# Optimistic Concurrency Control in a Distributed NameNode Architecture for Hadoop Distributed File System

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### Motivation

#### Industrial Standard in Big Data Era

Apache Hadoop Ecosystem

#### Limits to growth in HDFS

Number of Files	Memory Requirement	Physical Storage	
1 million	0.6 GB	0.6 PB	
100 million	60 GB	60 PB	
1 billion	600 GB	600 PB	

#### Hop-HDFS and Its Limitation

Distributed NameNode Architecture Maintain HDFS Strong Consistency Semantics Concurrency Restricted



### Problem Statement

#### **HDFS**

System-level Lock

#### Hop-HDFS v1

System-level Lock

#### Hop-HDFS v2

Row-level Lock

#### MySQL Cluster

Read Committed / Anomalies

#### Contribution

#### Architectures and Namespace Concurrency Control

GFS, HDFS, Hop-HDFS and MySQL Cluster

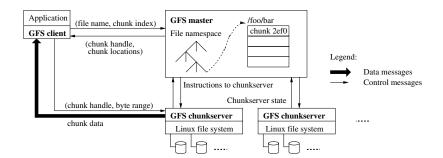
#### Performance Accessment and Limitation Analysis

HDFS v.s. Hop-HDFS v2 (PCC version)

#### Solution for Hop-HDFS

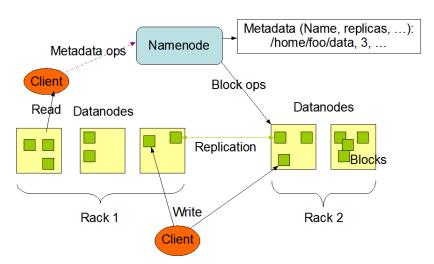
Optimistic Concurrency Control with Snapshot Isolation on Semantic Related Group

### **GFS** Architecture



### **HDFS** Architecture

#### **HDFS Architecture**



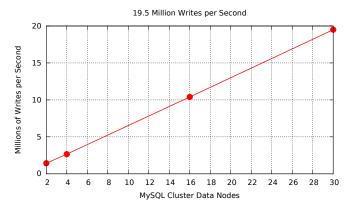
### **Isolation Level**

Berenson, Hal, et al. "A Critique of ANSI SQL Isolation Levels." ACM SIGMOD Record 24.2 (1995): 1-10.

Isolation Level	Lost	Fuzzy	Phantom	Read	Write
	Up-	Read		Skew	Skew
	date				
Read Uncommitted	✓	✓	<b>√</b>	<b>√</b>	<b>√</b>
Read Committed	✓	✓	✓	✓	<b>√</b>
Cursor Stability	some-	some-	<b>√</b>	✓	some-
	times	times			times
Repeatable Read	Χ	Χ	✓	Χ	Χ
Snapshot	Χ	X	sometimes	Χ	<b>√</b>
Serializable	Χ	X	Χ	Χ	Χ

# MySQL Cluster

- Distributed, in-memory, replicated database
- Supports only Read Committed
- High throughput:

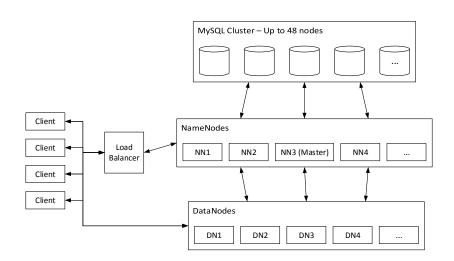


# Hop-HDFS

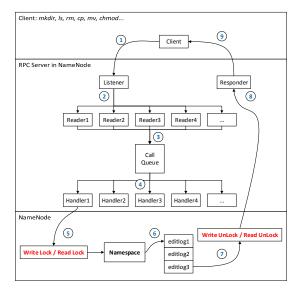
#### Overcome Limitations in HDFS NameNode

- Scalability of the Namespace
- Throughput Problem
- Failure Recovery

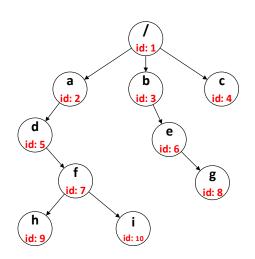
# Hop-HDFS Architecture



# Limitations in HDFS Namespace Concurrency Control



# **HDFS Namespace Structure**



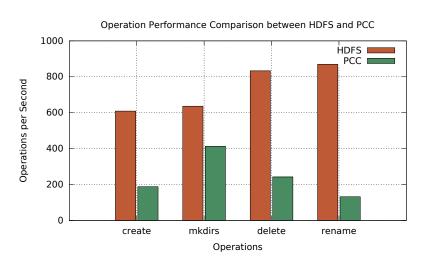
id	parent_id	name
1	0	/
2	1	a
3	1	b
4	1	С
5	2	d
6	3	е
7	5	f
8	6	g
9	7	h
10	7	i

# Limitations in Hop-HDFS Namespace Concurrency Control

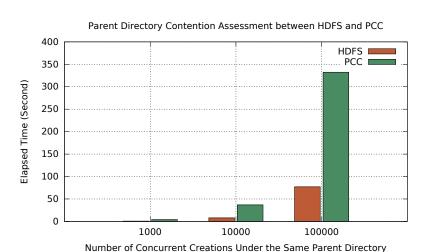
- Duplicated Round Trips
- Implicit Parent Locks:

id	parent_id	name	Locks by Tx1	Locks by Tx2
1	0	/	R	R
2	1	а	R	R
3	1	b		
4	1	С		
5	2	d	R	R
6	3	е		
7	5	f	W	W (Block)
8	6	g		
9	7	h (Tx1)	W (Implicit)	W (Implicit) (Block)
10	7	i (Tx2)	W (Implicit)	W (Implicit) (Block)

