Logical Database Design: Mapping ER to Relational

Chapter 3, Section 3.5

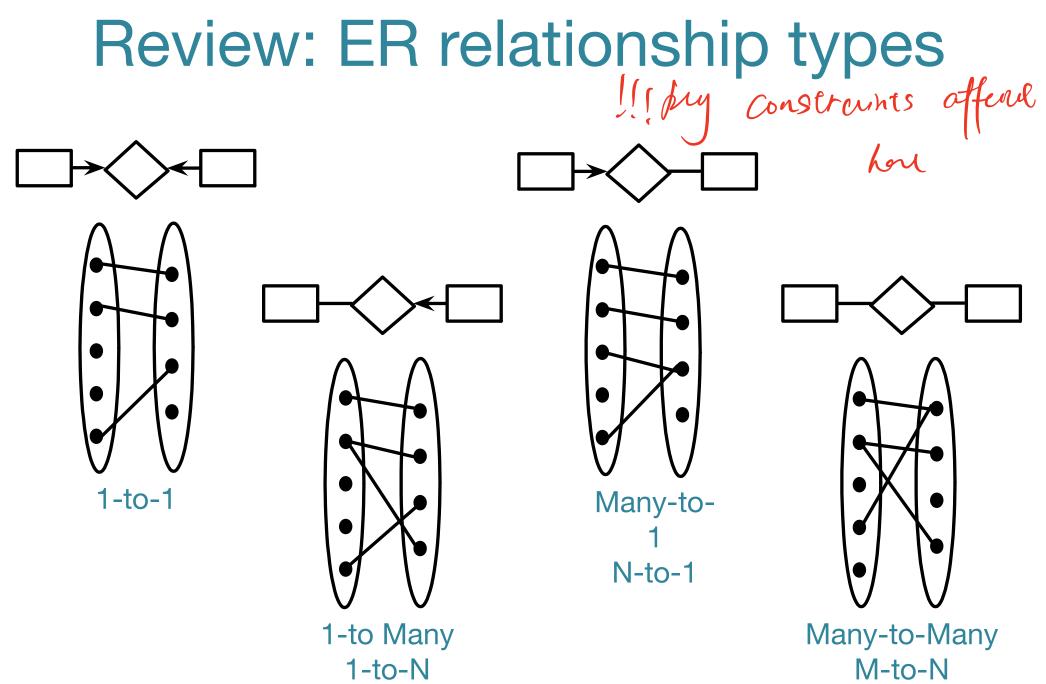
ER Model vs. Relational Model

- ER Model used for conceptual design
- Relational Model implemented by modern DBMS
- Important Step: Translate ER diagram to Relational schema

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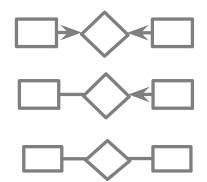
Recall ER Constructs

- Basic Constructs
 - Entity Sets
 - Relationship Sets
 - Attributes (of entities and relationships)
- Additional Constructs
 - ISA Hierarchies
 - Weak Entities
 - Aggregation
- Integrity Constraints
 - Key constraints
 - Participation constraints
 - Overlap / Covering constraints for ISA hierarchies



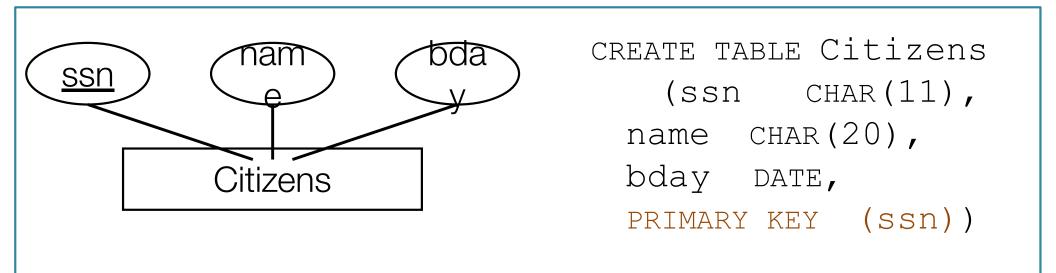
ER to Tables: Basics

- Strong entities:
 - key = primary key
- Relationships: keys come from the participating entities. E.g. for binary relationships:
 - 1-to-1: either key (other = candidate key)
 - 1-to-N: the key of the 'many' (N) part
 - M-to-N: both keys



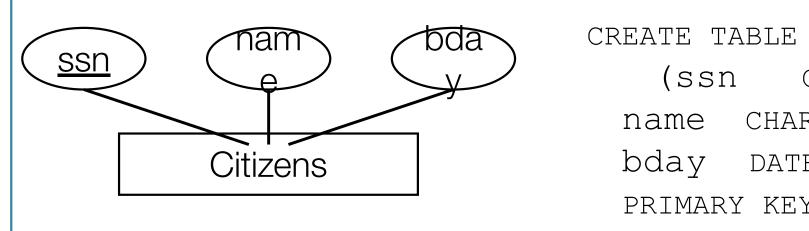
Entity Sets to Tables





Question??

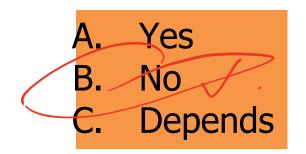




CREATE TABLE Citizens (ssn CHAR(11),name CHAR(20), bday DATE, PRIMARY KEY (SSn))

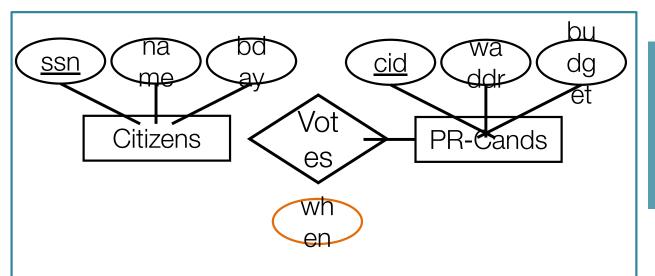
Can ssn have a null value?





Relationship Sets to Tables





- Foreign key: keys from participating entity sets
- Descriptive attributes

```
CREATE TABLE Votes (
ssn CHAR(11),
cid INTEGER,
when DATE,
PRIMARY KEY (ssn, cid),
FOREIGN KEY (ssn) REFERENCES Citizens,
FOREIGN KEY (cid) REFERENCES PR-Cands)
```

Generalizes to n-ary relationships

```
(we will see example later)

came cut = ) is read

rame

Sperify
```

Relationship Sets to Tables

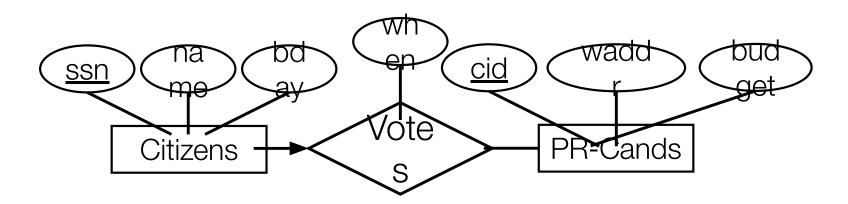
Same entity participate in tus different roles

```
Citizens
               CREATE TABLE Represents (
                                           (s'sn, ssn) > make no
sense!!
                   elected ssn CHAR(11),
                   cons ssn CHAR (11),
               PRIMARY KEY (elected ssn, cons ssn),
               FOREIGN KEY (elected ssn) REFERENCES Citizens (ssn)
               FOREIGN KEY (cons_ssn) REFERENCES Citizens (ssn))
Repre
```

Note that you need to specify the column that you are referring to in the Citizens table, as Citizens does not have "elected ssn" nor "cons ssn".

Key Constraints





Approach 1: Three Tables (Citizens, Votes, PR-Cands)

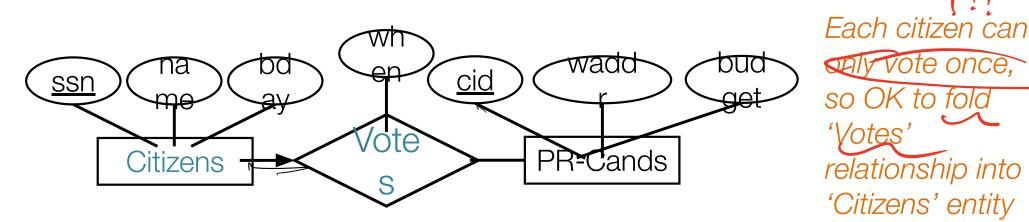
```
CREATE TABLE Votes

( ssn CHAR(11),
 cid INTEGER,
 when DATE,

PRIMARY KEY (ssn)

FOREIGN KEY (ssn) REFERENCES Citizens,
 FOREIGN KEY (cid) REFERENCES PR-Cands)
```

