A Two Level Linear Dependent Type Theory

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1 Syntax

2 Static Fragment

Sort Order

$$\overline{\mathbf{U} \sqsubseteq s}$$
 $\overline{\mathbf{L} \sqsubseteq \mathbf{L}}$

Static Context

$$\frac{\Gamma \vdash \qquad \Gamma \vdash A : s \qquad x \not\in \mathit{fresh}(\Gamma)}{\Gamma, x : A \vdash}$$

Static Typing

atic Typing
$$\frac{\Gamma \vdash \Gamma}{\Gamma \vdash s : U} \frac{\Gamma, x : A \vdash \Gamma}{\Gamma, x : A \vdash x : A} \frac{\Gamma \vdash A : s \quad \Gamma, x : A \vdash B : r}{\Gamma \vdash \Pi_t(x : A).B : t} \frac{\Gamma \vdash A : s \quad \Gamma, x : A \vdash B : r}{\Gamma \vdash \Pi_t\{x : A\}.B : t}$$

$$\frac{\Gamma, x : A \vdash m : B}{\Gamma \vdash \lambda_t(x : A).m : \Pi_t(x : A).B} \frac{\Gamma \vdash n : A}{\Gamma \vdash \lambda_t\{x : A\}.m : \Pi_t\{x : A\}.B} \frac{\Gamma \vdash m : \Pi_t(x : A).B}{\Gamma \vdash m n : B[n/x]} \frac{\Gamma \vdash m : \Pi_t(x : A).B}{\Gamma \vdash m n : B[n/x]} \frac{S \sqsubseteq t \quad r \sqsubseteq t \quad \Gamma \vdash A : s \quad \Gamma, x : A \vdash B : r}{\Gamma \vdash \Sigma_t(x : A).B : t}$$

$$\frac{S \sqsubseteq t \quad \Gamma \vdash A : s \quad \Gamma, x : A \vdash B : r}{\Gamma \vdash \Sigma_t\{x : A\}.B : t} \frac{\Gamma \vdash \Sigma_t(x : A).B}{\Gamma \vdash m : A} \frac{\Gamma \vdash n : B[m/x]}{\Gamma \vdash (m, n)_t : \Sigma_t(x : A).B}$$

$$\frac{\Gamma \vdash \Sigma_t\{x : A\}.B}{\Gamma \vdash \{m, n\}_t : \Sigma_t\{x : A\}.B} \frac{\Gamma \vdash m : A}{\Gamma \vdash n : B[m/x]}$$

$$\frac{\Gamma, z : \Sigma_t(x : A).B \vdash C : s \qquad \Gamma \vdash m : \Sigma_t(x : A).B \qquad \Gamma, x : A, y : B \vdash n : C[\langle x, y \rangle_t/z]}{\Gamma \vdash \mathbf{R}^{\Sigma}_{[z]C}(m, [x, y]n) : C[m/z]}$$

$$\frac{\Gamma,z:\Sigma_t\{x:A\}.B\vdash C:s\qquad \Gamma\vdash m:\Sigma_t\{x:A\}.B\qquad \Gamma,x:A,y:B\vdash n:C[\{x,y\}_t/z]}{\Gamma\vdash \mathbf{R}^\Sigma_{[z]C}(m,[x,y]n):C[m/z]}$$

$$\frac{\Gamma \vdash A : s \quad \Gamma \vdash B : r}{\Gamma \vdash A \&_t B : t} \qquad \frac{\Gamma \vdash m : A \quad \Gamma \vdash n : B}{\Gamma \vdash (m, n)_t : A \&_t B} \qquad \frac{\Gamma \vdash m : A \&_t B}{\Gamma \vdash \pi_1 m : A} \qquad \frac{\Gamma \vdash m : A \&_t B}{\Gamma \vdash \pi_2 m : B}$$

$$\frac{\Gamma \vdash A : s \quad \Gamma \vdash m : A \quad \Gamma \vdash n : A}{\Gamma \vdash m \equiv_A n : U} \qquad \frac{\Gamma \vdash m : A \&_t B}{\Gamma \vdash \pi_1 m : A}$$

$$\frac{\Gamma, x: A, y: A, p: x \equiv_A y \vdash B: s}{\Gamma \vdash \mathbf{R}^{\equiv}_{[x,y,p]B}([z]H, P): B[m/x, n/y, P/p]} \qquad \Gamma \vdash P: m \equiv_A n$$

$$\frac{\Gamma \vdash B : s \qquad \Gamma \vdash m : A \qquad A \simeq B}{\Gamma \vdash m : B}$$

