

EE360T/382V Software Testing

khurshid@ece.utexas.edu

February 19, 2020

Overview

Today – Logic coverage

Last class – Graph coverage for designs and specs

Next class – Review for Exam 1

Problem Set 1 reminder – Due: 2/23 11:59pm

Exam 1 reminder – February 26, i.e., only **1 week** away

- In-class, closed book, no cheat-sheet
- Practice questions posted on Canvas

EE360T/382V Software Testing

khurshid@ece.utexas.edu

Chapter 3*: Logic Coverage

*Introduction to Software Testing by Ammann and Offutt

Criteria based on structures

The textbook focuses on four kinds of structures to define criteria:

- Graphs
 - E.g., control-flow graphs (CFGs)
- Logical expressions
 - E.g., if-conditions
- Input domain characterization
 - E.g., sorted array
- Syntactic structures
 - E.g., mutation

3.1 Overview: logical expressions

Predicate – expression that evaluates to a boolean value, e.g., “ $((a > b) \parallel C) \&\& p(x)$ ”

- May contain variables (boolean or non-boolean) and methods
- Internal structure defined by logical operators, e.g., “!”, “&&”, “||”

Clause – predicate that contains no logical operator, e.g., “ $a > b$ ”, “ C ”, and “ $p(x)$ ”

Logical expressions come from various sources, e.g., program source-code

3.2 Logic expression coverage criteria

P – set of predicates

C – set of clauses in predicates in P

C3.12 Predicate coverage (PC) – for each $p \in P$, TR contains two requirements:

- p evaluates to *true*; and
- p evaluated to *false*

C3.13 Clause coverage (CC) -- for each $c \in C$, TR contains two requirements:

- c evaluates to *true*; and
- c evaluated to *false*

PC and CC relation

Consider the predicate “ $((a > b) \parallel C) \&\& p(x)$ ”

Predicate coverage satisfied by two tests:

1. $a = 5, b = 4, C = \text{true}, p(x) = \text{true}$
 2. $a = 5, b = 4, C = \text{true}, p(x) = \text{false}$
- The two tests do not satisfy clause coverage

Clause coverage satisfied by two tests:

1. $a = 5, b = 5, C = \text{false}, p(x) = \text{true}$
 2. $a = 5, b = 4, C = \text{true}, p(x) = \text{false}$
- The two tests do not satisfy predicate coverage!

PC does not subsume CC and CC does not subsume PC

Combinatorial coverage

C_p – set of clauses in predicate p

C3.14 Combinatorial coverage (CoC) – for each $p \in P$, TR has test requirements for the clauses in C_p to evaluate to each possible combination of truth values

- Also called multiple condition coverage

Example: “!a || b”

| a | b | $!a b$ |
|-----|-----|-----------|
| F | F | T |
| F | T | T |
| T | F | F |
| T | T | T |

Predicate with n clauses has 2^n possible assignments
Often impractical for predicates with > a few clauses

Determination

Motivation – need criteria that capture the effect of each clause using a reasonable number of tests

- Would like to exercise conditions where flipping a clause flips the predicate

Major clause – clause c_i that is our focus

Minor clause – clause c_j where $j \neq i$

D3.42 Determination – major clause c_i *determines* predicate p if the minor clauses have values so that changing the truth value of c_i changes the value of p

Active clause coverage

D3.43 Active clause coverage (ACC)* – TR has two requirements for each active clause $c_i \in C_p$ for each $p \in P$: c_i evaluates to *true*; and c_i evaluates to *false*

Example – consider predicate $p = a \mid \mid b$

- a determines p iff b is false; likewise for b
- 4 requirements:
 - $c_i = a: \{ \langle a = T, b = f \rangle, \langle a = F, b = f \rangle \}$
 - $c_i = b: \{ \langle a = f, b = T \rangle, \langle a = f, b = F \rangle \}$
 - 2 of these are identical, so 3 in total

Key question in ACC – do minor clauses have constant values when major clause c_i is and when c_i is false?

