ITCT Homework 2

● R07944007網媒碩一林良翰

1. [Prob. 3.1.]

a. Markov's Inequality

• By the definition of probability

$$E[X] = \int_{-\infty}^{\infty} x P(x) dx$$

• Since r.v. *X* is non-negative

$$E[X] = \int_0^\infty x P(x) dx$$

• By the property of integrals

$$E\left[X
ight] = \int_{0}^{t}xP\left(x
ight)dx + \int_{t}^{\infty}xP\left(x
ight)dx \geq \int_{t}^{\infty}xP\left(x
ight)dx \geq \int_{t}^{\infty}tP\left(x
ight)dx = tP\left(X\geq t
ight)$$

Thus

$$P(X \ge t) \le \frac{E[X]}{t}$$

b. Chebyshev's Inequality

• By the definition of mean and variance

$$\circ$$
 Mean : $\mu = E[Y]$

• Variance :
$$\sigma^2 = E\left[(Y - E[Y])^2 \right] = E\left[(Y - \mu)^2 \right] = E[X]$$

By Markov's enequality

$$P\left(X \geq \epsilon^2\right) \leq rac{E[X]}{\epsilon^2}$$

Thus

$$P\left(|Y-\mu| \geq \epsilon
ight) = P\left(\left(Y-\mu
ight)^2 \geq \epsilon^2
ight) \leq rac{E\left[\left(Y-\mu
ight)^2
ight]}{\epsilon^2} = rac{\sigma^2}{\epsilon^2}$$

c. Weak Law of Large Numbers

Chebyshev's inequality

A r.v.
$$X$$
 with its mean $E\left(X\right)=\mu$ and variance $Var\left(X\right)=\sigma^{2}$, for any $\epsilon\geq0$

$$P(|X - \mu| \ge \epsilon) \le \frac{\sigma^2}{\epsilon^2}$$

which is same as

$$P(|X - E(X)| \ge \epsilon) \le \frac{Var(X)}{\epsilon^2}$$

ullet Find $E\left(ar{Z}_n
ight)$ and $Var\left(ar{Z}_n
ight)$

$$\begin{array}{l} \circ \ E\left(\bar{Z}_{n}\right) = E\left(\frac{1}{n}\sum_{i=1}^{n}Z_{i}\right) = \frac{1}{n}E\left(\sum_{i=1}^{n}Z_{i}\right) = \frac{1}{n}\sum_{i=1}^{n}E\left(Z_{i}\right) = \frac{1}{n}(n\mu) = \mu \\ \circ \ Var\left(\bar{Z}_{n}\right) = Var\left(\frac{1}{n}\sum_{i=1}^{n}Z_{i}\right) = \frac{1}{n^{2}}Var\left(\sum_{i=1}^{n}Z_{i}\right) = \frac{1}{n^{2}}\sum_{i=1}^{n}Var\left(Z_{i}\right) = \frac{1}{n^{2}}\left(n\sigma^{2}\right) = \frac{\sigma^{2}}{n} \end{array}$$

ullet Thus, by Chebyshev's inequality and the mean, variance of $ar{Z}_n$, for any $\epsilon \geq 0$

$$P\left(\left|\bar{Z}_{n} - E\left(\bar{Z}_{n}\right)\right| \geq \epsilon\right) \leq rac{Var(\bar{Z}_{n})}{\epsilon^{2}}$$

 $\Rightarrow P\left(\left|\bar{Z}_{n} - \mu\right| \geq \epsilon\right) \leq rac{\sigma^{2}}{n\epsilon^{2}}$

2. [Prob. 3.4.]

a.

• By AEP and the property of typical set $A_{\epsilon}^{(n)}$ with respect to $P\left(X\right)$ which is the set of sequence $(X_1,X_2,\ldots,X_n)\in X^n$

$$2^{-n(H(X)+\epsilon)} \leq P\left(X_1,X_2,\ldots,X_n
ight) \leq 2^{-n(H(X)-\epsilon)}$$

• Derive A^n

$$egin{aligned} A^n &= \left\{ x^n \in X^n : \left| -rac{1}{n} \log P\left(x^n
ight) - H\left(X
ight)
ight| \leq \epsilon
ight\} \ &= \left\{ x^n \in X^n : -\epsilon \leq -rac{1}{n} \log P\left(x^n
ight) - H\left(X
ight) \leq \epsilon
ight\} \ &= \left\{ x^n \in X^n : 2^{-n(H(X) + \epsilon)} \leq P\left(x^n
ight) \leq 2^{-n(H(X) - \epsilon)}
ight\} \ &= A_{\epsilon}^{(n)} \end{aligned}$$

 ${\cal A}^n$ is just the same as typical set

$$P\left\{X^n\in A^n\right\}=P\left\{X^n\in A_\epsilon^{(n)}\right\}\to 1-\epsilon, \forall \epsilon>0 \text{ as } n\to\infty$$
 Thus, $P\left\{X^n\in A^n\right\}\to 1$

b.

