solving method : dynamic programming

Let T[i] be the length of a longest subsequence

   starting at c[i].

initialization: T[i]=1 for i=1 to n

Optimal Substructure:

T[i]= 1 + max {T[j]: c[i] and c[j] have a match and i<j }

Optimal solution= max {T[i]}.

Complexity: