

# 5.

# PACKETS, FRAMES, AND ADDRESSING



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传输介质

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局域通信

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布线、拓扑、接口硬件



# PART II Packet Transmission

## Ch 13 Local Area Networks: Packets, Frames, and Addressing

局域网：分组、帧与编址



# 分组交换机

怎么分组，怎么交换



# 7.1 Introduction

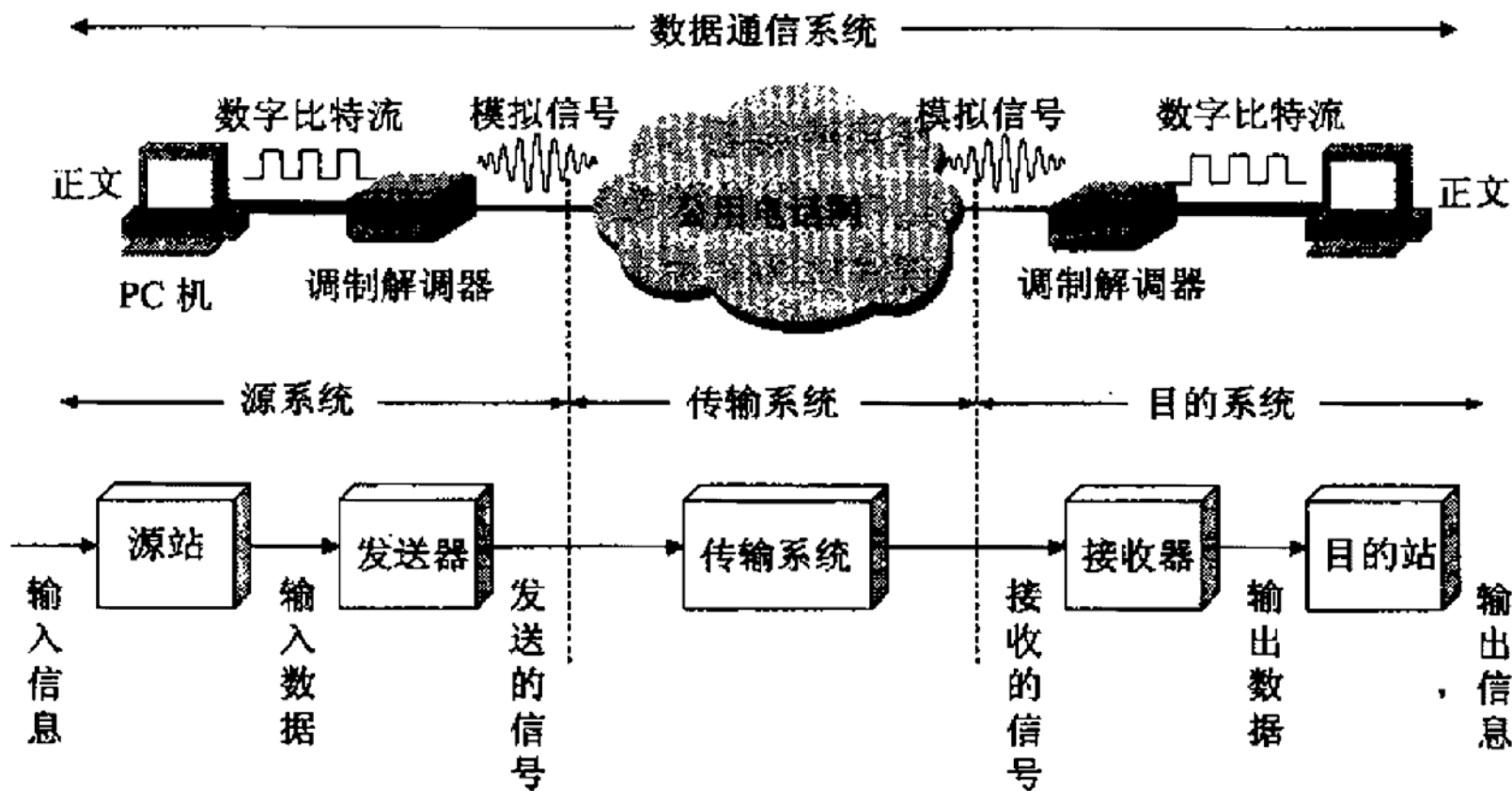


图 3-1 数据通信系统的模型

# 13.2 Circuit Switching

- Circuit switching (线路交换)

- 指的是在发送方和接收方之间建立通路的交互机制
- 收发方的路径与其他收发方是相互独立的
- 类似于电话技术 (associated with telephone technology)

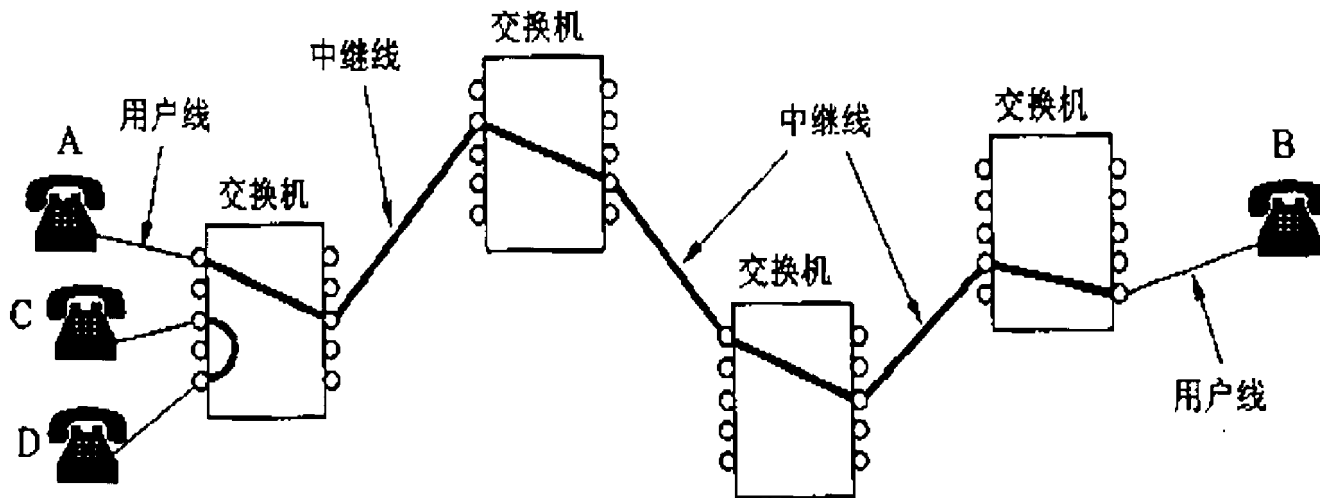


图 1-4 电路交换的示意图

# 13.2 Circuit Switching (线路交换)

- 线路交换 网络使用电子设备建立线路
  - Instead of having each circuit correspond to a physical path
    - multiple circuits are multiplexed over shared media
    - and the result is known as a **virtual circuit** (虚拟通路)

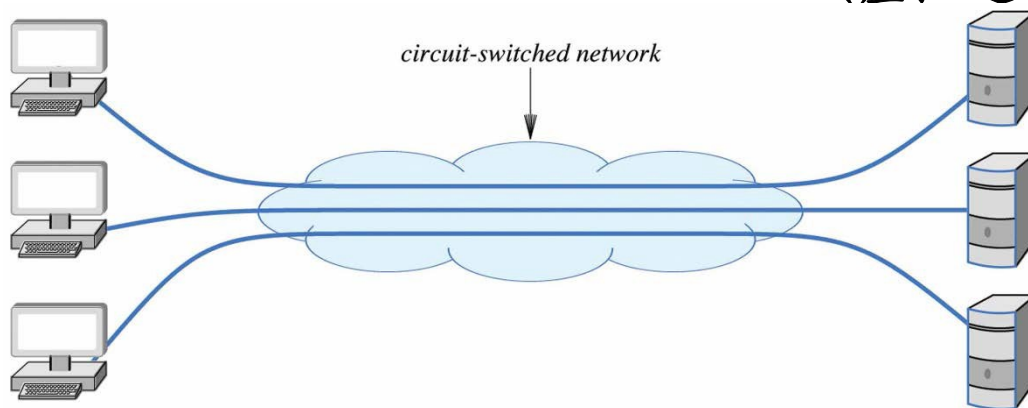


Figure 13.1 A circuit-switched network that provides a direct connection between each pair of communicating entities.





