# Optimal parallelization strategy on MiniMax search tree based on Kalah game

基於kalah之最大最小搜尋樹優化平行策略

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#### Outline

- Introduction- Kalah
- Minimax tree
- Pruned Minimax tree
- New Proposed Method: Advanced parallel method on game tree
- Result and Evaluation
- Conclusion

#### Introduction -Kalah

House score(計分洞)

- ●Kalah、Kalaha(播棋)
  - Assume Beans = 7
  - Assume Bean holes = 6
- 6 options on each round

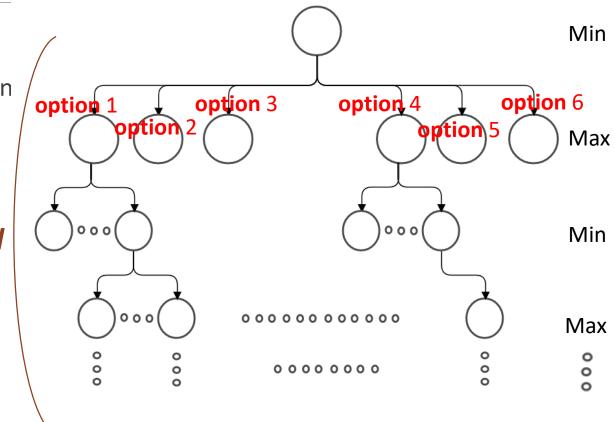


Player1

bean holes (opt. = 6)

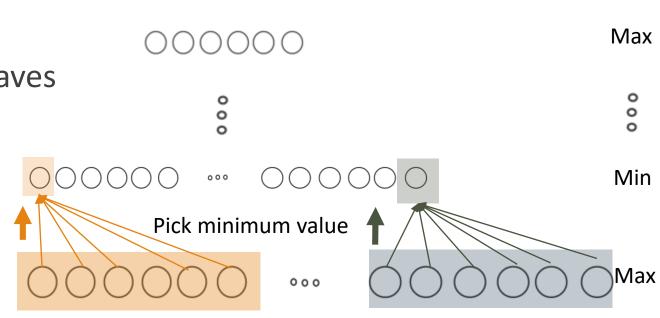
#### Minimax tree

- Concept:
  - Best option to win opponent's best step
  - Expand further d steps, and predict best option
- •complexity:  $O(6^d)$ 
  - Maxmum game length (d)
  - As far as we know, length can be more than 50
- Minimax prediction generally use recursion



## Experiment- (1)

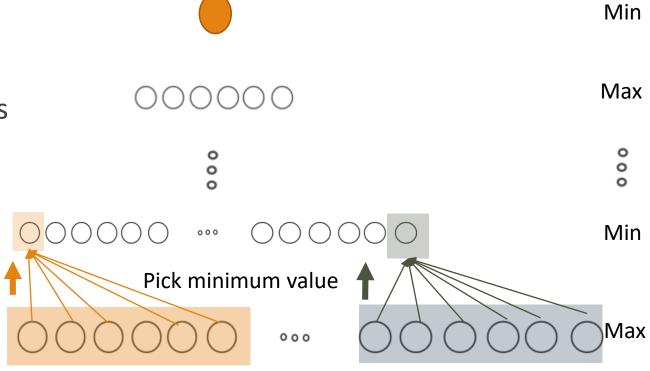
- Time overhead:
  - Reduction
  - Traverse
- assume d =10, yield 362,797,056 leaves
  - Need  $6^{10} + \cdots + 6^{1}$  comparison
  - How much parallelization can improve?



Min

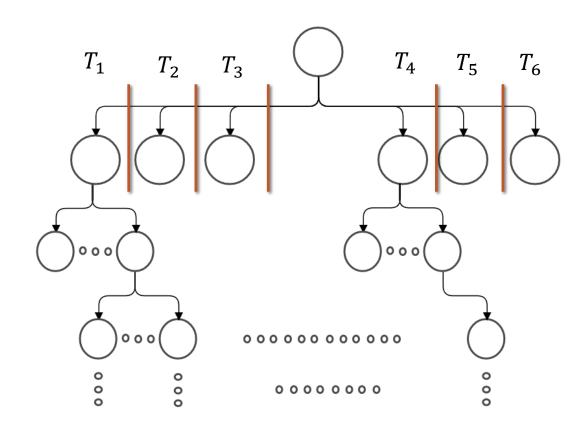
# Experiment- (1)