• By the strong law of large numbers

$$P\left(\lim_{n\to\infty} \frac{1}{n} \sum_{i=1}^{n} X_i = \mu\right) = 1$$

which can be rewrite as

$$orall \epsilon>0$$
 , $P\left(\left|\left(rac{1}{n}\sum_{i=1}^nX_i
ight)-\mu
ight|\leq\epsilon
ight) o 1$, as $n o\infty$

$$ullet \ B^n = \left\{ x^n \in X^n : \left| \left(rac{1}{n} \sum_{i=1}^n X_i
ight) - \mu
ight| \leq \epsilon
ight\}$$

$$P\left\{X^n\in B^n\right\} o 1$$

• By $P\left\{X^n\in A^n\right\} o 1$, there exists $\epsilon>0$ and N_1 such that $P\left\{X^n\in A^n\right\}>1-\frac{\epsilon}{2}$ for all $n>N_1$ By $P\left\{X^n\in B^n\right\} o 1$, there exists $\epsilon>0$ and N_2 such that $P\left\{X^n\in B^n\right\}>1-\frac{\epsilon}{2}$ for all $n>N_2$

ullet So, for all $n>N=\max{(N_1,N_2)}$ and for any $\epsilon>0$

$$P\{X^n \in A^n \cap B^n\} = P\{X^n \in A^n\} + P\{X^n \in B^n\} - P\{X^n \in A^n \cup B^n\}$$
$$> 1 - \frac{\epsilon}{2} + 1 - \frac{\epsilon}{2} - 1$$
$$= 1 - \epsilon$$

Thus

$$P\left\{X^n\in A^n\cap B^n
ight\}
ightarrow 1-\epsilon$$
 , for all $n>N$ and for any $\epsilon>0$

C.

• By the property of probability

$$\sum_{x^{n}\in A^{n}\cap B^{n}}P\left(x^{n}\right) \leq1$$

ullet By the property of typical set, for $x^n \in A^n$

$$P(x^n) > 2^{-n(H(X)+\epsilon)}$$

• Combine above two inequations

$$1 \geq \sum\limits_{x^n \in A^n \cap B^n} P\left(x^n
ight) \geq \sum\limits_{x^n \in A^n \cap B^n} 2^{-n(H(X) + \epsilon)} \geq \left|A^n \cap B^n
ight| 2^{-n(H(X) + \epsilon)}$$

Therefore

$$|A^n \cap B^n| \leq 2^{n(H(X)+\epsilon)}$$

d.

• From b.

$$P\left\{X^n\in A^n\cap B^n
ight\}
ightarrow 1-\epsilon$$
 , for all $n>\max\left(N_1,N_2
ight)$ and for any $\epsilon>0$

Thus, there exists a number N such that $P\left\{X^n\in A^n\cap B^n
ight\}\geq rac{1}{2}$, for all n>N

ullet By the property of typical set, for $x^n \in A^n$

$$P(x^n) \leq 2^{-n(H(X)-\epsilon)}$$

• Combine above two inequations

$$\textstyle \frac{1}{2} \leq P\left\{X^n \in A^n \cap B^n\right\} = \sum\limits_{x^n \in A^n \cap B^n} P\left(x^n\right) \leq \sum\limits_{x^n \in A^n \cap B^n} 2^{-n(H(X) - \epsilon)} = \left|A^n \cap B^n\right| 2^{-n(H(X) - \epsilon)}$$

Therefore

$$|A^n \cap B^n| \ge \left(\frac{1}{2}\right) 2^{n(H(X) - \epsilon)}$$

3. [Prob. 3.10.]

