

COA - May 2024

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Q.11]

Ans-c]

Given number = -20.25

Converting to IEEE Single precision floating point standard.

Sign bit = 1 (-ve)

→ Step 1:  
Converting the given number to binary form

$$\begin{array}{r|l} 2 & 20 \\ \hline 2 & 10 \\ 2 & 5 \\ 2 & 2 \\ & 1 \end{array}$$

0  
0  
1  
0  
1

↑

$$\begin{array}{l} 0.25 \times 2 = 0.5 \quad 0 \\ 0.5 \times 2 = 1 \quad 1 \end{array}$$

$$[-20.25] = \underbrace{10100}_{20} \underbrace{01}_{0.25}$$

• Step 2:  
Normalizing the binary number

$$[10100.01]_2$$

$$[1.010001 \times 2^4]$$

Mantissa

• Step 3:  
Finding the value of e

$$\begin{aligned} e' &= E + 127 \\ &= 4 + 127 \end{aligned}$$



$$\therefore E' = 131$$

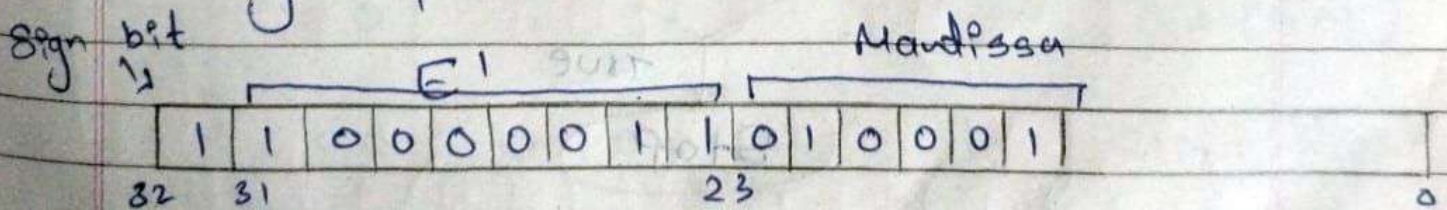
Note:  $E$  = Exponent & we got the value of exponent by shifting the decimal to the left while normalising. Then we got the value  $[24]$

- Step 4: Converting the value of  $E'$  to binary format

2	131	
2	65	1
2	32	100
2	16	00
2	8	01
2	4	0
2	2	0
	1	0

$$\therefore E' = 10000011$$

- Step 5: Single precision format (32-bit)