3 Dynamic Fragment

Dynamic Context

$$\frac{\Gamma;\Delta \vdash \qquad \Gamma \vdash A:s \qquad x \in \mathit{fresh}(\Gamma)}{\Gamma,x:A;\Delta,x:_s A \vdash} \qquad \qquad \frac{\Gamma;\Delta \vdash \qquad \Gamma \vdash A:s \qquad x \in \mathit{fresh}(\Gamma)}{\Gamma,x:A;\Delta \vdash}$$

Context Merge

$$\frac{\Delta_1 \cup \Delta_2 = \Delta \quad x \in \mathit{fresh}(\Delta)}{(\Delta_1, x :_U A) \cup (\Delta_2, x :_U A) = (\Delta, x :_U A)}$$

$$\frac{\Delta_1 \cup \Delta_2 = \Delta \quad x \in \mathit{fresh}(\Delta)}{(\Delta_1, x :_L A) \cup \Delta_2 = (\Delta, x :_L A)}$$

$$\frac{\Delta_1 \cup \Delta_2 = \Delta \quad x \in \mathit{fresh}(\Delta)}{\Delta_1 \cup (\Delta_2, x :_L A) = (\Delta, x :_L A)}$$

Context Constraint

$$\frac{\Delta \triangleright \mathbf{U}}{\Delta, x :_{\mathbf{U}} A \triangleright \mathbf{U}} \qquad \qquad \frac{\Delta \triangleright \mathbf{L}}{\Delta, x :_{s} A \triangleright \mathbf{L}}$$

Dynamic Typing

4 Erasure

Erasure Relation

$$\frac{\Gamma, x : A; \Delta, x :_s A \vdash \Delta \rhd \cup}{\Gamma, x : A; \Delta, x :_s A \vdash x \sim x : A} \qquad \frac{\Gamma, x : A; \Delta, x :_s A \vdash m \sim m' : B}{\Gamma; \Delta \vdash \lambda_t(x : A).m \sim \lambda_t(x : \Box).m' : \Pi_t(x : A).B}$$

$$\frac{\Gamma, x : A; \Delta, h = m \sim m' : B}{\Gamma; \Delta \vdash \lambda_t(x : A).m \sim \lambda_t(x : \Box).m' : \Pi_t(x : A).B} \qquad \frac{\Gamma; \Delta \vdash m \sim m' : \Pi_t(x : A).B}{\Gamma; \Delta \vdash \lambda_t(x : A).m \sim m' : n' : B[n/x]}$$

$$\frac{\Gamma; \Delta \vdash m \sim m' : \Pi_t(x : A).B}{\Gamma; \Delta \vdash m \sim m' : B[n/x]} \qquad \frac{\Gamma; \Delta \vdash m \sim m' : \Pi_t(x : A).B}{\Gamma; \Delta \vdash m \sim m' : B[n/x]}$$

$$\frac{\Gamma \vdash \Sigma_t(x : A).B : t}{\Gamma; \Delta \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta_t \vdash m \sim m' : B[m/x]}{\Gamma; \Delta_t \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta_t \vdash m \sim m' : B[m/x]}{\Gamma; \Delta_t \vdash m \sim m' : A}$$

$$\frac{\Gamma \vdash \Sigma_t(x : A).B : t}{\Gamma; \Delta \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta_t \vdash n : B[m/x]}{\Gamma; \Delta \vdash \{m, n\}_t \sim \{m', n'\}_t : \Sigma_t(x : A).B}$$

$$\frac{\Gamma; \Delta \vdash m \sim m' : \Sigma_t(x : A).B}{\Gamma; \Delta_t \vdash m \sim m' : \Sigma_t(x : A).B} \qquad \frac{\Gamma, x : A, y : B; \Delta_t(x : A).B}{\Gamma; \Delta_t \vdash m \sim m' : \Sigma_t(x : A).B}$$

$$\frac{\Gamma; \Sigma \vdash \Sigma_t(x : A).B \vdash C : s}{\Gamma; \Delta_t \vdash m \sim m' : \Sigma_t(x : A).B} \qquad \Gamma, x : A, y : B; \Delta_t(x : A).B \vdash C : S$$

$$\frac{\Gamma; \Delta \vdash m \sim m' : \Sigma_t(x : A).B}{\Gamma; \Delta_t \vdash m \sim m' : \Sigma_t(x : A).B} \qquad \Gamma, x : A, y : B; \Delta_t(x : x).T$$

$$\frac{\Gamma; \Delta \vdash m \sim m' : A}{\Gamma; \Delta \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta \vdash m \sim m' : A}{\Gamma; \Delta \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta \vdash m \sim m' : A \&_t B}{\Gamma; \Delta \vdash m \sim m' : B}$$

$$\frac{\Gamma; \Delta \vdash m \sim m' : A}{\Gamma; \Delta \vdash m \sim m' : A} \qquad \frac{\Gamma; \Delta \vdash m \sim m' : A \&_t B}{\Gamma; \Delta \vdash m \sim m' : B}$$

