# Digital search structure

Ch. 12

## Ch. 12: Digital search structure

Digital Search Tree



Reduce number of comparisons

Trie



Reduce required space for index

Compressed trie



Reduce required space for data

Patricia

## 1. Digital search tree

A binary tree. Each node contains one data pair.

- Keys are binary bit strings, indicating the path of search.
  - For previously introduced structures, key values are meaningful, and we do compare based on key values.
  - Fixed length: 0110, 0010, 1010, 1011.
  - Variable length: 01, 00, 101, 1011.

- Applications
  - Networks, e.g., hardware IP routing

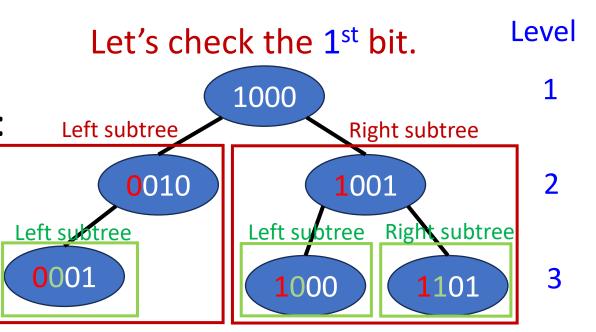
### Digital search tree

- Assume that the key is fixed length.
- Number the bit from left to right.

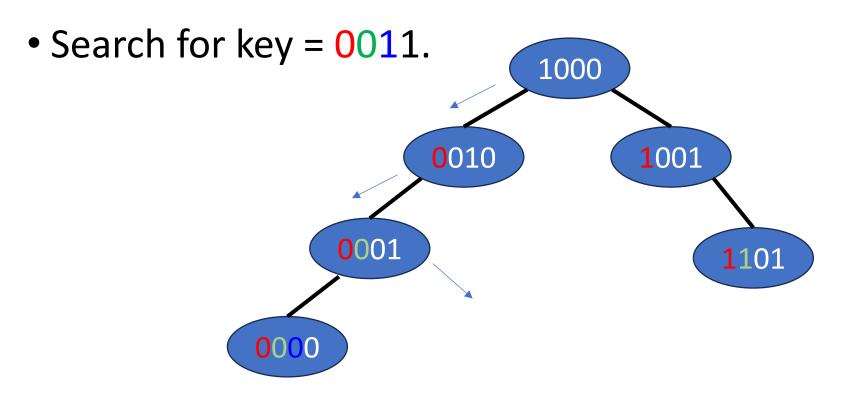
Bit 1 2 3 4 0 1 1 0

- For a node at level *i*:
  - Keys of pairs in its left subtree:
     the *i*-th bit = 0.
  - Keys of pairs in its right subtree:
     the *i*-th bit = 1.

Let's check the 2nd bit.

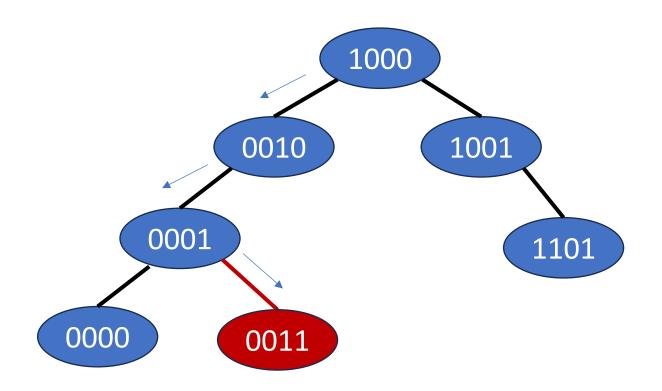


# Operation: Search



- Compare with the key in the root (1000).  $\rightarrow$  different
- 1st bit is 0. Compare with the key in the left child (0010).  $\rightarrow$  different
- $2^{nd}$  bit is 0. Compare with the key in the left child (0001).  $\rightarrow$  different
- $3^{rd}$  bit is 1. Compare with the key in the right child.  $\rightarrow$  No right child.
- 0011 is not in the search tree.