• Define the notataions

$$\log x = \log_2 x$$

$$\ln x = \log_e x$$

ullet The volume $V_n = \prod\limits_{i=1}^n X_i$

We can derive

$$\log\left(V_n
ight)^{rac{1}{n}} = rac{1}{n}\mathrm{log}\left(V_n
ight) = rac{1}{n}\sum_{i=1}^{n}\mathrm{log}\left(X_i
ight)$$

ullet By the strong law of large numbers, and X is uniform over $\left[0,1\right]$

$$rac{1}{n}\sum_{i=1}^{n}\log\left(X_{i}
ight)
ightarrow E\left[\log\left(X
ight)
ight] = \int_{0}^{1}\log\left(x
ight)dx = -rac{1}{\ln2}$$

• Finally we can derive $\lim_{n \to \infty} (V_n)^{\frac{1}{n}}$

$$\lim_{n o\infty}(V_n)^{rac{1}{n}}=e^{\lim_{n o\infty}\ln{(V_n)^{rac{1}{n}}}}=e^{\ln{2\left(\lim_{n o\infty}\log{(V_n)^{rac{1}{n}}}
ight)}}=e^{\ln{2\left(-rac{1}{\ln{2}}
ight)}}=rac{1}{e}$$

$$\bullet \ \left(E\left[V_n\right]\right)^{\frac{1}{n}} = \left(E\left[\prod_{i=1}^n X_i\right]\right)^{\frac{1}{n}} = \left(\prod_{i=1}^n E\left[X_i\right]\right)^{\frac{1}{n}} = \left(\left(\frac{1}{2}\right)^n\right)^{\frac{1}{n}} = \frac{1}{2} \neq \frac{1}{e}$$

Thus, the expected edge length does not capture the idea of the volume of the box.

4. [Prob. 3.11.]

a.

• By definition and the property of probability

$$P(A) > 1 - \epsilon_1$$

$$P(B) > 1 - \epsilon_2$$

$$P(A \cup B) \le 1 \Rightarrow -P(A \cup B) \ge -1$$

ullet We can obtain the probablility of the intersection of A and B

$$P(A \cap B) = P(A) + P(B) - P(A \cup B) \ge 1 - \epsilon_1 + 1 - \epsilon_2 - 1 = 1 - \epsilon_1 - \epsilon_2$$

b.

• (a) from a.

$$1 - \delta - \epsilon \le P\left(A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)}\right)$$

• (b) by the definition of probability of a set

$$P\left({A_{\epsilon}}^{(n)}\cap{B_{\delta}}^{(n)}
ight) = \sum\limits_{x^n\in{A_{\epsilon}}^{(n)}\cap{B_{\delta}}^{(n)}}P\left(x^n
ight)$$

ullet (c) by the property of typical set that $x^n \in A_{\epsilon}^{\;(n)}$

 $P\left(x^{n}\right)$ has upper bound $2^{-n(H-\epsilon)}$

$$\textstyle \sum_{x^n \in A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)}} P\left(x^n\right) \leq \sum_{x^n \in A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)}} 2^{-n(H-\epsilon)}$$

• (d) by the definition of the cardinality of a set

$$\sum_{x^n \in A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)}} 2^{-n(H-\epsilon)} = \left| A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)} \right| 2^{-n(H-\epsilon)}$$

ullet (e) by the property that ${A_\epsilon}^{(n)}\cap {B_\delta}^{(n)}\subseteq {B_\delta}^{(n)}$

$$\left|A_{\epsilon}^{\,(n)}\cap B_{\delta}^{\,(n)}
ight|2^{-n(H-\epsilon)}\leq \left|B_{\delta}^{\,(n)}
ight|2^{-n(H-\epsilon)}$$

C.

• We apply similar steps in b.

$$egin{aligned} 1-\delta & \leq P\left(B_{\delta}^{(n)}
ight) \ & = \sum_{x^n \in B_{\delta}^{(n)}} P\left(x^n
ight) \ & \leq \sum_{x^n \in B_{\delta}^{(n)}} 2^{-n(H-\delta)} \ & = \left|B_{\delta}^{(n)}
ight| 2^{-n(H-\delta)} \end{aligned}$$

• Let $\delta' = \delta - \frac{1}{n} \log (1 - \delta) > 0$

Finally we can get

$$\frac{1}{n}\log\left|B_{\delta}^{(n)}\right| \geq H - \delta'$$

5. [Prob. 4.6.]

a.

