

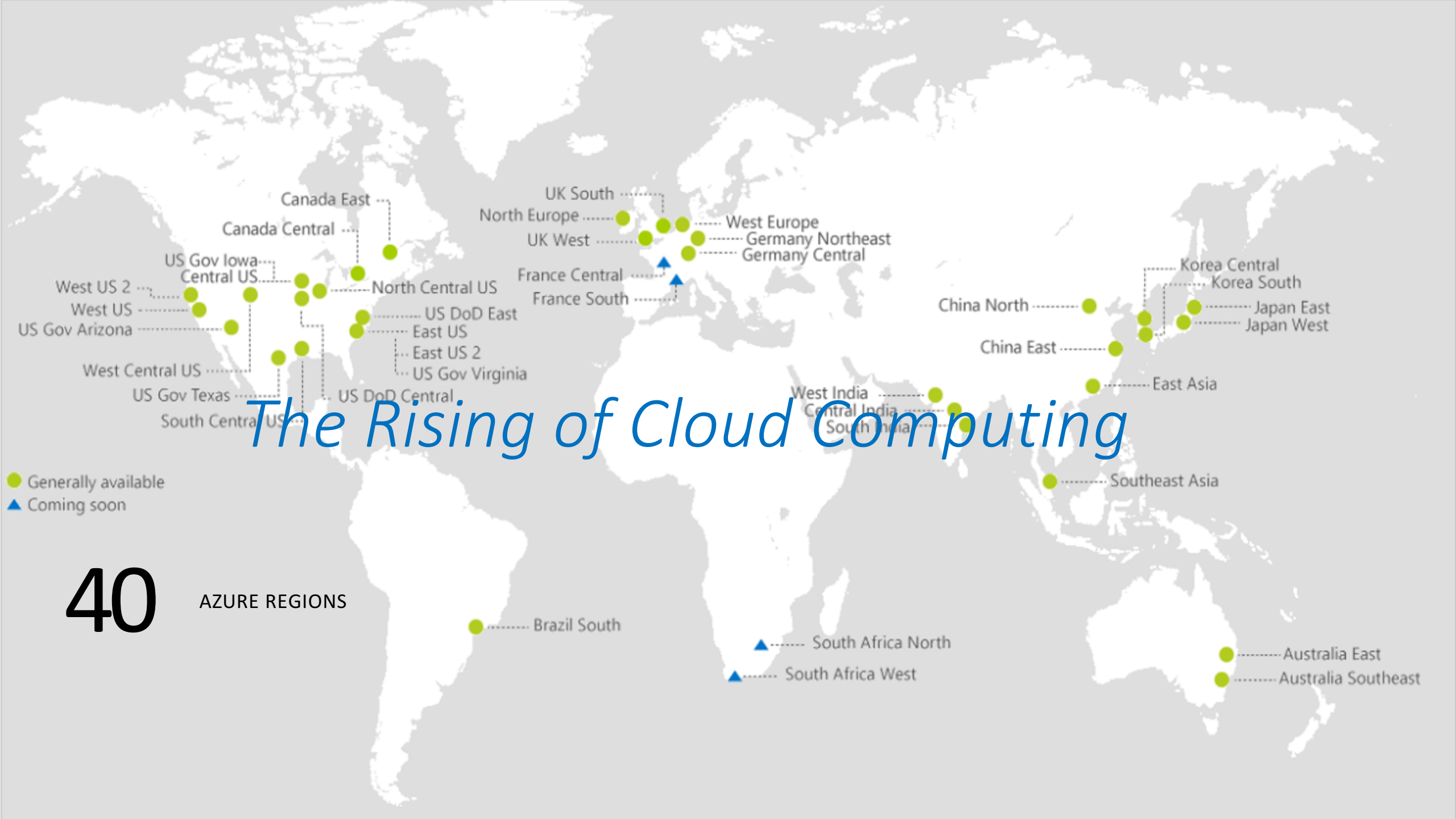
RDMA in Data Centers: Looking Back and Looking Forward

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The Rising of Cloud Computing



40

AZURE REGIONS

An aerial photograph of a large data center complex. The facility consists of several long, white, rectangular buildings with flat roofs. In the center, there is a large area filled with rows of white server racks or storage units. The complex is surrounded by green fields and dense forests. A road runs along the bottom right, and a small pond is visible in the lower left. The text "Data Centers" is overlaid in the center of the image.

Data Centers

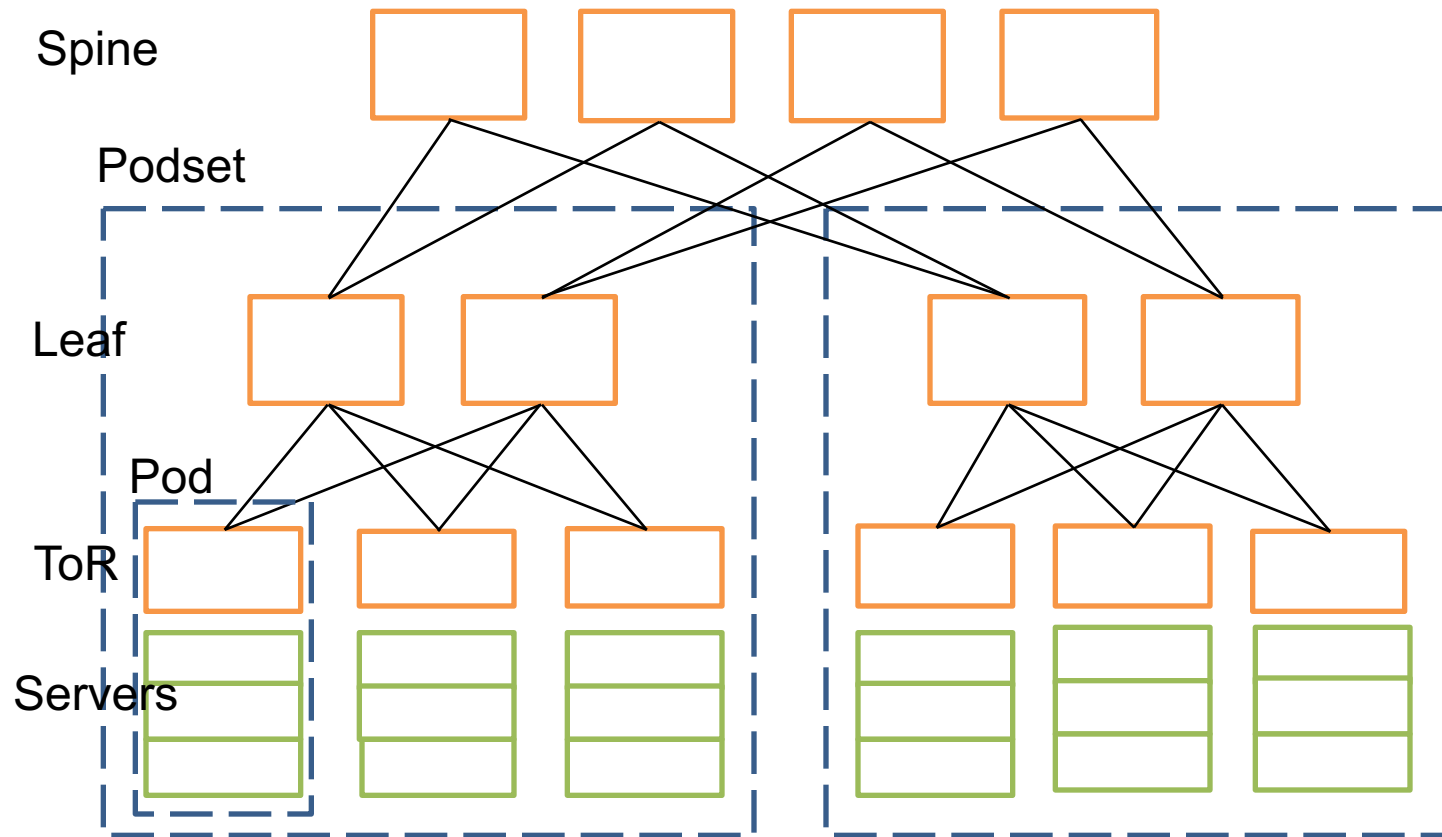


Data Centers

Data center networks (DCN)

- Cloud scale services: IaaS, PaaS, Search, BigData, Storage, Machine Learning, Deep Learning
- Services are latency sensitive or bandwidth hungry or both
- Cloud scale services need cloud scale computing and communication infrastructure

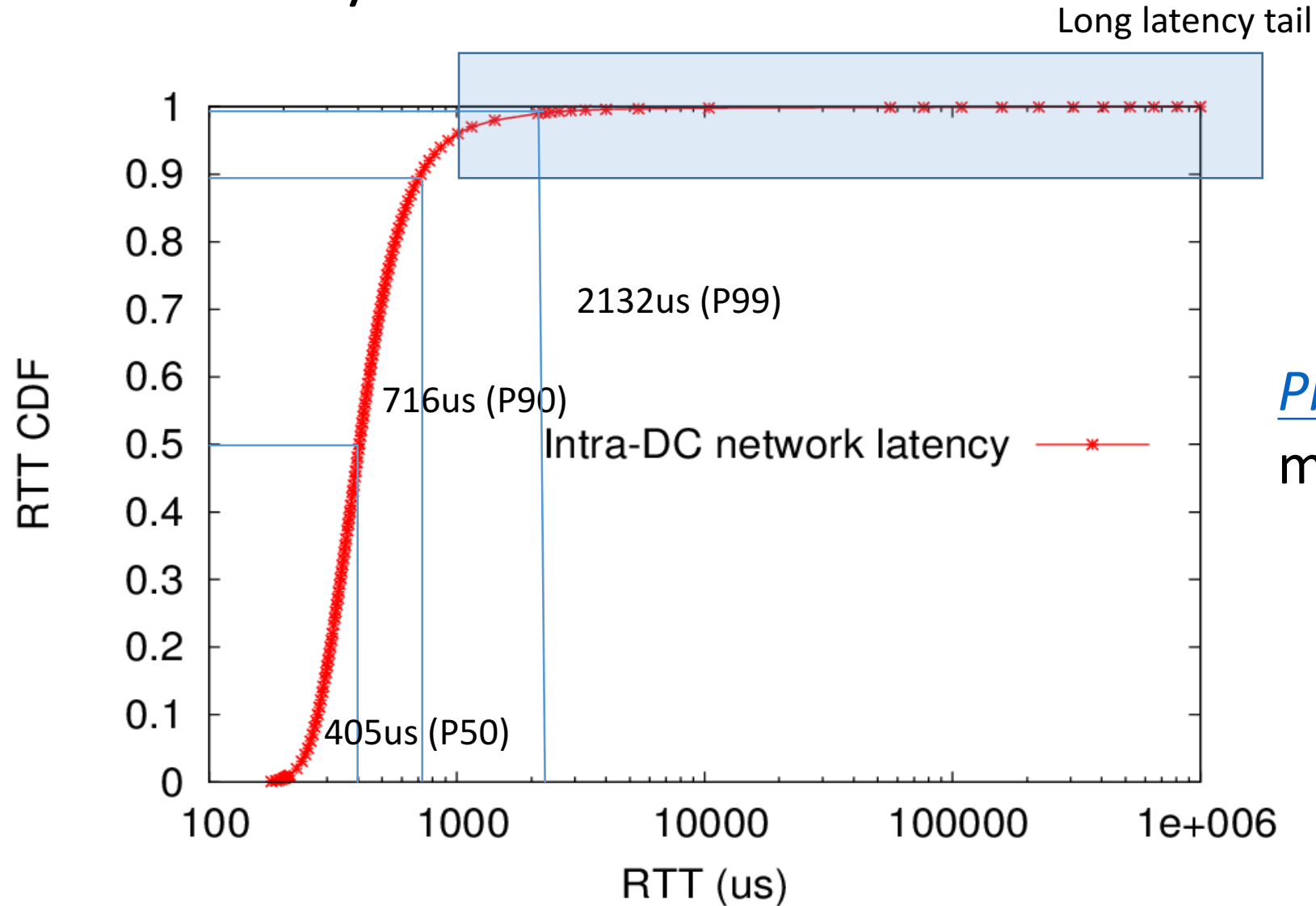
Data center networks (DCN)



- Single ownership
- Large scale
- High bisection bandwidth
- Commodity Ethernet switches
- TCP/IP protocol suite

But TCP/IP is not doing well

TCP latency



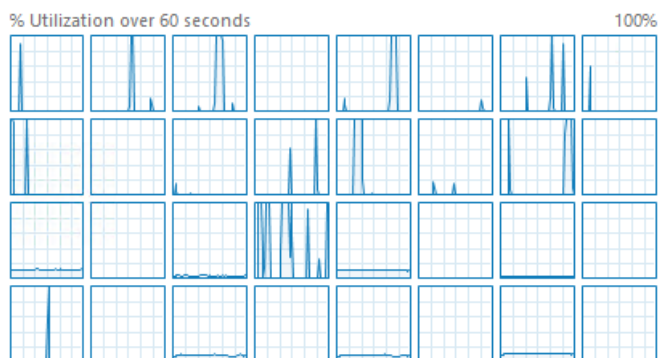
Pingmesh
measurement results

TCP processing overhead (40G)

Sender

CPU

Intel(R) Xeon(R) CPU E5-2690 0 @ 2.90GHz



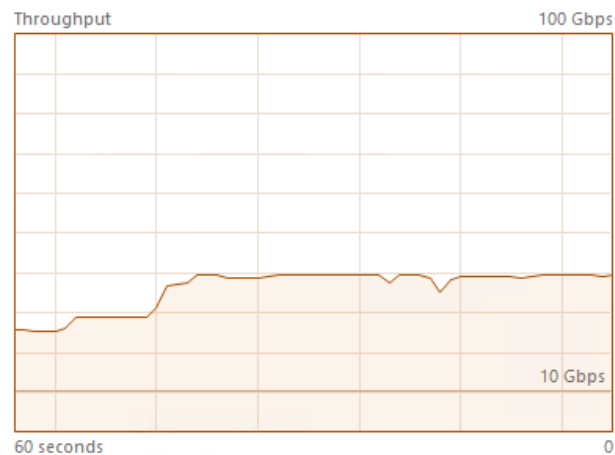
Utilization	Speed	Maximum speed:	2.90 GHz
6%	3.28 GHz	Sockets:	2
Processes	Threads	Cores:	16
36	799	Logical processors:	32
Up time		Virtualization:	Enabled
9:07:46:53		L1 cache:	1.0 MB
		L2 cache:	4.0 MB
		L3 cache:	40.0 MB

8 tcp connections

40G NIC

Ethernet

Mellanox ConnectX-3 Pro Ethernet A...

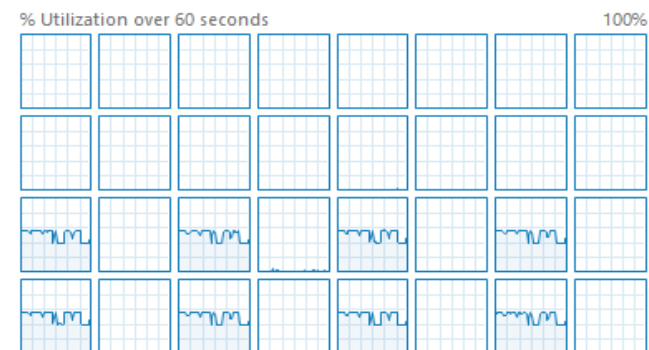


Send	Adapter name:	SLOT 4 4
70.0 Mbps	Connection type:	Ethernet
	IPv4 address:	169.254.232.5
Receive	IPv6 address:	fe80::e10c:b1bb:664c:e805%30
39.4 Gbps		

Receiver

CPU

Intel(R) Xeon(R) CPU E5-2690 0 @ 2.90GHz



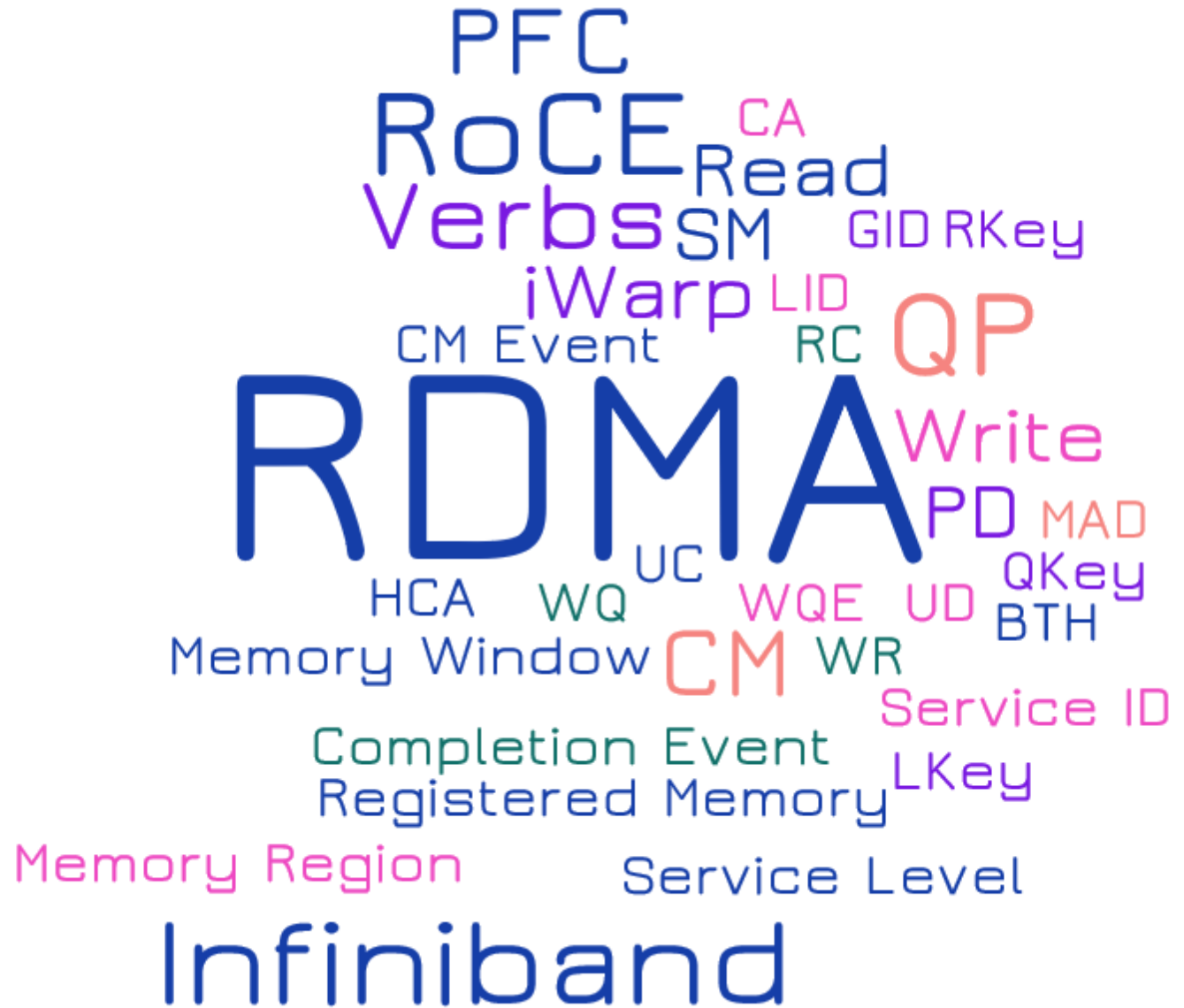
Utilization	Speed	Maximum speed:	2.90 GHz
12%	3.28 GHz	Sockets:	2
Processes	Threads	Cores:	16
36	793	Logical processors:	32
Up time		Virtualization:	Enabled
9:07:50:19		L1 cache:	1.0 MB
		L2 cache:	4.0 MB
		L3 cache:	40.0 MB

An RDMA renaissance story

Virtual Interface
Architecture Spec
1.0 1997

*Infiniband Architecture
Spec*

1.0 2000
1.1 2002
1.2 2004
1.3 2015
RoCE 2010
RoCEv2 2014



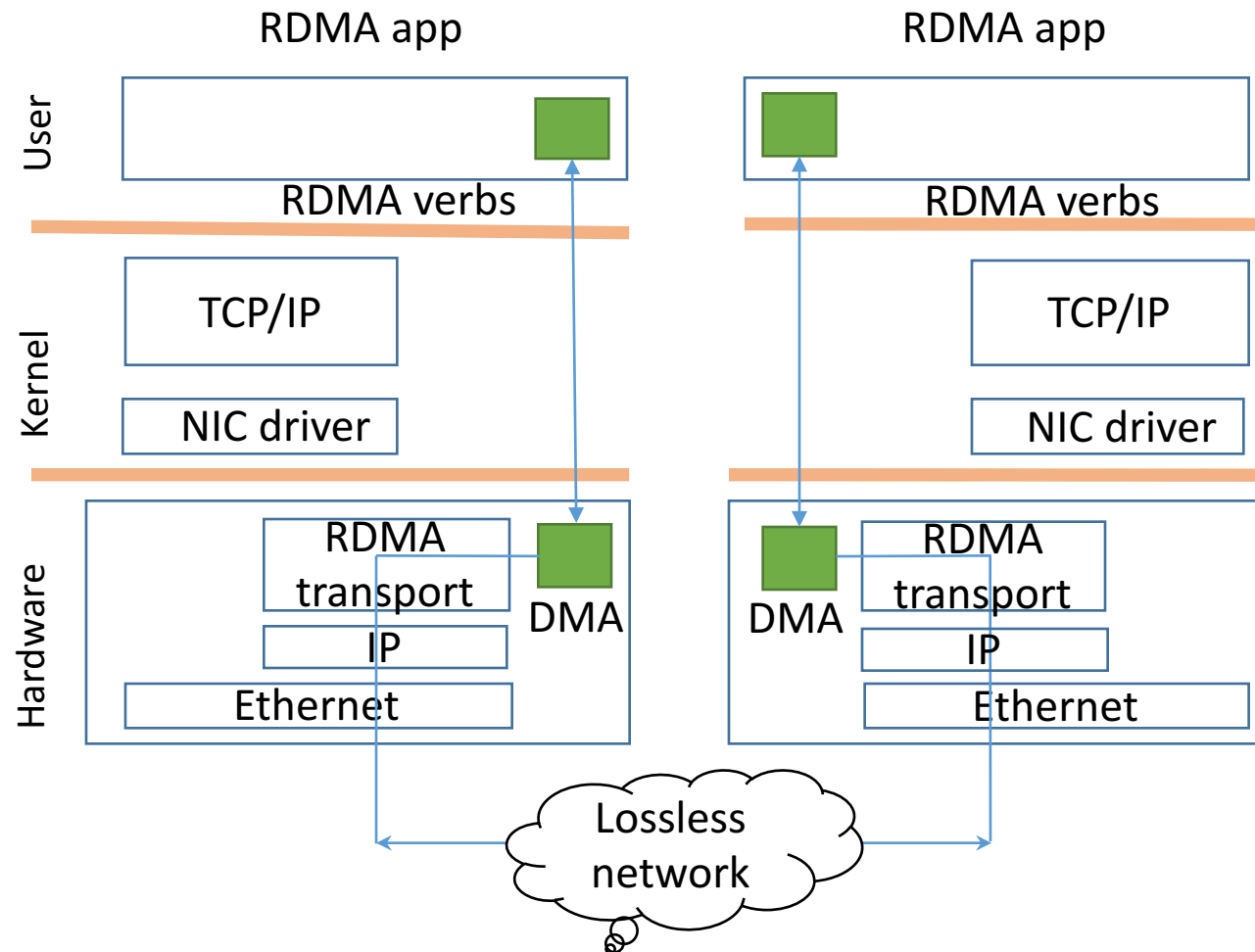
A word cloud of Infiniband terminology. The words are arranged in a roughly circular shape, with 'RDMA' being the largest and most central. Other prominent words include 'PFC', 'RoCE', 'Verbs', 'iWarp', 'CM', 'Event', 'HCA', 'WQ', 'Memory Window', 'Completion Event', 'Registered Memory', 'Memory Region', 'Service Level', 'Infiniband', 'QP', 'Write', 'PD', 'MAD', 'QKey', 'UD', 'BTH', 'Service ID', 'LKey', 'WR', 'WQE', 'UC', 'RC', 'LID', 'GID', 'RKey', 'CA', 'Read', 'SM', and 'Read'. The words are in various colors (blue, green, orange, pink, purple) and sizes, creating a dense, circular arrangement.

PFC
RoCE
Verbs
iWarp
CM Event
HCA WQ
Memory Window
Completion Event
Registered Memory
Memory Region
Service Level
Infiniband
QP
Write
PD MAD
QKey
UD BTH
Service ID
LKey
WR
WQE
UC
RC
LID
GID RKey
CA
Read
SM
Read

RDMA

- Remote Direct Memory Access (RDMA): Method of accessing memory on a remote system ***without*** interrupting the processing of the CPU(s) on that system
- RDMA offloads packet processing protocols to the NIC
- RDMA in Ethernet based data centers

RoCEv2: RDMA over Commodity Ethernet

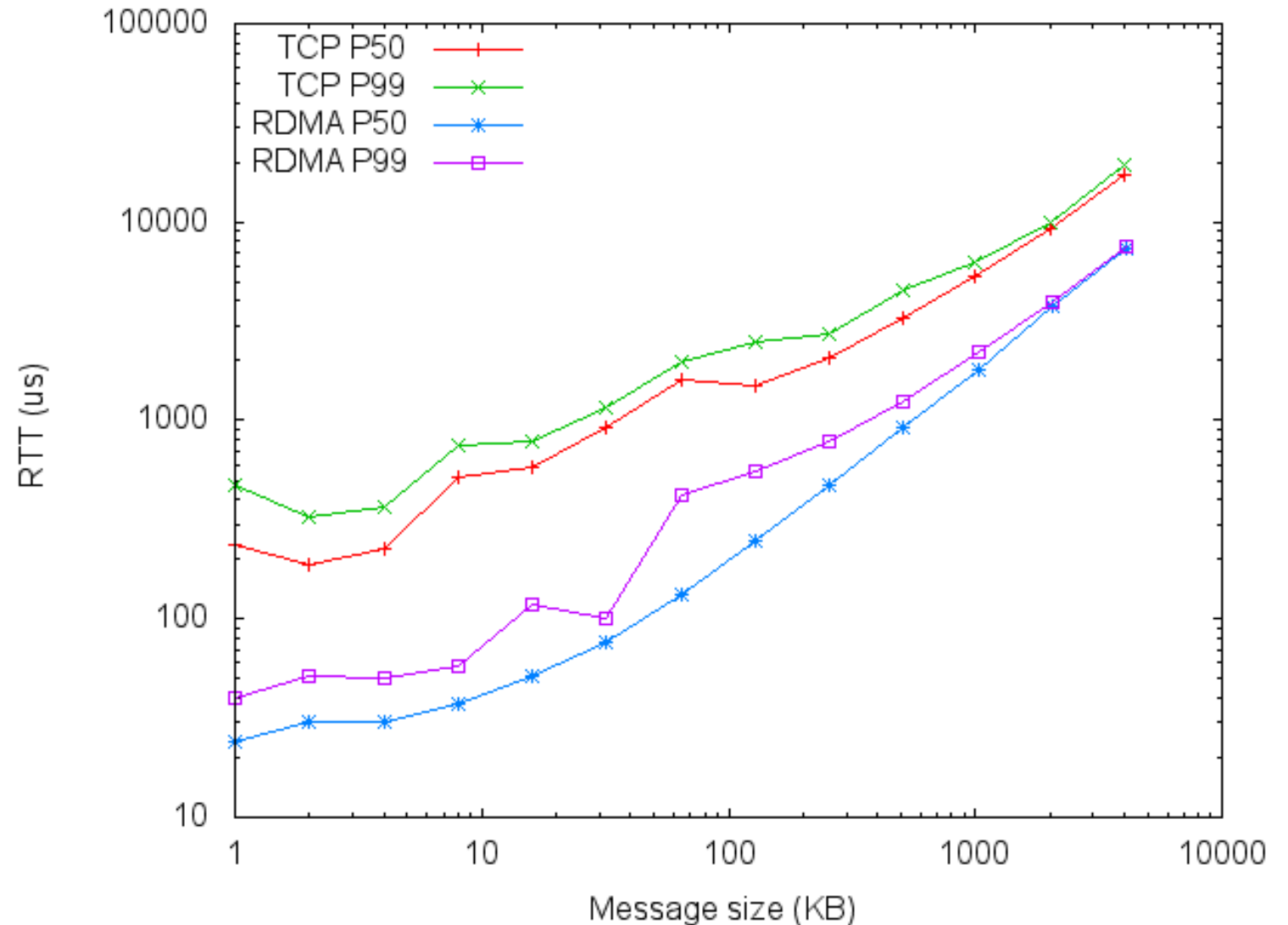


- RoCEv2 for Ethernet based data centers
- RoCEv2 encapsulates packets in UDP
- OS kernel is not in data path
- NIC for network protocol processing and message DMA

RDMA benefit: latency reduction

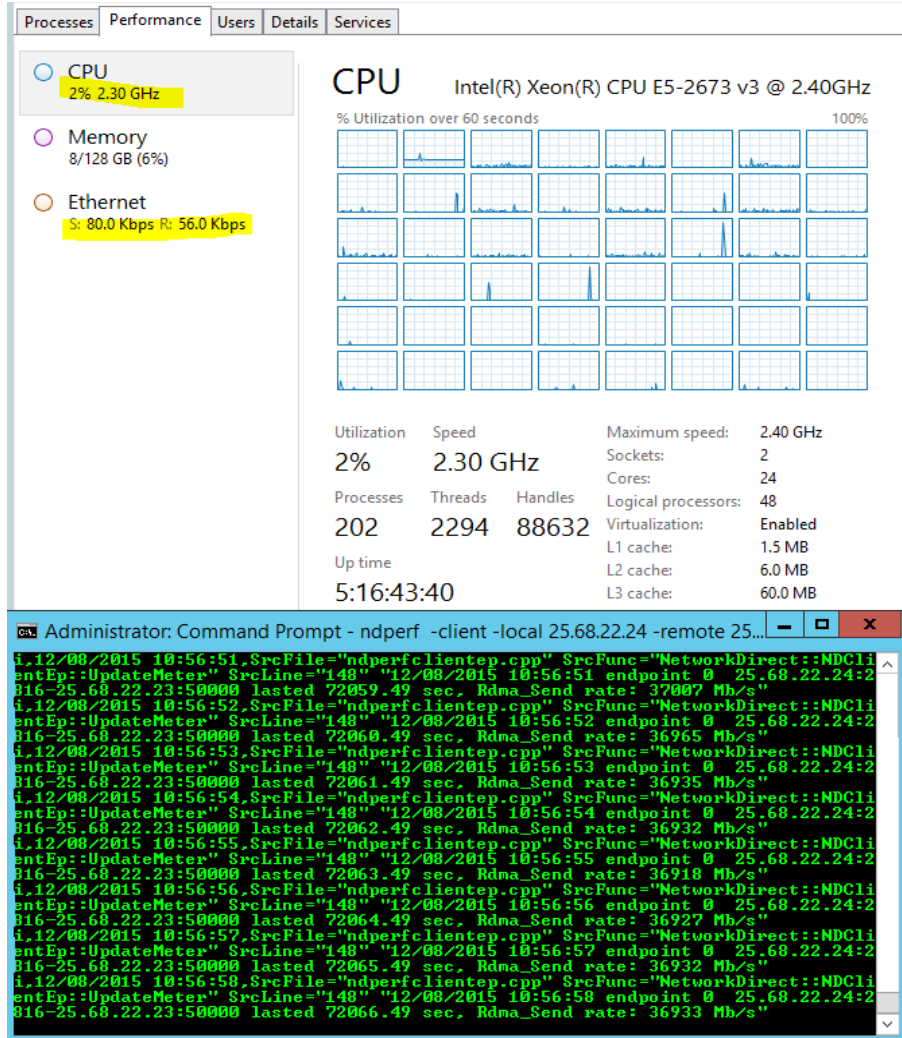
Msg size	TCP P50 (us)	TCP P99 (us)	RDMA P50 (us)	RDMA P99 (us)
1KB	236	467	24	40
16KB	580	788	51	117
128KB	1483	2491	247	551
1MB	5290	6195	1783	2214

- For small msgs (<32KB), OS processing latency matters
- For large msgs (100KB+), speed matters

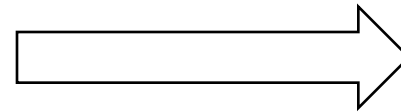


RDMA benefit: CPU overhead reduction

Sender



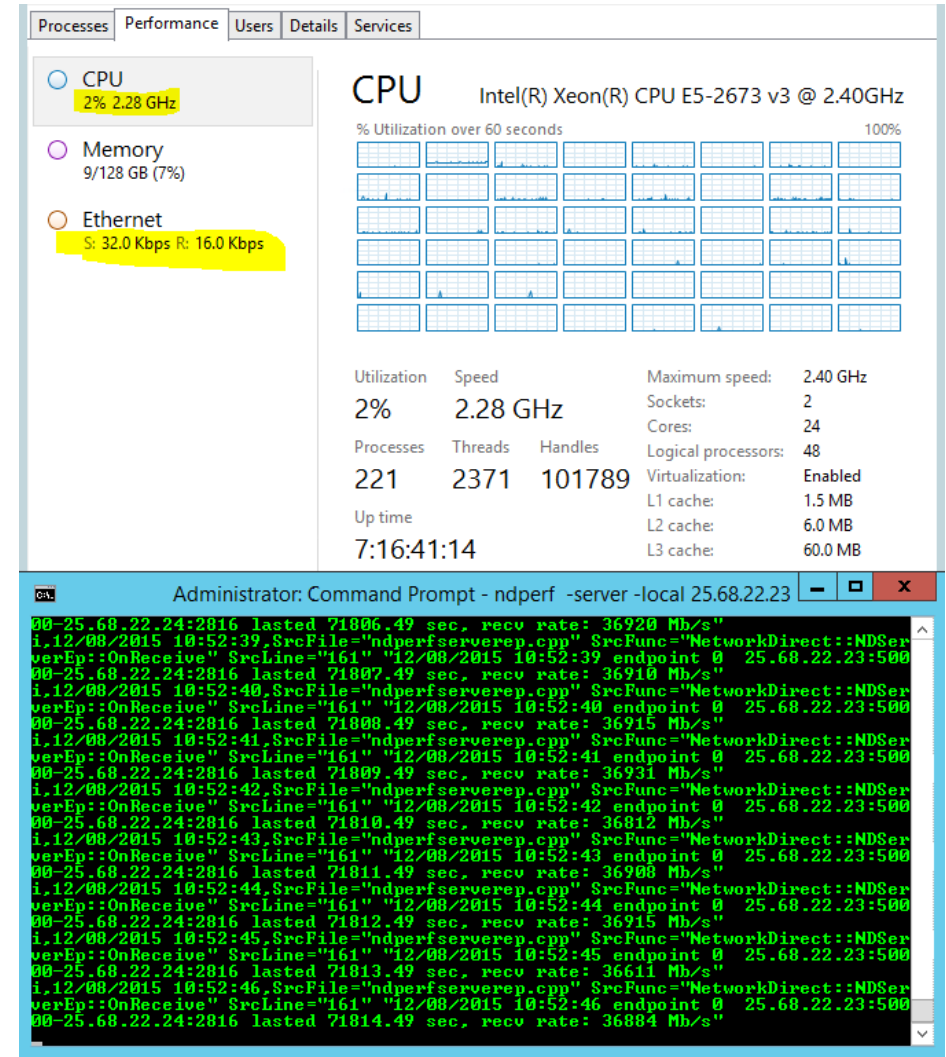
One ND connection



40G NIC

37Gb/s goodput

Receiver



RDMA benefit: CPU overhead reduction

Intel(R) Xeon(R) CPU E5-2690 v4 @ 2.60GHz, two sockets 28 cores

```
OnWrite rate = 88551 mbps
OnWrite rate = 87825 mbps
OnWrite rate = 88364 mbps
OnWrite rate = 87896 mbps
OnWrite rate = 87437 mbps
OnWrite rate = 87527 mbps
OnWrite rate = 86992 mbps
OnWrite rate = 87257 mbps
OnWrite rate = 87884 mbps
OnWrite rate = 87851 mbps
OnWrite rate = 88063 mbps
OnWrite rate = 87444 mbps
OnWrite rate = 88320 mbps
OnWrite rate = 87506 mbps
OnWrite rate = 87827 mbps
```

```
top - 18:16:47 up 21 days, 12:44, 0 users, load average: 5.38, 3.46, 3.37
Tasks: 8 total, 1 running, 7 sleeping, 0 stopped, 0 zombie
%Cpu(s): 15.0 us, 3.0 sy, 0.0 ni, 81.1 id, 0.0 wa, 0.3 hi, 0.7 si, 0.0 st
KiB Mem: 52827267+total, 64229080 used, 46404358+free, 496340 buffers
KiB Swap: 67108860 total, 2940 used, 67105920 free. 50241336 cached Mem
```

PID	USER	PR	NI	VIRT	RES	SHR	S	%CPU	%MEM	TIME+	COMMAND
1888	root	20	0	178048	4504	3528	S	1.7	0.0	0:31.01	rfwork
1	root	20	0	18200	3340	2836	S	0.0	0.0	0:00.02	bash
1785	root	20	0	24244	2572	2320	S	0.0	0.0	0:00.00	tmux
1787	root	20	0	24768	3532	2676	S	0.0	0.0	0:00.23	tmux
1788	root	20	0	18228	3464	2936	S	0.0	0.0	0:00.00	bash
1799	root	20	0	18228	3460	2932	S	0.0	0.0	0:00.00	bash
1813	root	20	0	18228	3436	2908	S	0.0	0.0	0:00.00	bash
1824	root	20	0	21956	2512	2180	R	0.0	0.0	0:00.09	top

```
76 recv rate: 30012423744
77 recv rate: 29785573024
78 recv rate: 32537844160
79 recv rate: 34104756640
80 recv rate: 32433151744
81 recv rate: 47936439424
82 recv rate: 47009696672
83 recv rate: 45762704800
84 recv rate: 36870458528
85 recv rate: 44697521312
86 recv rate: 47689360512
87 recv rate: 52305616256
88 recv rate: 36145854304
89 recv rate: 43588678304
90 recv rate: 50147339616
91 recv rate: 38479512416
92 recv rate: 49721405056
93 recv rate: 50277126944
94 recv rate: 46907554080
95 recv rate: 45712658208
96 recv rate: 50521826912
97 recv rate: 47425490240
```

```
top - 06:36:10 up 5 days, 1:59, 0 users, load average: 1.44, 0.80, 0.65
Tasks: 9 total, 1 running, 8 sleeping, 0 stopped, 0 zombie
%Cpu(s): 0.4 us, 4.5 sy, 0.0 ni, 92.6 id, 0.0 wa, 0.4 hi, 2.1 si, 0.0 st
KiB Mem: 52827267+total, 18587868 used, 50968480+free, 147664 buffers
KiB Swap: 67108860 total, 0 used, 67108860 free. 14185556 cached Mem
```

PID	USER	PR	NI	VIRT	RES	SHR	S	%CPU	%MEM	TIME+	COMMAND
1657	root	20	0	1212108	13460	2708	S	4.3	0.0	2:00.42	psrver
1	root	20	0	18200	3276	2776	S	0.0	0.0	0:00.01	bash
114	root	20	0	18192	3336	2840	S	0.0	0.0	0:00.00	bash
1620	root	20	0	24244	2572	2320	S	0.0	0.0	0:00.00	tmux
1622	root	20	0	24504	3080	2456	S	0.0	0.0	0:00.06	tmux
1623	root	20	0	18228	3468	2940	S	0.0	0.0	0:00.00	bash
1634	root	20	0	18228	3416	2888	S	0.0	0.0	0:00.00	bash
1645	root	20	0	18228	3420	2892	S	0.0	0.0	0:00.00	bash
1675	root	20	0	21956	2500	2176	R	0.0	0.0	0:00.01	top

RDMA: Single QP, 88 Gb/s, 1.7% CPU

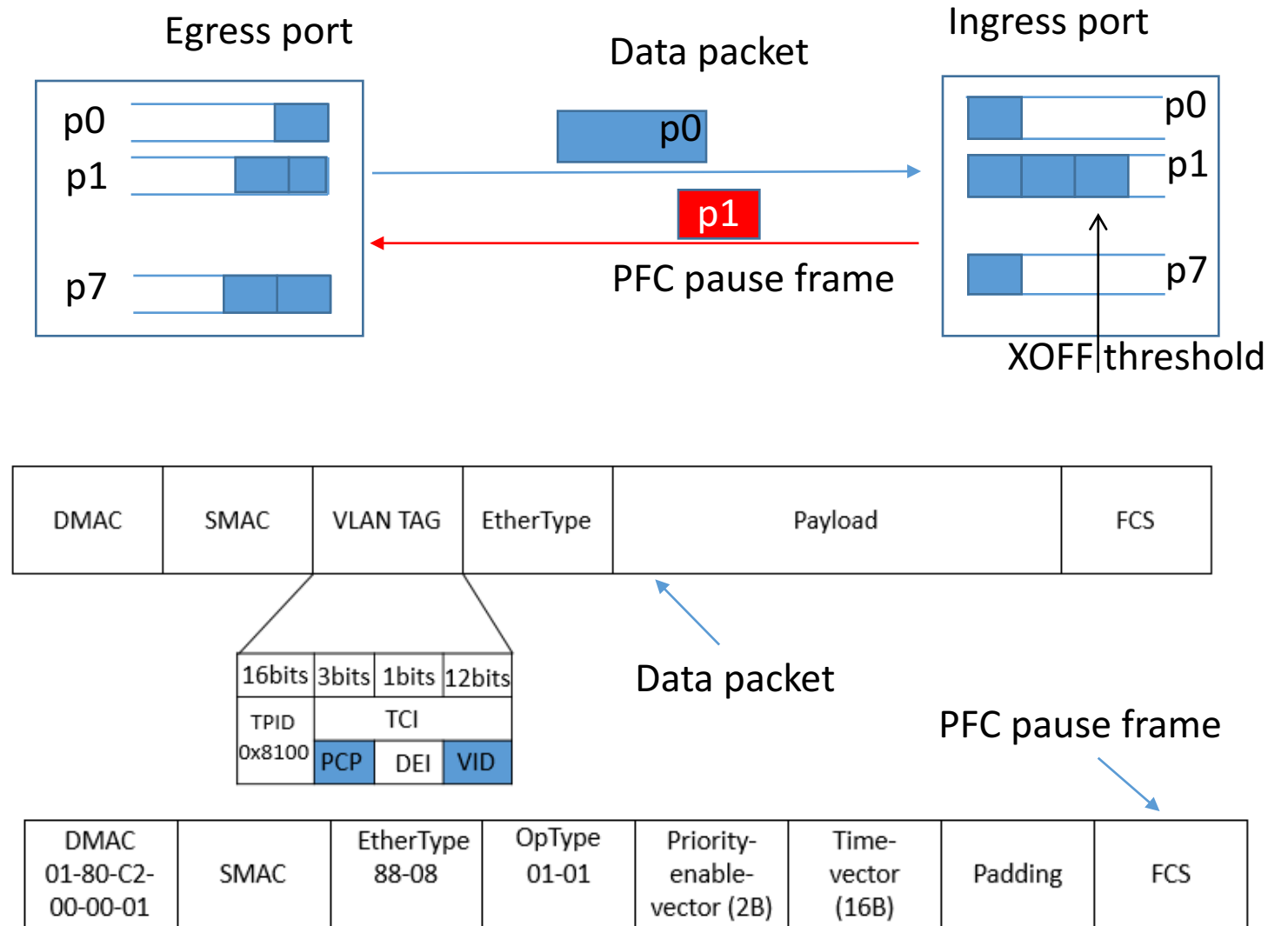
TCP: Eight connections, 30-50Gb/s,
Client: 2.6%, Server: 4.3% CPU

RoCEv2 needs a lossless Ethernet network

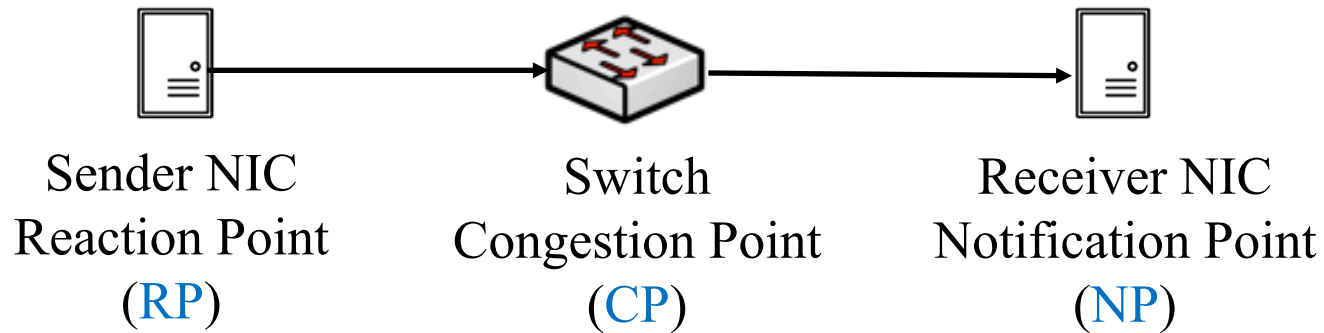
- PFC for hop-by-hop flow control
- DCQCN for connection-level congestion control

Priority-based flow control (PFC)

- Hop-by-hop flow control, with eight priorities for HOL blocking mitigation
- The priority in data packets is carried in the VLAN tag or DSCP
- PFC pause frame to inform the upstream to stop
- PFC causes HOL and collateral damage