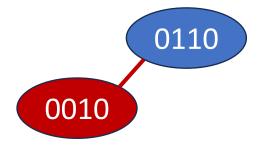
• Insert a data pair whose key = 0011.



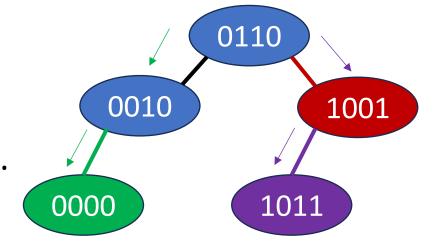
- Start with an empty digital search tree.
- Insert a pair whose key is **0110**.



- Then, insert a pair whose key is 0010.
  - The first bit is 0. Move to the left child.



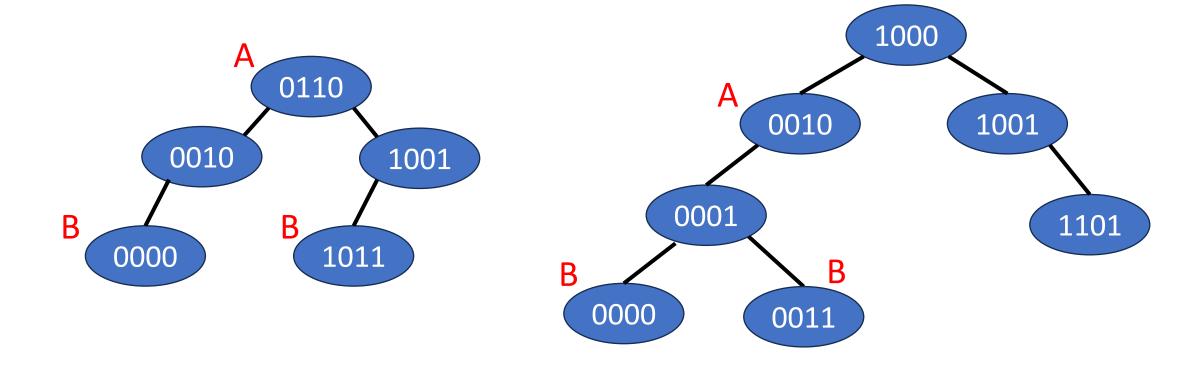
- Then, insert a pair whose key is 1001.
  - The first bit is 1. Move to the right child.
- Then, insert a pair whose key is 1011.
  - The first bit is 1. Move to the right child.
  - The second bit is 0. Move to the left child.



- Then, insert a pair whose key is 0000.
  - The first bit is 0. Move to the left child.
  - The second bit is 0. Move to the left child.

## Operation: Deletion

- Delete node A.
  - Find any leaf B in the subtree rooted at A.
  - Then replace A with B.



### Exercise

• Q5: Given the following keys 0000, 0011, 0001, 0010, 1000, 1001, 1100, 0110. Following this order, please insert the keys into an empty digital search tree. What will be the keys in the nodes of the 3rd level of the resultant tree?

Q6: (Continue Q5) When deleting the node whose key is 0011, what

is(are) the key(s) to replace key 0011?

Please reply your answers of Q5-Q6 via the following link:

https://forms.gle/NyCVWyzZ6eXwVdYg9

Group members: 2~4 people

### Discussion

Complexity of each operation is O(#bits in a key).

- #key comparisons = O(height).
  - Starting from root, we compare the key with each of the node along the path.
  - The height is at most 1 + #bits in a key.
- Expensive when keys are very long.
  - If keys are computed by SHA-1 hash function with 160 bits, the cost of a key comparison is high.

