An Intuition for Propagators

George Wilson

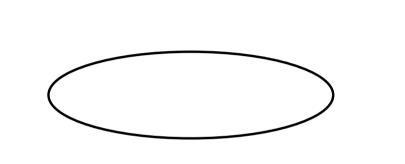
CSIRO's Data61

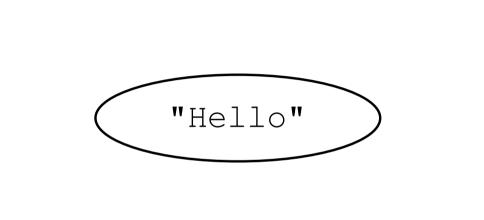
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2nd September 2019



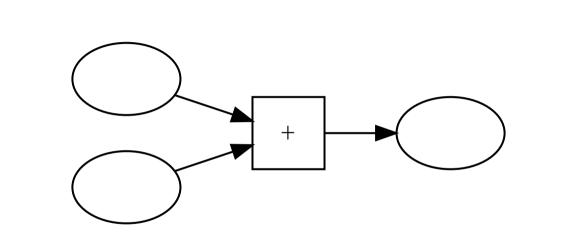
1970s, MIT

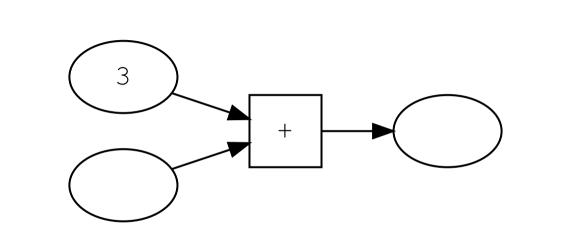


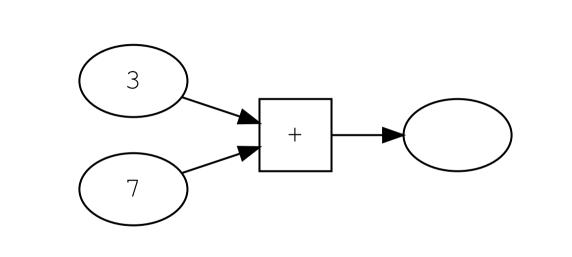


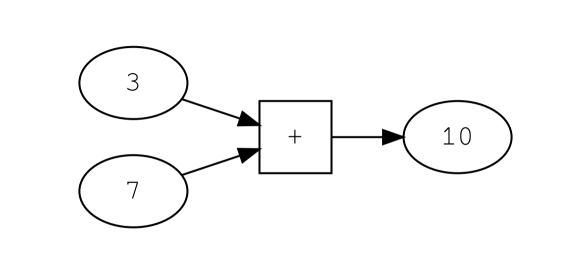


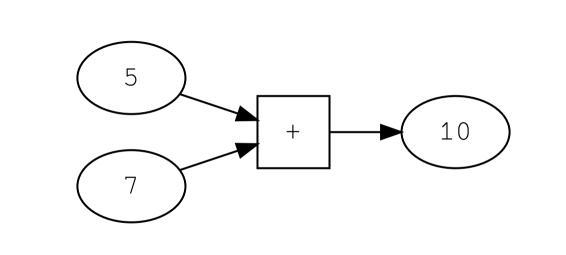
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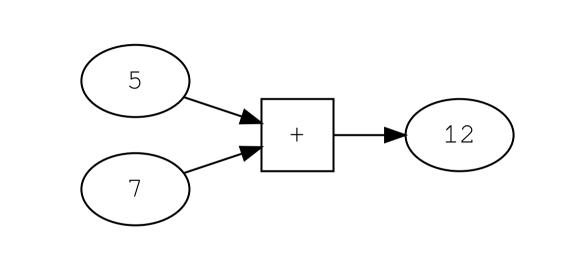












-- types data Cell a

data Par a
instance Monad Par

```
-- types
data Cell a
```

data Par a instance Monad Par

```
-- Creating a cell cell :: Par (Cell a)
```

```
-- types
data Cell a
data Par a
instance Monad Par
-- Creating a cell
cell :: Par (Cell a)
-- Working with Cells
content :: Cell a -> Par (Maybe a)
```

write :: Cell a -> a -> Par ()

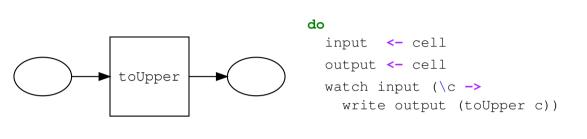
```
-- types
data Cell a
data Par a
instance Monad Par
-- Creating a cell
cell :: Par (Cell a)
-- Working with Cells
content :: Cell a -> Par (Maybe a)
write :: Cell a -> a -> Par ()
-- Creating a propagator
watch :: Cell a -> (a -> Par ()) -> Par ()
```