# 13.2 Circuit Switching

- **Three general properties:**
  - **Point-to-point communication**
  - **Separate steps for circuit creation, use, and termination**
  - **Performance equivalent to an isolated physical path**
- **交换线路使用三个步骤模拟拨打电话**
  - **建立线路**：a circuit is established between two parties
  - **线路交互**：the two parties use the circuit to communicate
  - **终止使用**：the two parties terminate use

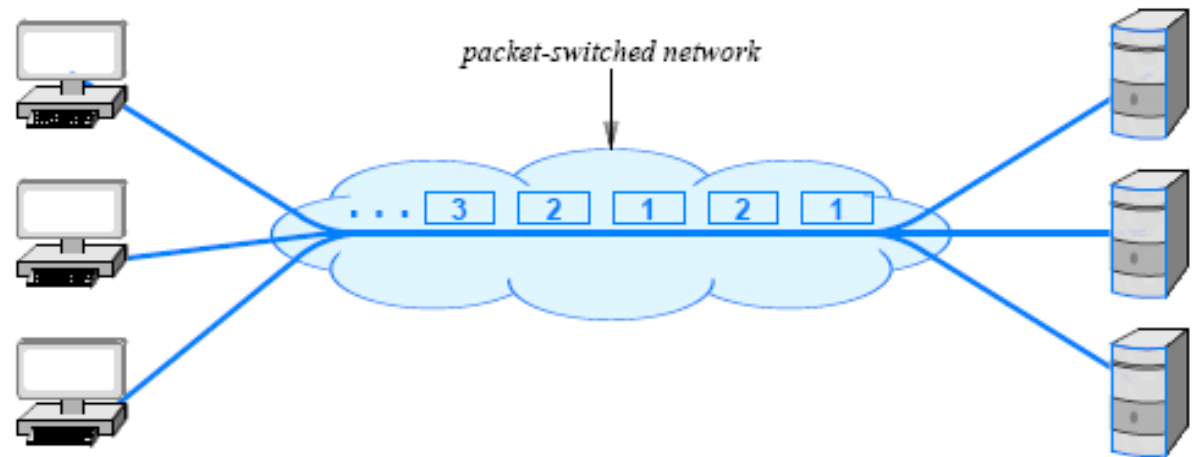


# 13.3 Packet Switching

- A packet switching system uses statistical multiplexing
  - multiple sources **compete** for the use of shared media
- It requires a sender to divide each message into **blocks** of data that are known as **packets (分组)**

- 特点

- 异步
- 无需建立
- 性能各异



# 13.3 Packet Switching

- 优点：To provide communication among  $N$  computers
  - 线路交换：at least  $N/2$  independent paths
  - 分组交换：only requires one path that is shared

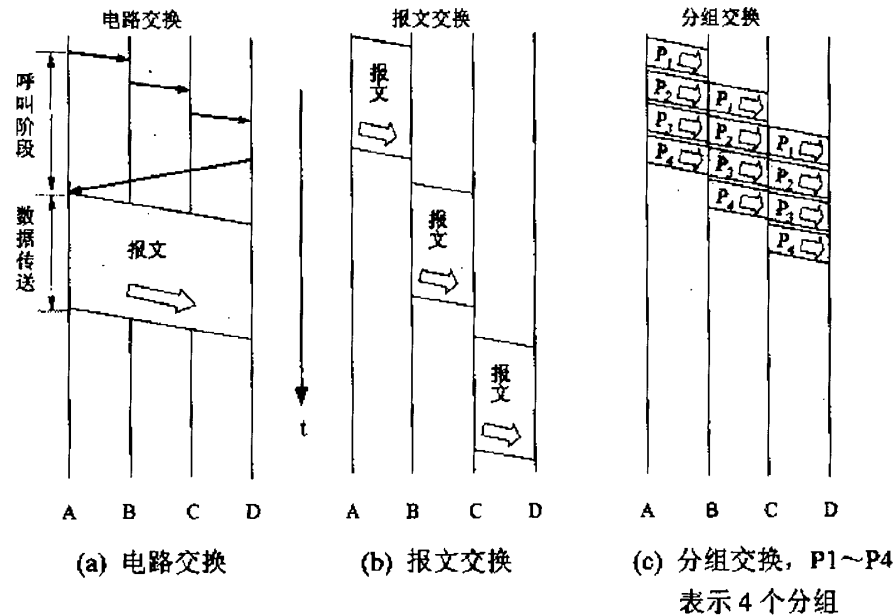


图 1-7 3 种交换的比较



# 分组交换

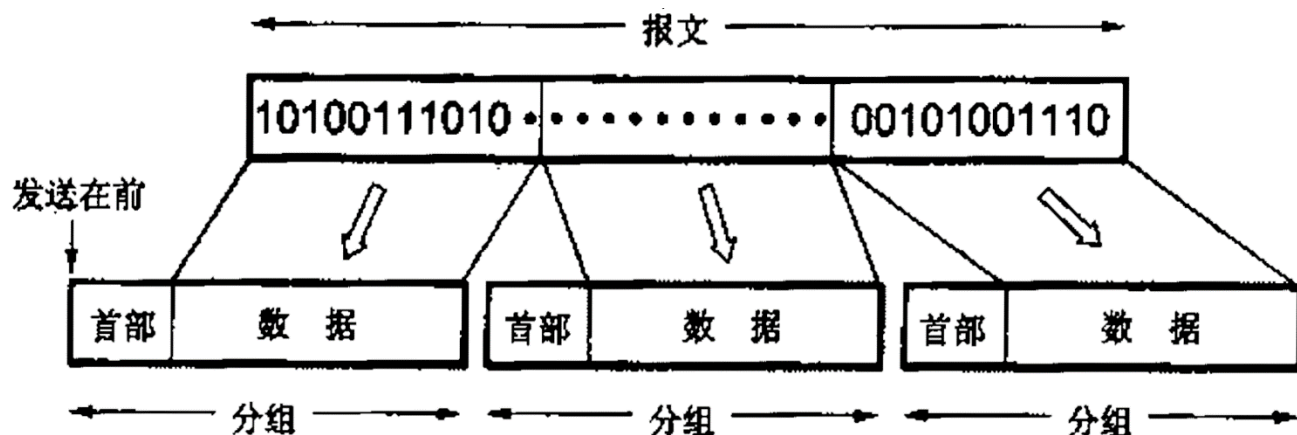


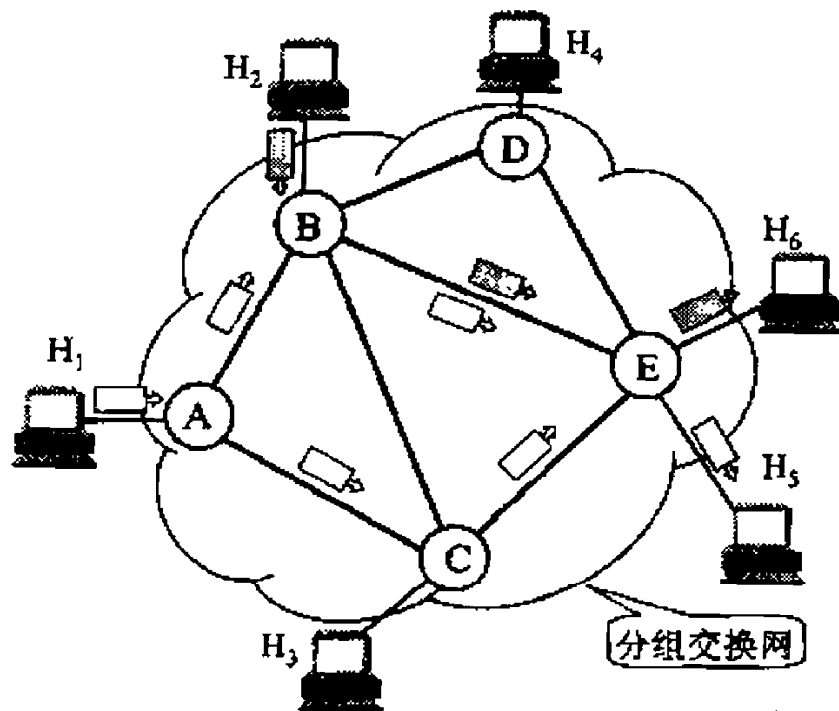
图 1-5 分组的概念

表 1-1 分组交换的优点

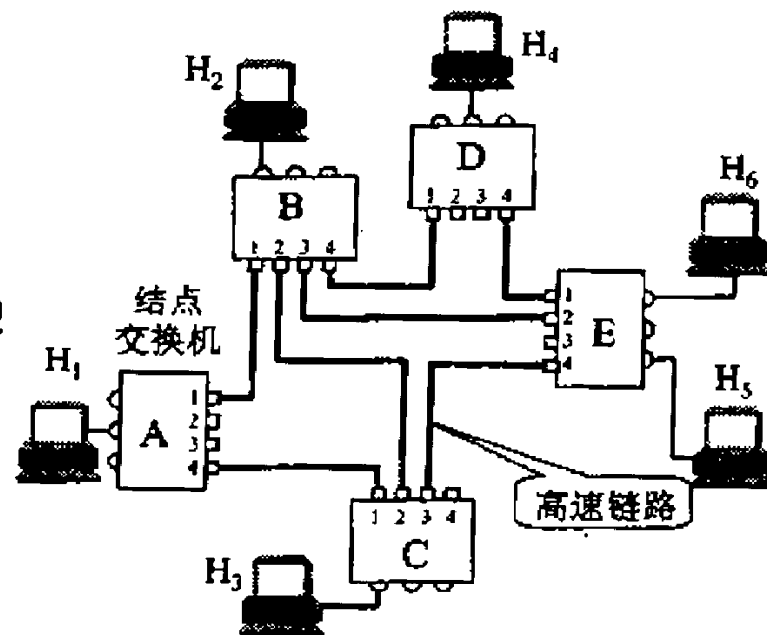
优点	所采用的手段
高效	在分组传输的过程中动态分配传输带宽
灵活	每个结点均有智能, 可根据情况决定路由和对数据做必要的处理
迅速	以分组作为传送单位, 在每个结点存储转发; 网络使用高速链路
可靠	完善的网络协议; 分布式多路由的通信子网



# 分组交换



(a) 通信子网和主机



(b) 结点交换机的两组端口

图 1-6 分组交换网的示意图

# 7.2 The Concept of Packets

- The network system divides data into small blocks called packets, which it sends individually.
  - Computer networks are often called packet networks or packet switching networks because they use packet technology.
- 为何分组
  - 收发双方需要协调传输
  - 多个电脑共用底层硬件



Figure 7.1 An illustration of one reason computer networks use packets. While one pair of computers communicate, others must wait.