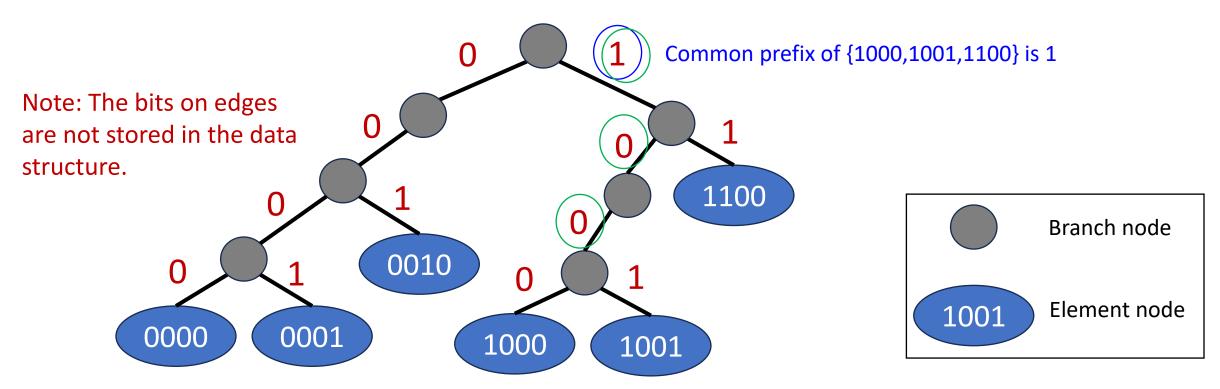
Can we reduce the number of key comparisons done during a search?

## 2. Binary Trie

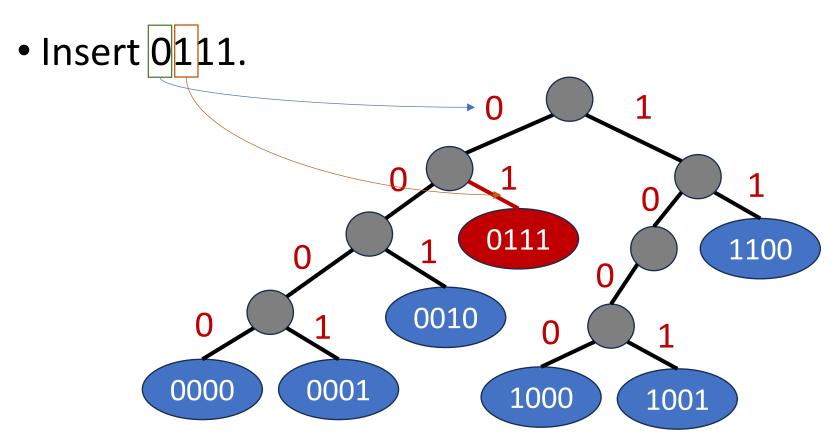
- Information Retrieval.
- At most one key comparison per operation.
- Fixed length keys.
- Two types of nodes:
  - Branch nodes
    - Left and right child pointer
    - No data field
  - Element nodes
    - No child pointer
    - Data field to hold dictionary data pair

# Example

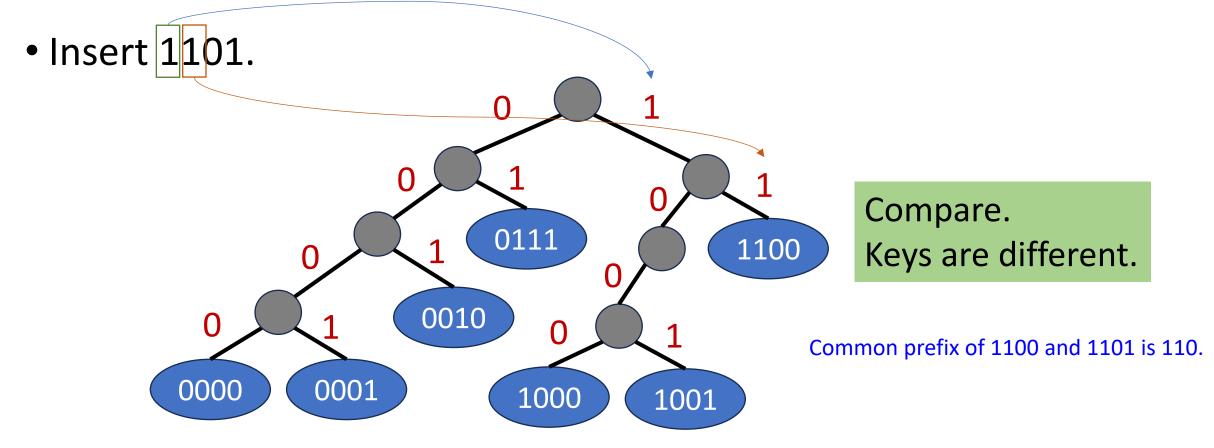
- Use bits to guide the direction to traverse top-down the tree.
- At most one key comparison.
  - The traversal does not introduce comparisons.
  - When finally reaching an element node, comparing with its key.