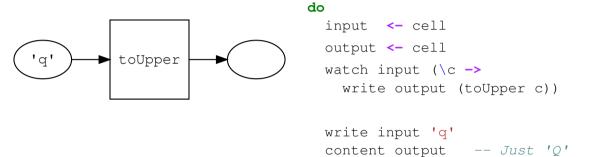
do

input <- cell



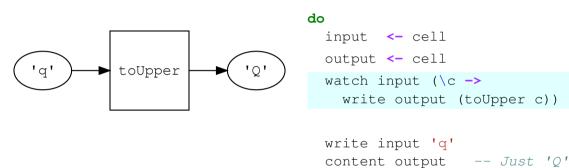






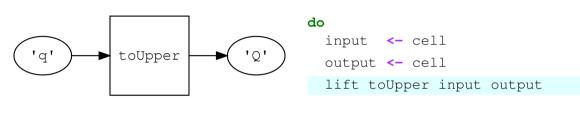
```
input <- cell
output <- cell
output (\c ->
write output (toUpper c))

write input 'q'
content output -- Just 'Q'
```



```
lift :: (a -> b) -> Cell a -> Cell b -> Par ()
lift f input output =
  watch input (\a ->
```

write output (f a))



write input 'q'
content output -- Just 'Q'



```
+
```

inL <- cell
inR <- cell
out <- cell

do

watch inL (\x -> do
 maybeY <- content inR
 case maybeY of
 Nothing -> pure ()
 Just y -> write out (x+y)

```
+
```

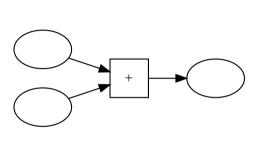
```
inL <- cell
inR <- cell
out <- cell
watch inL (\x -> do
 maybeY <- content inR
 case maybeY of
   Nothing -> pure ()
   Just y -> write out (x+y)
watch inR (\y -> do
 maybeX <- content inL
 case maybeX of
   Nothing -> pure ()
   Just x -> write out (x+v)
```

do

```
with :: Cell a -> (a -> Par ()) -> Par ()
with theCell callback = do
  maybeA <- content theCell</pre>
```

case maybeA of

Nothing -> pure ()
Just a -> callback a

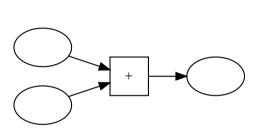


do

inL <- cell
inR <- cell
out <- cell</pre>

watch inL ($x \rightarrow$ with inR ($y \rightarrow$

write out (x+y)



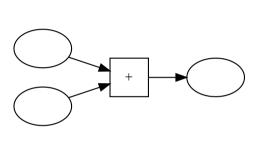
```
inL <- cell
inR <- cell
out <- cell</pre>
```

```
watch inL (\x ->
  with inR (\y ->
    write out (x+y)
```

```
watch inR (\y ->
  with inL (\x ->
  write out (x+y)
```

```
lift2 :: (a -> b -> c)
-> Cell a -> Cell b -> Cell c
-> Par ()
```

```
lift2 :: (a \rightarrow b \rightarrow c)
       -> Cell a -> Cell b -> Cell c
       -> Par ()
lift2 f inL inR out = do
  watch inL (\a ->
    with inR (b \rightarrow
      write out (f a b)))
  watch inR (\b ->
    with inL (\a ->
      write out (f a b)))
```



```
inL <- cell
inR <- cell
out <- cell
adder inL inR out</pre>
```

adder 1 r o = do lift2 (+) 1 r o lift2 (-) o 1 r lift2 (-) o r 1

where





How can we fix this?

```
data WriteOnce a
 = None
   Written a
   TooMany
tryWrite :: a -> WriteOnce a -> WriteOnce a
tryWrite a w = case w of
 None -> Written a
 Written b -> TooMany
 TooMany -> TooMany
```

```
data WriteOnce a
 = None
  | Written a
   TooMany
tryWrite :: (Eq a) => a -> WriteOnce a -> WriteOnce a
tryWrite a w = case w of
 None -> Written a
 Written b -> if a == b then Written b else TooMany
 TooMany -> TooMany
```

Now every network will give a **deterministic answer** in **finite time**

Mutability is **chaos**WriteOnce is **rigid**

Accumulate information about a value

Accumulate information about a value

merge :: a -> a -> a

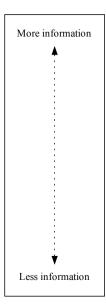
marace milorination about a value

data WriteOnce a

- -- I have heard contradictory answers!
- TooMany
- -- I know the answer exactly
- | **Written** a
- -- I don't know anything
- = None

data WriteOnce a

- -- I have heard contradictory answers!
- = TooMany
- -- I know the answer exactly
- | Written a
- -- I don't know anything
- = None



mergeWrites (Written a) (Written b) =

if a == b then Written a else TooMany

mergeWrites TooMany b = TooMany
mergeWrites a TooMany = TooMany
mergeWrites (Written a) (Written b) =

if a == b then Written a else TooMany

```
mergeWrites :: WriteOnce a -> WriteOnce a -> WriteOnce a
mergeWrites
                          None
                                     = a
                          b
                                     = b
```

mergeWrites None mergeWrites **TooMany** b = TooMany

TooMany = TooMany

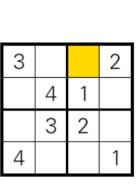
mergeWrites a mergeWrites (Written a) (Written b) =