# 7.3 Packets and TDM

- A network that permits many sources to take turns accessing a shared communication resource is providing a form of time-division multiplexing.

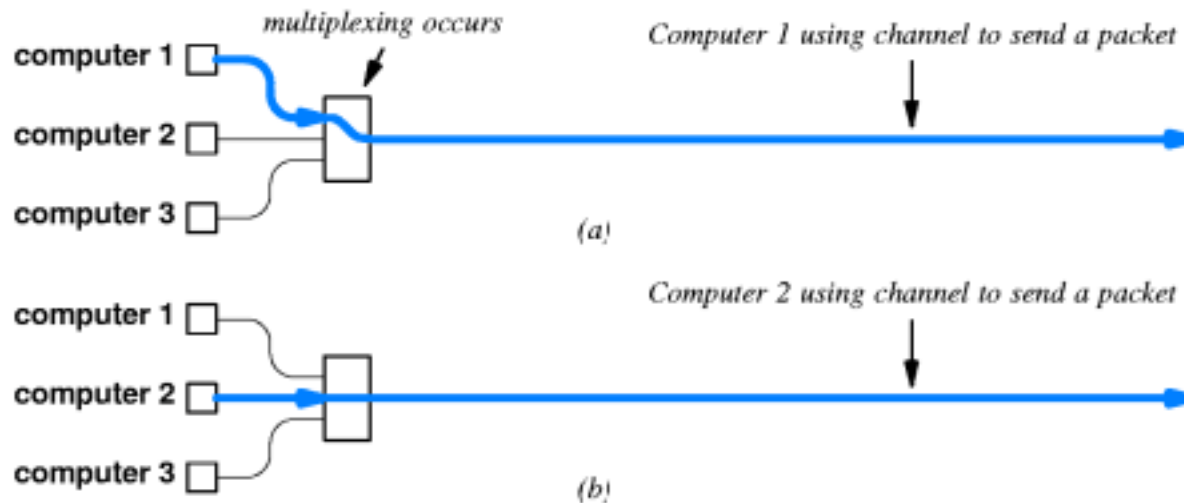


Figure 7.2 Illustration of multiplexing with packets. The sources take turns using the shared communication channel. (a) Computer 1 uses the resource to send a packet, and then (b) computer 2 uses the resource to send a packet.

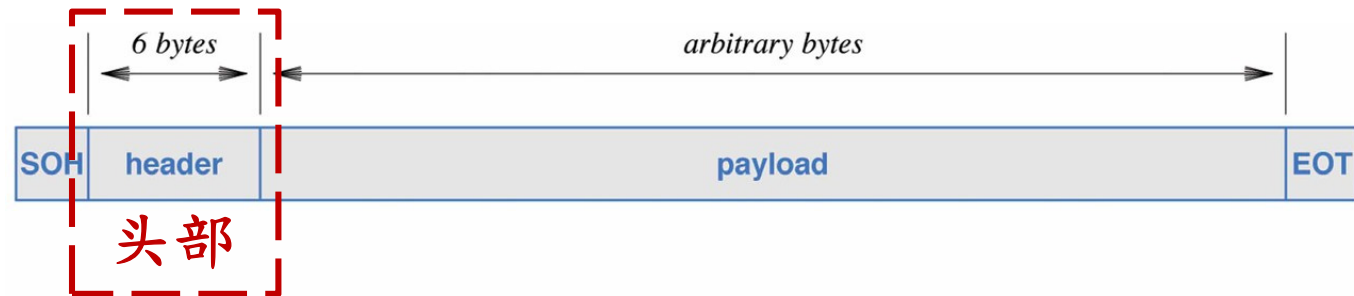
# 7.4 Packets and Hardware Frames

- **Packet (分组)**

- refers to the general concept of a small block of data
  - there is no universal agreement on the exact format of packet.
- 具体硬件技术定义其格式及如何发送

- **Frame (帧)**

- to denote the definition of a packet used with a specific type of network.



**Figure 13.11** An example frame format that uses SOH and EOT characters to delineate a frame.

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# 从双方到多方

## RS-232：不需要共享介质



## 13.9 Packet Identification, Demultiplexing, MAC Addresses

- **Demultiplexing (解复用)** to refer to the **separation** of a combination back into separate information streams
  - Demultiplexing uses **an identifier** known as an **address**
    - Each packet that travels across the shared medium is intended for a **specific recipient**
      - and only the intended recipient should process the packet
- **IEEE** has created a standard for **addressing**
  - Each computer is assigned a **unique** address
    - and each packet contains the address of the intended recipient



# 9.1 Introduction

- **Physically, any signal sent across a shared network reaches all attached stations (所有连接的工作站).**
- **At a given station, the network interface hardware detects the electrical signal and extracts (提取) a copy of the frame.**
  - **Communication does not usually involve all stations.**
  - **Most communication occurs because an application program on one computer sends data directly to an application program on another.**



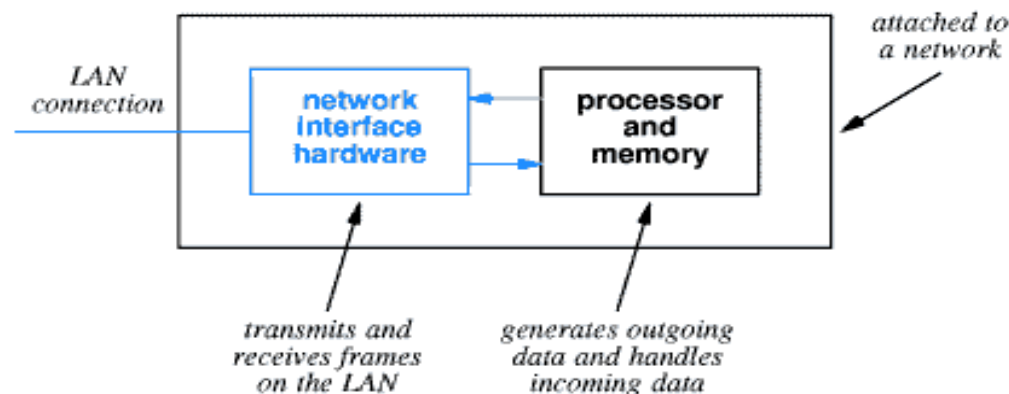
## 9.2 Specifying A Recipient

- Most LAN technologies use an addressing scheme to provide direct communication.
- Each station on the LAN is assigned a unique numeric value called a physical address (物理地址), hardware address (硬件地址), or media access address (MAC).
- A frame sent across a LAN contains the address of the sending computer, called a **source address**, and the address of the intended recipient, called the **destination address**.



## 9.3 How LAN Hardware Uses Addresses to Filter Packets

- LAN硬件完全与CPU、内存分离
  - 完全处理共享介质的收发细节，无需借助CPU
- LAN接口硬件使用物理地址防止收到LAN所有数据包
  - 一旦获得完整帧，接口硬件将其目的地址与该站的物理地址比较。



**Figure 9.1** Organization of the hardware in a computer attached to a LAN. Because it is powerful and independent, the network interface hardware does not use the CPU when transmitting or receiving bits of a frame.

# Network Interface Card

网卡



# 网卡

- 网卡的任务：
  - 检测帧是否存在
  - 有效性检验
  - 判断是否接收
- CPU的任务
  - 判断是否传给上层处理
- 分层处理的目的是：减少CPU的负荷



# 网卡

- 数据包过滤是通信接收方的动作
- 通过检测数据的标识，对符合接收条件的：
  - 复制帧
  - 中断CPU，将复制后的帧交给CPU
  - 继续等待下一个帧
- 对不符合接收条件的：放弃接收，继续等待下一个帧





# Media Access Control Address

MAC地址

如何识别收发双方



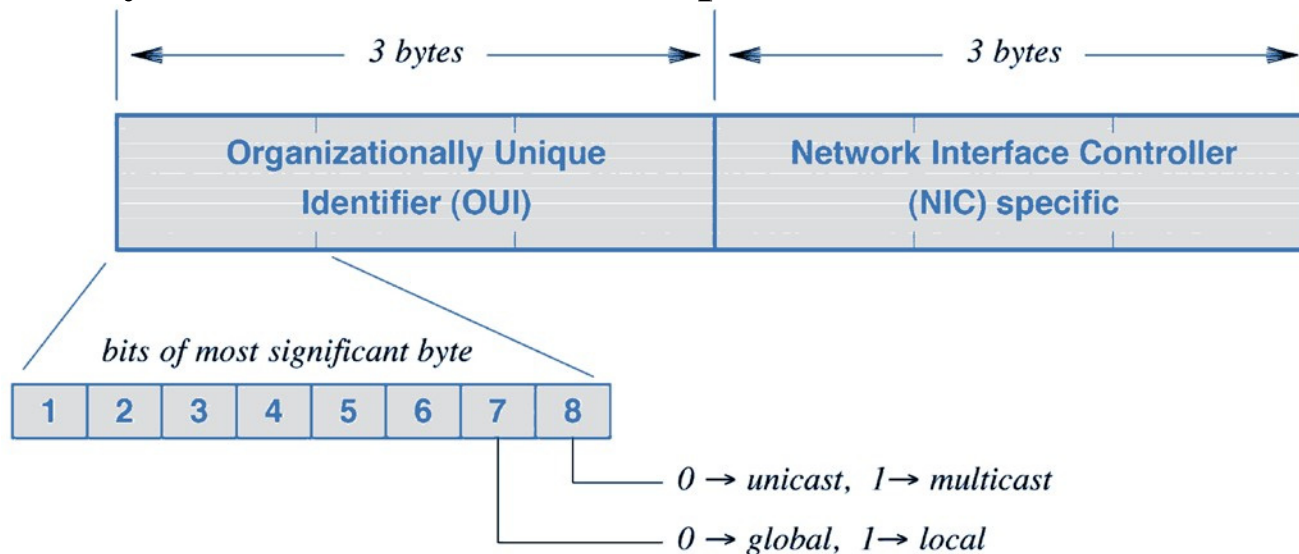
# IEEE Addressing

- In the IEEE addressing scheme, each address consists of **48 bits: Media Access Control (MAC) address**
  - often use the term **Ethernet address (以太网地址)**
- IEEE allocates a **unique** address for each piece of interface
  - Each **Network Interface Card (NIC)** contains a unique IEEE address assigned when the device was manufactured