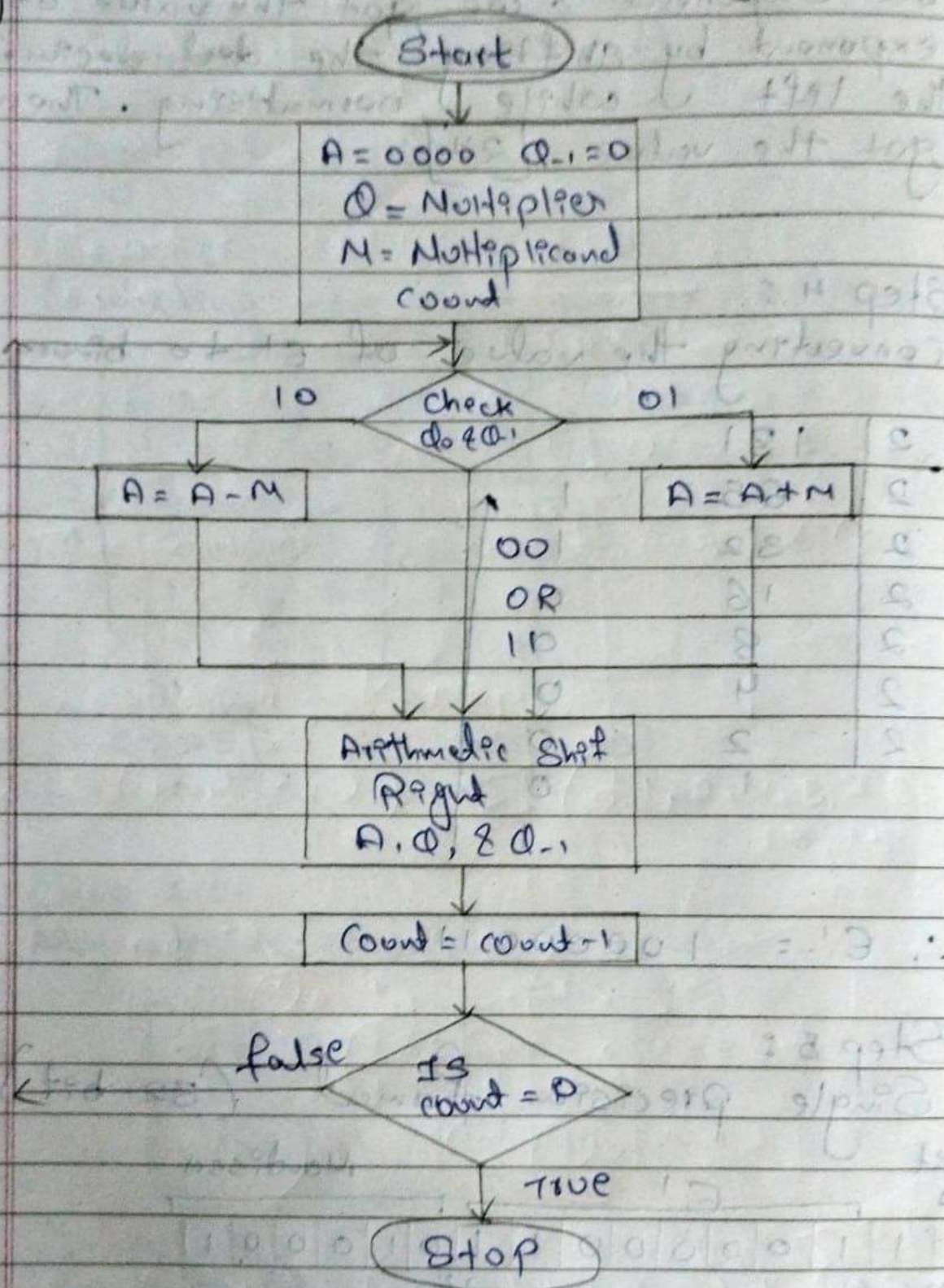




Q.2]

Ans-a]

## Flow-chart of Booths Algorithm





→ Multiplying  $(-5)$  and  $(6)$  using Booth's Algorithm.

$[+5]$  in binary bits  $(M) = \boxed{0101}$

$0101$

$1010$

→ 1's complement

+

$1011$

→ 2's complement

$[-5]$  in binary bits  $(-M) = \boxed{1011}$

$Q$  or  $(6)$  in binary bits  $= \boxed{0110}$

$A = 0000$

$Q_{-1} = 0$

Count = 4

A	$Q_s$	$Q_{-1}$	Count	Remark
0000	0110	0	4	Initialize
↓	0011	0	3	Shift Right
+ 1011				
1011				$A + (-M)$
↓	1001	0	2	Shift Right
+ 1011				
1000				$A + (-M)$
↓	0100	0	1	Shift Right
↓	0010	0	0	Shift Right



## \* Process Control Block

- Contains the process elements
- Created and managed by OS
- Allows support for multiple process.

Result: 1110 0010  
(A ⊕)

$$\begin{array}{r}
 00011101 \rightarrow 13 \text{ complement} \\
 + 11100010 \rightarrow 25 \text{ complement} \\
 \hline
 00011110 \rightarrow 14 \text{ complement}
 \end{array}$$