- Time overhead:
  - Reduction
  - Traverse
- assume d =10, yield 362,797,056 leaves Need  $6^{10} + \cdots + 6^{1}$  comparison
- Serial Reduction
  - $Time_{reduction}$  < 1% of Total time
  - Most of time spend on traverse

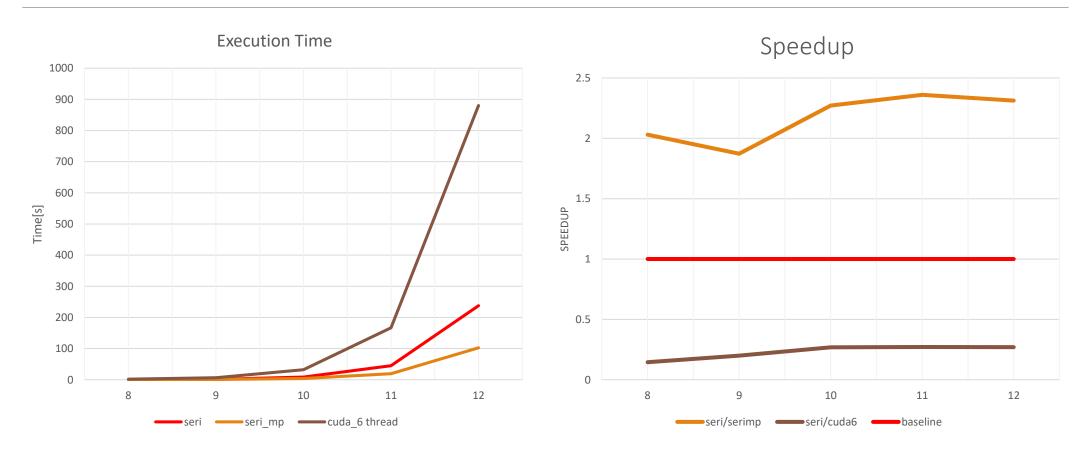


## Minimax tree - traverse parallelization

- 6 threads parallel
  - CUDA, openMP respectively



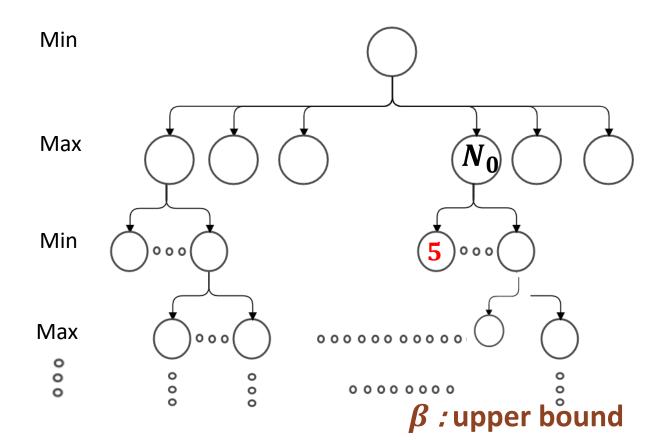
# Seri & Seri\_Mp



Complexity improvement:

 $\bullet O(6^d) \rightarrow O(6^{d/2})$ 

Step1: got min value 5



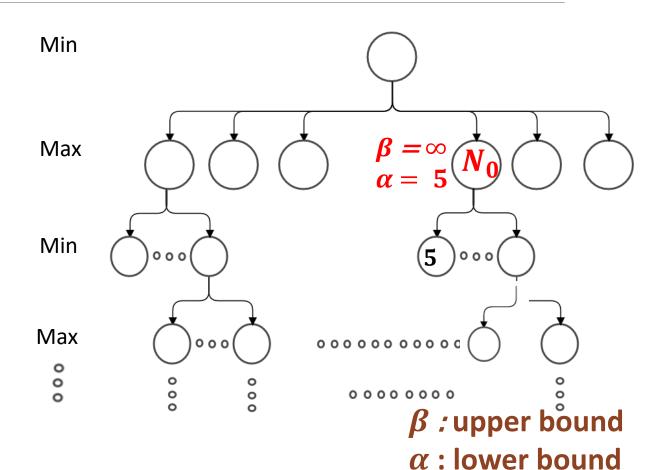
 $\alpha$  : lower bound

Complexity improvement:

