

MANDATORY ASSIGNMENT 4

MATINF4130

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October 23, 2017

Introduction

In this assignment we take a look at the PageRank algorithm, that helped establish GOOGLE as a powerful search engine. The algorithm assigns to each web page its “popularity” in a way that mimics how a human would define a popular web page. Before discussing the algorithm itself, we start with a mathematical intermezzo.

Exercise 1 We first establish some notation. Let \mathcal{S} denote the unit simplex

$$\mathcal{S} := \left\{ \mathbf{x} \in \mathbb{R}^n \mid x_i \geq 0 \text{ for } i = 1, \dots, n \text{ and } \sum_{i=1}^n x_i = 1 \right\}. \quad (1)$$

For later, we note that this is a closed and bounded set, which in \mathbb{R}^n is equivalent to compact. We let \mathbf{A} be a real $n \times n$ matrix with non-negative elements $a_{ij} \geq 0$, whose columns sum to one, and refer to this as a *stochastic matrix*. The image of \mathcal{S} under \mathbf{A} is denoted

$$\mathbf{A}(\mathcal{S}) := \{ \mathbf{Ax} \mid \mathbf{x} \in \mathcal{S} \}. \quad (2)$$

a) If $\mathbf{y} \in \mathbf{A}(\mathcal{S})$ then $\mathbf{y} = \mathbf{Ax}$ for some \mathbf{x} . Note that since both x_i and a_{ij} are non-negative, we must have y_i non-negative for $i = 1, \dots, n$. The sum

$$\sum_{i=1}^n y_i = \sum_{j=1}^n x_j \left(\sum_{i=1}^n a_{ij} \right) = \sum_{j=1}^n x_j = 1 \quad (3)$$

tells us that $\mathbf{y} \in \mathcal{S}$ and consequently $\mathbf{A}(\mathcal{S}) \subseteq \mathcal{S}$.

- b) Considering $\mathbf{A}: \mathcal{S} \rightarrow \mathcal{S}$ as a linear operator, it suffices to show that it is bounded to show continuity. We have that

$$\|\mathbf{A}\|_1 = \max_{\|x\|=1} \|\mathbf{A}x\|_1 = 1 \quad (4)$$

so \mathbf{A} is bounded, and therefore also continuous in the $\|\cdot\|_1$ norm.

- c) Assume that (λ, \mathbf{v}) is an eigenpair for \mathbf{A} . Since $\mathbf{A}\mathbf{v} = \lambda\mathbf{v} \in \mathcal{S}$, we must have $|\lambda| \leq 1$. Since \mathcal{S} is closed and bounded it is compact, and since \mathcal{S} is continuous, it follows by Brouwer's fixed-point theorem that there exists a \mathbf{w} such that

$$\mathbf{A}\mathbf{w} = \mathbf{w}. \quad (5)$$

Consequently, $(1, \mathbf{w})$ is a right eigenpair for \mathbf{A} .

From now on, we assume that the matrix entries a_{ij} are all strictly positive. Denote by \mathcal{S}^* the *interior* of \mathcal{S} :

$$\mathcal{S}^* := \{\mathbf{x} \in \mathcal{S} \mid x_i > 0 \text{ for } i = 1, \dots, n\}. \quad (6)$$

- d) Let $\mathbf{x} \in \mathcal{S}$ and set $\mathbf{y} = \mathbf{A}\mathbf{x}$. Since at least one of the elements x_i are non-negative, and all a_{ij} are strictly positive, we have $y_i > 0$ for all $i = 1, \dots, n$. This means that $\mathbf{y} \in \mathcal{S}^*$, so \mathbf{A} maps \mathcal{S} to its interior.
- e) Let $\mathbf{x}, \mathbf{y} \in \mathcal{S}$ be two distinct vectors. Since the components of \mathbf{x} and the components of \mathbf{y} sum to one, we have that the components of $\mathbf{z} := \mathbf{x} - \mathbf{y}$ sum to zero. This means that since \mathbf{x} and \mathbf{y} are different, \mathbf{z} is non-zero, hence $z_j < 0$ for at least one j . We will need this fact to achieve a strict inequality. We have that

$$\begin{aligned} \|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{y}\|_1 &= \|\mathbf{A}\mathbf{z}\|_1 = \sum_{i=1}^n \left| \sum_{j=1}^n a_{ij} z_j \right| \\ &< \sum_{j=1}^n \sum_{i=1}^n a_{ij} |z_j| \\ &= \sum_{j=1}^n |z_j| = \|\mathbf{z}\|_1 = \|\mathbf{x} - \mathbf{y}\|_1 \end{aligned} \quad (7)$$

Consequently, $\mathbf{A}: \mathcal{S} \rightarrow \mathcal{S}$ is a *contraction* in the $\|\cdot\|_1$ norm. Assume that $\mathbf{w}_1 \neq \mathbf{w}_2$ are two distinct eigenvectors with eigenvalue one. Then

$$\|\mathbf{w}_1 - \mathbf{w}_2\|_1 = \|\mathbf{A}\mathbf{w}_1 - \mathbf{A}\mathbf{w}_2\|_1 < \|\mathbf{w}_1 - \mathbf{w}_2\|_1, \quad (8)$$

which is a contradiction. We can therefore conclude that the geometric multiplicity $g(\lambda)$ of the eigenvalue $\lambda = 1$ is one.