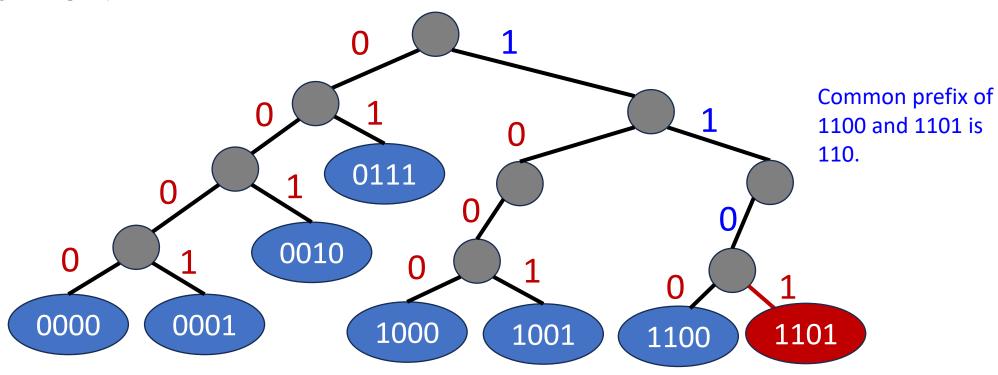
Common prefix of {1000,1001} is 100



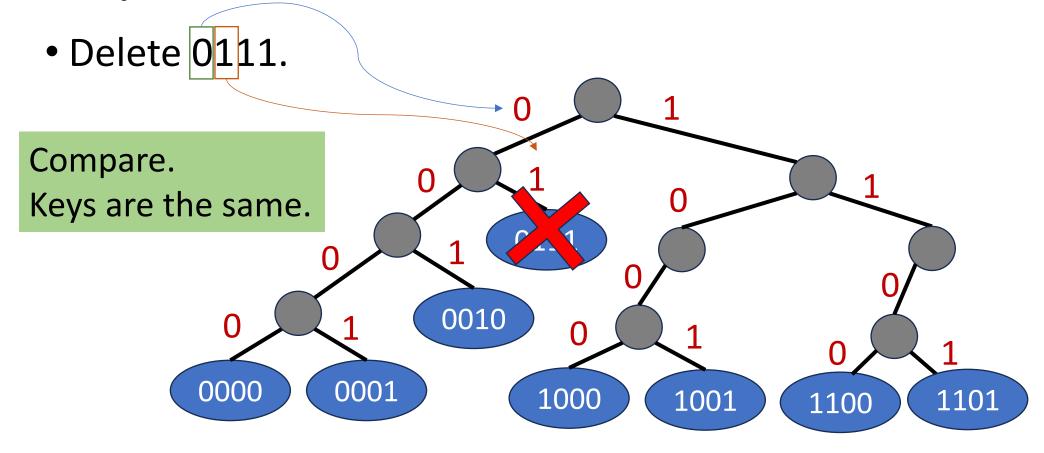
Zero compare.



