

# EE360T/382C-16 Software Testing

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## Lecture 2

# Today and next time

## Discrete math basics

- Review some material from MIT's 6.042 text
  - <https://courses.csail.mit.edu/6.042/spring17/mcs.pdf>
- Focus on Section I Proofs in the text
  - **Propositions, predicates**
  - **Logical formulas**
  - Mathematical data types
  - Induction
  - State machines

# Propositions

A proposition is a statement that is either true or false

[Prop. 1.1.1]  $2 + 3 = 5$

[Prop. 1.1.2]  $1 + 1 = 3$

Can you think of a sentence that is not a proposition?

# Propositions

[Claim 1.1.3] For every non-negative integer  $n$  the value of  $n^2 + n + 41$  is prime

Let  $p(n) ::= n^2 + n + 41$

Example values:  $p(0) = 41$ ,  $p(1) = 43$ ;  $p(2) = 47$ ;  
 $p(3) = 53$ , ...,  $p(20) = 461$  are all prime

But  $p(40) = 40 \cdot 40 + 40 + 41 = 41 \cdot 41$  is not a prime

# Propositions

[Euler's conjecture, 1769]  $a^4 + b^4 + c^4 = d^4$  has no solution when  $a, b, c, d$  are positive integers

Proved false 218 years later [Elkies]

- $A = 95800, b = 217519, c = 414560, d = 422481$

# Propositions

[Fermat's last theorem, 1630] There are no positive integers  $x$ ,  $y$ , and  $z$  such that  $x^n + y^n = z^n$  for some integer  $n > 2$

Fermat claimed to have a proof but not enough space to fit it in a margin

Over the years, shown to hold for all  $n \leq 4,000,000$

In 1994, Andrew Wiles gave a proof after working on it for 7 years

# Propositions

[Goldbach conjecture, 1742] Every even integer greater than 2 is a sum of 2 primes

Known to hold for all numbers up to  $10^{18}$

But we do not know if it true or false

# Predicates

A predicate is a proposition whose truth depends on the value of one or more variables

$P(n) ::= \text{"n is a perfect square"}$

- Truth depends on the value of  $n$

$P(4)$  is true but  $P(5)$  is false

A predicate is analogous to a function

- A predicate is a boolean function



# Logical formulas

Natural language sentences can have ambiguity

- “You may have cake, or you may have ice cream.”
- “If you can solve any problem we come up with, then you get an A for the course.”

# Boolean operators

P	NOT(P)
True	False
False	True

P	Q	P AND Q	P OR Q
T	T	T	T
T	F	F	T
F	T	F	T
F	F	F	F

# Boolean operators

P	Q	P if and only iff Q
T	T	T
T	F	F
F	T	F
F	F	T

P and Q have the same truth value

Ex: for any real number  $x$ ,  $x^2 - 4 \geq 0$  IFF  $|x| \geq 2$

# Boolean operators

P	Q	P IMPLIES Q
T	T	T
T	F	F
F	T	T
F	F	T

Is the following proposition true or false:

- If Goldbach's conjecture is true, then  $x^2 \geq 0$  for every real number  $x$

# Propositional logic in code

Example Java statement:

```
if (x > 0 || (x <= 0 && y > 100)) ...
```

Let A be the expression “ $x > 0$ ” and B be “ $y > 100$ ”

The condition is “A OR (NOT(A) AND B)”

# Propositional logic in code

A	B	A OR (NOT(A) AND B)			A OR B
T	T	<b>T</b>	F	F	<b>T</b>
T	F	<b>T</b>	F	F	<b>T</b>
F	T	<b>T</b>	T	T	<b>T</b>
F	F	<b>F</b>	T	F	<b>F</b>

“A OR (NOT(A) AND B)” is equivalent to “A OR B”

- Simpler, easier to comprehend
- Can be used to simplify the original program and possibly make it run faster

# Notation

English	Symbolic Notation
NOT( $P$ )	$\neg P$ (alternatively, $\overline{P}$ )
$P$ AND $Q$	$P \wedge Q$
$P$ OR $Q$	$P \vee Q$
$P$ IMPLIES $Q$	$P \longrightarrow Q$
if $P$ then $Q$	$P \longrightarrow Q$
$P$ IFF $Q$	$P \longleftrightarrow Q$
$P$ XOR $Q$	$P \oplus Q$

source: page 54 of <https://courses.csail.mit.edu/6.042/spring17/mcs.pdf>

# Equivalence

Do the following two sentences say the same thing?

- If I am hungry, then I am grumpy (S1)
- If I am not grumpy, then I am not hungry (S2)

Let P be “I am hungry” and Q be “I am grumpy”

- S1 is  $P \text{ IMPLIES } Q$
- S2 is  $\text{NOT}(Q) \text{ IMPLIES } \text{NOT}(P)$



# Equivalence

P	Q	P IMPLIES Q	NOT(Q) IMPLIES NOT(P)		
T	T	<b>T</b>	F	<b>T</b>	F
T	F	<b>F</b>	T	<b>F</b>	F
F	T	<b>T</b>	F	<b>T</b>	T
F	F	<b>T</b>	T	<b>T</b>	T

NOT(Q) IMPLIES NOT(P) is called the contrapositive of P IMPLIES Q

- An implication and its contrapositive are always equivalent

# Validity

A formula is valid if it is always true regardless of the values of its variables

Ex:  $P \text{ or } \text{not}(P)$

Ex:  $(P \text{ implies } Q \text{ and } Q \text{ implies } R) \text{ implies } (P \text{ implies } R)$

# Satisfiability

A formula is satisfiable if there is some assignment of values to its variables such that the formula is true

Ex:  $P \text{ and } Q$  is satisfiable because for  $P = T$  and  $Q = T$ ,  
 $P \text{ and } Q = T$

$P$  is satisfiable if and only if its negation  $\text{not}(P)$  is not valid

# The SAT problem

Is the given formula satisfiable?

- $(p \vee q \vee r) \wedge (\neg p \vee \neg q) \wedge (\neg p \vee \neg r) \wedge (\neg r \vee \neg q)$

Can construct a truth table to check satisfiability

- Size of table grows exponentially

Unknown whether there a polynomial-time solution

- “P versus NP” problem

# Quantifiers

Universal – for all:

- Ex:  $\forall x \in \mathbb{R}. x^2 \geq 0$

Existential – there exists:

- Ex:  $\exists x \in \mathbb{R}. 5x^2 - 7 = 0$

$\forall x \in \mathbb{R}. 5x^2 - 7 = 0$  is false

# Mixing quantifiers

Recall Goldbach's conjecture: Every even integer greater than 2 is a sum of 2 primes

Let *Evens* be the set of all evens  $> 2$  and *Primes* be the set of all primes

$$\forall n \in \textit{Evens} \exists p \in \textit{Primes} \exists q \in \textit{Primes} . n = p + q$$

# Order of quantifiers

Swapping the order of different types of quantifiers usually changes the meaning of the formula

Ex: the following is false:

$$\exists p \in \text{Primes} \exists q \in \text{Primes} \forall n \in \text{Evens} . n = p + q$$

# Negating quantifiers

Ex: the following sentences mean the same thing

- Not everyone likes ice cream
- There is someone who does not like ice cream

In general

- $\neg \forall x. P(x)$  is equivalent to  $\exists x. \neg P(x)$
- $\neg \exists x. P(x)$  is equivalent to  $\forall x. \neg P(x)$



?/!