DCQCN



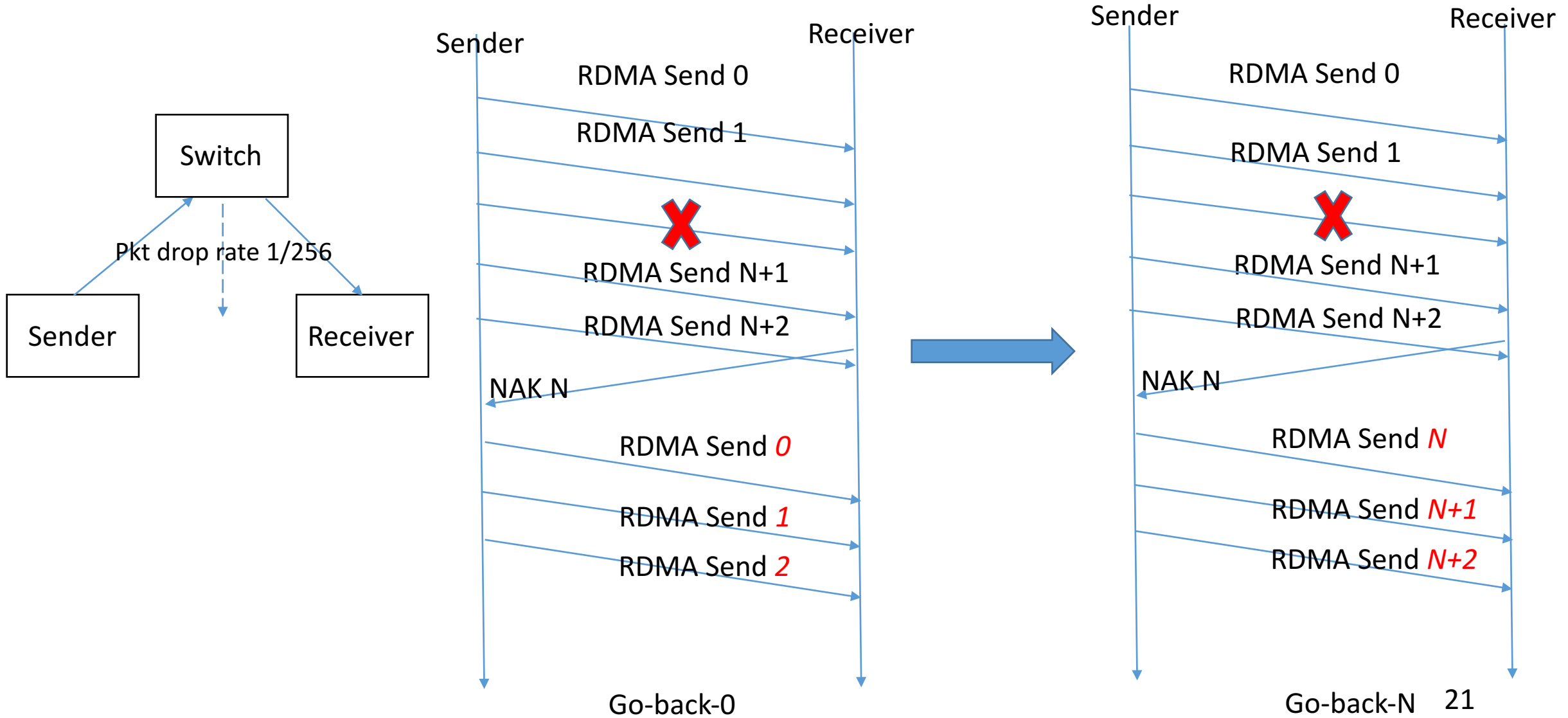
DCQCN = Keep PFC + Use ECN +
hardware rate-based congestion control

- **CP:** Switches use ECN for packet marking
- **NP:** periodically check if ECN-marked packets arrived, if so, notify the sender
- **RP:** adjust sending rate based on NP feedbacks

The lossless requirement causes safety and performance challenges

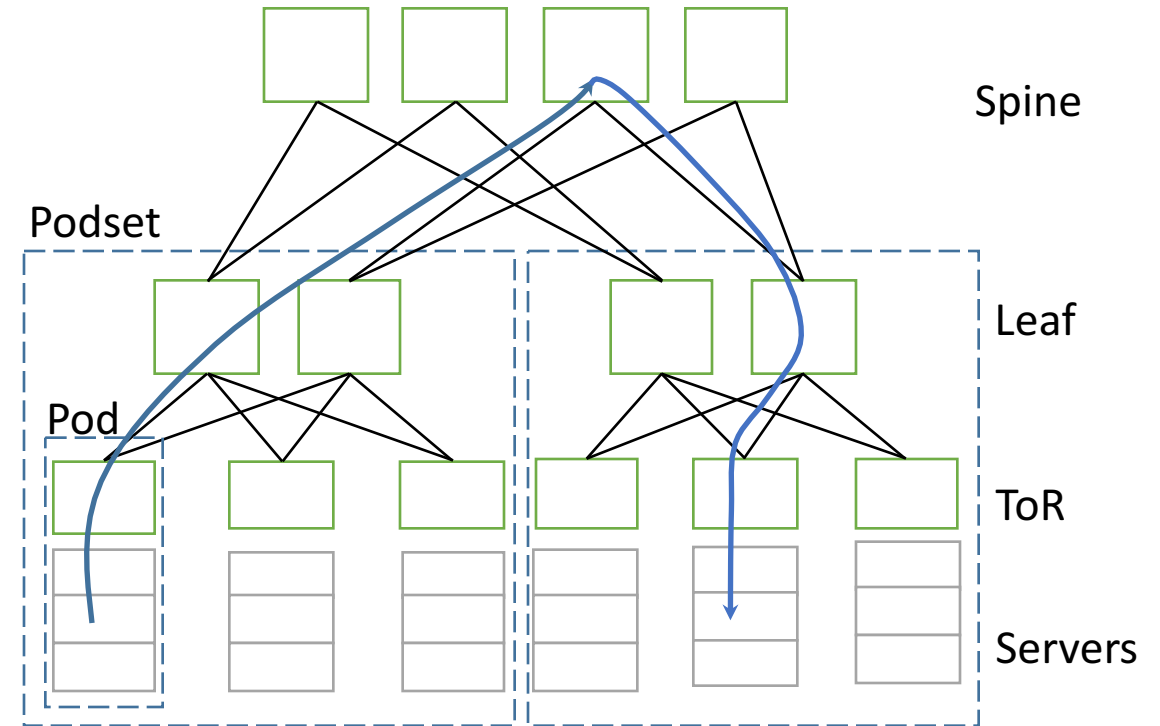
- RDMA transport livelock
- PFC pause frame storm
- PFC deadlock
- Slow-receiver symptom

RDMA transport livelock



PFC deadlock

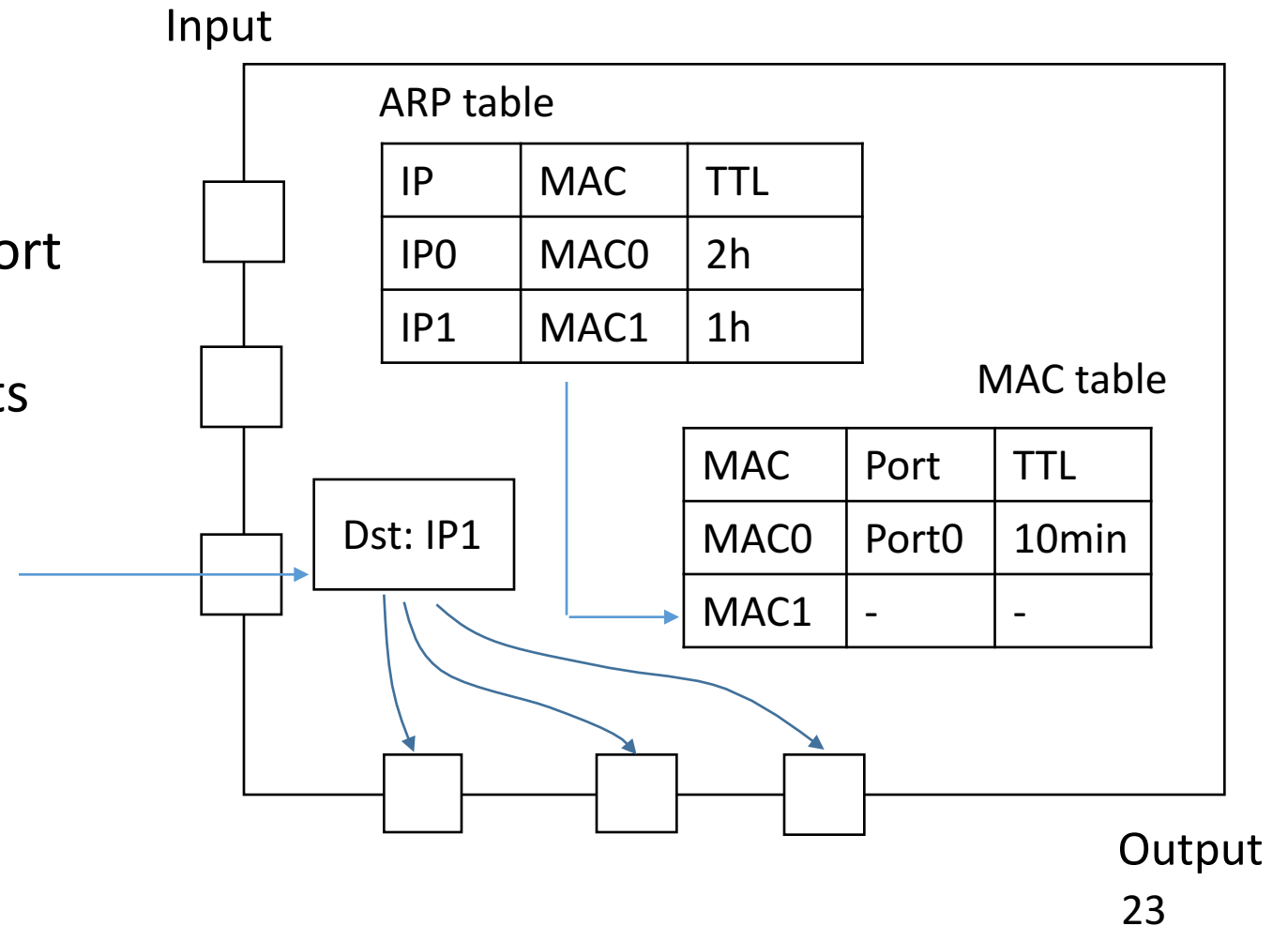
- Our data centers use Clos network
- Packets first travel up then go down
- No cyclic buffer dependency for up-down routing -> no deadlock
- But we did experience deadlock!