5 Static Semantics

Static Reduction

$$\begin{array}{c} A \hookrightarrow A' \\ \hline \Pi_t(x:A).B \leadsto \Pi_t(x:A').B \\ \hline \Pi_t(x:A).B \leadsto \Pi_t(x:A').B \\ \hline \Pi_t(x:A).B \leadsto \Pi_t(x:A).B' \\ \hline \Pi_t\{x:A).B \leadsto \Pi_t(x:A).B' \\ \hline \Pi_t\{x:A).B \leadsto \Pi_t\{x:A).B' \\ \hline \Pi_t\{x:A).B \leadsto \Pi_t\{x:A).B' \\ \hline \Lambda_t(x:A).M \leadsto \Lambda_t(x:A').m \\ \hline \Lambda_t(x:A).M \leadsto \Lambda_t(x:A).M' \\ \hline \Lambda_t(x:A).M \leadsto \Lambda_t(x:A).M \\ \hline \Lambda_t(x:A).M \leadsto \Lambda_t(x:A).M' \\ \hline \Lambda_t(x:$$

Conversion

$$\frac{A \simeq B \qquad B \leadsto C}{A \simeq A} \qquad \frac{A \simeq B \qquad C \leadsto B}{A \simeq C}$$

6 Dynamic Semantics

Value

$$\frac{1}{x \ value} \qquad \frac{1}{\lambda_t(x:A).m \ value} \qquad \frac{1}{\lambda_t(x:A).m \ value} \qquad \frac{u \ value \quad v \ value}{\langle u,v\rangle_t \ value} \qquad \frac{v \ value}{\{v,m\}_t \ value} \qquad \frac{v \ value}{(m,n)_t \ value}$$

Dynamic Reduction

$$\frac{m \rightsquigarrow m'}{m \ n \rightsquigarrow m' \ n} \qquad \frac{n \rightsquigarrow n'}{m \ n \rightsquigarrow m \ n'} \qquad \frac{v \ value}{(\lambda_t(x:A).m) \ v \rightsquigarrow m[v/x]} \qquad \frac{(\lambda_t\{x:A\}.m) \ n \rightsquigarrow m[n/x]}{(\lambda_t\{x:A\}.m) \ n \rightsquigarrow m[n/x]}$$

$$\frac{m \rightsquigarrow m'}{\langle m, n \rangle_t \rightsquigarrow \langle m', n \rangle_t} \qquad \frac{n \rightsquigarrow n'}{\langle m, n \rangle_t \rightsquigarrow \langle m, n' \rangle_t} \qquad \frac{m \rightsquigarrow m'}{\{m, n\}_t \rightsquigarrow \{m', n\}_t} \qquad \frac{R^{\Sigma}_{[z]A}(m, [x, y]n) \rightsquigarrow R^{\Sigma}_{[z]A}(m', [x, y]n)}{R^{\Sigma}_{[z]A}(w, v)_t, [x, y]n) \rightsquigarrow n[u/x, v/y]} \qquad \frac{m \rightsquigarrow m'}{R^{\Sigma}_{[z]A}(\{v, m\}_t, [x, y]n) \rightsquigarrow n[v/x, m/y]} \qquad \frac{m \rightsquigarrow m'}{\pi_1 \ m \rightsquigarrow \pi_1 \ m'}$$

$$\frac{m \rightsquigarrow m'}{\pi_2 \ m \rightsquigarrow \pi_2 \ m'} \qquad \frac{\pi_1 \ (m, n)_t \rightsquigarrow m}{\pi_1 \ (m, n)_t \rightsquigarrow m} \qquad \frac{\pi_2 \ (m, n)_t \rightsquigarrow n}{\pi_2 \ (m, n)_t \rightsquigarrow n}$$

Erasure-Dynamic Reduction

$$\Gamma; \Delta \vdash m \sim m' : A$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

$$\Gamma; \Delta \vdash n \sim n' : A$$