*Almost identical to MCDC in literature

CACC and RACC

C3.16 Correlated active clause coverage (CACC) – *TR* has two requirements for each major clause $c_i \in C_p$ for each $p \in P$: c_i evaluates to *true*; and c_i evaluates to *false*. **The values chosen for minor clauses c_j ($j \neq i$) must cause p to be true for one value of c_i , and false for the other**

C3.17 Restricted active clause coverage (RACC) – *TR* has two requirements for each major clause $c_i \in C_p$ for each $p \in P$: c_i evaluates to *true*; and c_i evaluates to *false*. **The values chosen for minor clauses c_j ($j \neq i$) must be the same when c_i is *true* as when c_i is *false***

CACC and RACC Example

a determines the predicate “ $a \ \&\& \ (b \ || \ c)$ ” when “ $(b \ || \ c)$ ” is true

- $\langle b = T, c = T \rangle$
- $\langle b = T, c = F \rangle$
- $\langle b = F, c = T \rangle$

| | a | b | c | $a \ \&\& \ (b \ \ c)$ |
|---|-----|-----|-----|---------------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

CACC and RACC Example

a determines the predicate “ $a \ \&\& \ (b \ || \ c)$ ” when “ $(b \ || \ c)$ ” is true

- $\langle b = T, c = T \rangle$
- $\langle b = T, c = F \rangle$
- $\langle b = F, c = T \rangle$

CACC – pick one of $\{ 1, 2, 3 \}$
and one of $\{ 5, 6, 7 \}$

- 9 choices

| | a | b | c | $a \ \&\& \ (b \ \ c)$ |
|---|-----|-----|-----|---------------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

CACC and RACC Example

a determines the predicate “ $a \ \&\& \ (b \ || \ c)$ ” when “ $(b \ || \ c)$ ” is true

- $\langle b = T, c = T \rangle$
- $\langle b = T, c = F \rangle$
- $\langle b = F, c = T \rangle$

CACC – pick one of $\{ 1, 2, 3 \}$ and one of $\{ 5, 6, 7 \}$

- 9 choices

RACC – pick one of $\langle 1, 5 \rangle$, $\langle 2, 6 \rangle$, and $\langle 3, 7 \rangle$

- 3 choices

| | a | b | c | $a \ \&\& \ (b \ \ c)$ |
|---|-----|-----|-----|---------------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

Inactive clause coverage

Basis – check that changing a clause that should not affect the predicate does not, in fact, affect it

D3.44 Inactive clause coverage (ICC) – for each $p \in P$ and each major clause $c_i \in C_p$, choose minor clauses c_j ($j \neq i$) so that c_i does **not** determine p . *TR* has four test requirements for c_i :

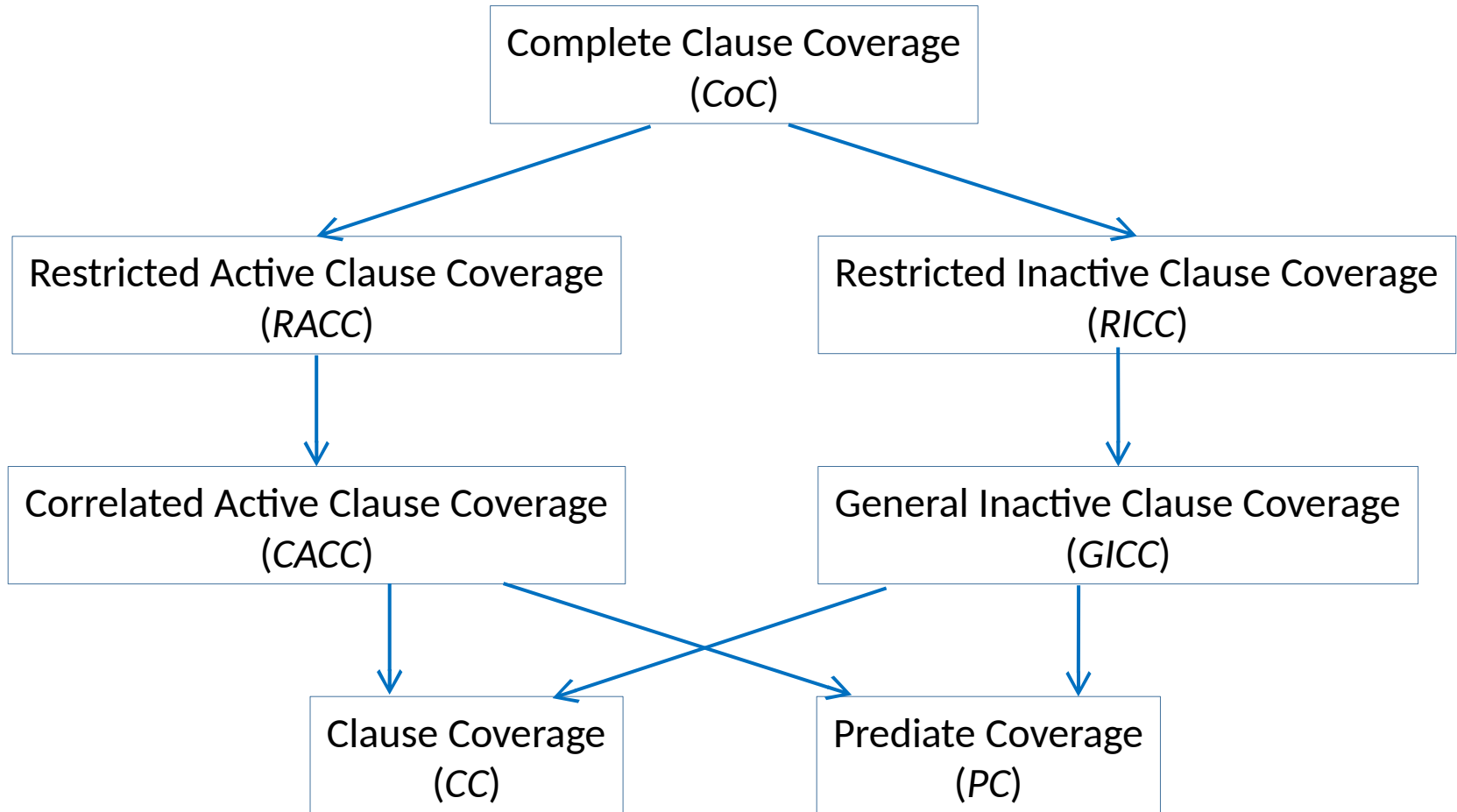
1. c_i evaluates to *true* with p *true*;
2. c_i evaluates to *false* with p *true*;
3. c_i evaluates to *true* with p *false*; and
4. c_i evaluates to *false* with p *false*;

