# A Two Level Dependent Session Type Theory

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# 1 Syntax

```
variables
                     x, y, z
channels
                     c, d
                                                              U \mid L
       sorts
                     s, r, t
       roles
                                                             + | -
                    m, n, A, B, C
     _{\rm terms}
                                                   ::= x \mid c \mid s
                                                              \Pi_t(x:A).B \mid \Pi_t\{x:A\}.B \mid \Sigma_t(x:A).B \mid \Sigma_t\{x:A\}.B
                                                              \lambda_t(x:A).m \mid \lambda_t\{x:A\}.m \mid \langle m,n \rangle_t \mid \{m,n\}_t
                                                             m \ n \mid \mathcal{R}^{\Sigma}_{[z]A}(m,[x,y]n) \mid \mu(x:A).m \\ 1 \mid () \mid 2 \mid \text{true} \mid \text{false} \mid \mathcal{R}^{2}_{[z]A}(m,n_{1},n_{2})
                                                              TA \mid \text{return } m \mid \text{let } x \leftarrow m \text{ in } n
                                                              proto \mid \operatorname{end} \rho \mid \rho\{x:A\}.B \mid \rho(x:A).B \mid \rho \operatorname{Ch} A
                                                              fork (x:A).m \mid \text{recv } m \mid \text{recv } \{m\} \mid \text{send } m \mid \text{send } \{m\}
                                                              close m \mid \text{wait } m
                                                              x \mid c \mid \lambda_t(x:A).m \mid \lambda_t\{x:A\}.m \mid \langle u,v \rangle_t \mid \{v,m\}_t
     values u, v
                                                               () | true | false | return v | let x \leftarrow v in m
                                                               fork (x:A).m \mid \text{recv } v \mid \text{recv } \{v\} \mid \text{send } v \mid \text{send } \{v\}
                                                               \operatorname{send} u v \mid \operatorname{send} \{v\} \ m \mid \operatorname{close} v \mid \operatorname{wait} v
```

### 2 Static Fragment

#### Sort Order

$$\overline{\mathbf{U} \sqsubseteq s}$$
  $\overline{\mathbf{L} \sqsubseteq \mathbf{L}}$ 

#### **Static Context**

$$\frac{\Gamma \vdash \qquad \Gamma \vdash A : s \qquad x \in \mathit{fresh}(\Gamma)}{\Gamma, x : A \vdash}$$

### Core Typing

### **Data Typing**

$$\frac{\Gamma \vdash}{\Gamma \vdash 1 : \mathbf{U}} \qquad \frac{\Gamma \vdash}{\Gamma \vdash () : 1} \qquad \frac{\Gamma \vdash}{\Gamma \vdash 2 : \mathbf{U}} \qquad \frac{\Gamma \vdash}{\Gamma \vdash \text{true} : 2} \qquad \frac{\Gamma \vdash}{\Gamma \vdash \text{false} : 2}$$

$$\frac{\Gamma, z : 2 \vdash A : s \qquad \Gamma \vdash m : 2 \qquad \Gamma \vdash n_1 : A[\text{true}/z] \qquad \Gamma \vdash n_2 : A[\text{false}/z]}{\Gamma \vdash \mathbf{R}^2_{[z]A}(m, n_1, n_2) : A[m/z]}$$

## Monadic Typing

$$\frac{\Gamma \vdash A : s}{\Gamma \vdash \Tau A : \mathsf{L}} \qquad \frac{\Gamma \vdash m : A}{\Gamma \vdash \mathsf{return} \ m : \Tau A} \qquad \frac{\Gamma \vdash B : s}{\Gamma \vdash \mathsf{let} \ x : A \leftarrow m \ \mathsf{in} \ n : \Tau B}{\Gamma \vdash \mathsf{let} \ x : A \leftarrow m \ \mathsf{in} \ n : \Tau B}$$

### Session Typing

### 3 Dynamic Fragment

### **Dynamic Context**

$$\frac{\Gamma;\Delta \vdash \qquad \Gamma \vdash A:s \qquad x \in \mathit{fresh}(\Gamma)}{\Gamma,x:A;\Delta,x:_s A \vdash} \qquad \qquad \frac{\Gamma;\Delta \vdash \qquad \Gamma \vdash A:s \qquad x \in \mathit{fresh}(\Gamma)}{\Gamma,x:A;\Delta \vdash}$$

#### Context Merge

$$\frac{\Delta_1 \cup \Delta_2 = \Delta \qquad x \in \mathit{fresh}(\Delta)}{(\Delta_1, x :_U A) \cup (\Delta_2, x :_U A) = (\Delta, x :_U A)}$$
 
$$\frac{\Delta_1 \cup \Delta_2 = \Delta \qquad x \in \mathit{fresh}(\Delta)}{(\Delta_1, x :_L A) \cup \Delta_2 = (\Delta, x :_L A)}$$
 
$$\frac{\Delta_1 \cup \Delta_2 = \Delta \qquad x \in \mathit{fresh}(\Delta)}{\Delta_1 \cup (\Delta_2, x :_L A) = (\Delta, x :_L A)}$$