 $\circ O(6^d) \to O(6^{d/2})$ 

Step1: got min value 5

Step2:  $N_0$  value at least bigger than 5



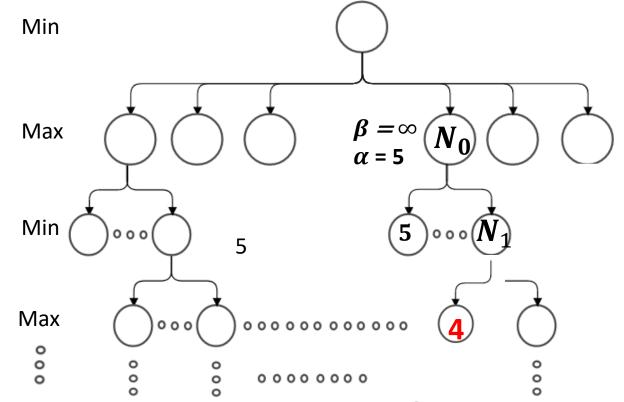
Complexity improvement:

 $0(6^d) \to 0(6^{d/2})$ 

Step1: got min value 5

Step2:  $N_0$  value at least bigger than 5

Step3: Max value 4 of  $N_1$  child



β: upper bound

 $\alpha$ : lower bound

Complexity improvement:

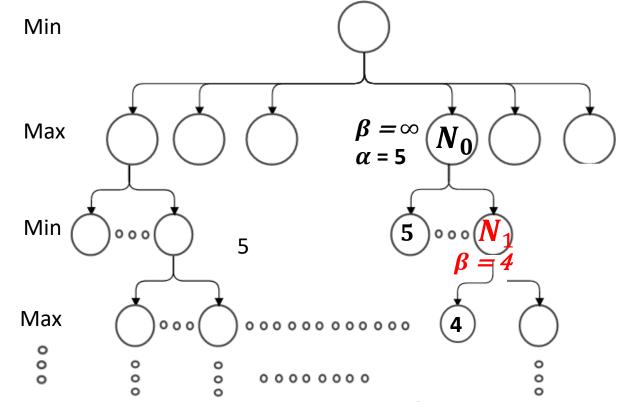
 $\circ$  O(6<sup>d</sup>)  $\to$  O(6<sup>d/2</sup>)

Step1: got min value 5

Step2:  $N_0$  value at least bigger than 5

Step3: Max value 4 of  $N_1$  child

Step4: N<sub>1</sub> value at most 5



β: upper bound

 $\alpha$ : lower bound

Complexity improvement:

 $\circ O(6^d) \to O(6^{d/2})$ 

Step1: got min value 5

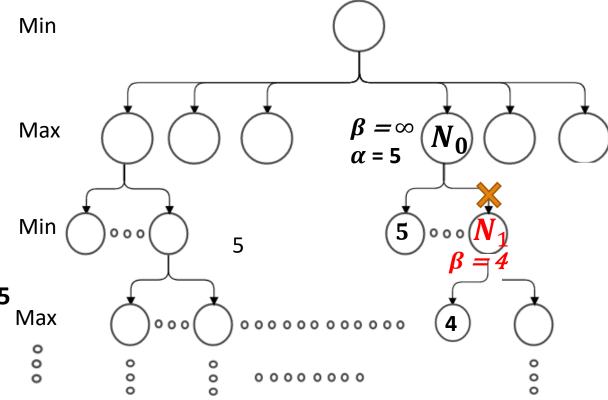
Step2:  $N_0$  value at least bigger than 5