Key Constraints

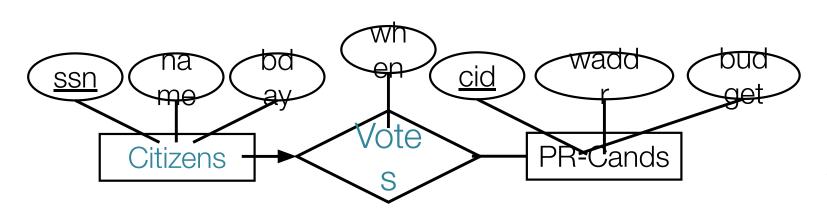


Approach 2: Two Tables (Citizen_Votes, PR-Cands)

```
CREATE TABLE
             Citizen Votes (
           CHAR (11),
   ssn
           CHAR (20),
   name
   bday
           DATE,
   when
           DATE,
   cid
           INTEGER,
PRIMARY KEY
              (ssn),
             (cid) REFERENCES PR-Cands)
FOREIGN KEY
```

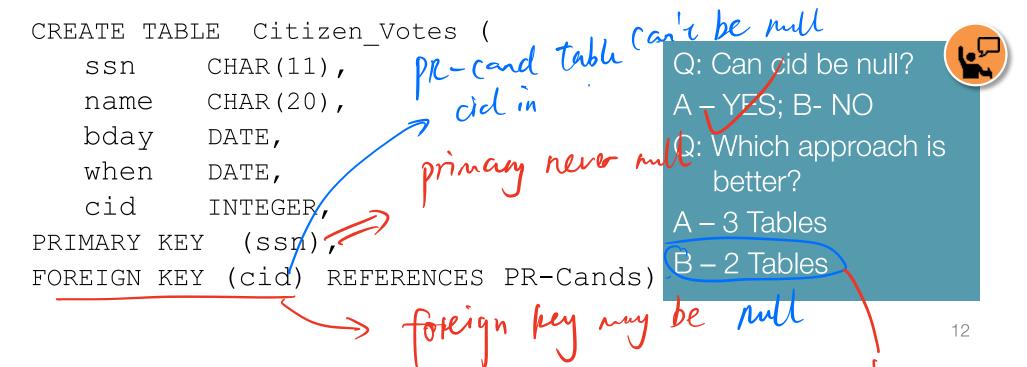
Question??

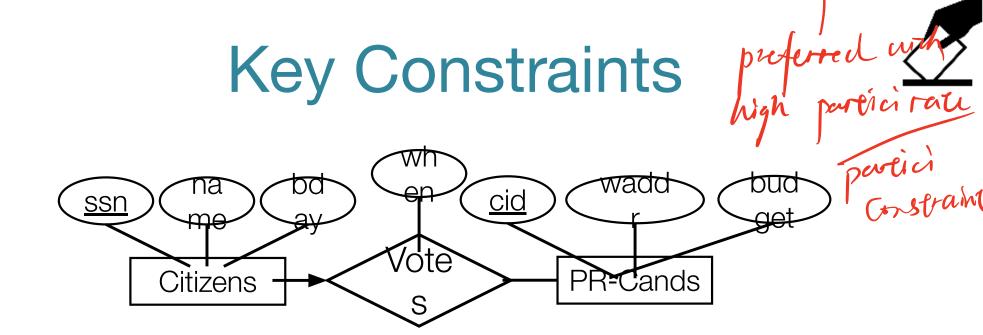




Each citizen can only vote once, so OK to fold 'Votes' relationship into 'Citizens' entity

Approach 2: Two Tables (Citizen_Votes, PR-Cands)





What about one table?