#### Context Constraint

$$\frac{\Delta \triangleright \mathbf{U}}{\Delta, x :_{\mathbf{U}} A \triangleright \mathbf{U}} \qquad \qquad \frac{\Delta \triangleright \mathbf{L}}{\Delta, x :_{s} A \triangleright \mathbf{L}}$$

#### Core Typing

### **Data Typing**

$$\frac{\Gamma; \Delta \vdash \Delta \triangleright \mathbf{U}}{\epsilon; \Gamma; \Delta \vdash (): 1}$$

$$\frac{\Gamma; \Delta \vdash \Delta \triangleright U}{\epsilon; \Gamma; \Delta \vdash \text{true} : 2}$$

$$\frac{\Gamma; \Delta \vdash \Delta \triangleright U}{\epsilon; \Gamma; \Delta \vdash \text{false} : 2}$$

$$\frac{\Gamma, z: 2 \vdash A: s \qquad \Theta_1; \Gamma; \Delta_1 \vdash m: 2 \qquad \Theta_2; \Gamma; \Delta_2 \vdash n_1: A[\text{true}/z] \qquad \Theta_2; \Gamma; \Delta_2 \vdash n_2: A[\text{false}/z]}{\Theta_1 \cup \Theta_2; \Gamma; \Delta_1 \cup \Delta_2 \vdash \mathbf{R}^2_{[z]A}(m, n_1, n_2): A[m/z]}$$

### Monadic Typing

$$\frac{\Theta; \Gamma; \Delta \vdash m : A}{\Theta; \Gamma; \Delta \vdash \text{return } m : \Tau A}$$

$$\frac{\Gamma \vdash B : s \qquad \Theta_1; \Gamma; \Delta_1 \vdash m : \Tau A \qquad \Theta_2; \Gamma, x : A; \Delta_2, x :_r A \vdash n : \Tau B}{\Theta_1 \cup \Theta_2; \Gamma; \Delta_1 \cup \Delta_2 \vdash \text{let } x \leftarrow m \text{ in } n : \Tau B}$$

### Session Typing

$$\frac{\Gamma; \Delta \vdash \quad \epsilon \vdash A : \operatorname{proto} \quad \Delta \rhd \mathbf{U}}{c :_{\mathbf{L}} \; \rho \mathsf{Ch} \, A; \Gamma; \Delta \vdash c : \rho \mathsf{Ch} \, A}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \operatorname{Ch} (\rho_2 \{x : A\}.B) \qquad \rho_1 \operatorname{xor} \rho_2 =?}{\Theta; \Gamma; \Delta \vdash \operatorname{recv} \{m\} : \operatorname{T} (\Sigma_L \{x : A\}.\rho_1 \operatorname{Ch} B)}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \mathrm{Ch} \left( \rho_2 \{ x : A \} . B \right) \qquad \rho_1 \operatorname{xor} \rho_2 = !}{\Theta; \Gamma; \Delta \vdash \operatorname{send} \{ m \} : \Pi_{\mathbf{L}} \{ x : A \} . T \left( \rho_1 \mathrm{Ch} \, B \right)}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \operatorname{Ch} \left( \operatorname{end} \rho_2 \right) \qquad \rho_1 \operatorname{xor} \rho_2 =?}{\Theta; \Gamma; \Delta \vdash \operatorname{wait} m : \operatorname{T} 1}$$

$$\frac{\Theta; \Gamma, x: + \operatorname{Ch} A; \Delta, x:_{\operatorname{L}} + \operatorname{Ch} A \vdash m: \operatorname{T} 1}{\Theta; \Gamma; \Delta \vdash \operatorname{fork} (x:A).m: \operatorname{T} (-\operatorname{Ch} A)}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \operatorname{Ch} (\rho_2(x : A).B) \qquad \rho_1 \operatorname{xor} \rho_2 =?}{\Theta; \Gamma; \Delta \vdash \operatorname{recv} m : \operatorname{T} (\Sigma_L(x : A).\rho_1 \operatorname{Ch} B)}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \operatorname{Ch} (\rho_2(x : A).B) \qquad \rho_1 \operatorname{xor} \rho_2 = !}{\Theta; \Gamma; \Delta \vdash \operatorname{send} m : \Pi_{\operatorname{L}}(x : A).\operatorname{T} (\rho_1 \operatorname{Ch} B)}$$

$$\frac{\Theta; \Gamma; \Delta \vdash m : \rho_1 \operatorname{Ch} (\operatorname{end} \rho_2) \qquad \rho_1 \operatorname{xor} \rho_2 = !}{\Theta; \Gamma; \Delta \vdash \operatorname{close} m : T 1}$$