The answer is correct  $(-8) \times (6) = (-30)$

## Q.3] Solving K-Map

Ans-3]  $f(A, B, C, D) = \sum (0, 1, 2, 4, 5, 7, 8, 10, 11)$

→ + HAB

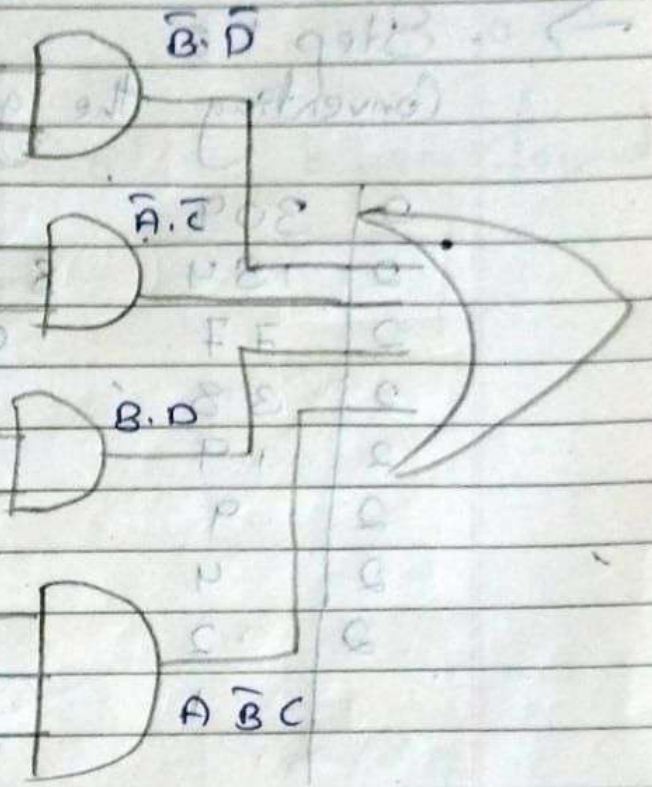
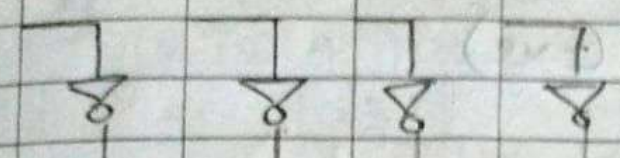
	00	01	11	10
00	1	1	1	1
01	1	1	1	1
11	1	1	1	1
10	1	1	1	1

$$f = (\bar{B} \cdot \bar{D}) + (\bar{A} \cdot \bar{C}) + (B \cdot D) + (A \bar{B} C)$$



# Logic Gates

A B C D



1100 = 12, 101011001 = 13

Non-overlapping binary numbers

100110010011001

100110010011001

10011001



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Q.1]

Ans-e]

Given number = -309.1875

In IEEE 754 "Double Precision Format"

Sign bit = 1 (ve)

→ Step 1:

Converting the given number to binary form

2 | 309

2 | 154

2 | 77

2 | 38

2 | 19

2 | 9

2 | 4

2 | 2

2 | 1

8 1

0

1

0

1

1

0

0

1

0.1875  $\times 2 = 0.375$  0

0.375  $\times 2 = 0.75$  0

0.75  $\times 2 = 1.5$  1

0.5  $\times 2 = 1$  1

309 = 100110101      0.1875 = 0011

• Step 2:

Normalising the binary number

[100110101.0011  $\times 2$ ]

[1.001101010011  $\times 2^8$ ] → [E/Exponent]