# IEEE Addressing

- **IEEE assigns a block of addresses to each vendor**
  - allows the vendor to assign a unique value to each device
  - there is a 3-byte **Organizationally Unique ID (OUI)**
    - OUI identifies the equipment vendor
    - a 3-byte block that identifies a particular NIC



**Figure 13.8** The division of a 48-bit IEEE MAC address.



# 9.4 Format of A Physical Address

- **物理地址 (Physical Address)**
  - 必须是所在局域网内的唯一标识
  - 局域网内要对物理地址提供检索的功能
- **Three broad categories:**
  - **Static address (静态地址)**
    - 出厂前由生产厂家一次性配置完成
    - 全球唯一硬件标识，不会出现地址冲突
    - 一次生成，永久使用
    - 地址较长，查询慢



# 9.4 Format of A Physical Address

- **Three broad categories:**

- **Configurable address (可配置地址)**

- 系统出厂后，用户动态分配，保持局域网内唯一
    - 地址短，查询快
    - 地址可以永久有效，也可以随需求改变

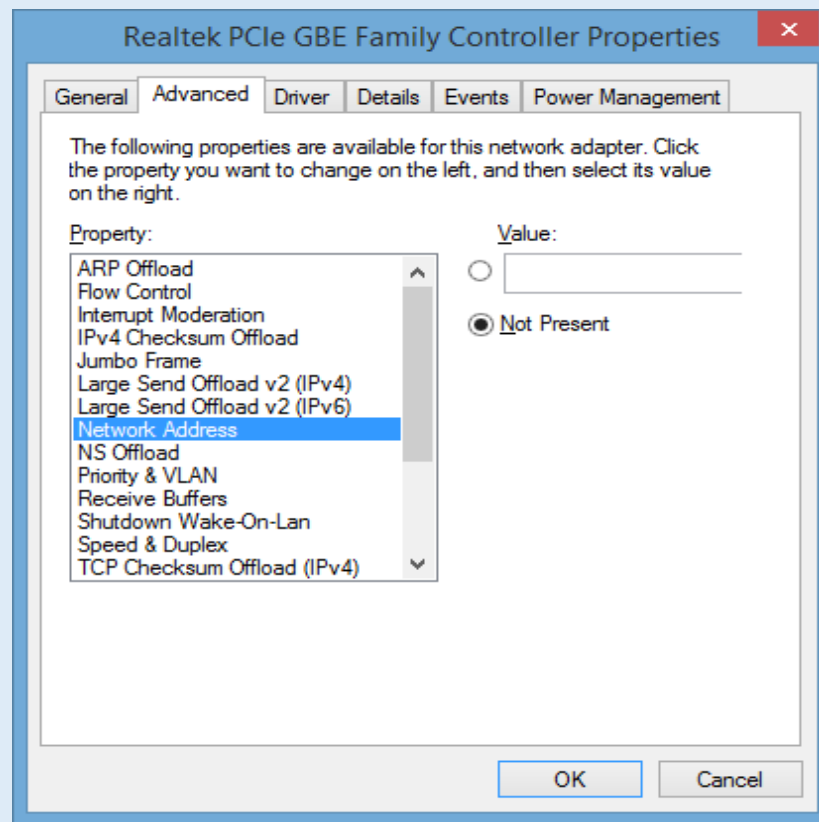
- **Dynamic address (动态地址)**

- 每当系统启动的时候动态分配，局域网内唯一
    - 地址短，查询快
    - 有可能出现地址冲突
    - 地址动态改变，不利于地址映射表的维护



# 思考题

## MAC地址能不能改



## 13.10 Unicast, Broadcast, and Multicast Addresses

- IEEE address format reserves a bit to distinguish between **unicast** and **multicast**
  - but does not provide a way to designate a broadcast address
  - specifies a **broadcast address** consists of **48** bits with all **1**s
    - Thus, a broadcast address has the multicast bit set
  - Broadcast can be viewed as a special form of multicast
    - Each multicast address corresponds to a group of computers
    - Broadcast address corresponds to a group that includes all computers on the network



## 13.11 Broadcast, Multicast, and Efficient Multi-Point Delivery

- **Broadcast and multicast addresses are useful in LANs**
  - because they permit efficient delivery to many computers
- **To understand the efficiency**
  - recall that a LAN transmits packets over a shared medium
- **In a typical LAN**
  - each computer on the LAN **monitors** the shared medium
  - **extracts** a copy of each packet
  - and then **examines** the address in the packet
  - determine whether the packet should be **processed** or **ignored**





# Algorithm

## Algorithm 13.1

Purpose:

Handle a packet that has arrived over a LAN

Method:

Extract destination address, D, from the packet;

if (D matches "my address") {

    accept and process the packet;

} else if (D matches the broadcast address) {

    accept and process the packet;

} else if (D matches one of the multicast addresses for a  
multicast group of which I am a member) {

    accept and process the packet;

} else {

    ignore the packet;

}

**Algorithm 13.1** Packet processing algorithm used in a LAN.



## 9.5 Broadcasting 广播

- **Broadcasting was originally applied to radio and television transmission.**
- **The network designers define a **special, reserved** address known as a broadcast address (广播地址).**
- **The hardware interface on a computer is built to recognize the special broadcast address as well as the station's physical address.**



## 9.5 Broadcasting 广播

- 当一台计算机需要向多台计算机广播信息时，若使用单播技术，必须向每一台计算机分别发送数据，将造成网络负载的成倍增长。
- 采用广播实现方式，发送方按照一个广播地址向共享介质发送数据帧，网上所有设备的网卡分别进行数据帧的拷贝，并交给CPU处理。需要接收数据帧的设备由CPU完成向上层的传递，不需要数据帧的设备由CPU丢弃该数据帧
  - 通过广播，实现了一对多的数据通信，比单点传输方式减少了负载，但所有计算机CPU都要决定接收还是丢弃数据帧，这样增加了无关设备的CPU负载。



## 9.6 Multicasting 组播

- A restricted form (限制形式) of broadcasting known as **multicasting**.
- A network interface does not automatically forward multicast frame to the CPU.
- The interface hardware must be programmed with specification of which multicast frame to accept and which to reject.
- The **interface hardware makes the decision** and only accepts those frames that match the specification.



## 9.7 Multicast Addressing 组播编址

- **When the computer boots, the interface is programmed to recognize only the computer's address and broadcast address.**
- **If an application on the computer wishes to receive multicast frame, the application must inform the network interface which multicast address to use.**
- **The interface adds the address to the set it will recognize, and begins accepting frame sent to that address.**



# 9.8 Identifying Packet Contents

- **A receiver cannot use data in the packet to determine what the packet contains.**
  - Each frame contains additional information that specifies the type of the contents.
- **Two methods to identify the contents of a frame.**
  - **Explicit frame type (显式帧类型)**
    - The bits of a frame used to identify the contents are called the frame type field (帧类型域), and the frame is self-identifying (自标识).
  - **Implicit frame type (隐式帧类型)**
    - 双方协商决定



# 帧与成帧

# Frame & Framing



# 13.12 Frames and Framing

- Chapter 9 introduces the concept of **framing**
  - – in synchronous communication systems it is used as a mechanism that allows a receiver to know where a message **begins** and **ends**
- In more general terms, **framing** refers to the structure added to a sequence of bits or bytes that allows a sender and receiver to agree on the exact format of the message





# 13.12 Frames and Framing

- In a packet-switched network, each frame corresponds to a packet
- A frame consists of two conceptual parts:
  - **Header** that contains **metadata**, such as an address
    - contains information used to process the frame
  - **Payload** that contains the data being sent
    - contains the message being sent
    - and is usually much larger than the frame header



# 13.12 Frames and Framing

- A message is **opaque** (不透明的)
  - that the network only examines the frame header
  - the payload can contain an arbitrary sequence of bytes that are only meaningful to the sender and receiver
- Some tech.'s **delineate** each frame by sending a short **prelude** before the frame and a short **postlude** after it

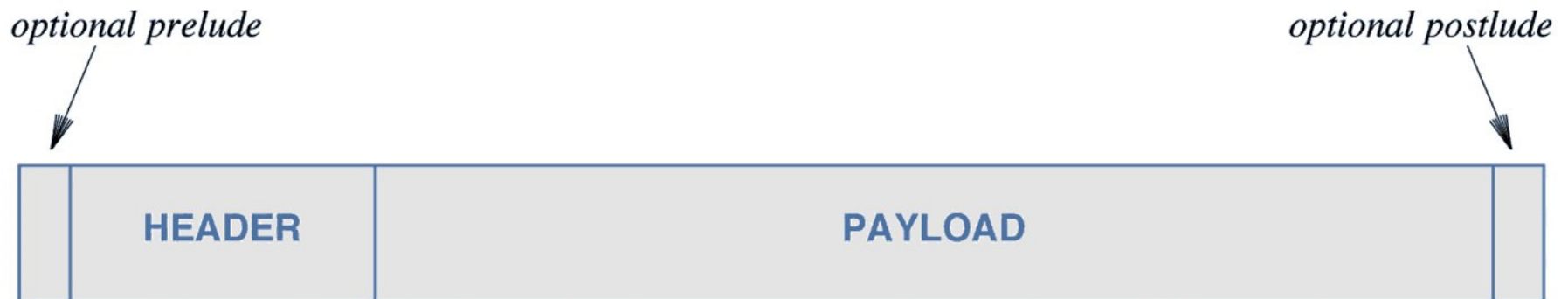


Figure 13.10 Typical structure of a frame in a packet-switched network.