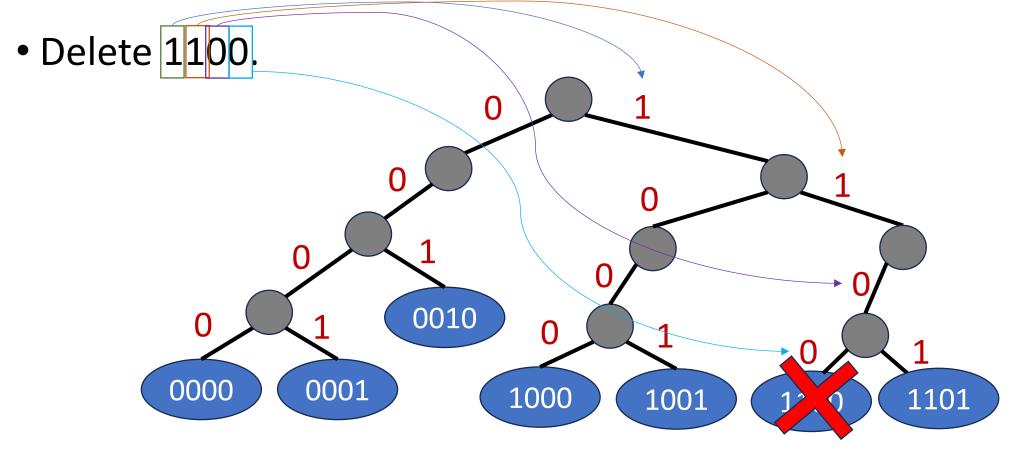
• Insert 1101.



# Operation: Delete



# Operation: Delete

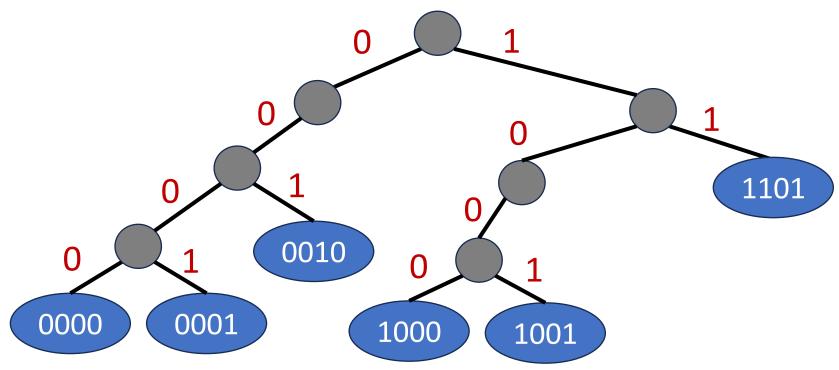


Compare.

Keys are the same.

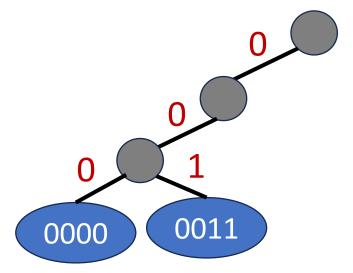
# Operation: Delete

• Delete 1100.



#### Exercise

• Q7: Given the keys 0010, 1000, 1100, 1101. Following this order, please insert the keys into the following binary trie. What will be the keys in the element nodes of the 5th level of the resultant tree?