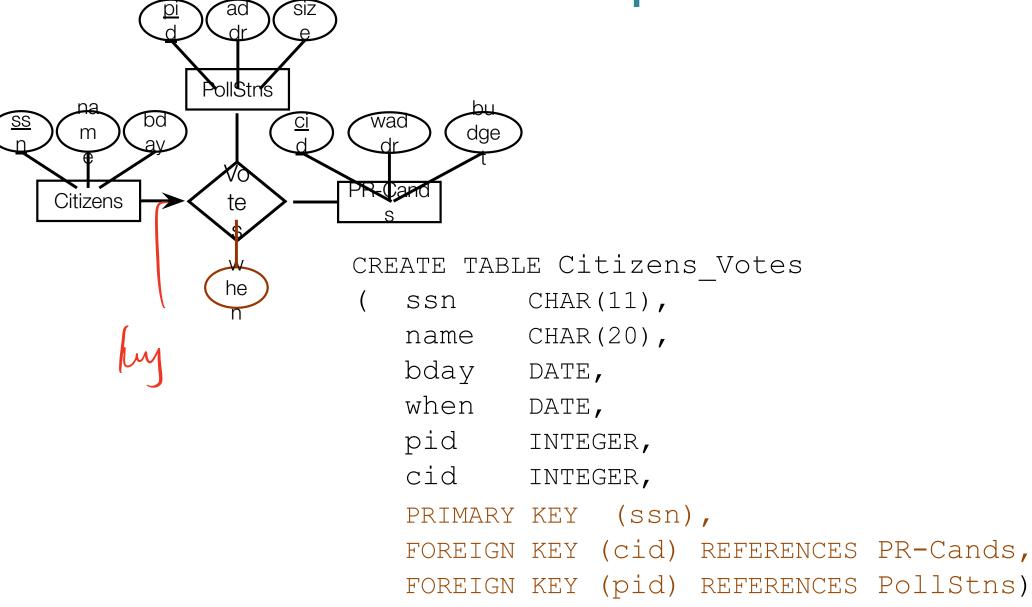


No! This is bad design.

e.g., For each citizen that votes for Candidate-X, we have to store Candidate-X information (cid, waddr, budget) => REDUNDANCY!

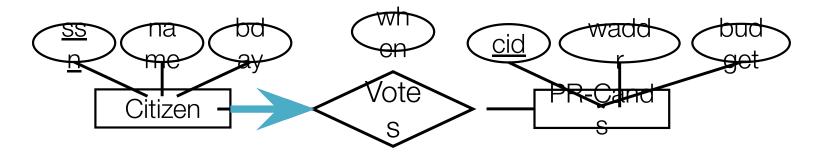
Key Constraints: N-ary relationships





Participation Constraints



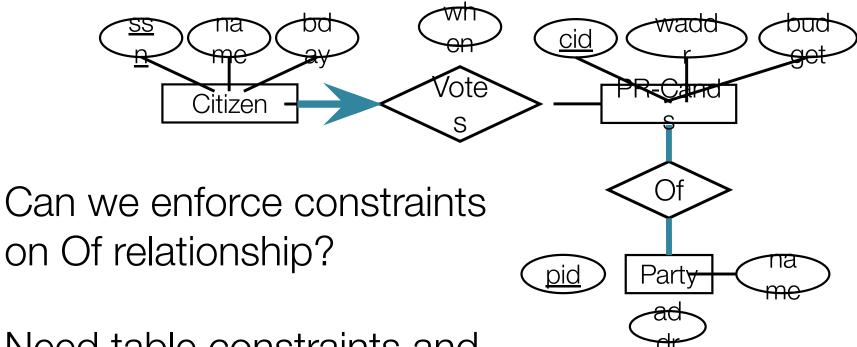


Using Approach 2

```
Citizen Votes(
CREATE TABLE
                                           Can we enforce the
            CHAR (11),
   ssn
                                          participation constraint
            CHAR (20),
   name
                                            using Approach 1
   bday
           DATE,
                                             (three tables)?
   when
            DATE,
   cid
            INTEGER NOT NULL,
                                     must vote.
              (ssn),
PRIMARY KEY
            (cid)
                     REFERENCES PR Cands
FOREIGN KEY
                         pon't allow (and airhdran)
   ON DELETE NO ACTION);
```

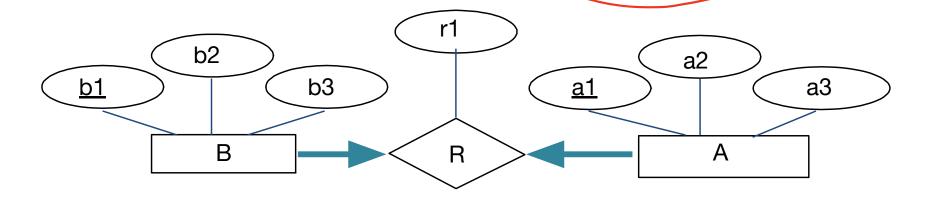
Participation Constraints





Need table constraints and assertions (later).

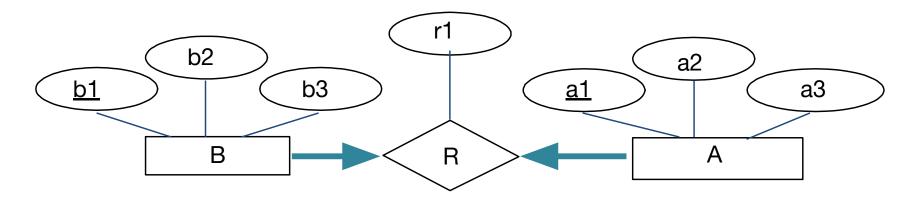
Mapping Participation Constraints (1-to-1)



```
CREATE TABLE RAB(
    r1 Integer,
    a1 Integer,
    a2 Integer,
    a3 Integer,
    b1 Integer,
    b2 Integer,
    b3 Integer ...)
```