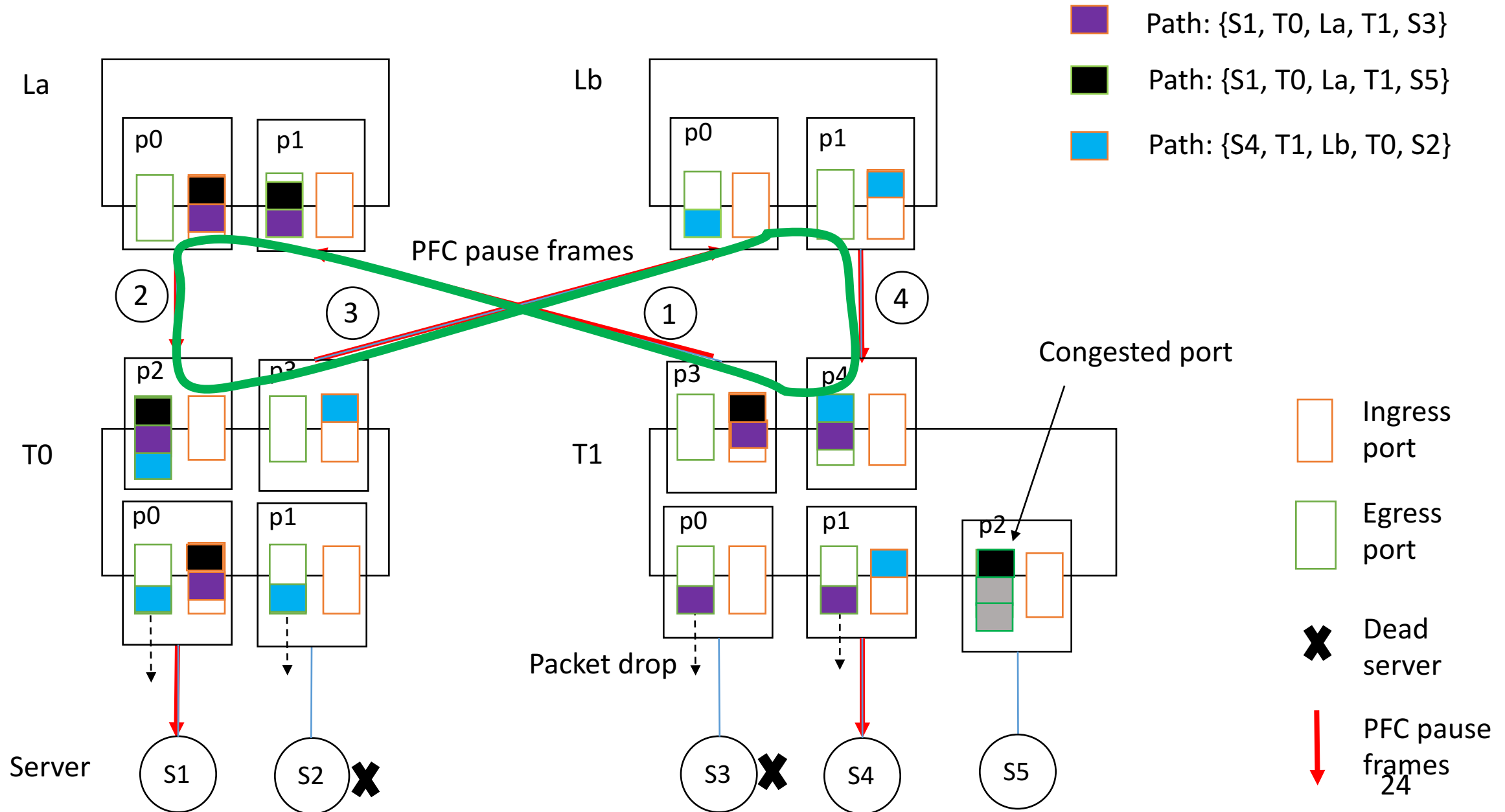
PFC deadlock

- Preliminaries

- ARP table: IP address to MAC address mapping
- MAC table: MAC address to port mapping
- If MAC entry is missing, packets are flooded to all ports



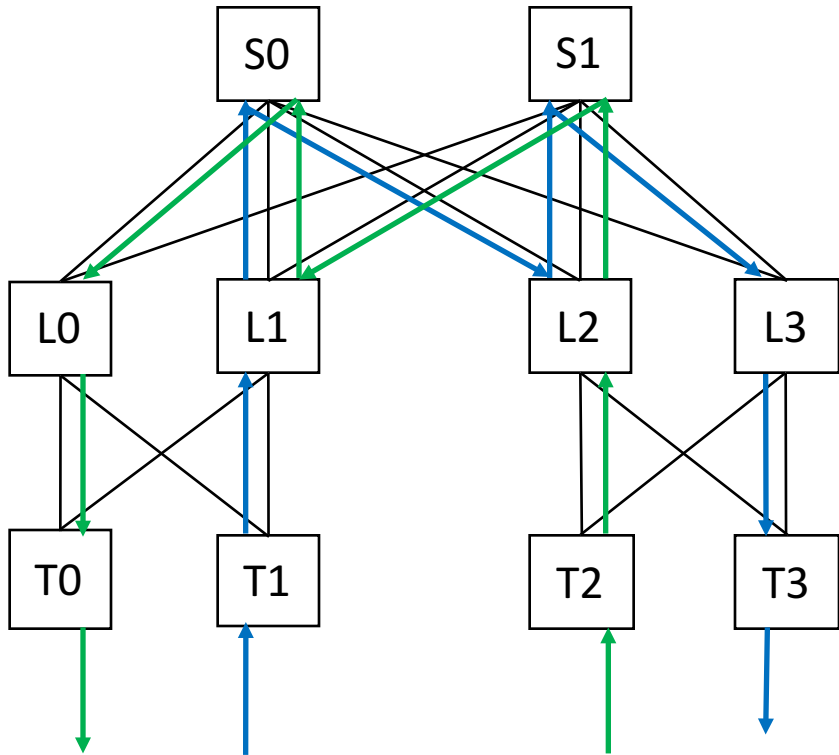
PFC deadlock



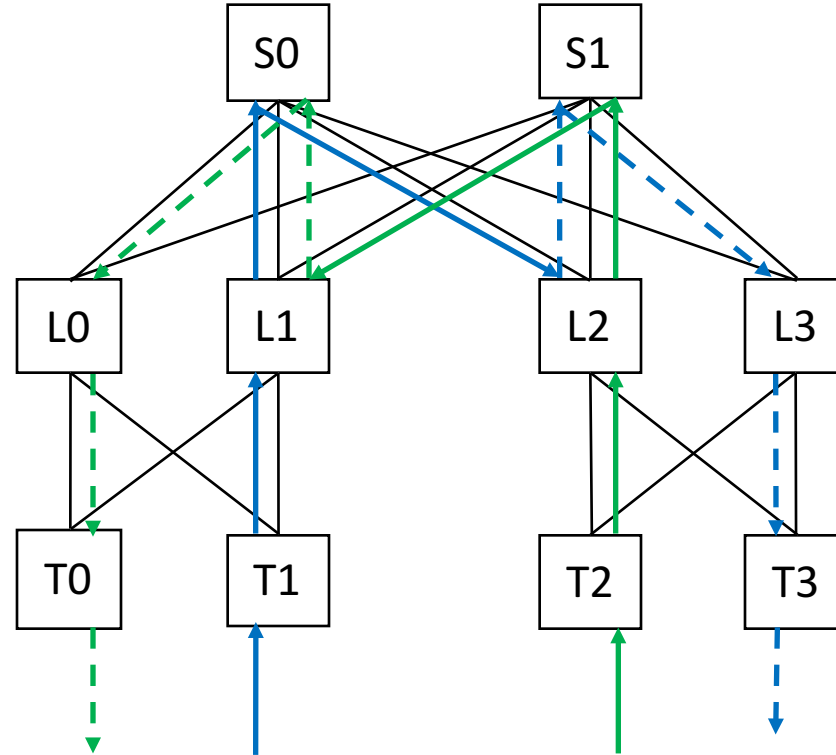
PFC deadlock

- The PFC deadlock root cause: the interaction between the PFC flow control and the Ethernet packet flooding
- Solution: drop the lossless packets if the ARP entry is incomplete
- Recommendation: do not flood or multicast for lossless traffic

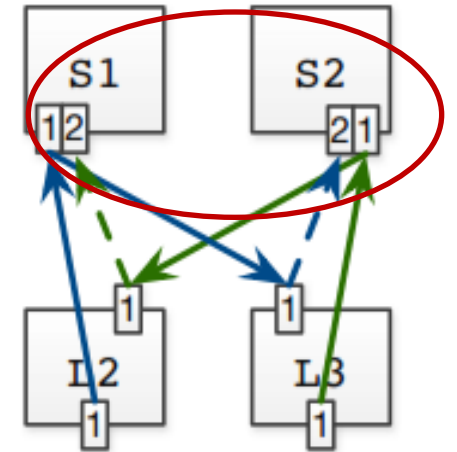
Tagger: practical PFC deadlock prevention



- Concept: Expected Lossless Path (ELP) to decouple Tagger from routing

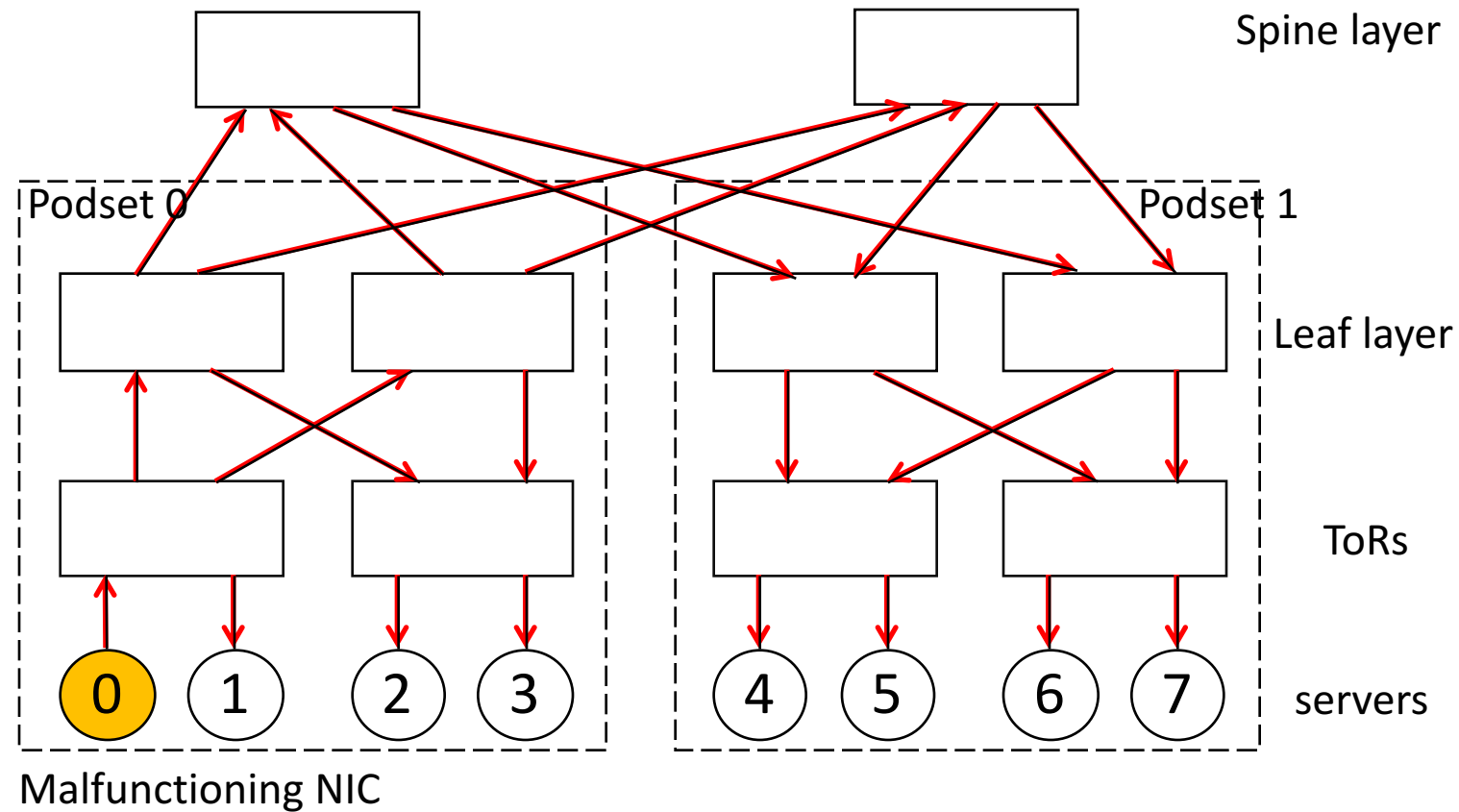


- Strategy: move packets to different lossless queue before CBD forming



- Tagger Algorithm works for general network topology
- Deployable in existing switching ASICs

