0,2,3}	{3,2,0,7,1}
4	{4}	{1,7,0,2,3,6}	{6,3,2,0,7,1}
-	{}	{1,7,0,2,3,6,4}	{4,6,3,2,0,7,1}
5	{5}	{1,7,0,2,3,6,4}	{4,6,3,2,0,7,1}
-	{}	{1,7,0,2,3,6,4,5}	{5,4,6,3,2,0,7,1}

topological Sorting Templates 1 -> 0 一直寻找入度为0 的节点 indegree = 0

它是一个 DAG 图,那么如何写出它的拓扑排序呢?这里说一种比较常用的方法:

- 1. 从 DAG 图中选择一个 没有前驱(即入度为0)的顶点并输出。
- 2. 从图中删除该顶点和所有以它为起点的有向边。
- 3.重复1和2直到当前的DAG图为空或当前图中不存在无前驱的顶点为止。后一种情况说明有向图中必然存在环。



★ DFS topologic Sorting

这个题目是完成事件1需要先完成事件 2,也就是走1 之前需要2,那么转化 成图,用topological sorting

```
bool canFinish(int n, vector<pair<int, int>>& pre)
    vector<vector<int>> adj(n, vector<int>());
    vector<int> degree(n, 0);
    for (auto &p: pre) {
        adj[p.second].push_back(p.first);
        degree[p.first]++;
    }
    queue<int> q;
    for (int i = 0; i < n; i++)
        if (degree[i] == 0) q.push(i);
    while (!q.empty()) {
        int curr = q.front(); q.pop(); n--;
        for (auto next: adj[curr])
            if (--degree[next] == 0) q.push(next);
    return n == 0;
}
```

★ 用<mark>双重map</mark>可以达到二维数组访问的效果,只不过index可以是字母!!!! unordered_map<string,unordered_map<string,double>> g;