Mantissa

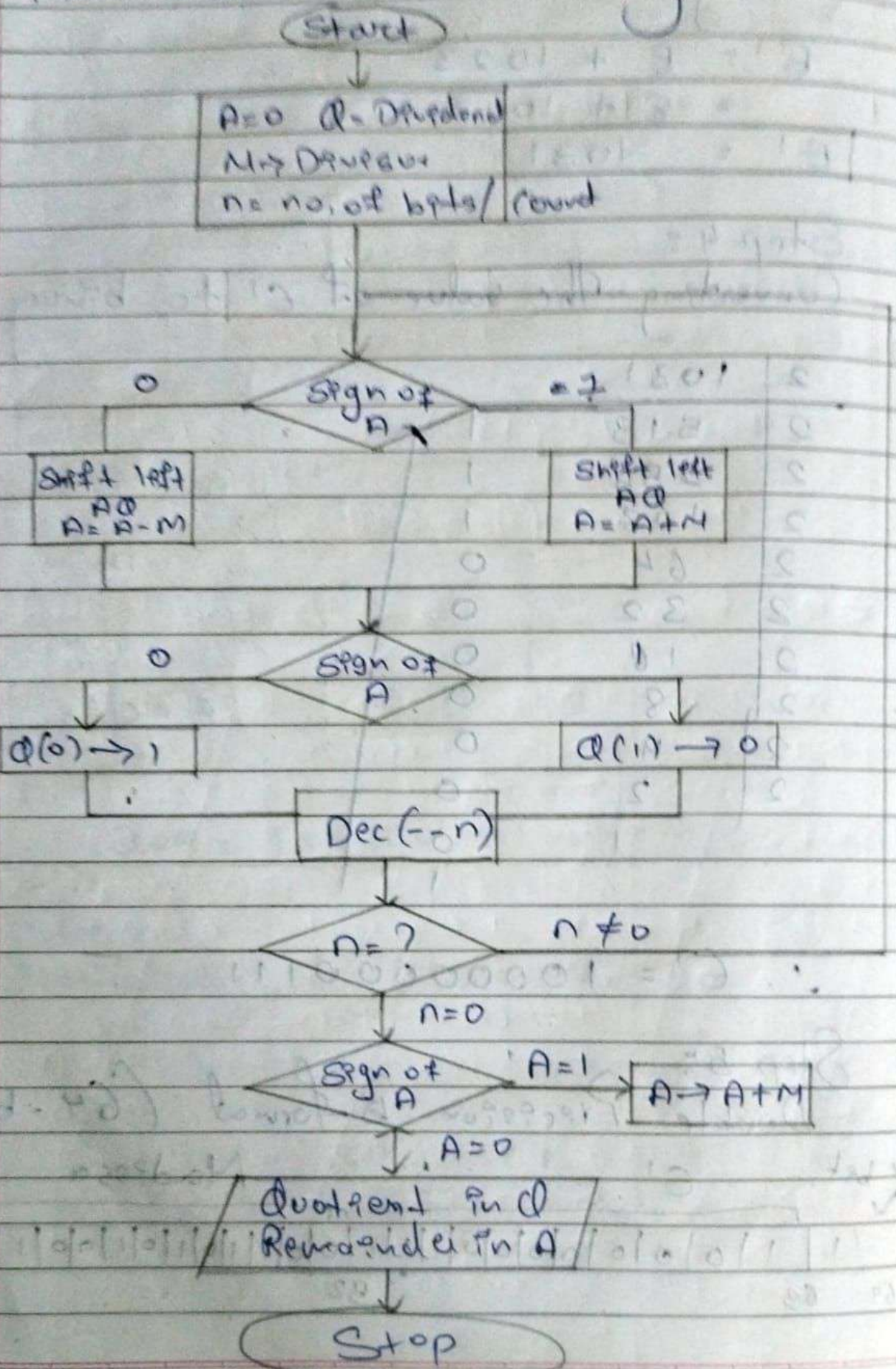






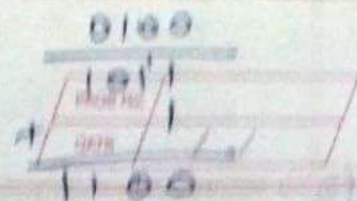
Q.3]  
Ans-a]

# Flow-chart of Non-Restoring Division





$12 \text{ (Dividend)} (D) = 1100$   
 $4 \text{ (Divisor)} (M) = 0100$   
 $4 \text{ (Divisor)} (C-M) = 1100$



→ Dividing 12 by 4

Operations	A	Dividend (D)	Count
Initialize	00000	1100	4
1. Shift left	00001	100	
A + M (-M)	11101	1000	3
2. Shift left	1011	000	
A + M	1111	0000	2
3. Shift left	1110	000	
A + M	1010	0001	1
4. Shift left	0100	001	
A + (-M)	0000	0011	0

Result :

A (Remainder) = 0000

Q (Quotient) = 0011 → 3

Ans: This matches  $12/4 = 3$



Q.5)

