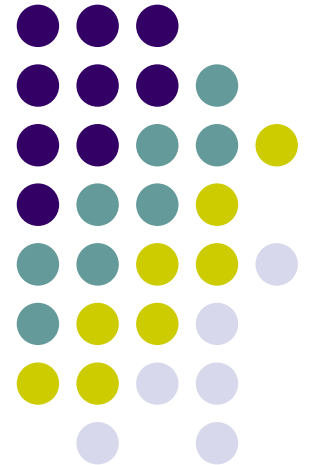


From Regular Expression to NFA

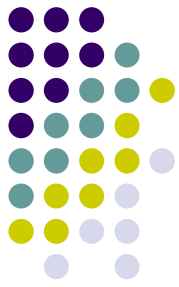
Thompson's Construction





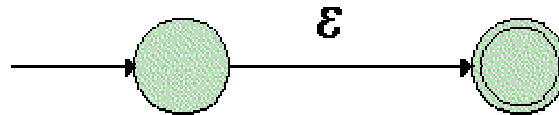
Thompson's Construction

- **Thompson's construction** is a technique for constructing a NFA from a regular expression.
- There are a few simple rules that can be combined to produce a NFA from any arbitrary regular expression.

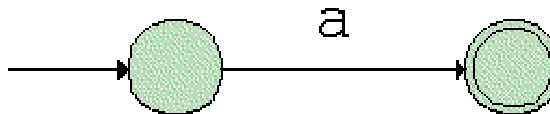


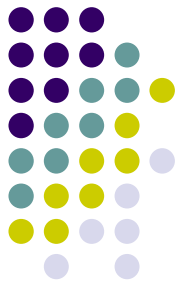
Thompson's Construction

- For ϵ , construct the NFA:



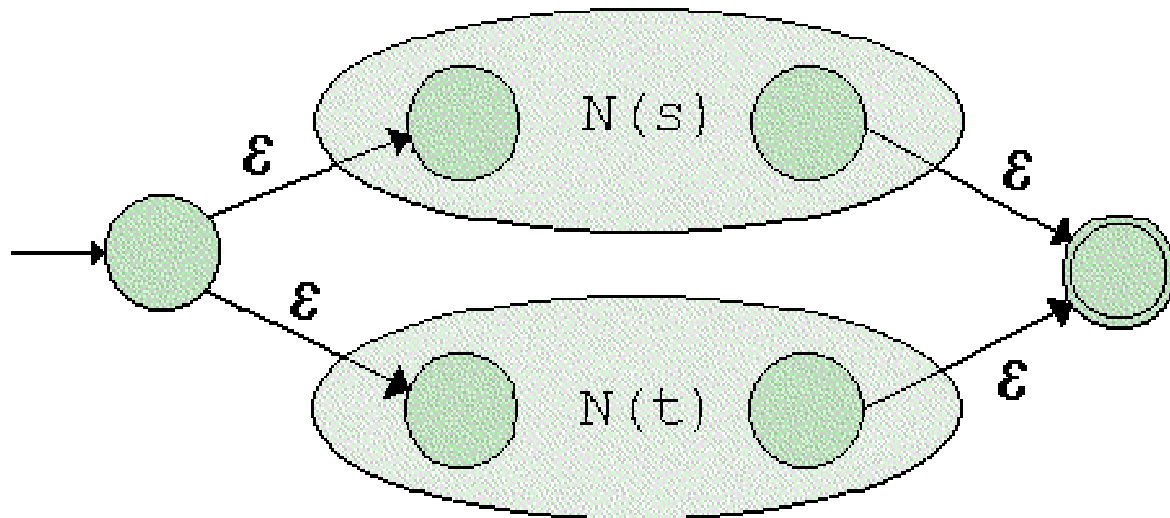
- For a , construct the NFA:





Thompson's Construction

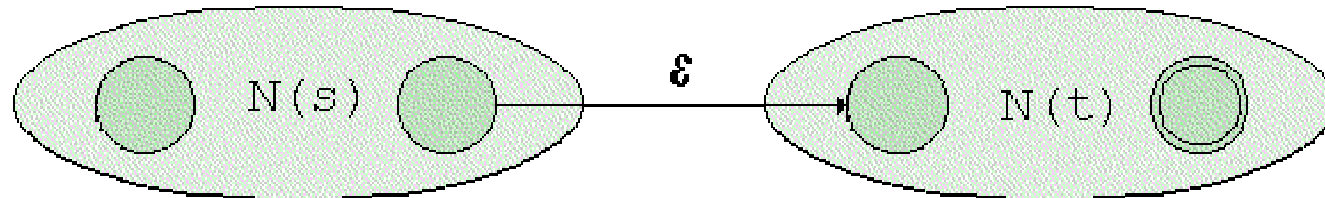
- Suppose $N(s)$ and $N(t)$ are NFAs for regular expressions s and t , respectively. For the regular expression $s \mid t$, construct the NFA:





Thompson's Construction

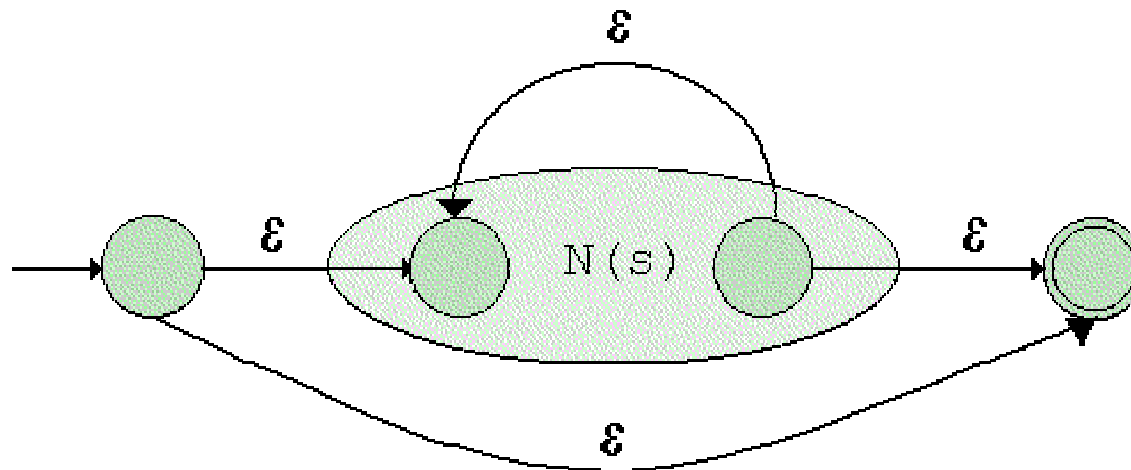
- Suppose $N(s)$ and $N(t)$ are NFAs for regular expressions s and t , respectively. For the regular expression st , construct the NFA:





Thompson's Construction

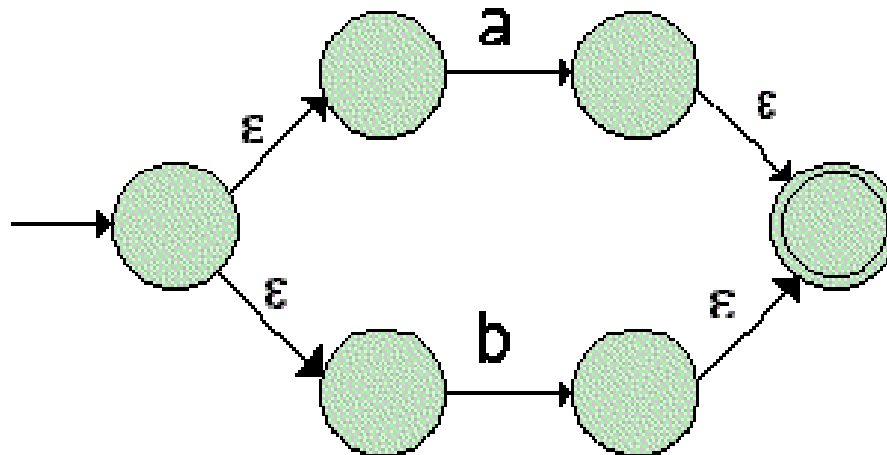
- Suppose $N(s)$ is the NFA for regular expression s . For the regular expression s^* , construct the NFA:



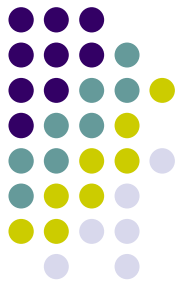


An example

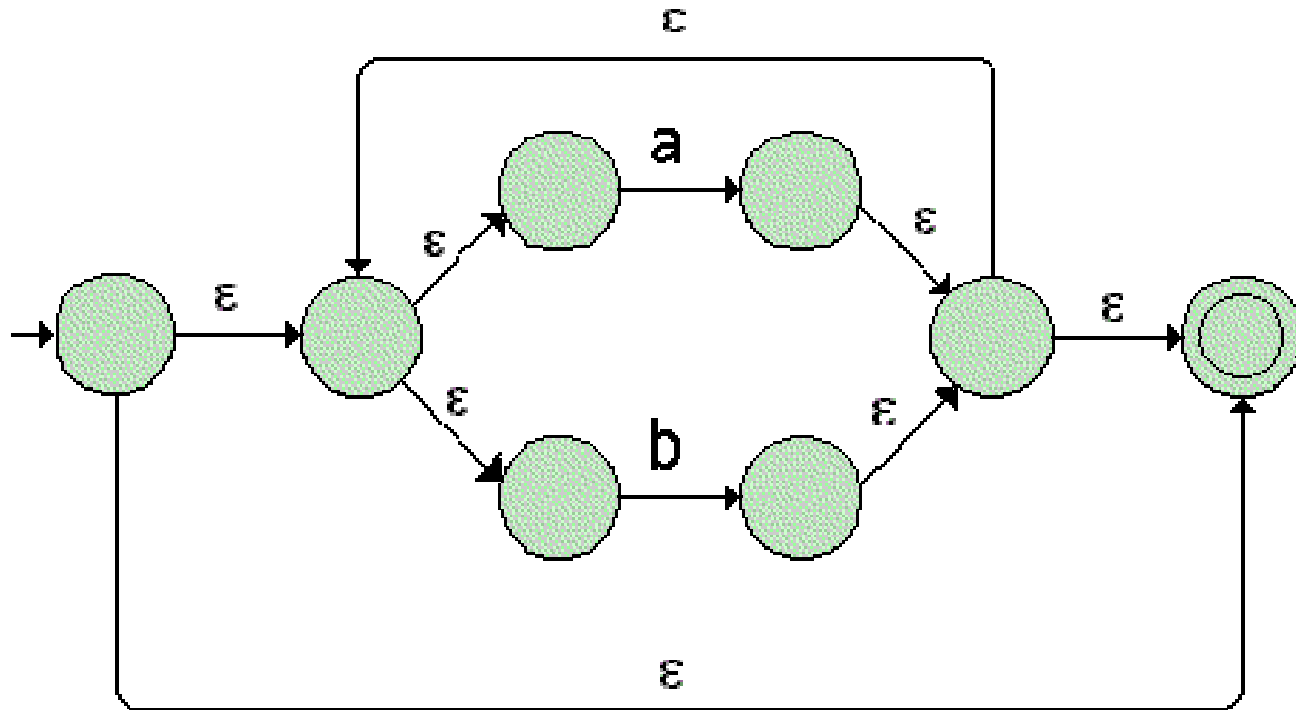
- To demonstrate, let's find an NFA for the regular expression $(a|b)^*aa$.
- First, the NFA for $a|b$ is