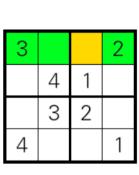
if a == b then Written a else TooMany

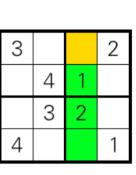
mergeWrites (Written a) (Written b) =

if a == b then Written a else TooMany

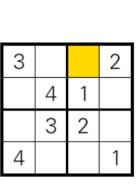
3			2
	4	1	
	Ω	2	
4			1

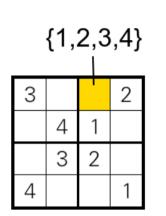


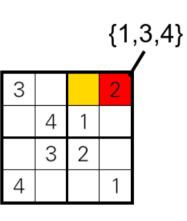


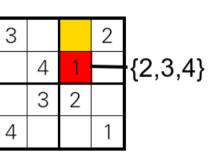


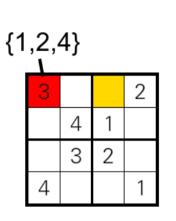
3			2	
	4	1		
	3	2		
4			1	











 $\{2,3,4\} \cap \{1,3,4\} \cap$

2 2 4 1

3		4	2
	4	1	
	$_{\odot}$	2	
4	·		1

The accumulation must:

- tolerate **reordering** of information
- tolerate **grouping** of information
- ignore **redundancy** of information

Join semilattice

Commutative:

$$x\vee y=y\vee x$$

Associative:

$$x \vee (y \vee z) = (x \vee y) \vee z$$

Idempotent:

$$x \lor x = x$$

```
class Semilattice a where
  (\/) :: a -> a -> a
```

bottom :: a

instance (Eq a) => SemiLattice (WriteOnce a) where

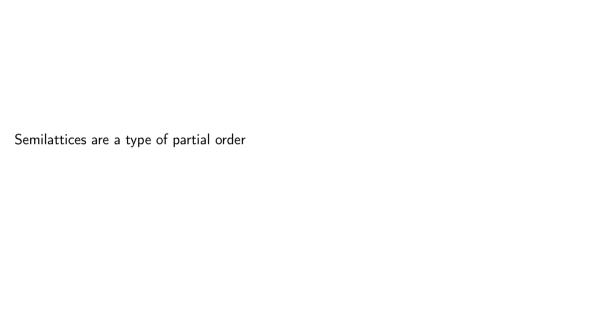
None \/ b **=** b

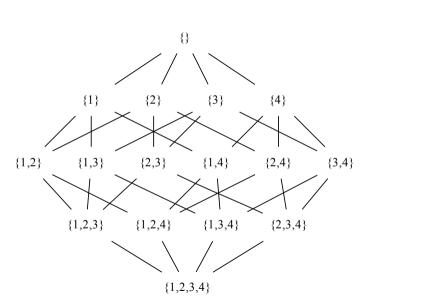
TooMany \/ x = TooMany Written a \/ None = Written a

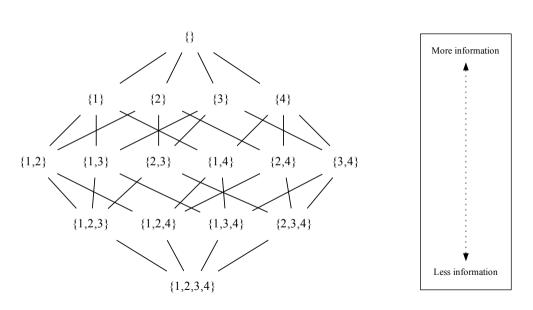
Written a \/ Written b = if a == b then Written a else TooMany Written a \/ TooMany = TooMany

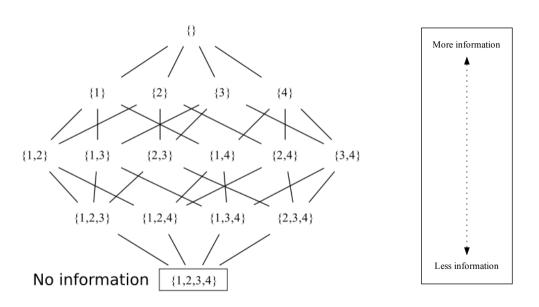
```
instance (Ord a) => Semilattice (Set a) where
p \/ q = Set.intersection p q
```