# 13.12 Frames and Framing

- Assume that a packet header consists of **6** bytes
  - the payload consists of an arbitrary number of bytes
- We can use ASCII character set
  - Start Of Header (**SOH**) marks the beginning of a frame
  - End Of Transmission (**EOT**) marks the end



**Figure 13.11** An example frame format that uses SOH and EOT characters to delineate a frame.



# 7.5 Byte Stuffing 字节填充

- 网络允许传输任意字节，因此需要对帧进行定界
  - 如何区分正常数据和定界？
- 回顾C语言的转义字符\ r \ n，右斜杠：\\
- 字节填充(data stuffing)：插入额外的比特或字节

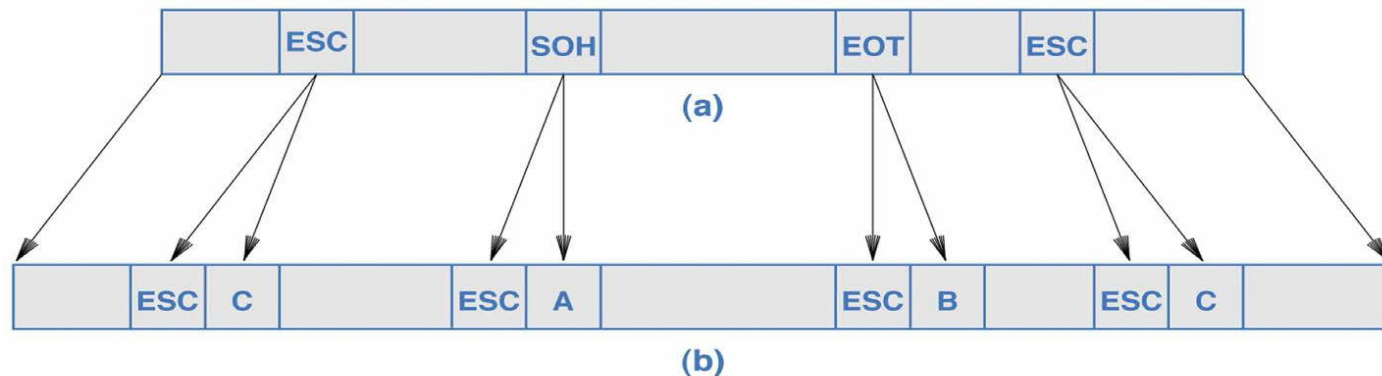
Byte In Payload	Sequence Sent
SOH	ESC A
EOT	ESC B
ESC	ESC C

**Figure 13.12** An example of byte stuffing that maps each special character into a 2-character sequence.



# 7.5 Byte Stuffing 字节填充

- The term byte stuffing and character stuffing refer to data stuffing used with character-oriented hardware, and bit stuffing refers to data stuffing used with bit-oriented hardware. For example:



**Figure 13.13** Illustration of (a) original data, and (b) a version after byte-stuffing has been performed.

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## • 字符填充

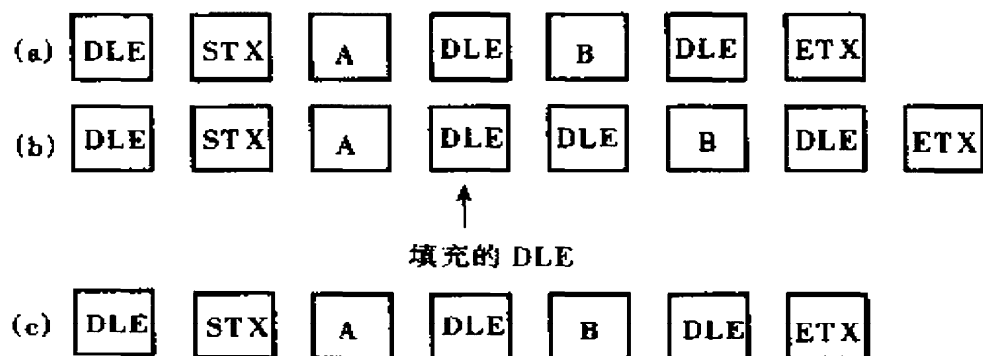


图 3-4 (a) 网络层发出的数据；(b) 经数据链路层填充后的数据；  
(c) 数据传送给接收方的网络层。

## • 位填充

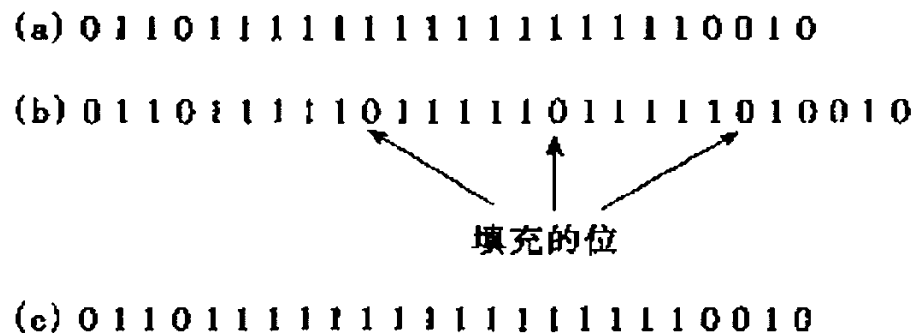


图 3-5 位填充

(a) 原始数据；(b) 线上数据；(c) 删除填充位后接收方存储器内的数据。



# 位填充

用这种方法成帧的主要缺点是要完全依赖于 8 位字符,特别是 ASCII 字符。随着网络的发展,这种在帧中嵌入字符的机制所带来的弊端变得越来越明显。于是,一种允许任意长短字符的新技术发展起来。

这种新技术允许数据帧包含任意个数的比特,而且也允许每个字符的编码包含任意个数的比特。它的工作方式如下:每一帧使用一个特殊的位模式,即 01111110 作为开始和结束标志(flag)字节。当发送方的数据链路层在数据中遇到 5 个连续的 1 时,它自动在其后插入一个 0 到输出比特流中。这种位填充技术类似于字符填充技术,后者是在数据中的 DLE 之前填充一个 DLE 到输出字符流中。

当接收方看到 5 个连续的 1 后面跟着一个 0 时,自动将此 0 删去。位填充(bit stuffing)技术和字符填充技术一样,对通信双方计算机的网络层来说都是完全透明的。如果用户数据包含着位模式 01111110,则将以 011111010 的形式传送出去,但是仍然以 01111110 的形式存放在接收方的存储器中。图 3-5 给出了一个位填充的例子。

采用位填充技术,两帧间的边界就可以通过位模式唯一地识别。因此,如果接收方失去同步,它只需在输入流中扫描标志序列,即可重新获得同步。因为这些标志序列只能是帧界,而决不应该在数据中。



# 帧格式





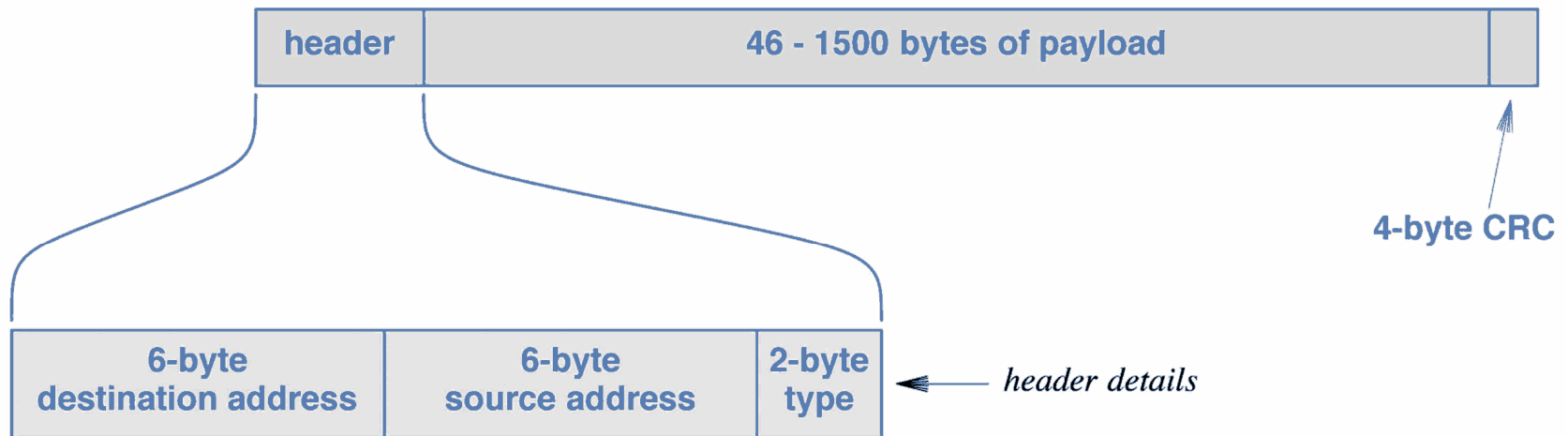
## 9.10 An Example Frame Format

- **An Ethernet frame begins with a header that contains three fields.**
- **The 64-bit preamble contains alternating 1s and 0s**
- **The first two fields of the header contain physical addresses (48-bit static addressing scheme).**
- **The third field of the header consists of a 16-bit Ethernet frame type.**



# Digital-Intel-Xerox (DIX)

- The DIX Ethernet standard specifies the value that can be used in the header fields and their meanings.
  - The address with all 48 bits set to 1 is reserved for broadcast, other addresses that start with a 1 bit are used for multicast.