Ans-a

Multiplying  $(-15)$  and  $(3)$  using booth's algorithm

$[-15]$  in binary bits  $(M) = 110001$

$[-15]$  in binary bits  $(-M) = 001111$

$[3]$  in binary bits  $(Q) = 000011$

$A = 000000$ ,  $Q_{-1} = 0$  Count = 6

A	$Q_{-1}$	Q	Count	Remarks
000000	000011	0	6	Initial
+001111				
001111	010110	0	5	$A + (-M)$
↓				
0001110	100001	1	5	Shift Right
↓				
000011	110000	1	4	Shift Right
+110001				
1101100	011011	0	3	$A + M$
↓				
111010	011000	0	3	Shift Right
↓				
111101	001100	0	2	Shift Right
↓				
11110	100110	0	1	Shift Right
↓				
1110110	010011	0	0	Shift Right



$$A = 111111$$

$$Q = 010011$$

$$\text{Result } (A+Q) = \boxed{111111010011}$$

$$000000101100 \rightarrow \text{complement}$$

$$+ \quad \quad \quad 1$$


---


$$\boxed{101101}$$

32 16 8 4 2 1

Since MSB = 1 (-ve)

Ans Final result is (-45)

Q.6] Given :

Ans. a-]  $F(A, B, C, D) = \sum m(0, 2, 7, 10, 15)$

$+ \sum d(3, 14)$

→

AB \ CD	00	01	11	10
00	1	1	X	1
01	1	1	1	1
11	1	1	X	1
10	1	1	1	1

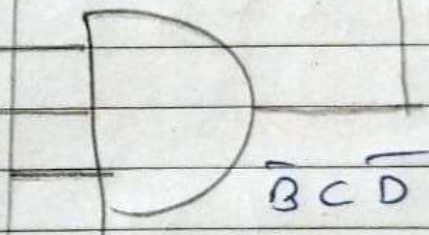
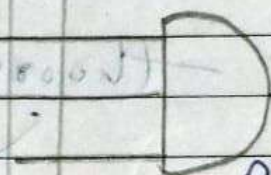
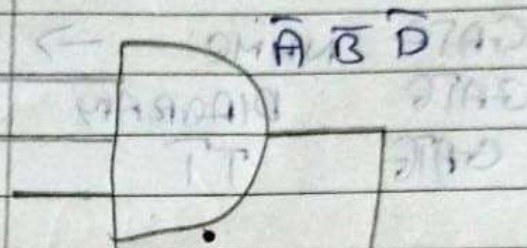
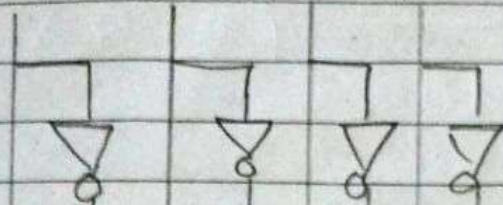
$$F = (A\bar{B}\bar{D}) + (BCD) + (\bar{B}C\bar{D})$$



Continuing COA May 2023

## Logic Gates

A B C D





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Q.1]

Given number is  $-10.125$

Ans-c]

IEEE 754 "Single Precision Method"

Sign bit = 1 (ve)

→

Step 1:

Converting the given numbers to binary form

2	10
2	5
2	2
	1

0 ↑  
1  
0

$$0.125 \times 2 = 0.25$$

0

$$0.25 \times 2 = 0.5$$

0

$$0.5 \times 2 =$$

1 ↓

$$101000001 = 10.125$$

$$[10.125] = \underbrace{1010}_{10} \underbrace{001}_{0.125}$$

ii.

Step 2:

Normalising the binary numbers

$$\rightarrow [1010.001]_2$$

$$\rightarrow [1.\underbrace{010001}_{\text{mantissa}}] \times 2^3 \rightarrow E$$

iii.