### NameNode Throughput Benchmark



# Parent Directory Contention Assessment





# Resolving the Semantic Related Group

Path: /a/d/f/h

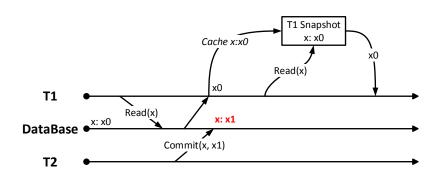
h: 
$${/->a->d->f}$$

	id	parent_id	name	other parameters
Related *	1	0	/	
Related *	2	1	a	
	3	1	b	
	4	1	С	
Related *	5	2	d	
	6	3	е	
Related *	7	5	f	
	8	6	g	
Selected ✓	9	7	h	
	10	7	i	

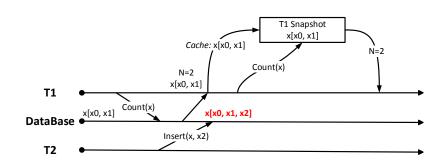
# Per-Transaction Snapshot Isolation

- Snapshot the whole Semantic Related Group
- Transaction performs on its own snapshot
- Preclude: Fuzzy Read & Phantom Read

# Snapshot Isolation Precludes Fuzzy Read



# Snapshot Isolation with Semantic Related Group Precludes Phantom Read



# Lock Mode in MySQL Cluster

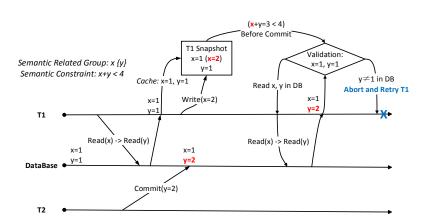
### Read\_Committed: Consistent nonlocking reads (based on MVCC)

Lock Type	Shared	Exclusive	Read_Committed
Shared	✓	Block	✓
Exclusive	Block	Block	✓
Read_Committed	✓	✓	✓

# Per-Transaction Snapshot Isolation

- Read Phase: Read\_Committed on snapshot
- Validation Phase: Shared on related Rows, Exclusive on modified rows / Compare versions / Abort and retry
- Preclude: Write Skew

# OCC with Snapshot Isolation on Semantic Related Group Precludes Write Skew



### Per-Transaction Snapshot Isolation

- Total order update modified rows by ids
- Abort and retry transactions if "new" rows already exists
- Increase versions for successful update

# Four Phases in Algorithm

- Read Phase
- Execution Phase
- Validation Phase
- Update Phase

# Experimental Testbed

#### MySQL Cluster

6 data nodes, 1 Gbps LAN, Intel Xeon X5660 CPU @ 2.80GHz, 6\*6=36 GB RAM, 2 data replicas

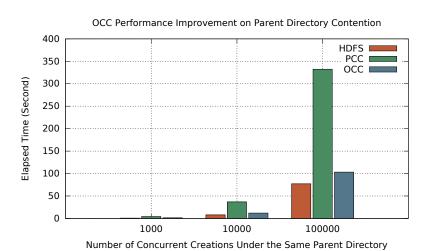
#### NameNode and Clients

Intel i7-4770T CPU at 2.50GHz and 16 GB RAM

#### MySQL Cluster and NameNode

100 Mbps LAN

### Parent Directory Contention Assessment

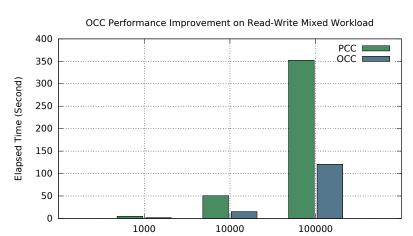




# Parent Directory Contention Assessment

Num. of Concurrent Creation	1000	10000	100000
HDFS	0.82s	7.83s	77.13s
PCC	4.35s	36.74s	332.36s
OCC	1.36s	12.01s	103.23s
PCC / HDFS	530.5%	469.2%	430.9%
OCC / HDFS	165.9%	153.4%	133.8%
OCC Improvement: (PCC-OCC) / PCC	68.7%	67.3%	68.9%

### Read-Write Mixed Workload



Total Number of Concurrent Operations Under the Same Parent Directory

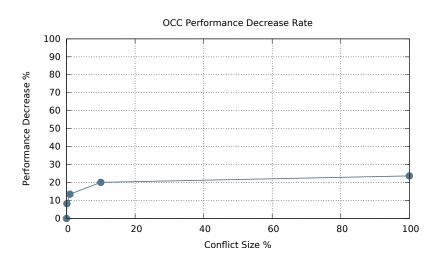
### Read-Write Mixed Workload

Concurrent Read+Creation	1000	10000	100000
PCC	4.92s	50.69s	352.25s
OCC	1.78s	15.31s	120.64s
OCC Improvement: (PCC-OCC) / PCC	63.8%	69.8%	65.8%

### OCC Performance with Different Size of Conflicts

Creations for	Conflict	Elapsed Time	Performance
10000 Operations	Size	(Second)	Decrease
1	100%	14.53	23.7%
10	10%	14.11	20.1%
100	1%	13.51	15.0%
1000	0.1%	12.72	8.23%
10000	0%	11.75	0%

### OCC Performance Decrease Rate



### Conclusion & Future Work

#### Conclusion

- Increase Performance up to 70 %
- Maintain HDFS Strong Consistency Semantics

#### Future Work

- OCC implementation on other operations
- OCC evaluation on multiple NameNodes

# Thank you!