# MAT-INF4130 MAT-INF4130

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## 1 Introduction

Given a set of points  $\{x_1, \ldots, x_n\} \subseteq \mathbb{R}^m$  and an integer  $k \leq m$ , we are interested in finding the k-dimensional subspace W that minimizes the squared distance:

$$\sum_{i=1}^{m} \operatorname{dist}(\boldsymbol{x}_i - W)^2. \tag{1}$$

To this end we first establish some notation. Let  $X := [x_1, \ldots, x_n]$  be the matrix with ijth element denoted  $x_{ij}$ . If W is any k-dimensional subspace of  $\mathbb{R}^m$  and  $\{w_j\}_{j=1}^k$  is an orthonormal basis for W, we may extend this basis to an orthonormal basis  $\{w_j\}_{j=1}^m$  for  $\mathbb{R}^m$  by the Grahm-Schmidt algorithm. We can therefore decompose the space  $\mathbb{R}^m$  into W and its orthogonal complement  $W^{\perp}$ :

$$\mathbb{R}^m = W \oplus W^{\perp}. \tag{2}$$

For a vector  $\boldsymbol{x} \in \mathbb{R}^m$ , denote by  $\operatorname{proj}_{W^{\perp}}(\boldsymbol{x})$ , its orthogonal projection onto  $W^{\perp}$ . Furthermore, recall that the Frobenius norm of a matrix  $\boldsymbol{A} \in \mathbb{C}^{m \times n}$  is defined as:

$$\|\mathbf{A}\|_{\mathrm{F}} \coloneqq \left(\sum_{i=1}^{m} \sum_{j=1}^{n} |a_{ij}|^2\right)^{1/2}$$
 (3)

and that this is equivalent to the 2-norm of the vector  $\text{Vec}(\boldsymbol{A})$  formed by stacking the columns of  $\boldsymbol{A}$  on top of each other. We denote by  $\langle \cdot, \cdot \rangle$  the standard inner product on  $\mathbb{R}^m$  defined by  $\langle \boldsymbol{x}, \boldsymbol{y} \rangle \coloneqq \boldsymbol{y}^T \boldsymbol{x}$ .

## Exercise 1

We start by showing that

$$\sum_{i=1}^{n} \left\| \operatorname{proj}_{W^{\perp}}(\boldsymbol{x}_{i}) \right\|_{2}^{2} = \left\| \boldsymbol{X}^{T} \boldsymbol{W} \right\|_{F}^{2}.$$

$$(4)$$

where  $\mathbf{W}$  is the  $m \times (m-k)$  matrix defined as  $\mathbf{W} := [\mathbf{w}_{k+1}, \dots, \mathbf{w}_m]$ . Since  $\{\mathbf{w}_{k+1}, \dots, \mathbf{w}_m\}$  is a basis for  $W^{\perp}$ , we can write  $\operatorname{proj}_{W^{\perp}}(\mathbf{x}_i) = \sum_{j=k+1}^m c_i \mathbf{w}_i$  where  $c_j = \langle \mathbf{x}_i, \mathbf{w}_j \rangle$ . Expanding the left hand side of eq. (4) yields:

$$\sum_{i=1}^{n} \|\operatorname{proj}_{W^{\perp}}(\boldsymbol{x}_{i})\|_{2}^{2} = \sum_{i=1}^{n} \langle \operatorname{proj}_{W^{\perp}}(\boldsymbol{x}_{i}), \operatorname{proj}_{W^{\perp}}(\boldsymbol{x}_{i}) \rangle$$

$$= \sum_{i=1}^{n} \sum_{j=k+1}^{m} \sum_{\ell=k+1}^{m} \langle \boldsymbol{x}_{i}, \boldsymbol{w}_{j} \rangle \langle \boldsymbol{x}_{i}, \boldsymbol{w}_{\ell} \rangle \delta_{j,\ell}$$

$$= \sum_{i=1}^{n} \sum_{j=k+1}^{m} |\langle \boldsymbol{x}_{i}, \boldsymbol{w}_{j} \rangle|^{2}.$$
(5)

Since we are working over the reals, we may write  $\langle \boldsymbol{x}_i, \boldsymbol{w}_j \rangle = \langle \boldsymbol{w}_j, \boldsymbol{x}_i \rangle$ , hence the equality follows in eq. (4).