• By the chain rule for entropy

$$\frac{1}{n}H(X_1, X_2, \dots, X_n) = \frac{1}{n} \sum_{i=1}^n H(X_i | X_{i-1}, \dots, X_1)$$

$$= \frac{1}{n} \left(H(X_n | X_{n-1}, \dots, X_1) + \sum_{i=1}^{n-1} H(X_i | X_{i-1}, \dots, X_1) \right)$$

$$= \frac{1}{n} (H(X_n | X_{n-1}, \dots, X_1) + H(X_1, X_2, \dots, X_{n-1}))$$

• From stationarity, for all $1 \le i \le n$

$$H(X_n|X_{n-1},...,X_1) \leq H(X_i|X_{i-1},...,X_1)$$

We can derive

$$H\left(X_{n}|X_{n-1},\ldots,X_{1}
ight) \leq rac{1}{n-1}\sum_{i=1}^{n-1}H\left(X_{i}|X_{i-1},\ldots,X_{1}
ight) = rac{1}{n-1}H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)$$

• Finally we can get

$$rac{1}{n}H\left(X_{1},X_{2},\ldots,X_{n}
ight)\leqrac{1}{n}\Big(rac{1}{n-1}H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)+H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)\Big)=rac{1}{n-1}H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)$$

b.

ullet By the stationarity, for all $1 \leq i \leq n$

$$H(X_n|X_{n-1},...,X_1) \leq H(X_i|X_{i-1},...,X_1)$$

• Thus we can derive

$$egin{aligned} H\left(X_{n}|X_{n-1},\ldots,X_{1}
ight) &= rac{1}{n}\sum_{i=1}^{n}H\left(X_{n}|X_{n-1},\ldots,X_{1}
ight) \ &\leq rac{1}{n}\sum_{i=1}^{n}H\left(X_{i}|X_{i-1},\ldots,X_{1}
ight) \ &= rac{1}{n}H\left(X_{1},X_{2},\ldots,X_{n}
ight) \end{aligned}$$

6. [Prob. 4.8.]

ullet By the definition of entropy and $p_1+p_2=1$

$$H(X) = -p_1 \log p_1 - p_2 \log p_2 = -p_1 \log p_1 - (1 - p_1) \log (1 - p_1)$$

• Define r.v. $T=\{t_1,t_2\}$ is the symbol duration of r.v. X

$$E[T] = p_1 t_1 + p_2 t_2 = p_1 + 2p_2 = 2 - p_1$$

• Define $f(p_1)$ is the source entropy per unit time controlled by p_1

$$f\left(p_{1}
ight)=rac{H(X)}{E[T]}=rac{-p_{1}\log p_{1}-(1-p_{1})\log \left(1-p_{1}
ight)}{2-p_{1}}$$

• Since f(0) = f(1) = 0, the maximum point occurs between $0 < p_1 < 1$, thus we take 1st derivative of f and set it to 0 to find optimal p_1^\star such that $f(p_1^\star)$ has maximum

$$rac{d}{dp_1}\Big(rac{-p_1\log p_1-(1-p_1)\log\left(1-p_1
ight)}{2-p_1}\Big)=rac{1}{\left(2-p_1
ight)^2}[\ln\left(1-p_1
ight)-2\ln p_1]=0$$

We can get

$$1 - p_1 = (p_1)^2 \Rightarrow p_1^\star = -\frac{1}{2} \pm \frac{\sqrt{5}}{2}$$

Since $0 < p_1 < 1$

$$p_1^\star = -rac{1}{2} + rac{\sqrt{5}}{2}$$

Organize the equation

$$\begin{split} f\left(p_{1}\right) &= \frac{-p_{1} \log p_{1} - \left(1 - p_{1}\right) \log \left(1 - p_{1}\right)}{2 - p_{1}} \\ &= \frac{-p_{1} \log p_{1} - \left(1 - p_{1}\right) \log \left(1 - p_{1}\right)}{1 + \left(1 - p_{1}\right)} \\ &= \frac{-p_{1} \log p_{1} - \left(p_{1}\right)^{2} \log \left(p_{1}\right)^{2}}{1 + \left(p_{1}\right)^{2}} \\ &= \frac{-\left(1 - \left(p_{1}\right)^{2}\right) \log p_{1} - 2\left(p_{1}\right)^{2} \log p_{1}}{1 + \left(p_{1}\right)^{2}} \\ &= \frac{-\left(1 + \left(p_{1}\right)^{2}\right) \log p_{1}}{1 + \left(p_{1}\right)^{2}} \\ &= -\log p_{1} \end{split}$$

Finally we can get

$$-\log p_1^* = -\log\left(-rac{1}{2} + rac{\sqrt{5}}{2}
ight) pprox 0.6942$$

7. [Prob. 4.12.]

a.

• The dog's walk is 2nd order Markov, because we need the previous 2 position to know the previous position and direction.

