Chun-Wei Chen CSE 351 Homework 4 05/26/13

1.

A. When reading x[1][0] after reading x[0][0], x[0][0] along with x[0][1], x[0][2], and x[0][3] will be evicted from the cache; therefore, when it will be a miss when reading x[0][1]. Consequently, miss rate is 100%.

B. If we double the cache size, we won't encounter thrashing in case 1, that is, when reading x[0][1] after reading x[1][0], we'll have a hit instead of miss since the block that contains x[0][1] isn't evicted by the block the contains x[1][0]; therefore, we'll only face cold misses in this case and the miss rate is 25%.

C. Two-way set associative cache eliminates the thrashing problem in case 1 and we'll only face cold misses in this case; therefore, the miss rate is 25%.

D. No, it won't reduce miss rate since increasing the cache size won't eliminate cold misses; in other words, we'll still encounter misses whenever we read x[0][4k] or x[1][4k] where $k \ge 0$ and $k \le 31$ because the block size is still 16 bytes. So the miss rate is still 25%.

E. Yes, it'll reduce the miss rate. For instance, if we increase the block size by a multiple of 2, we'll now encounter misses whenever we read x[0][8k] or x[1][8k] where $k \ge 0$ and $k \le 15$; therefore, the miss rate is reduced by a multiple of 2 because of increasing the block size.

2. Possible outputs: 243, 423, 432

3. n = 64 and P = 4 KB: # PTEs = 2^{64} / $(4*2^{10}) = <math>2^{64}$ / $2^{12} = 2^{52}$ (or 2^2 P) n = 64 and P = 4096 KB: # PTEs = 2^{64} / $(4096*2^{10}) = <math>2^{64}$ / $2^{22} = 2^{42}$ (or 2^2 T)

4. a 64-bit virtual address space and a 36-bit physical address

Р	# VPN bits	# VPO bits	# PPN bits	# PPO bits
4 KB	52	12	24	12
4096 KB	42	22	14	22

5. Virtual address: 0x027c

A.

13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	0	1	1	1	1	1	0	0

В.

Parameter	Value				
VPN	0x09				
TLB index	0x1				
TLB tag	0x02				

			hit? (Y			N N							
		Page	fault? PPN	(Y/IN)			0x17						
									ONLY				
	C							2	2	1	0		
11	10	9	8	7 1	6	5	4	3	2		0		
							<u> </u>	<u> </u>					
D.							Г						
			aramet							Value			
			yte offs che ind							0x0 0xF			
			ache ta							0xF			
			ne hit? (N			
			byte re										
6. Virtual address: 0x03a9 A.													
13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	0	1	0	1	0	0	1
В.													
		Pa	aramet	er			Value						
			VPN				0x0E						
			LB inde				0x2						
			TLB tag				0x03 N						
			hit? (Y fault?				N N						
		rage	PPN	(1/14)			0x11						
C.	4.0	•	•	_	_	_	_	2	•		•		
11	10	9	8	7	6	5 1	4	3	2	0	0	1	
0	1	U	U	U	<u> 1</u>	I	U	1	U	U	1		
D.	D.												
Parameter										Value			
			yte offs				0x1						
			che ind				0xA						
			ache ta				0x11						
Cache hit? (Y/N)							N						

Cache byte returned

7. Virtual address: 0x0040

A.

13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	1	0	0	0	0	0	0

B.

Parameter	Value				
VPN	0x01				
TLB index	0x1				
TLB tag	0x00				
TLB hit? (Y/N)	N				
Page fault? (Y/N)	Υ				
PPN					

C.

11	10	9	8	7	6	5	4	3	2	1	0

D.

Б.	
Parameter	Value
Byte offset	
Cache index	
Cache tag	
Cache hit? (Y/N)	
Cache byte returned	