## Exercise 2

Let  $X = U\Sigma V^T$  be a singular value decomposition of X.

a) We wish to show that

$$\|\boldsymbol{X}^T \boldsymbol{W}\|_{\mathrm{F}}^2 = \|\boldsymbol{\Sigma}^T \boldsymbol{U}^T \boldsymbol{W}\|_{\mathrm{F}}^2. \tag{6}$$

This can easily be shown by noting that for arbitrary unitary matrix  $\boldsymbol{B}$  and matrix  $\boldsymbol{A}$  we have

$$\|\boldsymbol{B}\boldsymbol{A}\|_{\mathrm{F}} = \|\boldsymbol{A}\|_{\mathrm{F}}.\tag{7}$$

This follows from the fact that the Frobenius norm on  $\mathbb{C}^{m\times n}$  is the same as the 2-norm on the vector space  $\mathbb{C}^{mn}$  and that  $\|\boldsymbol{B}\boldsymbol{x}\|_2 = \|\boldsymbol{x}\|_2$  for all  $\boldsymbol{x}$ . Indeed, we have

$$\|\boldsymbol{B}\boldsymbol{A}\|_{\mathrm{F}} = \sum_{j=1}^{n} \|\boldsymbol{B}\boldsymbol{a}_{:j}\|_{2} = \sum_{j=1}^{n} \|\boldsymbol{a}_{:j}\|_{2} = \|\boldsymbol{A}\|_{\mathrm{F}}.$$
 (8)

Setting  $\mathbf{B} \coloneqq \mathbf{V}$  and  $\mathbf{A} \coloneqq \mathbf{\Sigma}^T \mathbf{U}^T \mathbf{W}$  in eq. (7) and squaring proves the result.

b) We now wish to show that

$$\|\boldsymbol{X}^T \boldsymbol{W}\|_{\mathrm{F}} = \sum_{i=1}^m \sigma_i^2 \|\operatorname{proj}_{W^{\perp}}(\boldsymbol{u}_i)\|_2^2$$
(9)

Expanding the matrix  $U^TW$  tells us that the ijth entry is given by

$$(\boldsymbol{U}^T \boldsymbol{W})_{ij} = \langle \boldsymbol{u}_i, \boldsymbol{w}_{i+k} \rangle \tag{10}$$

for  $i=1,\ldots,m$  and  $j=1,\ldots,m-k$ . Pre-multiplying by  $\Sigma^T$  gives

$$(\mathbf{\Sigma}^T \mathbf{U}^T \mathbf{W})_{ij} = \sigma_i \langle \mathbf{u}_i, \mathbf{w}_{j+k} \rangle \tag{11}$$

We can therefore write the Frobenius norm of this matrix as

$$\left\|\boldsymbol{\Sigma}^T \boldsymbol{U}^T \boldsymbol{W}\right\|_{\mathrm{F}}^2 = \sum_{i=1}^m \sigma_i^2 \sum_{j=1}^{m-k} \langle \boldsymbol{u}_i, \boldsymbol{w}_{j+k} \rangle = \sum_{i=1}^m \sigma_i^2 \|\mathrm{proj}_{W^{\perp}}(\boldsymbol{u}_i)\|_2^2$$

as we wanted to show. Note that the singular values are ordered by convention  $\sigma_1 \geq \ldots \sigma_m \geq 0$ .

c) We also have the property that since the columns of W are orthonormal, and since there are m-k such columns, that

$$\|\boldsymbol{W}\|_{\mathrm{F}}^2 = m - k. \tag{12}$$

Since  $U^T$  is unitary, we have that

$$\|\boldsymbol{W}\|_{\mathrm{F}}^{2} = \|\boldsymbol{U}^{T}\boldsymbol{W}\|_{\mathrm{F}}^{2} = \sum_{i=1}^{m} \sum_{j=k+1}^{m} \langle \boldsymbol{u}_{i}, \boldsymbol{w}_{j} \rangle^{2} = \sum_{i=1}^{m} \|\operatorname{proj}_{W^{\perp}}(\boldsymbol{u}_{i})\|_{2}^{2}.$$
 (13)