Step3: Max value 4 of  $N_1$  child

Step4:  $N_1$  value at most 5

Step5:  $N_1$  would never find number better than 5

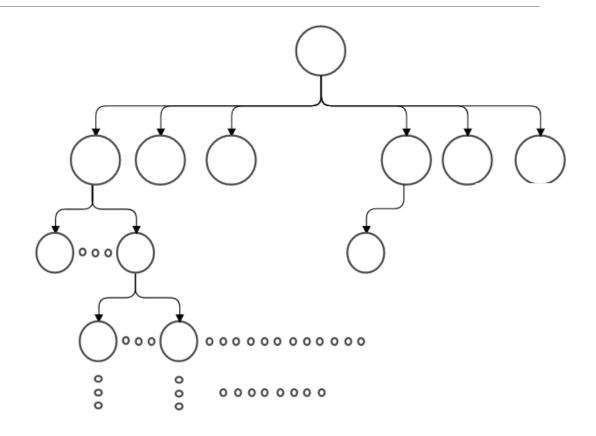
-> prune **N**<sub>1</sub>



 $\beta$ : upper bound  $\alpha$ : lower bound

- Complexity improvement:
  - $\circ$  O(6<sup>d</sup>)  $\to$  O(6<sup>d/2</sup>)

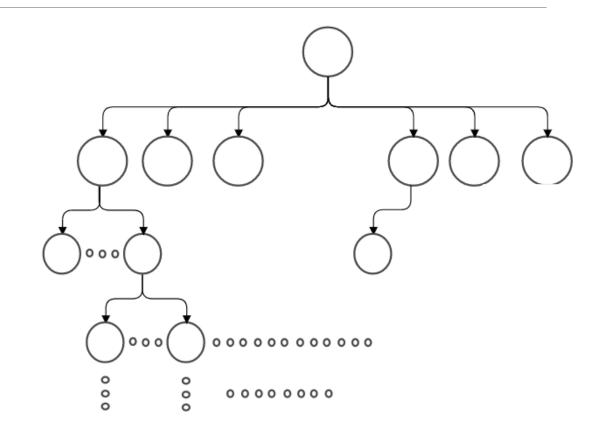
Get a pruned tree



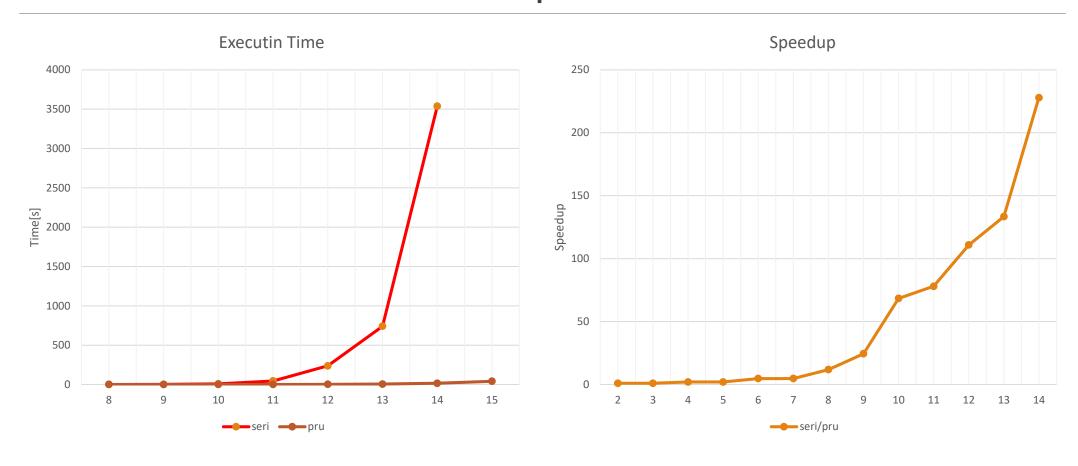
#### Complexity improvement:

$$\bullet O(6^d) \to O(6^{d/2})$$

Predict layer	original	pruned	save
1	6	6	0%
2	36	11	69%
3	216	42	81%
4	1,296	123	91%
5	7,776	507	93%
6	46,656	1,622	96%
7	279,936	6,348	98%
8	1,679,616	15,400	99%
9	10,077,696	79,274	99%