Mapping Participation Constraints (1-to-1)



```
CREATE TABLE RAB(

r1 Integer,

a1 Integer,

a2 Integer,

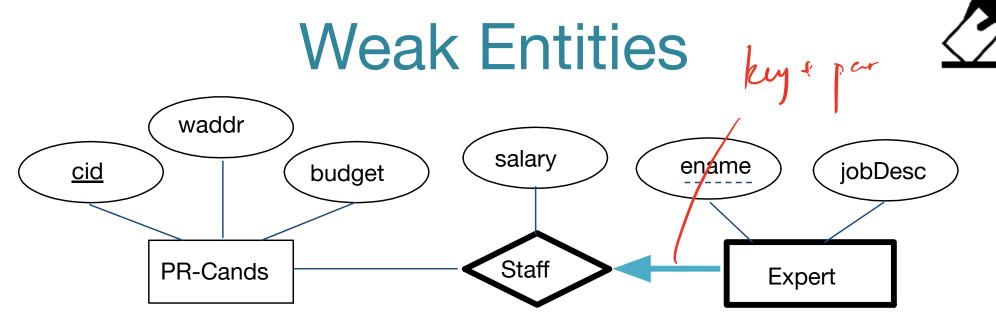
b1 Integer NOT NULL,

b2 Integer,

b3 Integer,

UNIQUE (b1), PRIMARY KEY (a1))

(adddate key
```



- Approach 2: Combine weak entity and owning relationship into one relation
 - Delete all weak entities when an owner entity is deleted.

```
create table Expert Staff ( ) the table design ename CHAR(20),
jobDesc CHAR(40),
salary REAL,
cid INTEGER,
PRIMARY KEY (ename, cid), part for weak entities
FOREIGN KEY (cid) REFERENCES PR-Cands
ON DELETE CASCADE)
```

ISA Hierarchies: General Approach separate table for powers and addition,

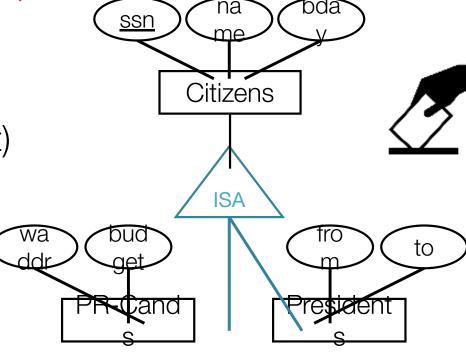
Three relations:

Citizens (<u>ssn</u>, name, bday)

PR-Cands (<u>ssn</u>, waddr, budget)

• Presidents (<u>ssn</u>, from, to)

- Queries:
 - Involving all citizens => Easy
 - Involving just PR-Cands => need to join PR-cands with Citizens to from Jase class get some attributes



ISA Hierarchies: Alternative

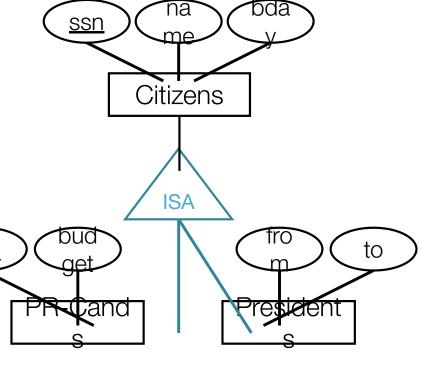


Two relations:

 PR-Cands (<u>ssn</u>, name, bday, waddr, budget)

 Presidents (<u>ssn</u>, name, bday, from, to)

+ citem (not pre/pre-can)



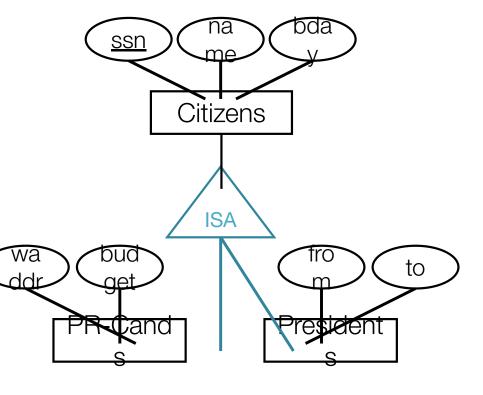
Question??