## Exercise 3

a) We can now find the "distance to the best subspace"  $\|\boldsymbol{X}^T\boldsymbol{W}\|_{\mathrm{F}}^2$  by solving the minimization problem

$$\min \sum_{i=1}^{m} \sigma_i^2 \kappa_i \tag{14}$$

where  $0 \le \kappa_i \le 1$  and  $\sum_{i=1}^m \kappa_i = m - k$ . Indeed, we have

$$\|\boldsymbol{X}^T \boldsymbol{W}\|_{\mathrm{F}}^2 = \sum_{i=1}^m \sigma_i^2 \|\mathrm{proj}_{W^{\perp}}(\boldsymbol{u}_i)\|_2^2$$
 (15)

where setting  $\kappa_i := \|\operatorname{proj}_{W^{\perp}}(\boldsymbol{u}_i)\|_2^2$  yields the above sum. Furthermore, we have that  $0 \le \kappa_i \le 1$  since the  $\boldsymbol{u}_i$ 's are orthonormal, and we showed above that  $\sum_{i=1}^m \kappa_i = m - k$ .

- b) The solution to this problem is obtained by setting  $\kappa_1 = \ldots = \kappa_k = 0$  and  $\kappa_{k+1} = \ldots = \kappa_m = 1$ , since the  $\sigma_i$ 's are ordered by magnitude. Indeed, the sum is minimized if we chose the vectors such that the columns of  $\boldsymbol{W}$  are perpendicular to  $\mathrm{Span}(\boldsymbol{u}_1,\ldots,\boldsymbol{u}_k)$ . Since  $0 \le \kappa_i \le 1$ , this means in order for  $\sum_{i=1}^m = m-k$ , we must have  $\kappa_{k+1} = \ldots = \kappa_m = 1$ .
- c) We can therefore choose the matrix W to consist of the last m-k vectors  $u_{k+1}, \ldots, u_m$  and let  $W^{\perp} = \operatorname{Span}(u_{k+1}, \ldots, u_m)$ . This set of vectors satisfies the minimization problem discussed above.
- d) We implement a routine for solving this minimization problem in Python.

```
def bestfit(X, k):
    U = np.linalg.svd(X)[0]
    B = U[:, :k]
    W = U[:, k:]
    dist = np.linalg.norm(X.T.dot(W))
    return B, dist
```

## Exercise 4

We can alternatively choose to implement a "vanilla" least squares, where we view the first k components of each column in  $\boldsymbol{X}$  as input data and the last m-k components of each column as output data. By defining  $\boldsymbol{Y} := \boldsymbol{X}_{1:k,1:n}$  and  $\boldsymbol{Z} := \boldsymbol{X}_{k+1:m,1:n}$ , we seek the matrix  $\boldsymbol{A}$  that minimizes

$$\|\boldsymbol{A}\boldsymbol{Y} - \boldsymbol{Z}\|_{\mathrm{F}}.\tag{16}$$

a) It can be shown that if  $\boldsymbol{A}$  solves the least squares problem, then it must satisfy the relation

$$\boldsymbol{A}^T = (\boldsymbol{Y}\boldsymbol{Y}^T)^{-1}\boldsymbol{Y}\boldsymbol{Z}^T, \tag{17}$$

however I was not successful in proving this.

b) We implement vanilla least squares in Python:

```
def vanleastsqr(X, k):
    Y = X[0:k, 0:]
    Z = X[k:, 0:]
    A = la.inv(Y.dot(Y.T)).dot(Y).dot(Z.T).T
    dist = la.norm(A.dot(Y) - Z)
    return A, dist
```

## Exercise 5

The two algorithms previously discussed are tested on two test  $2 \times 21$  matrices  $X_0$  and  $X_1$ . The first row of matrix  $X_1$  contains 21 uniformly spaced values between -1 and 1, while the second row contains uniformly random numbers between -1 and 1. All the elements in  $X_2$  are random.

As we see from fig. 1, the "bestfit" algorithm provides a better fit in both cases. This might be due to numerical instability in the computation of the inverse in the vanilla least squares.

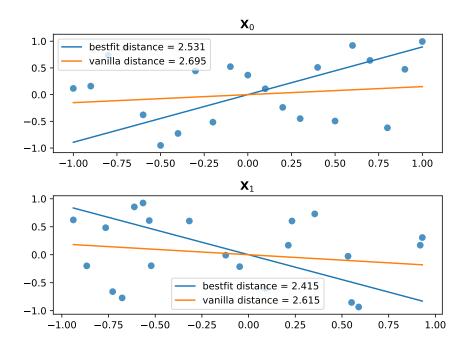


Figure 1: Two approaches to the least squares problem of finding the best subspace. As we see, the fit is better for the "bestfit" algorithm, in both cases.