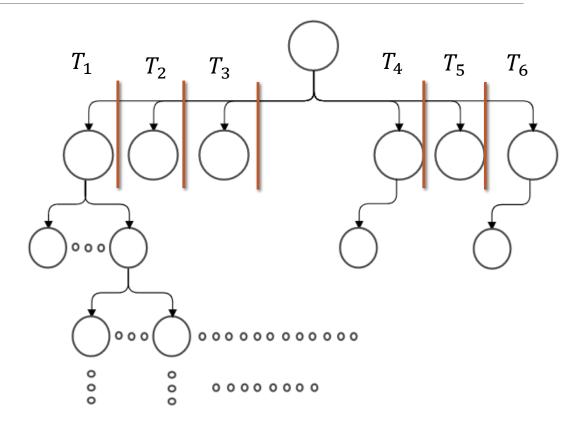


# Min-max &min-max prune

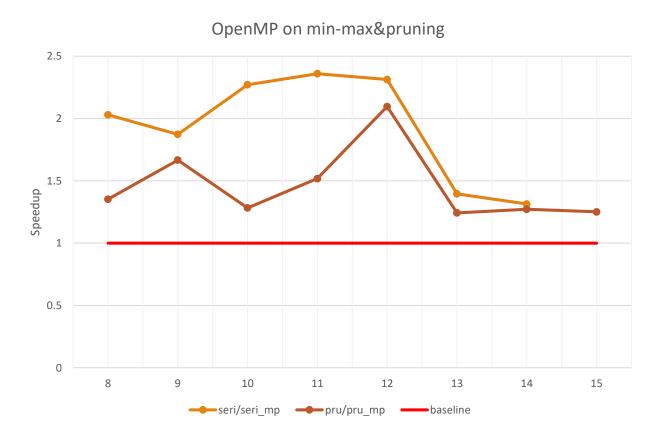


#### Pruned Minimax tree – Parallelization

- Hard to parallelize recursive algorithm
  - Thread number: 6



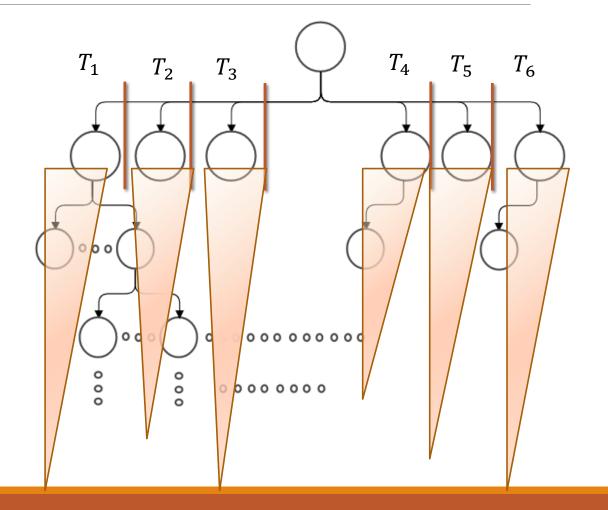
# Min-max\_mp &prune\_mp speedup



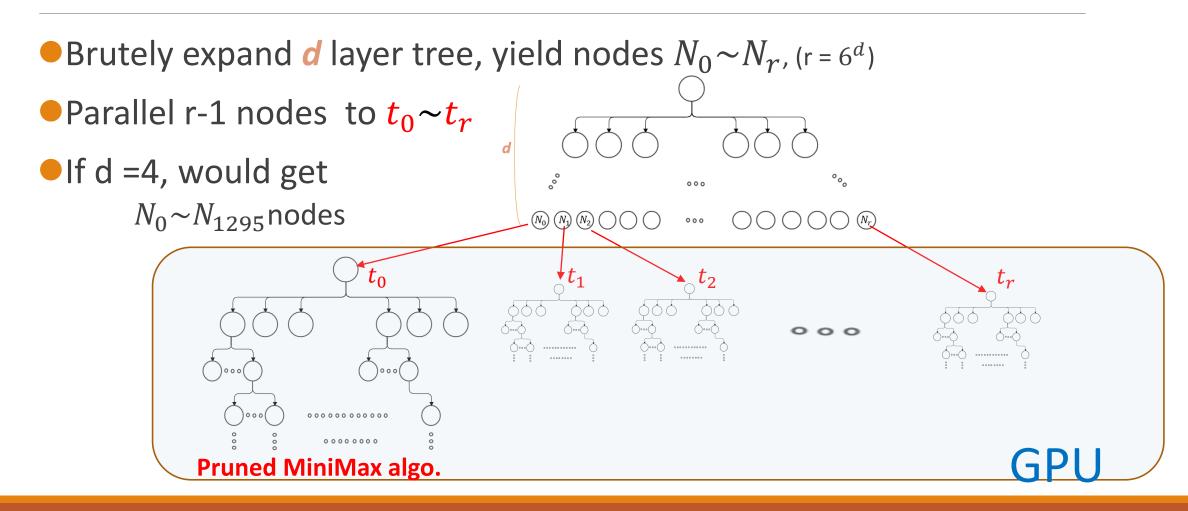
#### Pruned Minimax tree – Parallelization