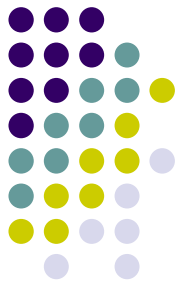
Example, continued



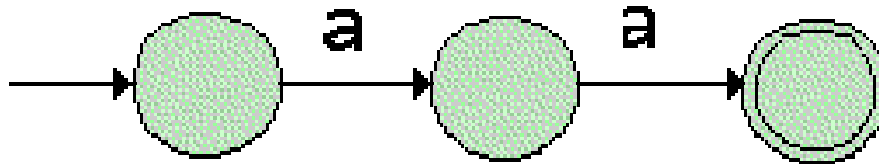
- Now, we draw the NFA for $(a|b)^*$:



Example, continued



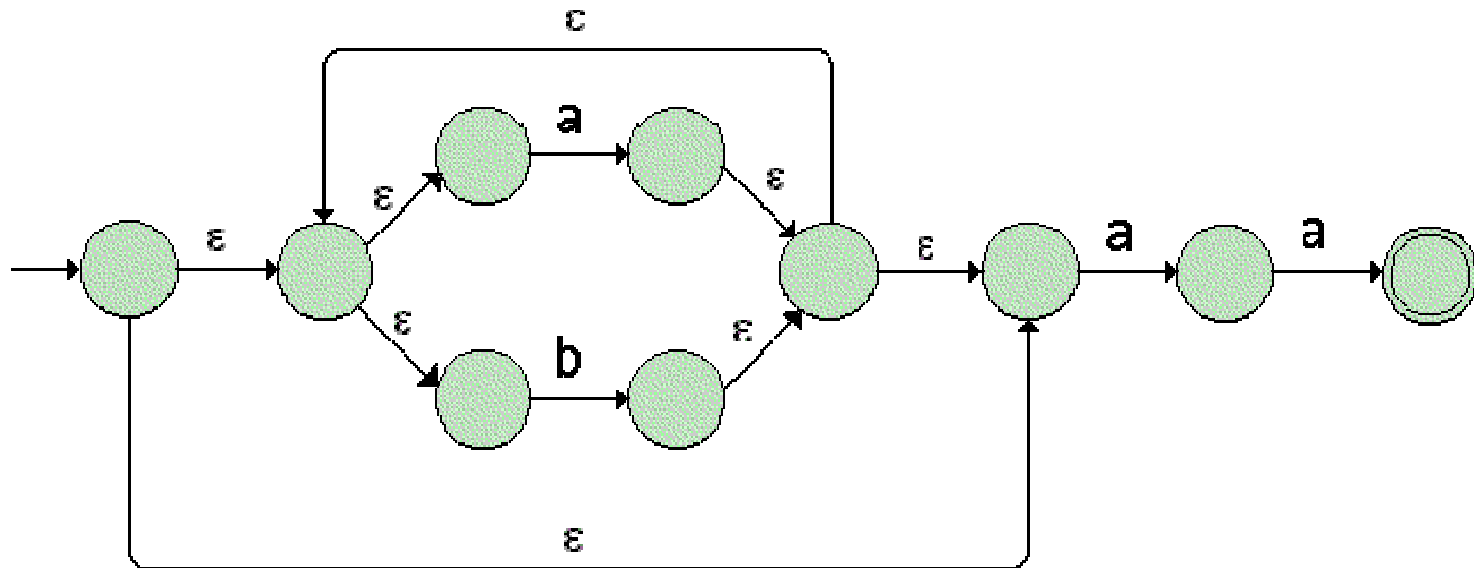
- The NFA for **aa**:

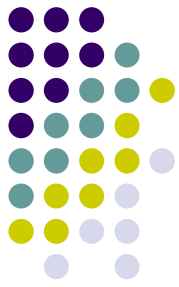


End of example



- Putting this all together, we get the NFA for $(a|b)^*aa$:





What now?

- We know how to produce an NFA that recognizes a string represented by a regular expression.
- Problem: it is difficult to write a program to implement an NFA
- Solution: transform the NFA into a DFA that recognizes the same language
- Next step: From NFA to DFA