• By the chain rule of entropy

$$H\left(X_{0},X_{1},\ldots,X_{n}
ight)=\sum\limits_{i=0}^{n}H\left(X_{i}|X_{i-1},\ldots,X_{0}
ight)=H\left(X_{0}
ight)+H\left(X_{1}|X_{0}
ight)+\sum\limits_{i=2}^{n}H\left(X_{i}|X_{i-1},\ldots,X_{0}
ight)$$

• The initial position is deterministic $\Rightarrow H\left(X_{0}\right)=0$

The first move is equally likely to positive or negative $\Rightarrow H(X_1|X_0) = H(\frac{1}{2},\frac{1}{2}) = 1$

The rest of the move follow the rule defined by the problem

$$\sum_{i=2}^{n} H\left(X_{i} | X_{i-1}, \dots, X_{0}
ight) = \sum_{i=2}^{n} H\left(rac{1}{10}, rac{9}{10}
ight) = (n-1) \, H\left(rac{1}{10}, rac{9}{10}
ight)$$

• Thus we have

$$H(X_0, X_1, \dots, X_n) = 1 + (n-1) H(\frac{1}{10}, \frac{9}{10})$$

b.

• The entropy rate

$$rac{H(X_0,X_1,\ldots,X_n)}{n+1}=rac{1+(n-1)H\left(rac{1}{10},rac{9}{10}
ight)}{n+1}
ightarrow H\left(rac{1}{10},rac{9}{10}
ight)$$
 as $n
ightarrow\infty$

c.?

ullet Let r.v. S be the number of steps taken between reversals, we have

$$E(S) = \sum_{s=1}^{\infty} s(\frac{9}{10})^{s-1} \frac{1}{10} = 10$$

• Thus the expected number of steps the dog takes before reversing direction is

$$10 + 1 = 11$$

8. [Prob. 4.19.]

• Define r.v. $V_n \in \{1, 2, 3, 4, 5\}$ is the vertex position at time n

• Define r.v.
$$X_n = \begin{bmatrix} P\left(V_n=1\right) \\ P\left(V_n=2\right) \\ P\left(V_n=3\right) \\ P\left(V_n=4\right) \\ P\left(V_n=5\right) \end{bmatrix}$$
 is the vertex position distribution at time n

• The state transition matrix *p*

$$P_{ij} = P\left(V_{n+1} = j | V_n = i
ight) = egin{bmatrix} 0 & rac{1}{3} & 0 & rac{1}{3} & rac{1}{3} \ rac{1}{3} & 0 & rac{1}{3} & 0 & rac{1}{3} \ 0 & rac{1}{3} & 0 & rac{1}{3} & rac{1}{3} \ rac{1}{3} & 0 & rac{1}{3} & rac{1}{3} \ rac{1}{4} & rac{1}{4} & rac{1}{4} & rac{1}{4} & 0 \ \end{bmatrix}$$

a.

• The stationary distribution of vertex position X_n^*

$$X_n^* = \lim_{n o \infty} X_n = \left[rac{3}{16} \quad rac{3}{16} \quad rac{3}{16} \quad rac{4}{16}
ight]^T$$
 such that $X_n^* = P_{ij} \cdot X_n^*$

b.

• By the definition of entropy rate

$$H\left(X
ight) = \lim_{n o \infty} rac{1}{n} H\left(X_1, X_2, \dots, X_n
ight)$$

ullet Since X_n is a stationary process and it's also a Markov chain

$$H\left(X
ight)=H'\left(X
ight)=\lim_{n
ightarrow\infty}H\left(X_{n}|X_{n-1},\ldots,X_{0}
ight)=H\left(X_{1}|X_{0}
ight)$$

$$H\left(X_{1}|X_{0}
ight) = -\sum_{i=1}^{5} \left(P\left(V_{0}=i
ight)\sum_{j=1}^{5} P\left(V_{1}=j|V_{0}=i
ight) \log P\left(V_{1}=j|V_{0}=i
ight)
ight) = 4 \cdot rac{3}{16} \log 3 + rac{4}{16} \log 4$$

C.

• By the definition of mutual information

$$\begin{split} I\left(X_{n+1}; X_n\right) &= H\left(X_{n+1}\right) - H\left(X_{n+1} | X_n\right) \\ &= H\left(X_{n+1}\right) - H\left(X_1 | X_0\right) \\ &= \left[4 \cdot \frac{3}{16} \log \frac{16}{3} + \frac{4}{16} \log \frac{16}{4}\right] - \left[4 \cdot \frac{3}{16} \log 3 + \frac{4}{16} \log 4\right] \\ &= \frac{3}{4} \log \frac{16}{9} \end{split}$$

9.