- Two relations:
 - PR-Cands (<u>ssn</u>, name, bday, waddr, budget)
 - Presidents (<u>ssn</u>, name, bday, from, to)

This design works best if we have:

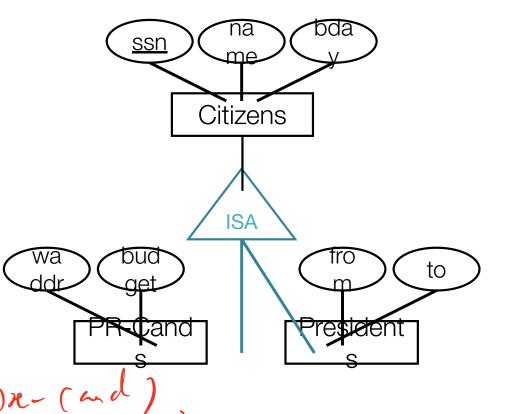
- Total Covering, Overlapping
- 3. Total Covering, Disjoint
- Partial Covering, Overlapping
- Partial Covering, Disjoint



ISA Hierarchies: Alternative



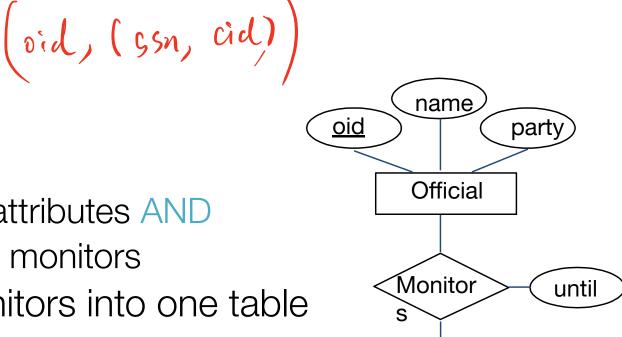
- Two relations:
 - PR-Cands (<u>ssn</u>, name, bday, waddr, budget)
 - Presidents (<u>ssn</u>, name, bday, from, to)
- Problems:
 - What if citizen is both?
 - Redundancy (pre + pre- (and
 - What if citizen is neither?
 - Use General approach

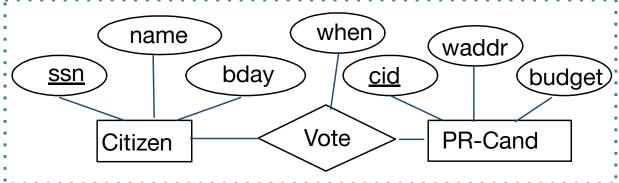


Aggregation



- Keep keys of all participating entity sets; decide primary
- Keys for Monitors
 - oid
 - (ssn, cid)
- Q: What if Vote:
 - has no descriptive attributes AND
 - total participation in monitors
- A: Fold Vote and Monitors into one table



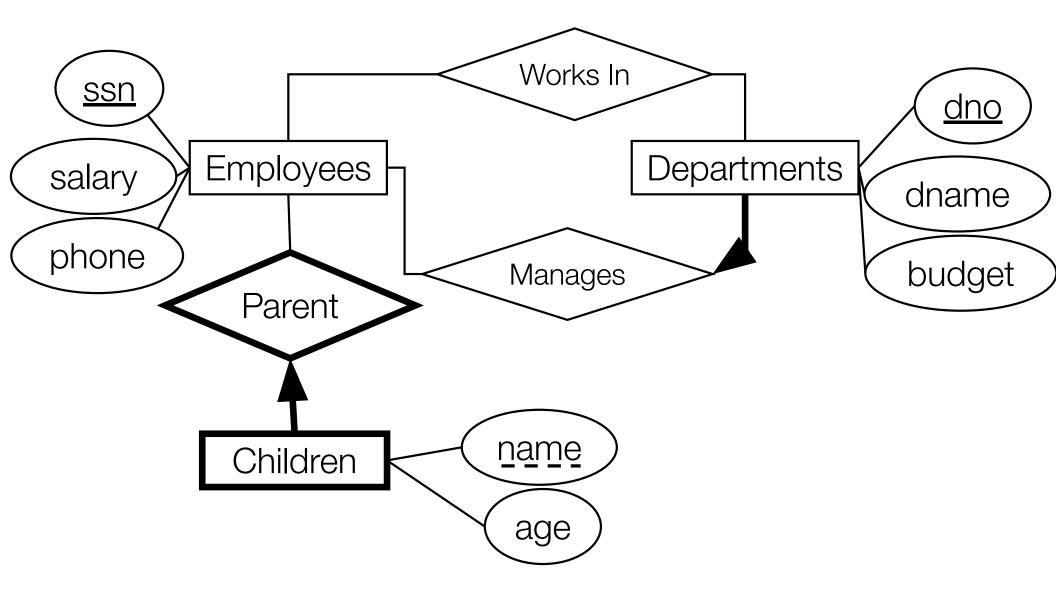


Exercise - Part 1

- A company database needs to store information about
 - employees (identified by <u>ssn</u>, with salary and phone attributes),
 - departments (identified by <u>dno</u>, with dname and budget attributes), and
 - children of employees (with name and age attributes).
- Employees work in (zero or more) departments
- Each department is managed by exactly one employee
- A child must be identified uniquely by name when the parent (who is an employee; assume only one parent works for the company) is known.
- We are not interested in information about a child once the parent leaves the company.
- Draw an ER diagram that captures this information