Step 3:

Finding the value of  $E'$

$$E' = E + 127$$

$$E' = 3 + 127$$

$$E' = 130$$



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$$\begin{array}{r} 2 \overline{) 130} \end{array}$$

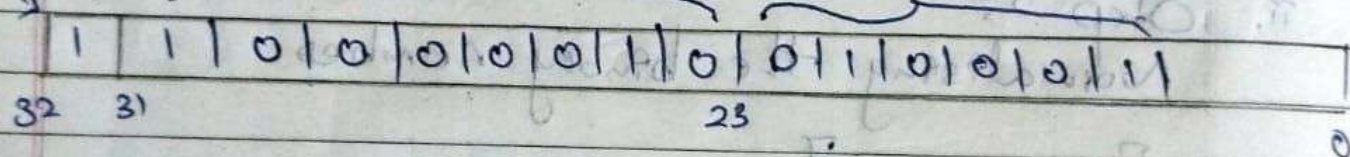
2	65	0
2	32	1
2	16	0
2	8	0
2	4	0
2	2	0
2	1	0
2	1	0

$$\therefore E' = 10000010$$

Single Precision format (32-bit)

६१

Handpass 9



Ans a)

$[-8]$  in binary format  $(M) = 01000$

1<sup>st</sup> complaint →

2<sup>nd</sup> complement

$N = 11000$

$$-M = 01000$$

Q [multiplier] = 00101

$D = 100000$

$d_1 = 0$       Count = 5



A	Q's	Qr	Count	Remarks
000000 +010000 ----- 010000	00101	0	5	Initialize
↓				A + (-M)
001000 +110000 ----- 111000	00010	1	4	Right Shift
↓				A + M
111100 +010000 ----- 001100	00001	0	3	Right Shift
↓				A + (-M)
000110 +110000 ----- 110110	00000	1	2	Right Shift
↓				A + M
111010 +100000 ----- 011010	10000	0	1	Right Shift
↓				Right Shift
111100	11000	0	0	

Result (A) = 11110      2(Q) = 11000  
(AQ) = 1111011000

Discarding the redundant and extra sign bit at MSB.

So the Result becomes = 111011000 → 9 bits

000100111 → 1<sup>st</sup> complement

+  
000101000

Note =  $32 + 8$   
= 40

The answer is correct for  $(-8) \times 5 = \underline{\underline{40}}$



Q.3]

Ans]