• Define the notation

$$X \in \{a, b, c\}$$

$$Y \in \{a',b',c'\}$$

and we have the following equation

$$\bar{P} + 2\bar{Q} = 1$$

• By the definition of channel capacity

$$C = \max_{P_{Y}} I\left(X; Y\right)$$

• Find the formula of mutual information

$$\begin{split} I\left(X;Y\right) &= H\left(X\right) - H\left(Y|X\right) \\ &= -\bar{P}\ln\bar{P} - 2\bar{Q}\ln\bar{Q} - \left[-\sum_{x \in X} \sum_{y \in Y} P\left(X = x, Y = y\right) \ln P\left(X = x|Y = y\right) \right] \\ &= -\bar{P}\ln\bar{P} - 2\bar{Q}\ln\bar{Q} + 2\bar{Q}q\ln q + 2\bar{Q}\left(1 - q\right) \ln\left(1 - q\right) \\ &= -\left(1 - 2\bar{Q}\right) \ln\left(1 - 2\bar{Q}\right) - 2\bar{Q}\ln\bar{Q} + 2\bar{Q}q\ln q + 2\bar{Q}\left(1 - q\right) \ln\left(1 - q\right) \end{split}$$

ullet To find the maximum value, we take 1st derivative of mutual information with $ar{Q}$ and set it to 0

$$\begin{split} \frac{dI\left(X;Y\right)}{d\bar{Q}} &= \frac{d\left[-\left(1-2\bar{Q}\right)\ln\left(1-2\bar{Q}\right)-2\bar{Q}\ln\bar{Q}+2\bar{Q}q\ln q+2\bar{Q}\left(1-q\right)\ln\left(1-q\right)\right]}{d\bar{Q}} \\ &= 2\ln\left(1-2\bar{Q}\right)-2\ln\bar{Q}+2\left[q\ln q+\left(1-q\right)\ln\left(1-q\right)\right]=0 \end{split}$$
 Let $\alpha = q^q(1-q)^{(1-q)}$

The optimal $ar{Q}^* = rac{lpha}{1+2lpha}$

ullet Set $ar Q=ar Q^*$

$$\begin{split} I_{\bar{Q}^*}\left(X;Y\right) &= -\left(1-2\bar{Q}^*\right)\ln\left(1-2\bar{Q}^*\right) - 2\bar{Q}^*\ln\bar{Q}^* + 2\bar{Q}^*q\ln q + 2\bar{Q}^*\left(1-q\right)\ln\left(1-q\right) \\ &= -\left(1-2\frac{\alpha}{1+2\alpha}\right)\ln\left(1-2\frac{\alpha}{1+2\alpha}\right) - 2\frac{\alpha}{1+2\alpha}\ln\frac{\alpha}{1+2\alpha} + 2\frac{\alpha}{1+2\alpha}q\ln q + 2\frac{\alpha}{1+2\alpha}(1-q)\ln\left(1-q\right) \\ &= \ln\left(1+2\alpha\right) \\ &= \ln\left(1+2q^q(1-q)^{(1-q)}\right) \end{split}$$

10.

a.

$$\begin{array}{l} \bullet \quad P\left(y_{0}\right) = P\left(x_{0}, y_{0}\right) + P\left(x_{1}, y_{0}\right) = \frac{1}{2} \frac{99}{100} + \frac{1}{2} \frac{1}{100} = \frac{1}{2} \\ \bullet \quad P\left(y_{1}\right) = P\left(x_{0}, y_{1}\right) + P\left(x_{1}, y_{1}\right) = \frac{1}{2} \frac{1}{100} + \frac{1}{2} \frac{99}{100} = \frac{1}{2} \end{array}$$

b.

•
$$H(Y) = \frac{1}{2}\log 2 + \frac{1}{2}\log 2 = 1$$

C.

• By the definition of mutual information

$$\begin{split} I\left(X;Y\right) &= H\left(X\right) - H\left(Y|X\right) \\ &= 1 - \left[\frac{1}{2}\frac{99}{100}\log\frac{100}{99} + \frac{1}{2}\frac{1}{100}\log\frac{100}{1} + \frac{1}{2}\frac{99}{100}\log\frac{100}{99} + \frac{1}{2}\frac{1}{100}\log\frac{100}{1}\right] \\ &= 1 - \frac{99}{100}\log\frac{100}{99} - \frac{1}{100}\log\frac{100}{1} \\ &\approx 0.9192 \end{split}$$