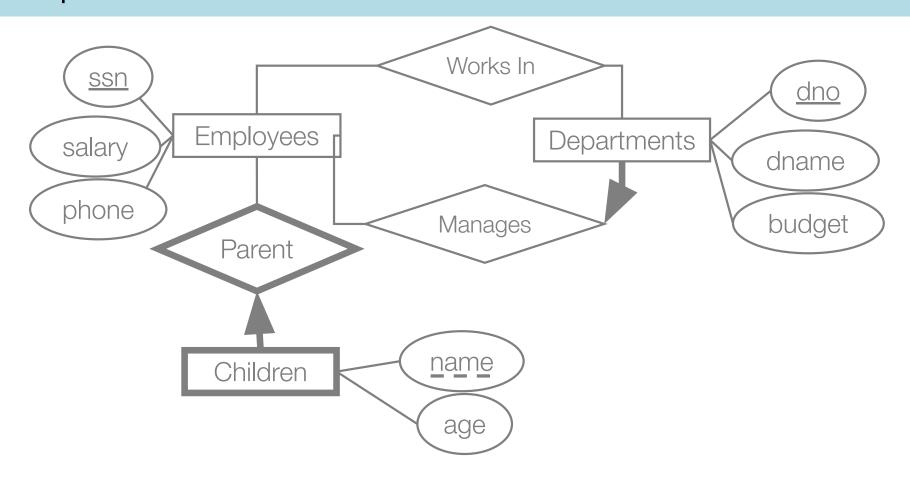
ER Diagram (one solution)





Exercise – Part 2

 Write SQL statements to create the corresponding relations, and to capture as many of the constraints as possible.



SQL DDL – One solution



```
CREATE TABLE Employees (
ssn INTEGER,
salary REAL,
phone CHAR(10),
PRIMARY KEY(ssn))
```

```
CREATE TABLE Works (
ssn INTEGER,
dno INTEGER,
PRIMARY KEY (ssn, dno),
FOREIGN KEY (ssn)
REFERENCES employees,
FOREIGN KEY (dno)
REFERENCES departments)
```

```
CREATE TABLE Departments (
dno INTEGER,
dname CHAR(20),
budget real,
manager INTEGER NOT NULL,
PRIMARY KEY (dno),
FOREIGN KEY (manager)
REFERENCES employees(ssn))
```

```
CREATE TABLE Children (
name CHAR(20),
age REAL,
parent INTEGER NOT NULL,
PRIMARY KEY(name, parent),
FOREIGN KEY(parent)
REFERENCES employees
ON DELETE CASCADE)
```

Integrity Constraints

- Describe conditions that must be satisfied by every legal instance
- Types of integrity constraints
 - Domain constraints
 - Primary key constraints
 - Foreign key constraints
 - General constraints

Table Constraints

- More general than key constraints
- Can use a query to express constraint
 - Constraints checked each time table updated
 - CHECK constraint always true for empty relation

```
CREATE TABLE Compete
CREATE TABLE Athlete
                          (aid INTEGER, oid INTEGER,
(aid INTEGER PRIMARY KEY,
                           PRIMARY KEY (aid, oid),
name CHAR(30),
                           FOREIGN KEY (aid)
age INTEGER,
                            REFERENCES Athlete
country CHAR(20),
                           CONSTRAINT noRussia
                                                      constrain
sport CHAR(20)),
                            CHECK (`Russia' <>
CHECK (age >= 18)
                                   (SELECT O.city
      AND age \leq 80)
                                  FROM Olympics O
                                   WHERE O.oid=oid)))
```



Try it out in sqlplus or sqlite

Do you get a constraint violation error?

Try it out in sqlplus or sqlite SQLite (Example in book 5.7)

Do you get a constraint violation error?