$$f(A, B, C, D) = \sum (0, 1, 2, 3, 4, 5, 8, 9, 10, 11, 12, 13)$$

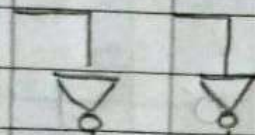
→

AB \ CD	00	01	11	10
00	1	0	1	1
01	1	4	1	5
11	1	2	1	3
10	1	6	1	9

$$f = \bar{C} + (\bar{B} \cdot C)$$

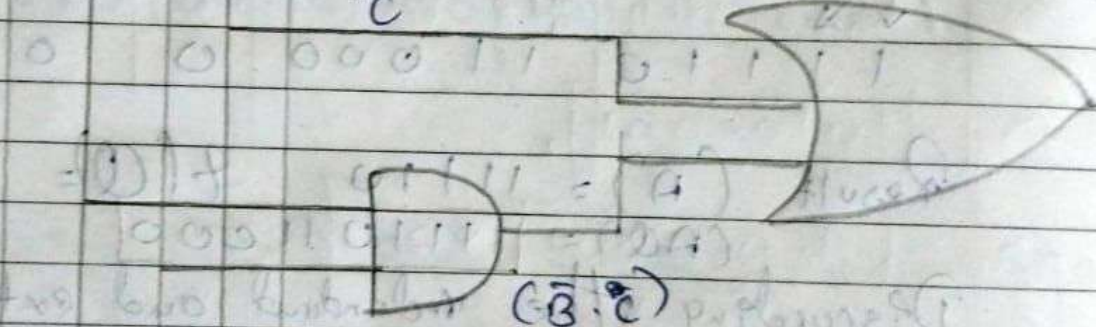
Logic Gates

A B C D



$\bar{C}$

$\bar{C} + (\bar{B} \cdot C)$



$(\bar{B} \cdot C)$

truth table for the function  $f(A, B, C, D)$

truth table for the function  $f(A, B, C, D)$

truth table for the function  $f(A, B, C, D)$

truth table for the function  $f(A, B, C, D)$

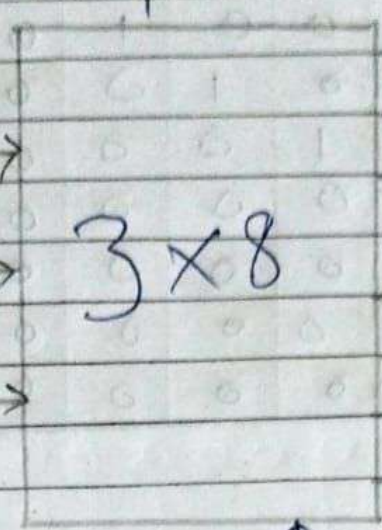


13)

Q.2]   
 Ans-b

# Decoder Example:

Input lines  
X  
Y  
Z



Decoder or output lines  
D<sub>0</sub>  
D<sub>1</sub>  
D<sub>2</sub>  
D<sub>3</sub>  
D<sub>4</sub>  
D<sub>5</sub>  
D<sub>6</sub>  
D<sub>7</sub>

$$n : 2^n$$

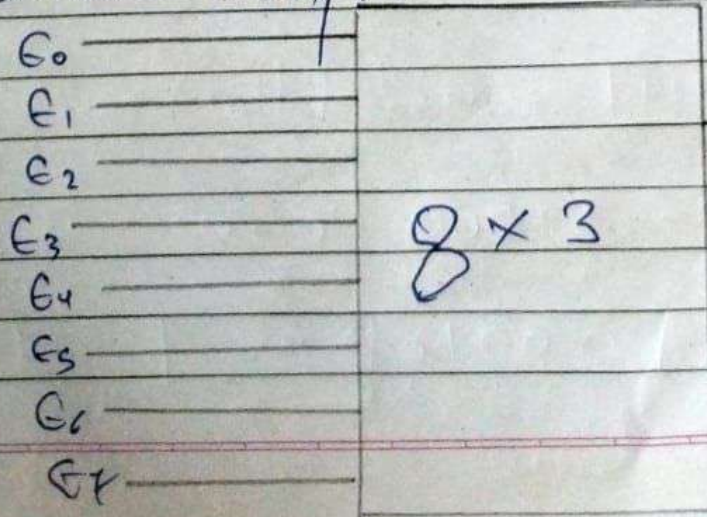
↑ Enabled = 1

Input lines

→ output lines

X	Y	Z	D <sub>7</sub>	D <sub>6</sub>	D <sub>5</sub>	D <sub>4</sub>	D <sub>3</sub>	D <sub>2</sub>	D <sub>1</sub>	D <sub>0</sub>
0	0	0	0	0	0	0	0	0	0	1
0	0	1	0	0	0	0	0	0	1	0
0	1	0	0	0	0	0	0	1	0	0
0	1	1	0	0	0	0	1	0	0	0
1	0	0	0	0	0	1	0	0	0	0
1	0	1	0	0	1	0	0	0	0	0
1	1	0	0	1	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0	0

# Encoder Example:



X

Y

Z



$E_7$	$E_6$	$E_5$	$E_4$	$E_3$	$E_2$	$E_1$	$E_0$	$X$	$Y$	$Z$
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

↑ output 1  
input 1

$X$	$Y$	$Z$	$D_7$	$D_6$	$D_5$	$D_4$	$D_3$	$D_2$	$D_1$	$D_0$
0	0	0	0	0	0	0	0	0	0	0
0	0	1	0	0	0	0	0	0	1	0
0	1	0	0	0	0	0	0	0	0	0
0	1	1	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0	0	0
1	1	1	0	0	0	0	0	0	0	0

Encoder (Example)

$E_0$   
 $E_1$   
 $E_2$   
 $E_3$   
 $E_4$   
 $E_5$   
 $E_6$   
 $E_7$