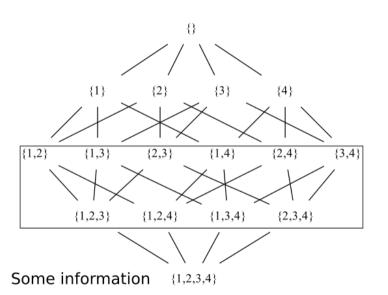
bottom = Set.empty

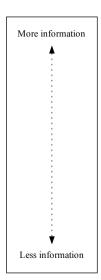


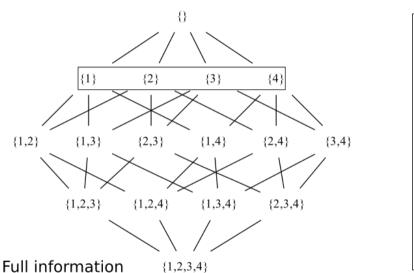




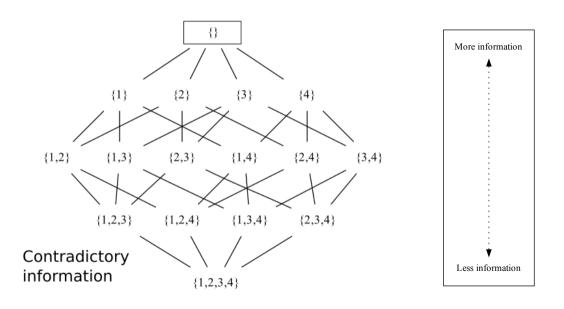


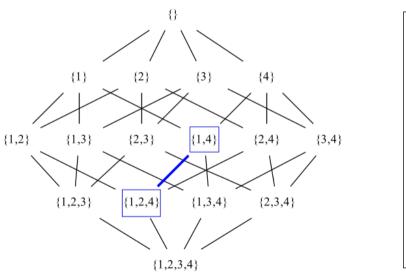






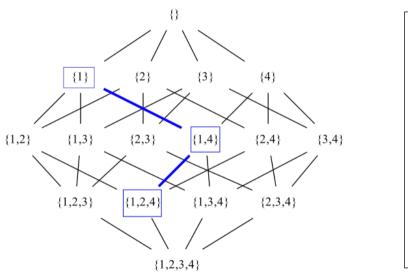






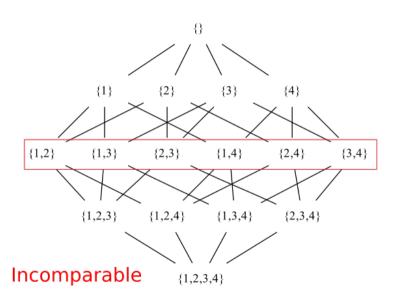
More information Less information

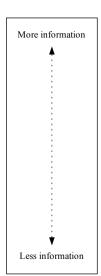
 $\{1,2,4\} < \{1,4\}$

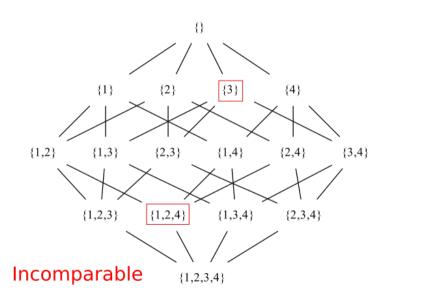


More information Less information

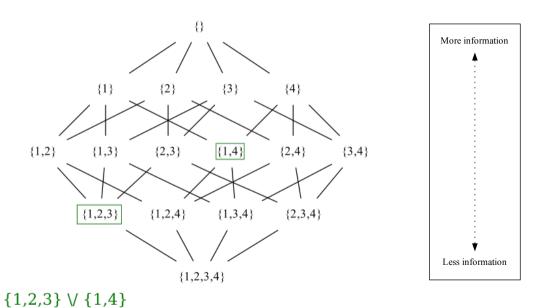
 $\{1,2,4\} < \{1,4\} < \{1\}$

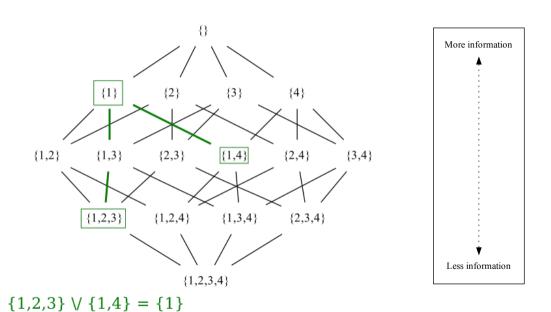












Monotonicity

 $x \le y \implies f(x) \le f(y)$

Thanks for listening!