> Error: CHECK constraint failed: Sailors

More general CHECK



- Interlake boats cannot be reserved.
- Note: these are not supported in Oracle or SQLite

Constraints Over Multiple Relations

 For general constraint over multiple tables, use an assertion

Number of boats plus number of sailors is < 100

```
CREATE ASSERTION smallClub
CHECK
((SELECT COUNT (S.sid) FROM Sailors S) +
```

(SELECT COUNT (B.bid) FROM Boats B) < 100)

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Practical Considerations

- CHECK with subqueries and ASSERTIONS
 - Part of SQL standard (since 1992?)
 - But, they are not supported in many major databases
 - Because it's only been 30 years lolol
 - Main concern: Performance issues; CHECK or ASSERTION constraints over multiple tables are very slow
- Instead: Triggers
 - Most major database systems support them
 - Triggers are procedural

Active Databases & Triggers (5.8 in the book)

Trigger: Procedure that starts automatically if specified changes occur to the DBMS

- Three parts:
 - Event (activates the trigger)
 - Condition (test that is run when the trigger is activated)
 - Action (what happens if the trigger runs)
 - Before and After Triggers
- Trigger Execution
 - Row-level Triggers: Once per modified row
 - Statement-level Triggers: Once per SQL statement

Triggers

- When condition is checked, success/failure can be used to trigger arbitrary actions.
- Used for many things:
 - Often used to fill out fields in a form
 - Check complex actions (such as credit limit in a shopping application)
 - Check preferred customer status
 - Generate logs for auditing and security checks.
- Must specify when to check condition
 - Usually just before or just after an update

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Try it out



- Modify athlete_create.sql so that it has the UPDATE and DELETE constraints in COMPETE relation as in the previous slide.
 - Modified file available in athlete_modified.sql.
- Try the following and check COMPETE:

```
sqlite> DELETE FROM Athlete WHERE name='Michael
Phelps';
sqlite> UPDATE Athlete SET aid=5 WHERE aid=4;
```

keys = on; "command to enforce foreign key constraints. By default, SQLite ignores them for backward compatibility.

Implementation Notes





- Oracle's sqlplus:
 - You cannot use NO ACTION constraints. They are the default and thus not needed.
 - String literals like 'USA' must use single quotes, not double quotes



- SQLite:
 - You need PRAGMA foreign_keys = ON; to enforce foreign key constraints. This is for backward compatibility.



Oracle Trigger Example

- First trigger executed <u>before</u> the activating statement, second executes <u>after</u> the activating statement.
- In combination with:
 - FOR EACH ROW execute once per modified record
 - (default) execute once per activating statement
- Activating statements:
 - INSERT
 - DELETE
 - UPDATE

```
statement-level
CREATE TRIGGER init count
BEFORE INSERT ON Student
                                  /* Event */
   DECLARE
    count INTEGER; -> dedare,
                                   /* Action */
    count := 0
   END;
CREATE TRIGGER incr count
   FOR EACH ROW por level by the Event */
WHEN (new.age >= 18)

REGIN
AFTER INSERT ON Student,
                             /* Action */
   BEGIN
    count := count + 1;
   END;
```

Recall CASCADE constraints

```
CREATE TABLE Athlete
                              CREATE TABLE Olympics
(aid INTEGER PRIMARY KEY,
                               (oid INTEGER PRIMARY KEY,
                               year INTEGER,
 name CHAR(30),
 country CHAR (20),
                               city CHAR(20));
 sport CHAR(20));
        CREATE TABLE Compete
            (aid INTEGER,
            oid INTEGER,
            PRIMARY KEY (aid, oid),
            FOREIGN KEY (aid) REFERENCES Athlete
                ON DELETE CASCADE
            FOREIGN KEY (oid) REFERENCES Olympics
            );
```

CASCADE Using Triggers

```
CREATE TABLE Compete
   (aid INTEGER,
    oid INTEGER,
    PRIMARY KEY (aid, oid),
    FOREIGN KEY (aid) REFERENCES Athlete
    FOREIGN KEY (oid) REFERENCES Olympics);
CREATE OR REPLACE TRIGGER cascade on delete
AFTER DELETE ON Athlete
FOR EACH ROW
BEGIN
  DELETE FROM Compete
  WHERE Compete.aid = :OLD.aid;
END;
```

Trying out triggers



- Try out the file athlete_trigger_cascade.sql in sqlplus
- Also try it out in sqlite to see if it supports triggers the same way
- Try removing a row from athlete. Does it cascade to Compete via the trigger?
- Try dropping the trigger:
 - DROP TRIGGER triggername;
- What happens now on deleting a row from Athlete?

Triggers: Pitfalls and Pain

- Triggers can be recursive!
 - Chain of triggers can be hard to predict, which makes triggers difficult to understand and debug
- Errors with "mutating" table
 - A table that is currently being modified by an UPDATE,
 DELETE, or INSERT statement, or a table that might be updated by the effects of a DELETE CASCADE constraint
 - The session that issued the triggering statement cannot query or modify a mutating table
 - Used to prevent a trigger from seeing inconsistent data