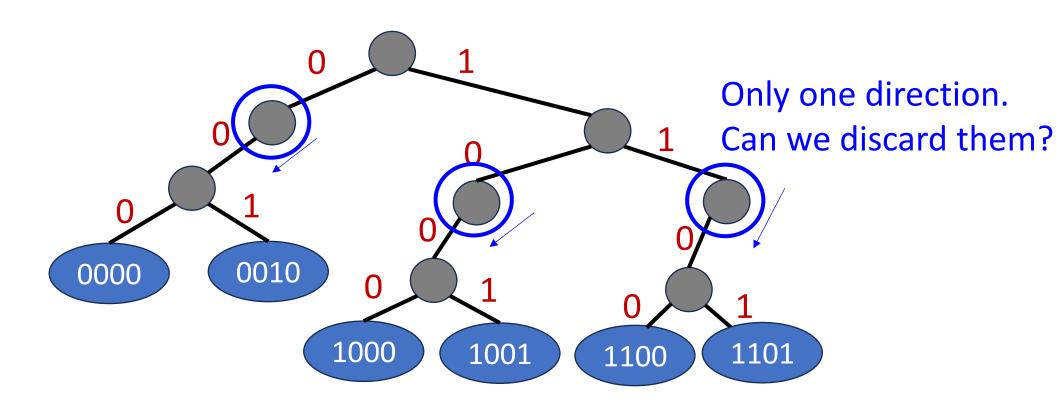


 Q8: (Continue Q7) After deleting the node whose key is 0011, what will be the keys in the element nodes of the 5th level of the resultant tree?



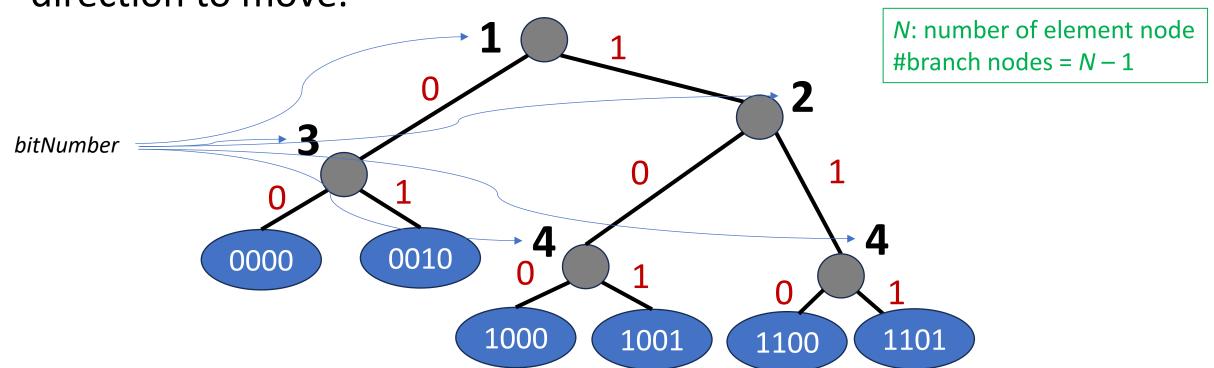
### Discussion

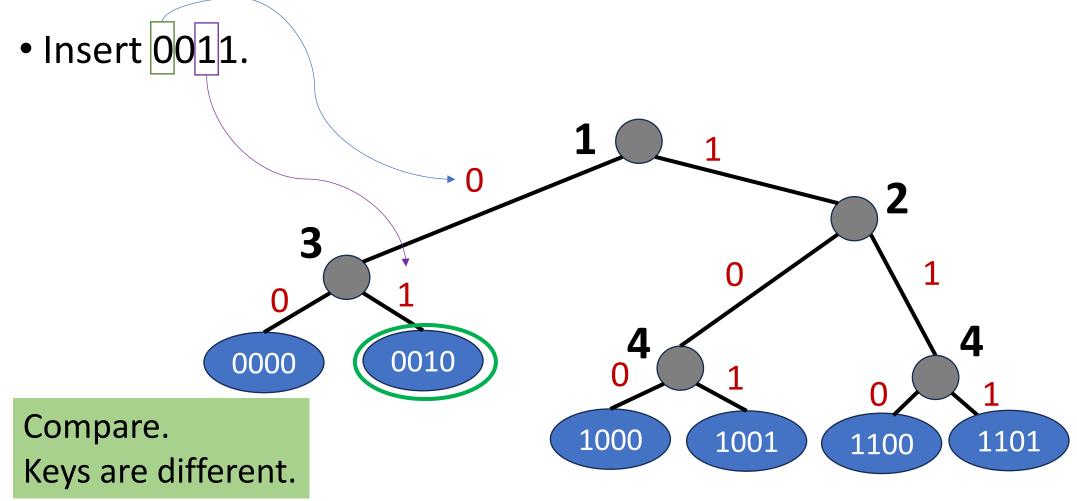
- 6 element nodes.
- 8 branch nodes.
- Too much overhead.