**Figure 15.1** Illustration of the Ethernet frame format and header details.

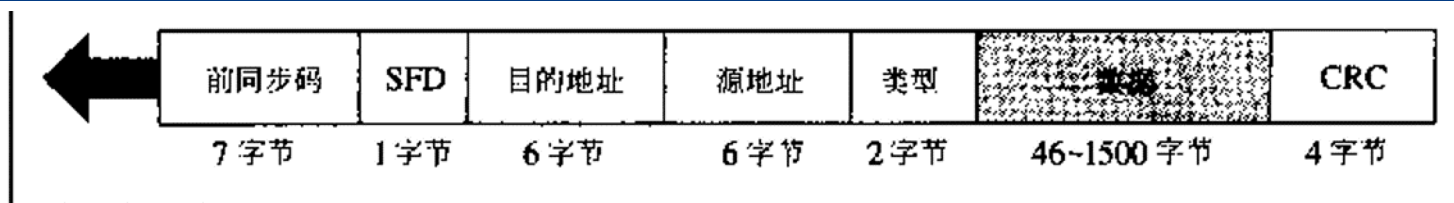
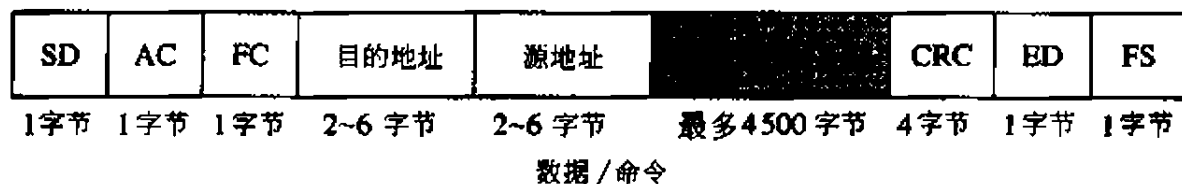


图 3.10 以太网的格式

以太网定义的帧共有 7 个字段：

- **前同步码**。这个字段有 7 个字节(56 比特)的交替出现的 0 和 1,其作用就是提醒接收系统有帧到来,以及使到来的帧与计时器同步。
- **帧首定界符(SFD)**。这里用 1 字节(10101011)作为标志,并指出帧的开始。
- **目的地址**。6 字节的物理地址。
- **源地址**。6 字节的源物理地址。
- **类型**。这个字段定义了封装在帧中的数据类型。
- **数据**。这个字段包含从上层来的数据。数据长度必须在 46 到 1500 字节之间。如果上层协议产生的数据长度小于 46 字节,则应将其填补到 46 字节。若数据长度超过 1500 字节,上层就必须将其进行分片。
- **循环冗余检验(CRC)**。这是一个 4 字节的字段用作差错检测,它使用 CRC-32



SD 起始定界符(标志)  
 AC 接入控制(优先级)  
 FC 帧控制(帧类型)  
 ED 结束定界符(标志)  
 FS 帧状态

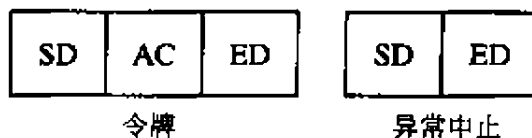


图 3.15 令牌环的帧格式

- **目的地址(DA)**。这个可变长度(2 至 6 字节)字段是下一站的物理地址。
- **源地址(SA)**。这个可变长度(2 至 6 字节)字段是前一站的物理地址。
- **数据**。这个字段是数据。数据可多到 4500 字节。
- **CRC**。这个字节为 4 字节长,包含 CRC-32 检错序列(见附录 D)。
- **结束定界符(ED)**。这个 1 字节字段指出发送器的数据结束,同时还包含更多的控制信息。
- **帧状态(FS)**。这个 FS 字段由接收器设置,指出帧已被读取,或由监督站设置,指出该帧已在环上转了一圈。

**令牌帧**,令牌帧包括三个字段:SD、AC 和 ED。

**异常中止帧**,异常中止帧只有两个字段:SD 和 ED。当出现一些问题时,监督站使用异常中止帧来中止令牌传递机制。



## 15.4 Ethernet Type Field and Demultiplexing

- The **type field** in an Ethernet frame provides multiplexing and demultiplexing
  - that allows a given computer to have multiple protocols operating simultaneously
- When a frame arrives at its destination
  - the receiver examines the type field, and it uses the value to determine which software module should process the frame



# 15.4 Ethernet Type Field and Demultiplexing

- The protocols used on the Internet send IP datagrams and ARP messages over Ethernet
  - Each is assigned a unique Eth. type

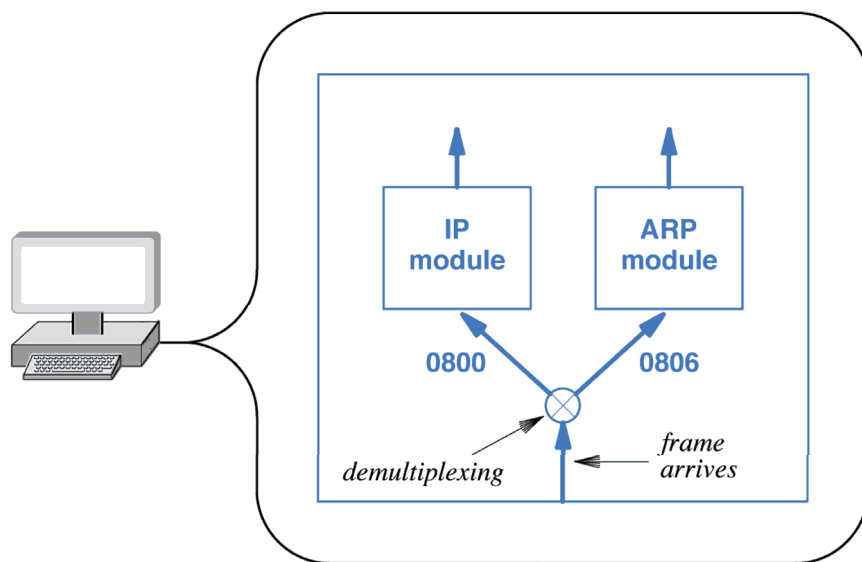


Figure 15.2 Illustration of using the frame type field for demultiplexing.

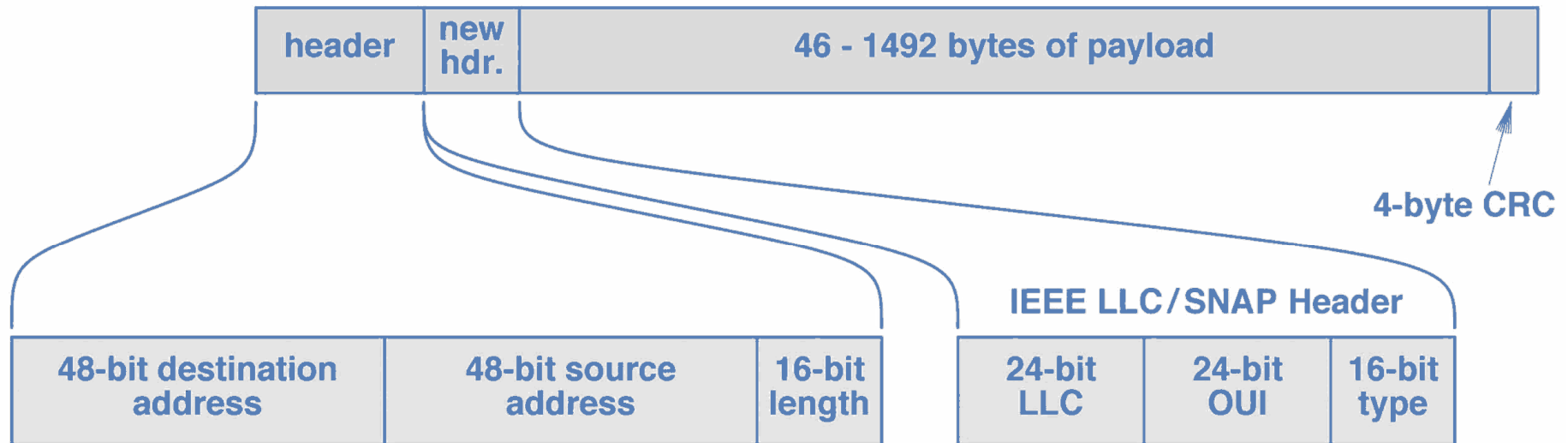
Value	Meaning
0000-05DC	Reserved for use with IEEE LLC/SNAP
0800	Internet IP Version 4
0805	CCITT X.25
0900	Ungermann-Bass Corporation network debugger
0BAD	Banyan Systems Corporation VINES
1000-100F	Berkeley UNIX Traller encapsulation
6004	Digital Equipment Corporation LAT
6559	Frame Relay
8005	Hewlett Packard Corporation network probe
8008	AT&T Corporation
8014	Silicon Graphics Corporation network games
8035	Internet Reverse ARP
8038	Digital Equipment Corporation LANBridge
805C	Stanford University V Kernel
809B	Apple Computer Corporation AppleTalk
80C4-80C5	Banyan Systems Corporation
80D5	IBM Corporation SNA
80FF-8103	Wellfleet Communications
8137-8138	Novell Corporation IPX
818D	Motorola Corporation
FFFF	Reserved