GICC and *RICC*

C 3.18 General inactive clause coverage (*GICC*) – ICC such that the values chosen for the minor clauses for a given major clause may vary among the four cases

C 3.19 Restricted inactive clause coverage (*RICC*) – ICC such that the values chosen for the minor clauses for a given major clause must be the same in cases (1) and (2), and also be the same in cases (3) and (4)

Subsumption



Making a clause determine a predicate

Let p be a predicate and c be a clause in p

Let $p_{c=true}$ be p with c set to *true*

Let $p_{c=false}$ be p with c set to *false*

Then, *solutions* to formula $p_{c=true} \oplus p_{c=false}$ give values for clauses $\neq c$ such that c determines p

- Each solution is assignment of values to clauses in $p_{c=true} \oplus p_{c=false}$ such that the formula is *true*

Example of making a clause active

$$p = a \mid \mid b$$

$$\begin{aligned}\text{To make } a \text{ active, solve } p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \mid \mid b) \oplus (\text{false} \mid \mid b) \\ &= \text{true} \oplus b \\ &= !b\end{aligned}$$

There is one solution to p_a , which is $b = \text{false}$

Thus, setting $b = \text{false}$ makes clause a active

By symmetry, $p_b = !a$

Another example

$$p = a \ \&\& \ b$$

$$\begin{aligned} \text{To make } a \text{ active, solve } p_a &= p_{a=true} \oplus p_{a=false} \\ &= (true \ \&\& \ b) \oplus (false \ \&\& \ b) \\ &= b \oplus false \\ &= b \end{aligned}$$

There is one solution to p_a , which is $b = true$

Thus, setting $b = true$ makes clause a active

By symmetry, $p_b = a$

An example with no constraint

$$p = a \Leftrightarrow b$$

To make a active, solve $p_a = p_{a=true} \oplus p_{a=false}$

$$\begin{aligned} &= (true \Leftrightarrow b) \oplus (false \Leftrightarrow b) \\ &= b \oplus !b \\ &= true \end{aligned}$$

Any value of b is a solution

Thus, setting b to any value makes clause a active

By symmetry, $p_b = true$

A degenerate case

$$p = a \ \&\& \ b \ || \ a \ \&\& \ !b$$

To make b active, solve $p_b = p_{b=true} \oplus p_{b=false}$

$$= (a \ \&\& \ true \ || \ a \ \&\& \ !true) \oplus (a \ \&\& \ false \ || \ a \ \&\& \ !false)$$

$$= (a \ || \ false) \oplus (false \ || \ a)$$

$$= a \oplus a$$

$$= false$$

There is no solution to p_b

Thus, there is no value for a that makes b active

An example with 3 clauses

$$p = a \ \&\& \ (b \ || \ c)$$

$$\begin{aligned}\text{To make } a \text{ active, solve } p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \ \&\& \ (b \ || \ c)) \oplus (\text{false} \ \&\& \ (b \ || \ c)) \\ &= (b \ || \ c) \oplus \text{false} \\ &= b \ || \ c\end{aligned}$$

There are three distinct solutions to p_a :

- $\langle b = T, c = T \rangle$, $\langle b = T, c = F \rangle$, and $\langle b = F, c = T \rangle$

Any of these three pairs makes clause a active

Exercise – make b active

?/!