

I use DPLL with BCP and PLP as the SAT solver in bonus.

1. DPLL

My DPLL is based on the pseudo-code from the class as following:

```
bool DPLL(CNF  $\phi$ , AssignMap A)
{
  1.  $\phi' = \text{BCP}(\phi, A)$ 
  2.  $\phi'' = \text{PLP}(\phi')$ 
  3. if( $\phi'' = \top$ ) then return SAT;
  4. else if( $\phi'' = \perp$ ) then return UNSAT;
  5.  $p = \text{choose\_var}(\phi'')$ ;
  6. if(DPLL( $\phi''$ , A[ $p \mapsto \top$ ])) then return SAT;
  7. else return (DPLL( $\phi''$ , A[ $p \mapsto \perp$ ]));
}
```

Like line (1), first call BCP. In order to improve the performance of DPLL, I judge whether cnf is sat after BCP. If it is, there is no need to call the PLP.

```
std::vector<std::vector<int>> modifiedCnf = BCP(cnf, &: assignMap);
if (modifiedCnf.empty()) {
    return true;
} else if (fomularIsFalse(modifiedCnf)) {
    return false;
}
```

Like line (2)-(4), call PLP and then judge whether cnf is sat. If cnf is empty, then it is sat. If cnf has a clause like [0], then it is unsat.

```
modifiedCnf = PLP(numVar, cnf: modifiedCnf, &: assignMap);
if (modifiedCnf.empty()) {
    return true;
} else if (fomularIsFalse(modifiedCnf)) {
    return false;
}
```

Like line (5), choose a variable that is not assigned a value.

```
int p = chooseVar(numVar, assignMap);
```

Like line (6)-(7), Call the DPLL recursively with p is true and p is false.

```
std::map<int, bool> assignMapPTrue = assignMap;
assignMapPTrue.insert( v: { &: p, u2: true});
std::map<int, bool> assignMapPFalse = assignMap;
assignMapPFalse.insert( v: { &: p, u2: false});
if (DPLL(numVar, cnf: modifiedCnf, &: assignMapPTrue)) {
    return true;
} else {
    return DPLL(numVar, cnf: modifiedCnf, &: assignMapPFalse);
}
```

2. BCP

My BCP is based on the pseudo-code from Piazza as following:

BCP maintains a queue or map A of (variable, value) pairs. This is a list of unit clauses and decisions.

- 1: for each pair (v, u) in A
- 2: for each clause C in ϕ
- 3: if $((v$ occurs positively in C and u is the value true) OR
 v occurs negatively in C and u is the value false)) then
 mark the clause C as satisfied and remove from ϕ
- 4: if $((v$ occurs positively in C and u is the value false) OR
 v occurs negatively and u is the value true)) then
 mark the literal as false, and shrink the clause C by dropping v
- 5: if C becomes unit, add it to the map A with appropriate value and remove it from ϕ
- 6: Return the modified formula //Note that the formula ϕ is being modified by either dropping clauses (step 3 and step 5) or dropping literals (step 4)

First use a list to record the order of keys in *assignMap*. Because line (5) will add pair to *assignMap* and the map in C++ will be automatically sorted based on the key. So, the adding will affect the for loop for the pair in *assignMap*. I use a list to avoid that.

```
std::list<int> assignMapOrder;
for (mapIter = assignMap.rbegin(); mapIter != assignMap.rend(); ++mapIter) {
    assignMapOrder.push_back( x: mapIter->first);
}
```

Like line (1)-(2), for loop the pair in *assignMap*, the clause in *cnf* and the literal in clause.

```
for (orderIter = assignMapOrder.begin(); orderIter != assignMapOrder.end(); ++orderIter) {
    for (int i = 0; i < modifiedCnf.size(); i++) {
        for (int j = 0; j < modifiedCnf[i].size(); j++) {
```

Like line (3), if the clause has a literal that is true, remove the clause from *cnf*.

```
if ((*orderIter == modifiedCnf[i][j] && assignMap[*orderIter]) ||
    (*orderIter == 0 - modifiedCnf[i][j] && !assignMap[*orderIter])) {
    modifiedCnf.erase( position: modifiedCnf.begin() + i);
```

Like line (4), if the clause has a literal that is false, remove the literal from the clause. And if the literal is the only literal in the clause, make the clause be [0], which means the clause is false.

```

else if ((*orderIter == modifiedCnf[i][j] && !assignMap[*orderIter]) ||
        (*orderIter == 0 - modifiedCnf[i][j] && assignMap[*orderIter])) {
if (modifiedCnf[i].size() == 1) {
    modifiedCnf[i][j] = 0;
} else {
    modifiedCnf[i].erase( position: modifiedCnf[i].begin() + j);
    j--;
}
}