# 15.5 IEEE's Version of Ethernet (802.3)

- **IEEE developed a standard for Ethernet (1983) and attempted to redefine the Ethernet frame format**
  - This IEEE working group is numbered 802.3
- **传统Ethernet与802.3 Ethernet区别于类型域的解释**
  - 802.3 standard interprets the original type field as a packet length, and adds 8-byte header that contains the packet type
  - The extra header is known as a Logical Link Control / Sub-Network Attachment Point (LLC/SNAP) header; most professionals simply call it a SNAP header



# 15.5 IEEE's Version of Ethernet (802.3)



**Figure 15.3** The IEEE 802.3 frame format with an LLC/SNAP header.





# 15.5 IEEE's Version of Ethernet (802.3)

- The overall frame size in 802.3 Ethernet remains the same as conventional Ethernet: **1514** bytes
  - IEEE reduced the max. payload from **1500** bytes to **1492**
    - SNAP header occupies the first 8 bytes of the payload
- To keep the two versions of Ethernet compatible:
  - If bytes **13-14** of a frame is less than **1500**
    - the field is interpreted as a packet length and the 802.3 standard applies
  - otherwise, the field is interpreted as a type field and the original Ethernet standard applies



## 15.6 LAN Connections and Network Interface Cards

- **NIC appears to be an I/O device**
  - it is connected in the same manner as a disk or video device
  - handles address recognition; CRC computation; frame **recognition**; sending/receiving frames
- **consists of a **circuit board** with a plug on one side.**
- **Most computers come with a NIC already installed**
- **The NIC is independent from the rest of the computer**
  - a user can replace the NIC without making other changes



- 可用的标识方式

- 把传送的数据看作有效数据，发送和接收者通过其它协议约定传输的数据类型标识
- 有效数据开始部分作为数据类型标识

- 应用层数据标识的问题

- 类型标识值可能被重复使用，很难保持全局唯一



- 数据链路层标识

- 全球唯一的标识的建立：IEEE802.2的LLC/SNAP
- 在帧的有效数据头附加一个唯一标识

- 逻辑链路层控制(LLC)：3bytes

- 子网连接点：SNAP

- 组织唯一标识(OUI)：3bytes，标准组织的编号
- 组织内部唯一类型标识：2bytes



## 9.12 Network Analyzers, Physical Addresses, Frame Type

- **A network analyzer or network monitor is a device that can be configured to count or display frame as they pass across a shared network.**
- **To read packets, analyzer software places the computer's network interface hardware into promiscuous mode (混合模式).**
- **An analyzer obtains a copy of each frame, then uses header fields such as the physical source address, physical destination address, or type information to determine how to process the frame.**



实践是检验真理的唯一标准



# 实验

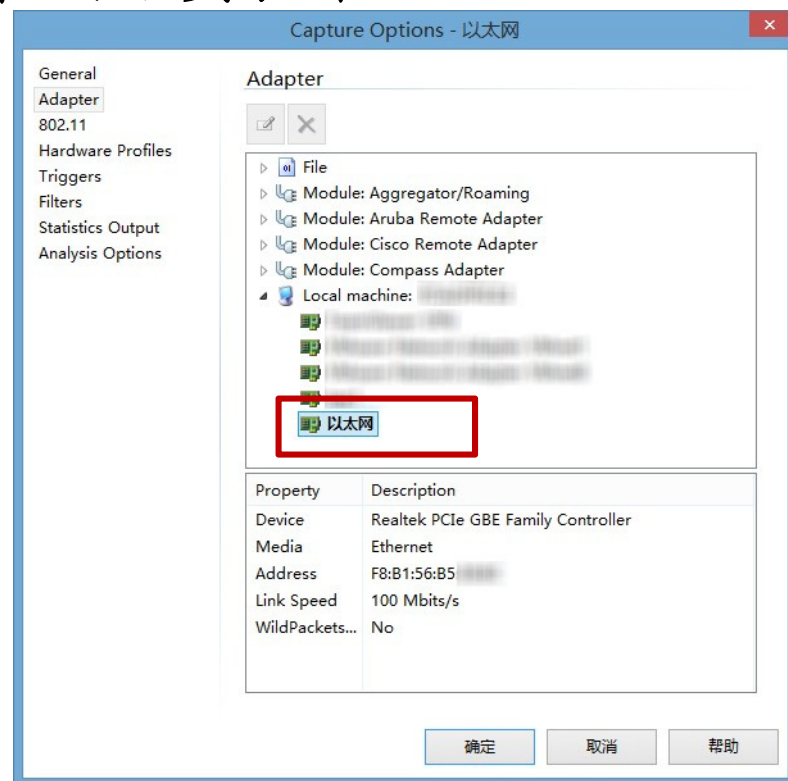
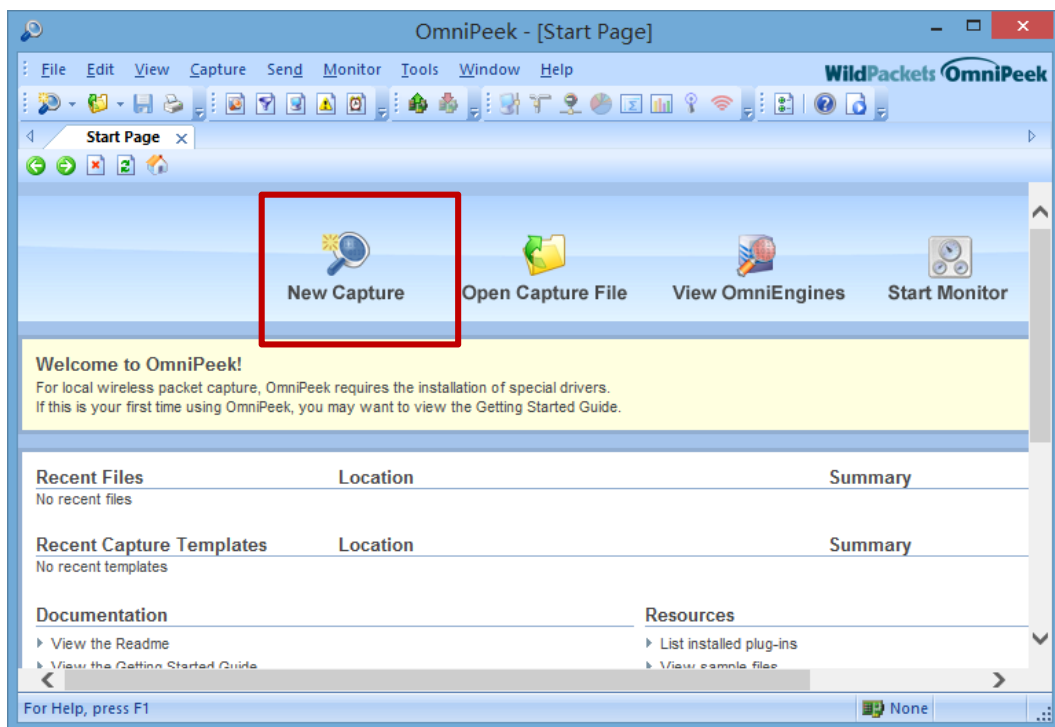
- 请用Omnipeek监听本机的收发包，观察MAC地址
  - 在宿舍里（设宿舍内不用路由），PING不同室友的电脑（或QQ传文件），观察包的MAC地址
  - PING公网上的不同主机，观察包的MAC地址
  - 以上实验也可以在机房完成
  - Omnipeek程序在FTP上，用法请上网搜教程。
- 得到了什么结论？



