

CSE 351: The Hardware/Software Interface

Section 5

Structs as parameters, buffer overflows, and lab 3

Structs as parameters

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- In the previous section, we looked at how integers and pointers were passed as arguments to functions
- If we were to pass a struct by value to a function (as in not passing a pointer), how would the compiler use assembly instructions to achieve this?
- Let's take a look at an example (see `pass_struct.c` from the section material)

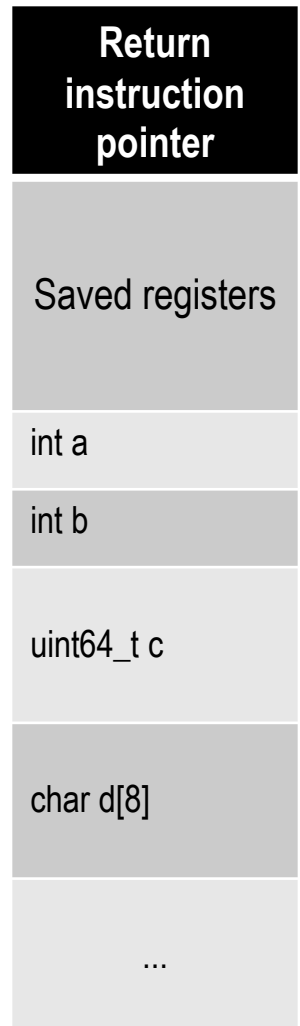
Buffer overflows

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- **C performs no bounds-checking on array accesses; this makes it fast but also unsafe**
 - What would we need to add to C to support checked array accesses?
- **For example: `int arr[10]; arr[15] = 3;`**
 - No compiler warning, just memory corruption
- **What symptoms are there when programs write past the end of arrays?**

Stack layout

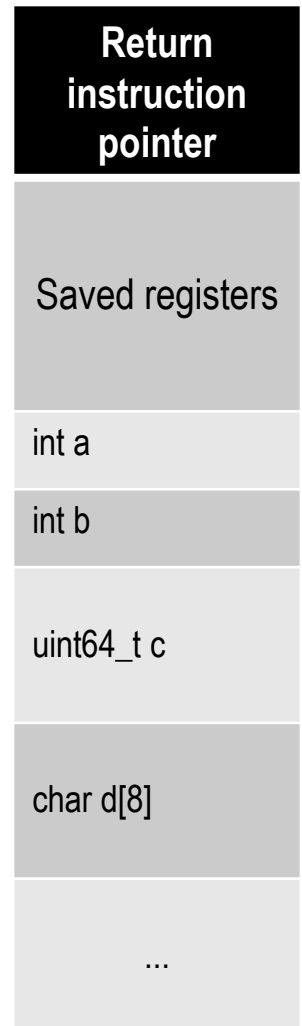
- As we've seen previously, when values are declared on the stack, the compiler shifts `%rsp` (in x86-64 assembly) to allocate space for them
- When a function returns, the return instruction pointer indicates where to begin executing again



Stack layout

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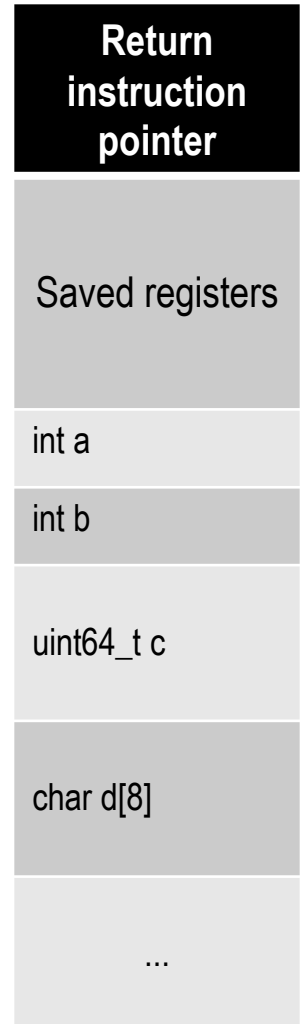
- Note that the top of the diagram represents higher addresses, and the bottom is lower addresses
- To which memory does `d[10]` refer in this example?



Buffer overflow attacks

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- In buffer overflow *attacks*, malicious users pass values to attempt to overwrite important parts of the stack or heap
- For example, an attacker could overwrite the return instruction pointer with the address of a malicious block of code



Buffer overflow attacks

- C has some inherently unsafe functions that facilitate buffer overflows, including `gets` and `strcpy`
- `gets(char* s)` reads from standard input until reaching a newline character (`'\n'`) or EOF (end of file)
 - How long should `s` be to contain the entire input string?
- `strcpy(char* dest, const char* src)` copies the contents of the `src` string into the `dest` string
 - What happens if `dest` is smaller than `src`?

Protecting against overflows

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- **As a programmer, you can protect against buffer overflow bugs/attacks by checking buffer lengths and using safer string-related functions**
 - `fgets(char* s, int size, FILE* stream)` takes a size parameter and will only read that many bytes from the given input stream
 - `strncpy(char* dest, const char* src, size_t n)` will copy at most `n` bytes from `src` to `dest`

Protecting against overflows

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■ Stack canaries

- At runtime, programs place a (pseudo-)random integer on the stack immediately before the return instruction pointer. If the integer value doesn't match when the function returns, the program generates a segmentation fault

■ Data execution prevention

- Some parts of memory (notably the stack) are marked as non-executable. The CPU will refuse to execute instructions from such locations and the program will terminate

Lab 3: Buffer overflows

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- The purpose of lab 3 is to become familiar with how buffer overflow attacks work
- The various stages of the lab require different types of attacks to achieve certain goals
- If you have become comfortable with GDB and understanding assembly instructions, you should have no problem

Lab 3: Buffer overflows

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- The exploitable function in lab 3 is called `Gets` (capital 'G') and is called from the `getbuf` function
- `getbuf` allocates a small array and reads user input into it via `Gets`. If the user input is too long, then certain values on the stack within the `getbuf` function will be overwritten...

Lab 3: Buffer overflows

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- The first thing to do is to become familiar with the provided tools for the lab
- To generate malicious strings for testing buffer overflows, use the provided `sendstring` tool. It takes a list of space-separated hex values and translates them to the corresponding Ascii characters
- Each lab is slightly different as determined by the username given to it; when you run the `bufbomb` binary, you have to pass in “`-u [cse-username]`”
- Let's take a look at how this works