- Hard to parallelize recursive algorithm
  - Thread number: 6
- Still face with load imbalance.
  - Note: htop



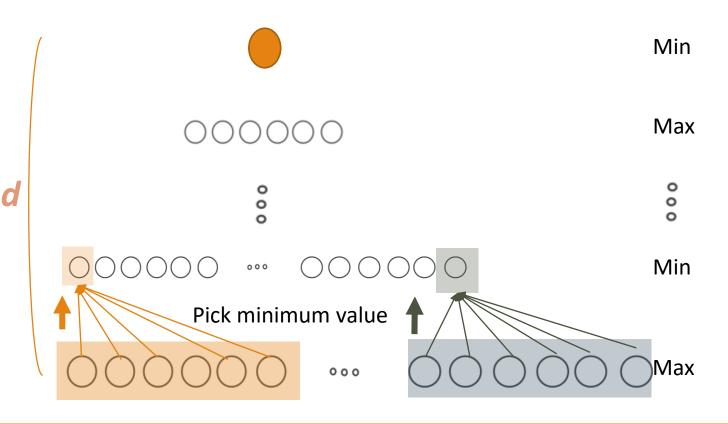


# Pruned Minimax tree parallel method with GPU

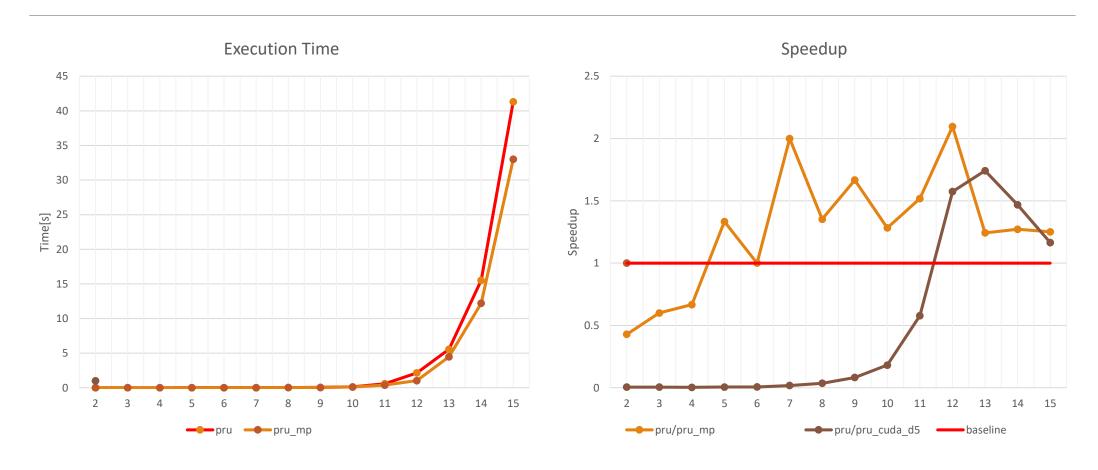


## Proposed Method: Advanced Pruned Minimax tree parallel method with GPU

- Brutely expand d layer tree, yield nodes  $N_0 \sim N_r$ , (r = 6<sup>d</sup>)
- ullet Parallel r-1 nodes to  $t_0 \sim t_r$
- Collect predictions and get best option.



# Pruned Minimax tree parallel method with GPU



#### Conclusion

 When we are design a parallel algorithm. The issues firstly come to our mind are synchronization problem and dependency problem.
However, picking a good algorithm to parallelize is also a big deal.
e.q. recursion.