

NATIONAL UNIVERSITY OF COMPUTER AND EMERGING SCIENCES Peshawar

OPERATING SYSTEMS LAB

Lab Manual Process Creation

fork, wait, exit, getpid and getppid Unix System Calls

15 COMPILE AND EXECUTE A C PROGRAM

Write down the C code in a text editor of your choice and save the file using **.c** extension i.e. *filename.c*. Now through your terminal move to the directory where *filename.c* resides.

gcc firstprogram . c

It will compile the code and by default an executable named *a.out* is created.

gcc –o firstprogram firstprogram . c

If your file is named firstprogram.c then type -o firstprogram as the parameter to gcc. It would create an executable by the name firstprogram for the source code named firstprogram.c.

To execute the program simply write the following:

./ a . out OR ./ firstprogram

In case you have written the code in C++ saved using .cpp extension, compile the code using g++ as:

g++ filename . cpp OR g++ -o exec_name filename . cpp

And execute as ./a.out or ./exec_name

16 Process Creation

When is a new process created?

- 1. System startup
- 2. Submit a new batch job/Start program



- 3. Execution of a system call by process
- 4. A user request to create a process

On the creation of a process following actions are performed:

- 1. Assign a unique process identifier. Every process has a unique process ID, a nonnegative integer. Because the process ID is the only well-known identifier of a process that is always unique, it is used to guarantee uniqueness.
- 2. Allocate space for the process.
- 3. Initialize process control block.
- 4. Set up appropriate linkage to the scheduling queue.

17 FORK SYSTEM CALL

An existing process can create a new process by calling the fork function.

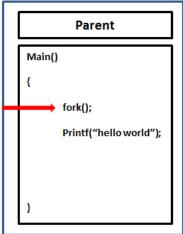
```
# include <unistd . h> pid_t fork (
void );
// Returns : 0 in child , process ID of child in parent , -1 on error
```

The definition of the pid_t is given in <sys/types> include file and <unistd.h> contain the declaration of *fork* system call.

IMPORTANT POINTS

- 1. The new process created by fork is called the child process.
- 2. This function is called once but returns twice. The only difference in the returns is that the return value in the child is 0, whereas the return value in the parent is the process ID of the new child.
- 3. Both the child and the parent continue executing with the instruction that follows the call to fork.
- 4. The child is a copy of the parent. For example, the child gets a copy of the parent's data space, heap, and stack. Note that this is a copy for the child; the parent and the child do not share these portions of memory.
- 5. In general, we never know whether the child starts executing before the parent, or vice versa. The order depends on the scheduling algorithm used by the kernel.
- [2][1] [The process will execute sequentially until it arrives at fork system call. This process will be named as parent process and the one created using fork() is termed as





child process.] [After fork system call both the child and the parent process continue executing with the instruction that follows the call to

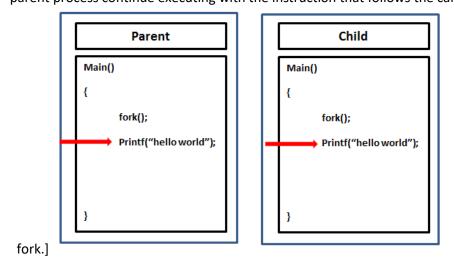


Figure 3: Parent and Child Process

21 [Fork system call returns child process ID (created using fork) in parent process and



zero in child process]

[We

can execute different portions of code in parent and child process based on the return value of fork that is either zero or greater than

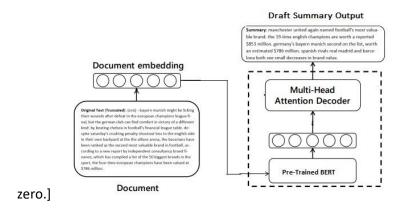


Figure 4: Return Value of fork system call

18 Examples Of Fork()

1. Fork()'s Return Value



2. Manipulating Local and Global Variables

```
# include
             <unistd . h>
# include
             <sys / types . h>
# include
             <errno . h>
# include
            <s t d i o . h>
# include <sys / wait . h> #
include < s t d l i b . h> int
         global =0;
     main ()
int
           int
                 status;
           pid_t child_pid ; int local =
           0;
           /* now create new process */ child_pid = fork ( );
           if (child pid
                           >=
                                     0){
                                              /* fork succeeded */
                      if ( child_pid
                                              0){
                                 /* fork ( ) returns 0 for the child process */ printf ( " child process !\n"
                                 // Increment the local and
                                                                 global
                                                                         variables local ++;
                                 global ++;
                                 printf ( "\n child ' s local = %d , child ' s global = %d\n" , local ,
                                                        global);
           }
                       else {
                                /* parent
                                              process */
                                   printf ( " parent
                                                        process !\n");
                                 printf ( " parent PID = %d , child pid = %d\n" , getpid ( ) , child_pid );
                                    int w=wait(& status);
                                 //The change in local and global variable //in child process
                                 should not reflect//here in parent process.
                                     printf ( "\n Parent ' z
                                                              local = %d, parent's
                                                                                              global
                                                       = %d\n" , local , global );
                                 printf ( " Parent says bye!\n" ); exit ( 0 ); /*
                                 parent exits */
           }
```



19 WAIT SYSTEM CALL

This function blocks the calling process until one of its child processes exits. wait() takes the address of an integer variable and returns the process ID of the completed process.

```
pid_t wait ( int * status )
Include <sys/wait.h> and <stdlib.h> for definition of wait().
```

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The execution of wait() could have two possible situations

- 1. If there are at least one child processes running when the call to wait() is made, the caller will be blocked until one of its child processes exits. At that moment, the caller resumes its execution
- 2. If there is no child process running when the call to wait() is made, then this wait() has no effect at all. That is, it is as if no wait() is there.

EXAMPLE

```
# include
               <s t d i o . h>
# include
                <unistd . h>
# include <sys / types . h> # include
<stdlib.h>
void
            main(void) { int
                    status =0;
            pid t
                       pid=fork();
             if ( pid ==
                                           check will
                                                           pass
                          0){ //This
                                                                       only for child
                                                                                            process
                        \boldsymbol{printf} ( " I am Child process with pid %d and i am not waiting\n" , getpid ( ) ) ;
                        exit (status);
            }
            else if (pid > 0){ // THIS Check will pass only for parent printf ("I am Parent process
                        with pid %d and i am waiting\n", getpid ());
                                      exitedChildId=wait(& status);
                        printf ( " I am Parent process and the child with pid %d is exited\n",
                                    exitedChildId);
            }
            else if ( pid < 0){ // if fork ( ) f a i l s printf ( " Error in Fork " ) ;
            }
```



20 EXIT SYSTEM CALL

A computer process terminates its execution by making an exit system call.

21 COMMAND LINE ARGUMENT TO C PROGRAM

To run:

./commandlineargument.out abc def 1 2