```

Like line (5), if the clause is unit, judge whether the literal is in *assignMap*. If it is in, like the red square, if the literal is false, then make the clause be [0]. Otherwise, remove the clause from *cnf*. If the literal is not in *assignMap*, like the blue square, if the literal > 0 , add it into *assignMap* with true; if the literal < 0 , add it into *assignMap* with false. And then remove the clause from *cnf*.

```

if (modifiedCnf[i].size() == 1 && modifiedCnf[i][0] != 0) {
    if (assignMap.count( k: modifiedCnf[i][0]) == 1) {
        if (!assignMap[modifiedCnf[i][0]]) {
            modifiedCnf[i][0] = 0;
        } else {
            modifiedCnf.erase( position: modifiedCnf.begin() + i);
            i--;
            break;
        }
    } else if (assignMap.count( k: 0 - modifiedCnf[i][0]) == 1) {
        if (assignMap[0 - modifiedCnf[i][0]]) {
            modifiedCnf[i][0] = 0;
        } else {
            modifiedCnf.erase( position: modifiedCnf.begin() + i);
            i--;
            break;
        }
    }
} else {
    if (modifiedCnf[i][0] > 0) {
        assignMap.insert( v: { &: modifiedCnf[i][0], u2: true});
        assignMapOrder.push_back( x: modifiedCnf[i][0]);
    } else {
        assignMap.insert( v: { u1: 0 - modifiedCnf[i][0], u2: false});
        assignMapOrder.push_back( x: 0 - modifiedCnf[i][0]);
    }
    modifiedCnf.erase( position: modifiedCnf.begin() + i);
    i--;
    break;
}
}
}

```

3. PLP

My PLP is based on the pseudo-code from Piazza as following:

```

PLP ( Formula F , Assignment & A ) {
    L = []
    for every variable v in the formula F {
        if ( v occurs only positively )
            A[v] = True
            L.append( v )
        else if ( v negatively in the entire formula F )
            A[v] = False
            L.append( v )
    }
    for each variable u in L {
        for each clause C in F
            if ( u occurs in C )
                drop C from F
    }
}

```

First, find and use *reverseList* to store all the non-pure literal in cnf. Perform a for loop on cnf, when *literal* > 0, *pureLiteralCount*++. When *literal* < 0, *pureLiteralCount*++. If *literal* > 0 and *pureLiteralCount* < 0 or *literal* < 0 and *pureLiteralCount* > 0, the literal is not pure literal, add it to *reverseList*.

```

for (int i = 0; i < modifiedCnf.size(); i++) {
    for (int j = 0; j < modifiedCnf[i].size(); j++) {
        if (modifiedCnf[i][j] > 0 && pureLiteralCount[modifiedCnf[i][j]] >= 0) {
            pureLiteralCount[modifiedCnf[i][j]]++;
        } else if (modifiedCnf[i][j] < 0 && pureLiteralCount[0 - modifiedCnf[i][j]] <= 0) {
            pureLiteralCount[0 - modifiedCnf[i][j]]--;
        } else {
            reverseList.push_back( x: abs(modifiedCnf[i][j]));
        }
    }
}

```

Then, all the variables that is not in *reverseList* is pure literal, add it to list and *assignMap*.

```

for (int i = 1; i <= numVar; i++) {
    if (std::find( first: reverseList.begin(), last: reverseList.end(), value_: i) == reverseList.end() && pureLiteralCount[i] > 0) {
        list.push_back( x: i);
        assignMap.insert( v: { &i, u2: true});
    } else if (std::find( first: reverseList.begin(), last: reverseList.end(), value_: i) == reverseList.end() && pureLiteralCount[i] < 0) {
        list.push_back( x: 0 - i);
        assignMap.insert( v: { u1: 0 - i, u2: false});
    }
}

```

Finally, for all the clauses that has pure literal, remove it from cnf.

```
std::list<int>::iterator iter;
for (iter = list.begin(); iter != list.end(); ++iter) {
    for (int i = 0; i < modifiedCnf.size(); i++) {
        for (int j = 0; j < modifiedCnf[i].size(); j++) {
            if (*iter == modifiedCnf[i][j]) {
                modifiedCnf.erase( position: modifiedCnf.begin() + i);
                i--;
                break;
            }
        }
    }
}
```