# Omnipeek 7.5教程

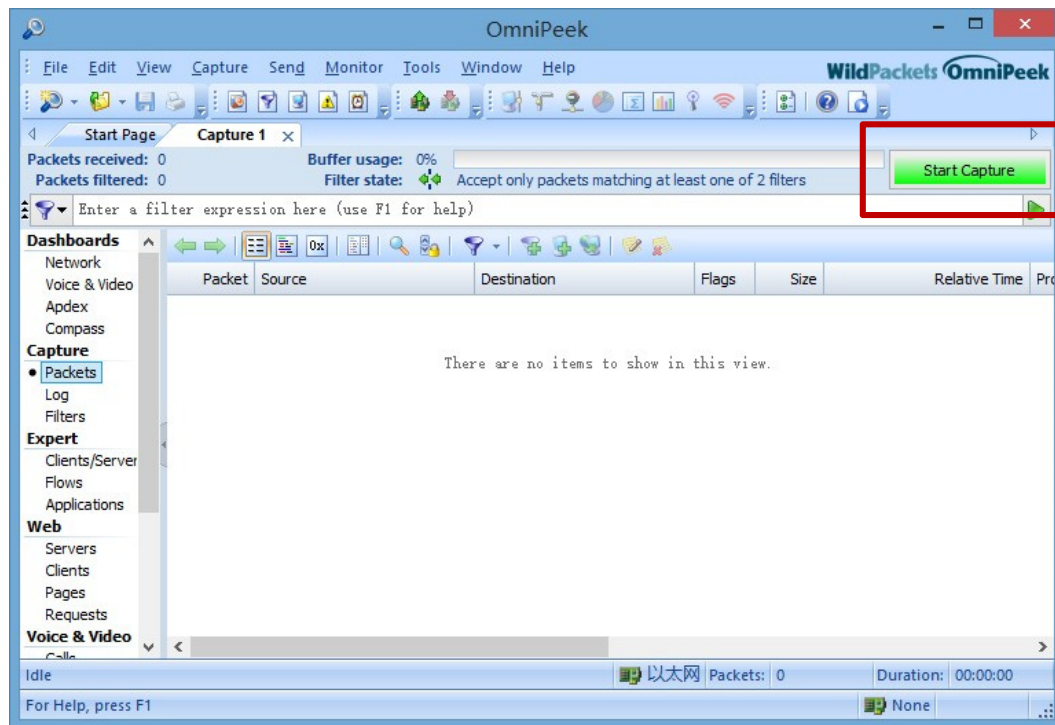
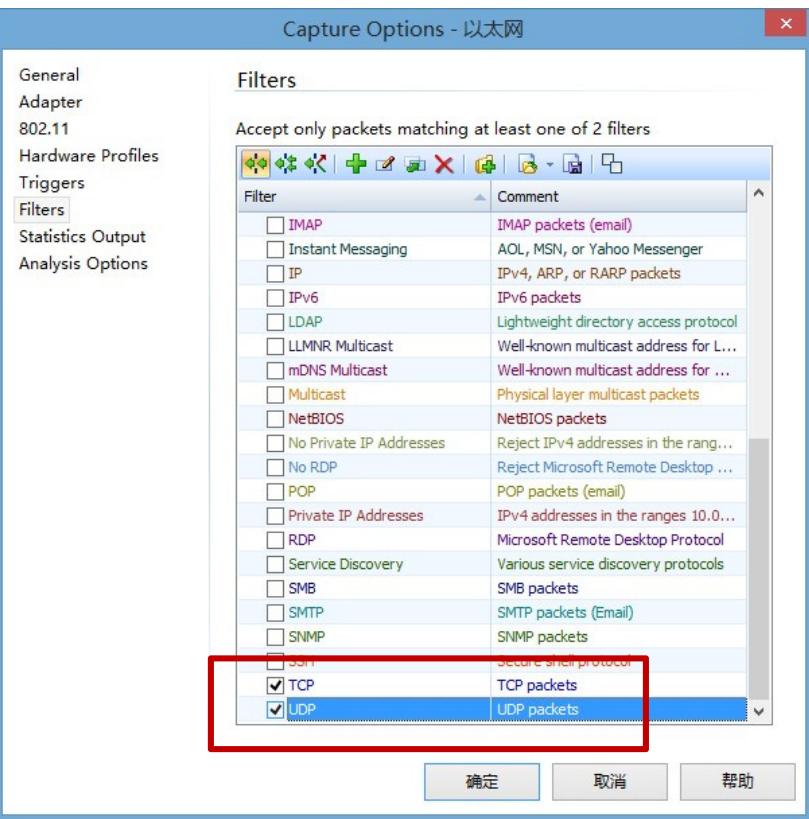
• 下载地址：

– FTP://XXXX/教学课件/林/计算机网络实验课/tools/

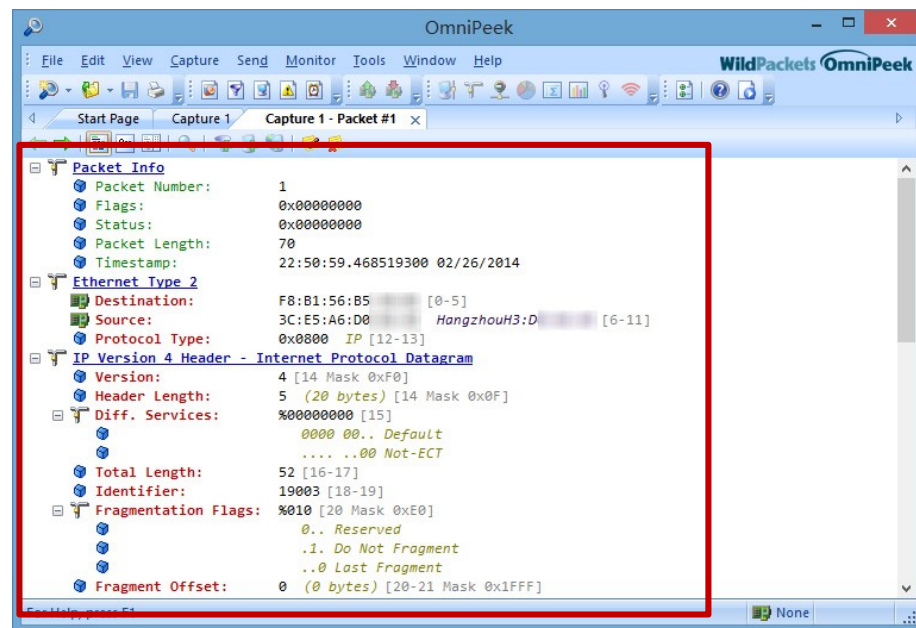
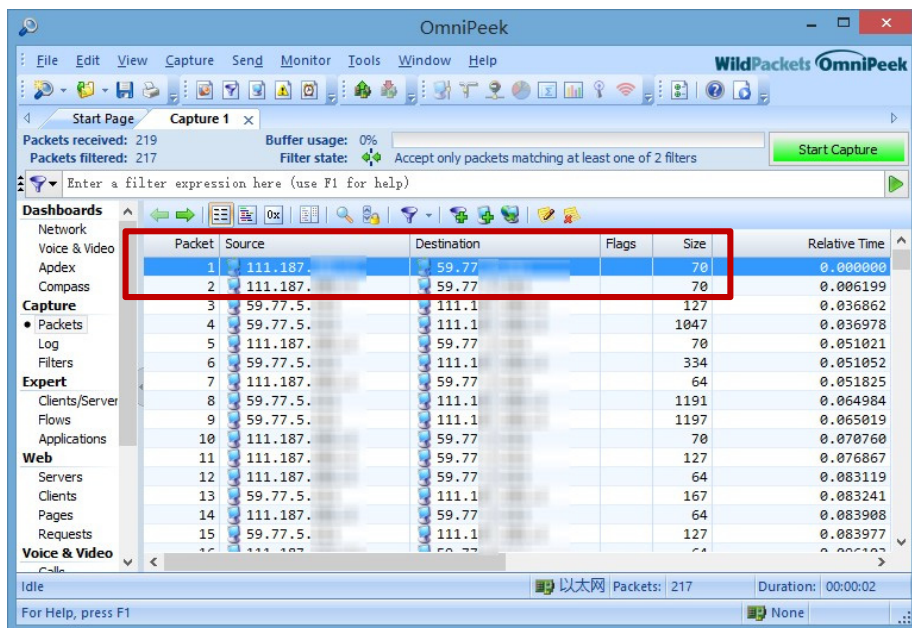




# Omnipeek 7.5教程



# Omnipeek 7.5教程



# 监听结果节选

## Packet #1

### Ethernet Type 2

**Destination:** FF:FF:FF:FF:FF:FF *Ethernet Broadcast* [0-5]  
**Source:** 00:0C:29:37:5A:1B *VMware:37:5A:1B* [6-11]  
**Protocol Type:** 0x0800 *IP* [12-13]

### IP Version 4 Header - Internet Protocol Datagram

**Version:** 4 [14 Mask 0xF0]  
**Protocol:** 17 *UDP* [23]  
**Source IP Address:** 0.0.0.0 [26-29]  
**Dest. IP Address:** 255.255.255.255 *IP Broadcast* [30-33]

### UDP - User Datagram Protocol

**Source Port:** 68 *bootpc* [34-35]  
**Destination Port:** 67 *bootps* [36-37]

### BootP - Bootstrap Protocol

**IP Address Known By Client:** 0.0.0.0 *IP Address Not Known By Client* [54-57]  
**Client IP Addr Given By Srvr:** 0.0.0.0 [58-61]  
**Server IP Address:** 0.0.0.0 [62-65]  
**Gateway IP Address:** 0.0.0.0 [66-69]  
**Client Hardware Addr:** 00:0C:29:37:5A:1B *VMware:37:5A:1B* [70-75]

### DHCP - Dynamic Host Configuration Protocol

#### Requested IP Address

**Address:** 192.168.7.132 [296-299]

#### Host Name Address

**String:** WIN-KG9CLM76UIA [302-316]



5.



**THANK YOU.**

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