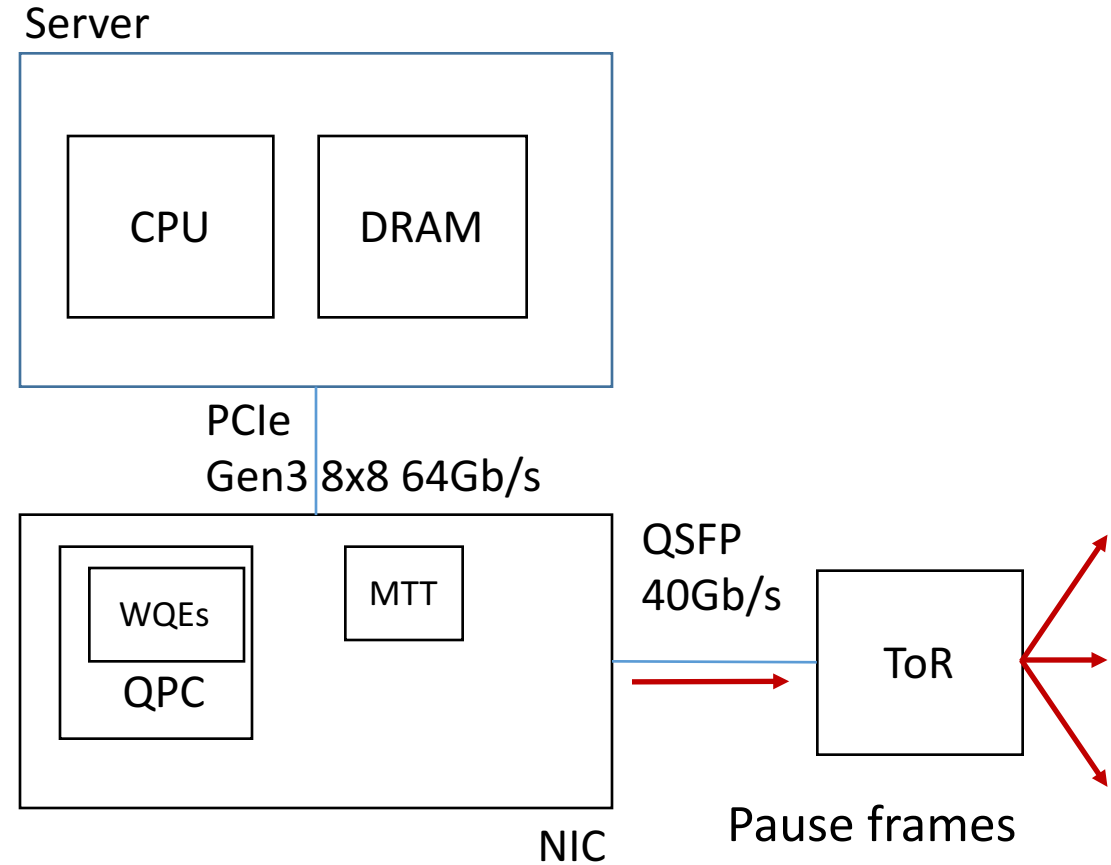
NIC PFC pause frame storm



- A malfunctioning NIC may block the whole network
- PFC pause frame storms caused several incidents
- Solution: watchdogs at both NIC and switch sides to stop the storm

The slow-receiver symptom

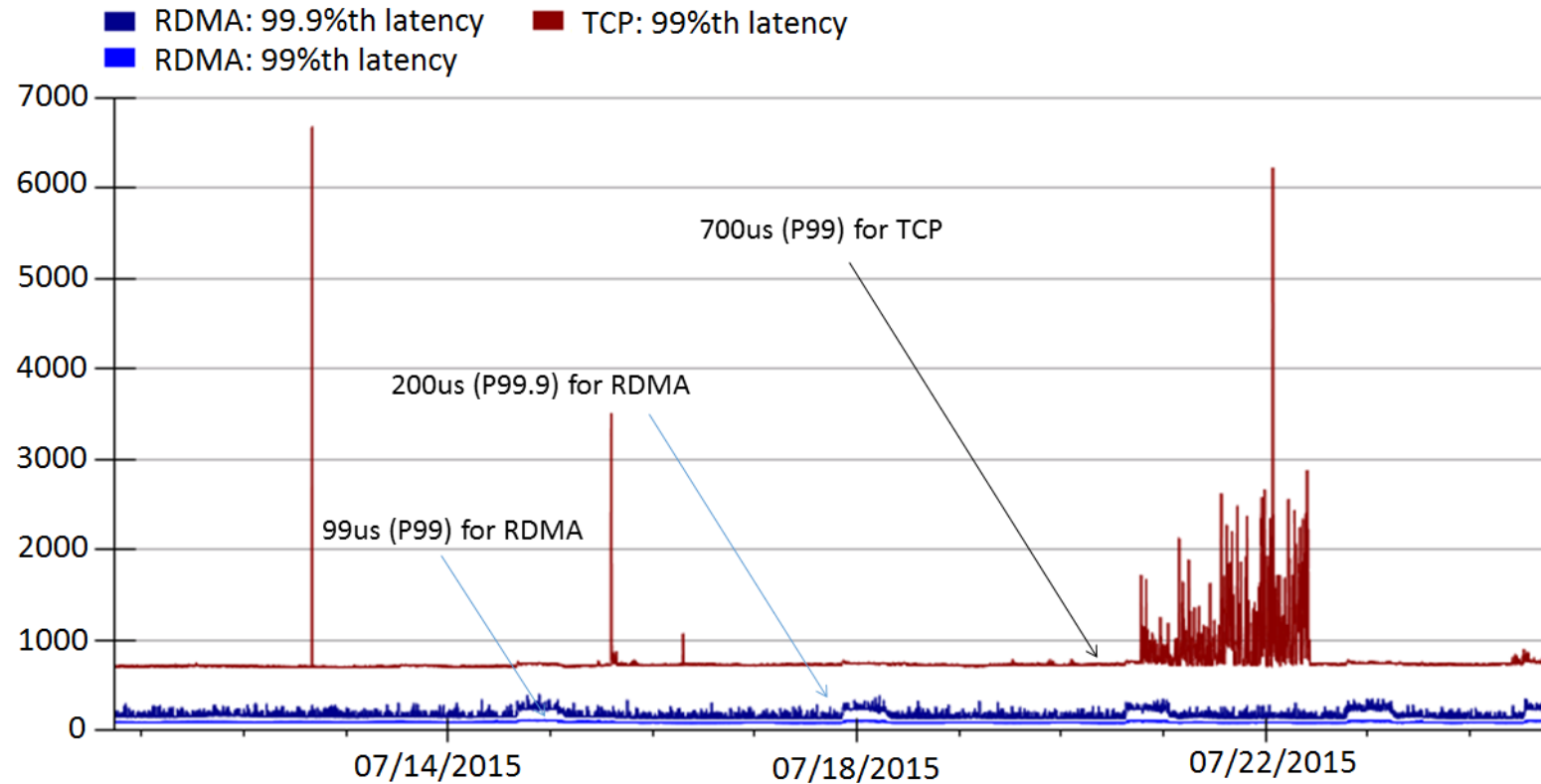
- ToR to NIC is 40Gb/s, NIC to server is 64Gb/s
- But NICs may generate large number of PFC pause frames
- Root cause: NIC is resource constrained
- Mitigation
 - Large page size for the MTT (memory translation table) entry
 - Dynamic buffer sharing at the ToR



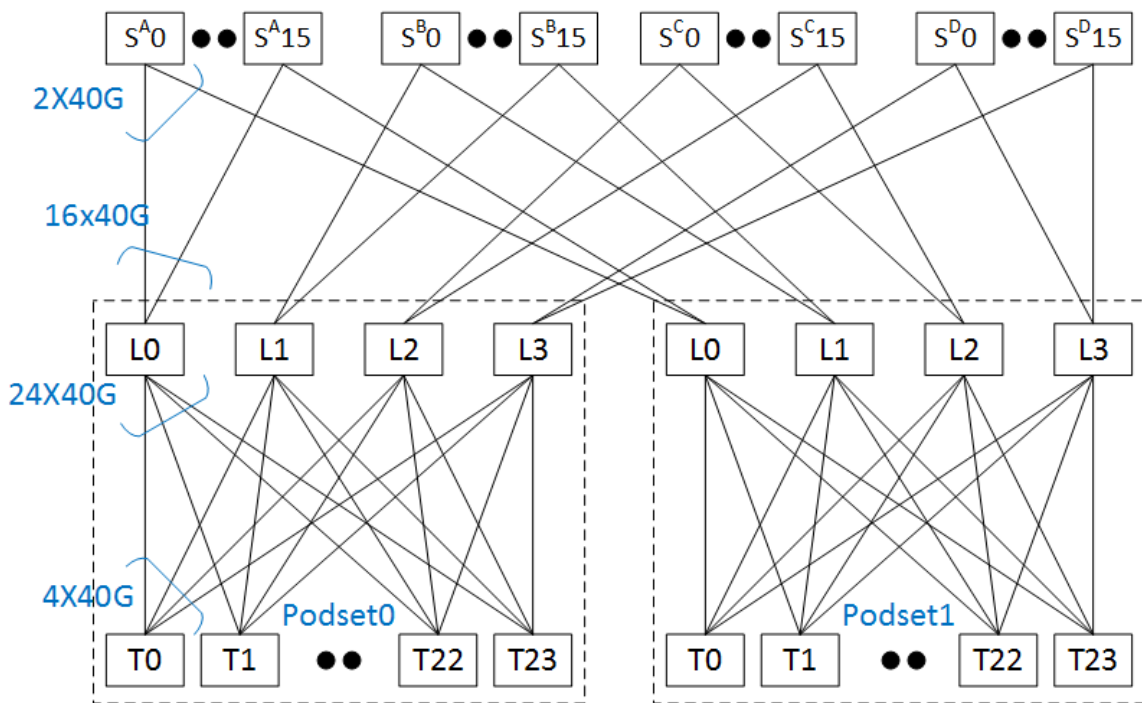
Deployment experiences and lessons learned

Latency reduction

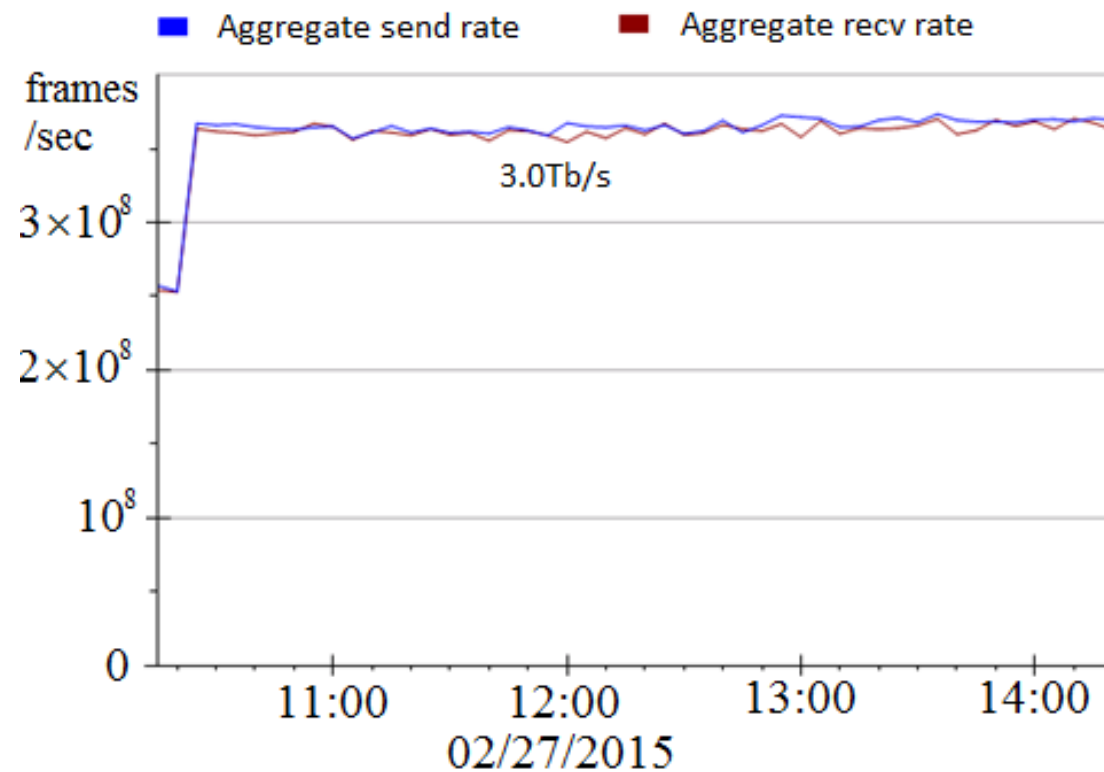
- RoCEv2 deployed in Bing world-wide for two and half years
- Significant latency reduction
- Incast problem solved as no packet drops