## 3. Compressed Binary Trie

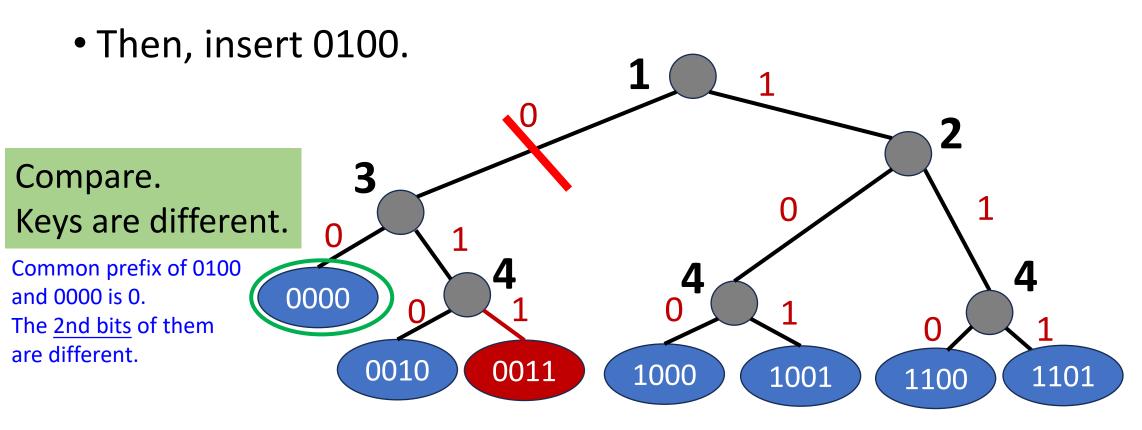
- No branch node whose degree is 1.
- Add a bitNumber field to each branch node
- bitNumber tells you which bit of the key to use to decide the direction to move.





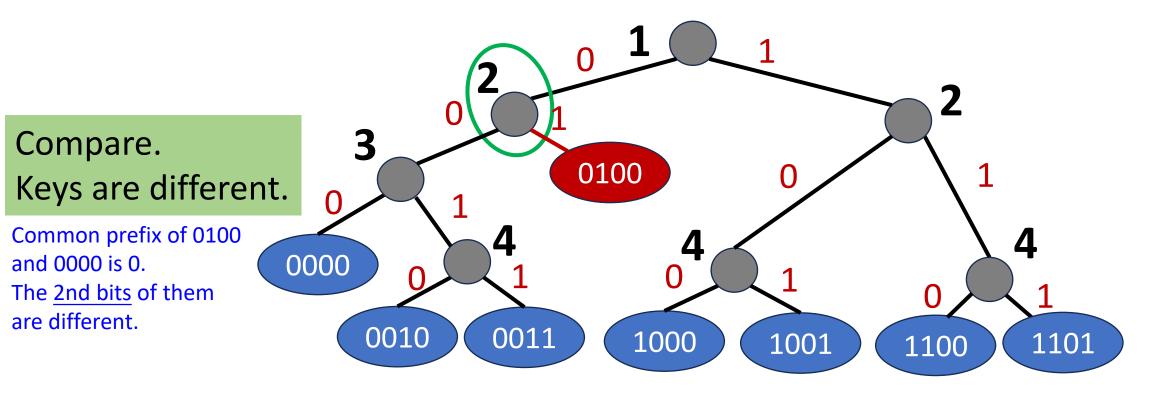
Common prefix of 0011 and 0010 is 001.

• Insert 0011.



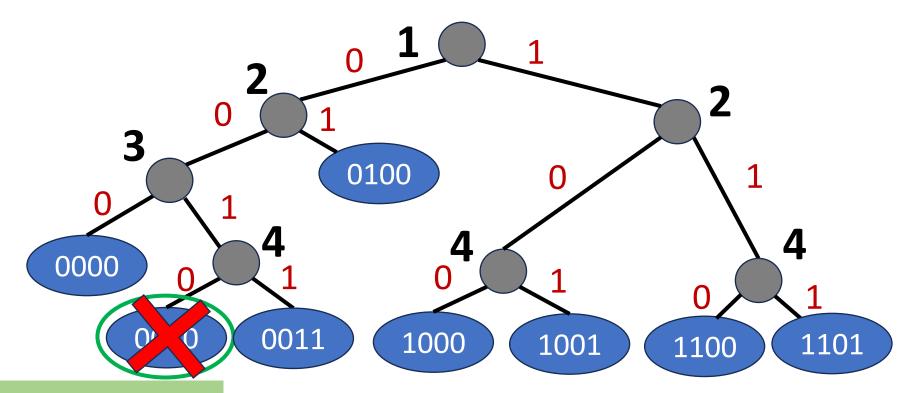
Common prefix of 0011 and 0010 is 001. The 4<sup>th</sup> bits of 0011 and 0010 are different.

• Insert 0100.



# Operation: Deletion

• Delete 0010.

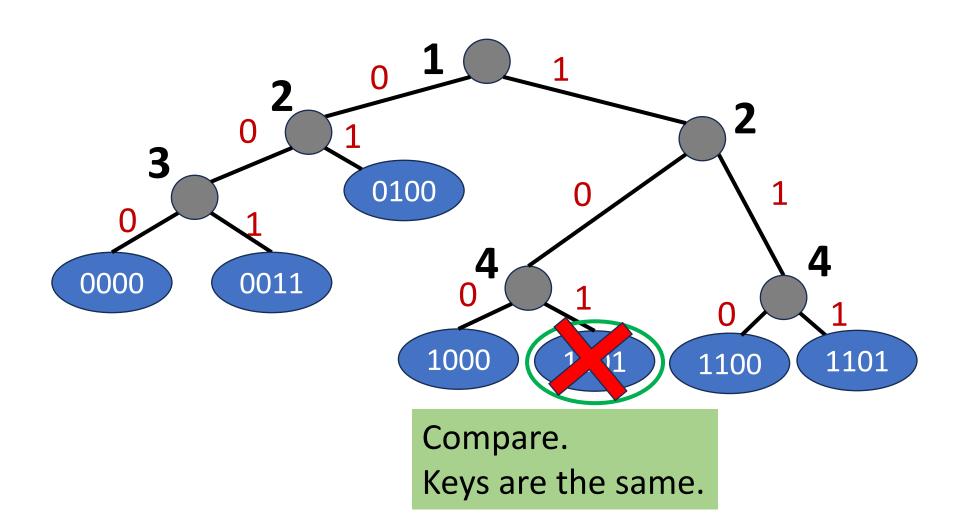


Compare.

Keys are the same.

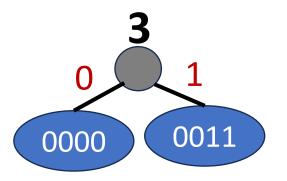
# Operation: Deletion

• Delete 1001.



#### Exercise

Q9: Given the keys 0010, 1000, 1100, 1101. Following this order, please insert the keys into the following compressed binary trie.
 What will be the keys in the element nodes of the 3rd level of the resultant tree?





### Exercise

• Q10: (Continue Q9) After deleting the node whose key is 0011, what will be the keys in the element nodes of the 3rd level of the resultant tree?



### Performance

Digital Search Tree ✓



Reduce number of comparisons (decrease time)

Performance is improved.

- Space
- Time

- Trie ✓

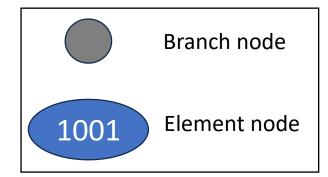
Reduce required space for index (branch node)

Compressed trie



Reduce required space for data (element nodes)

Patricia



# Summary

- Digital search tree
- Trie
- Compressed trie

- Operations:
  - Search
  - Insertion
  - Deletion