RDMA throughput

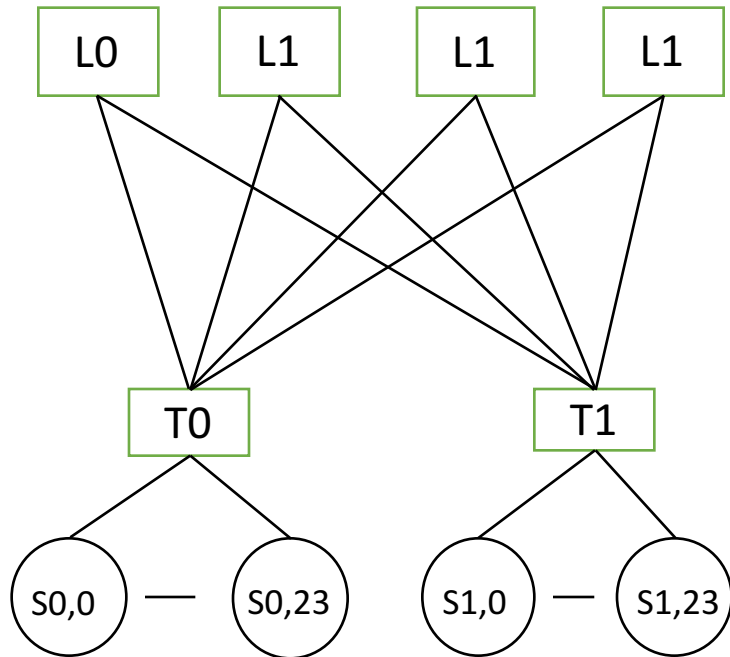


- Using two podsets each with 500+ servers
- 5Tb/s capacity between the two podsets

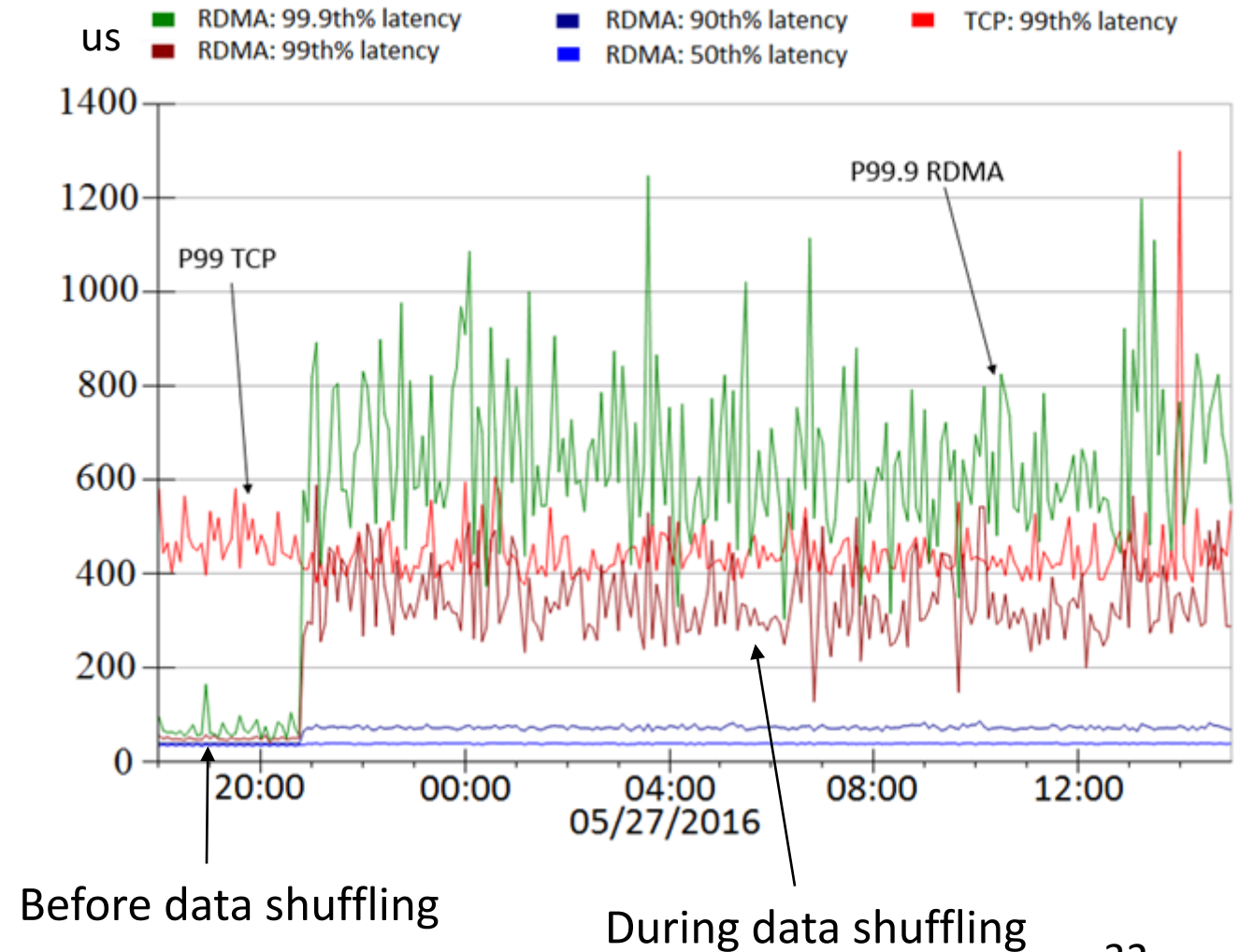


- Achieved 3Tb/s inter-podset throughput
- Bottlenecked by ECMP routing
- Close to 0 CPU overhead

Latency and throughput tradeoff



- RDMA latencies increase as data shuffling started
- Low latency vs high throughput



Lessons learned

- Providing lossless is hard!
- Deadlock, livelock, PFC pause frames propagation and storm did happen
- Be prepared for the unexpected
 - Configuration management, latency/availability, PFC pause frame, RDMA traffic monitoring
- NICs are the key to make RoCEv2 work

What's next?

Applications

- RDMA for X (Search, Storage, HFT, DNN, etc.)

Architectures

- Software vs hardware
- Lossy vs lossless network
- RDMA for heterogenous computing systems

Technologies

- RDMA programming
- RDMA virtualization
- RDMA security
- Inter-DC RDMA

Protocols

- Practical, large-scale deadlock free network
- Reducing colleterial damage

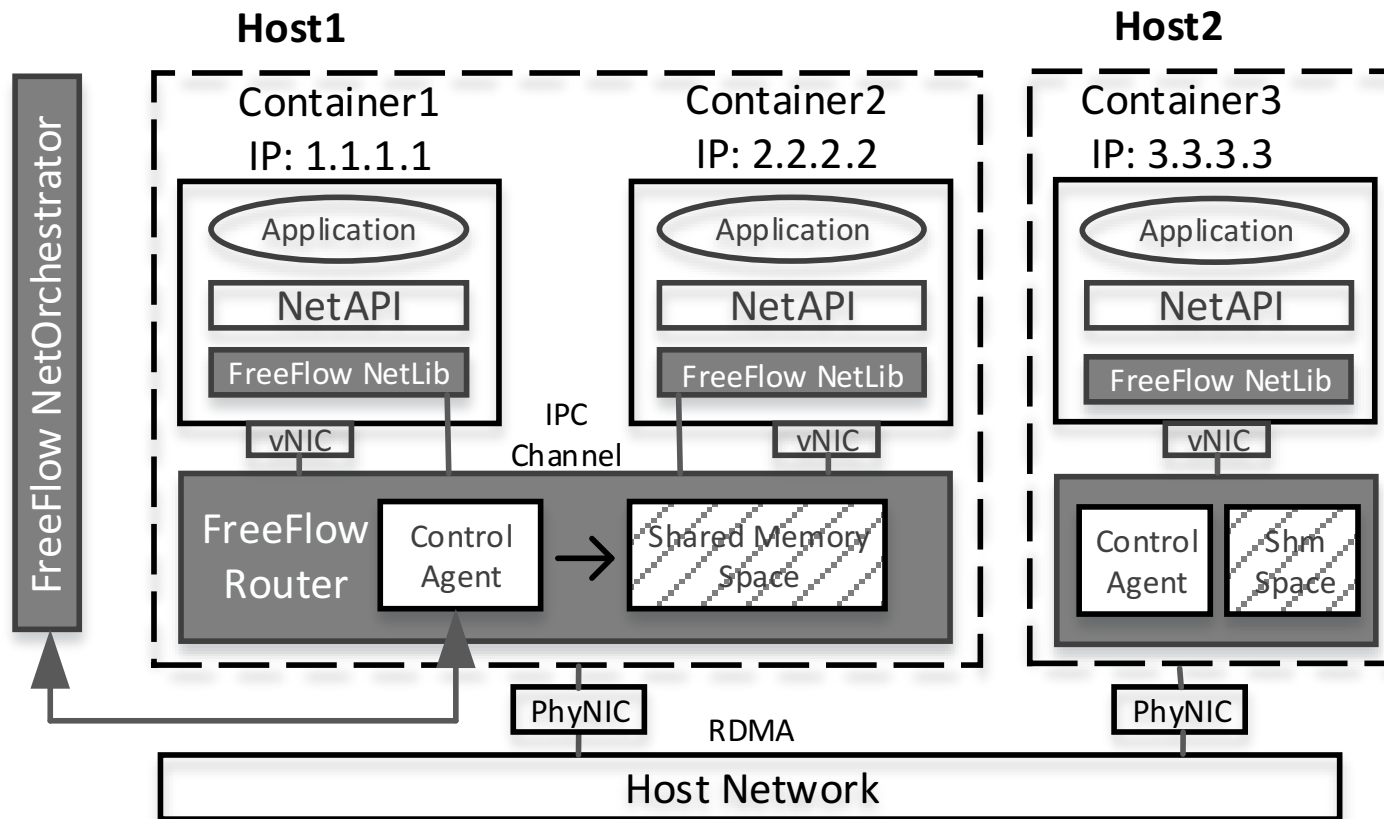
Will software win (again)?

- Historically, software based packet processing won (multiple times)
 - TCP processing overhead analysis by David Clark, et al.
 - Non of the stateful TCP offloading took off (e.g., TCP Chimney)
- The story is different this time
 - Moore's law is ending
 - Accelerators are coming
 - Network speed keep increasing
 - Demands for ultra low latency are real

Is lossless mandatory for RDMA?

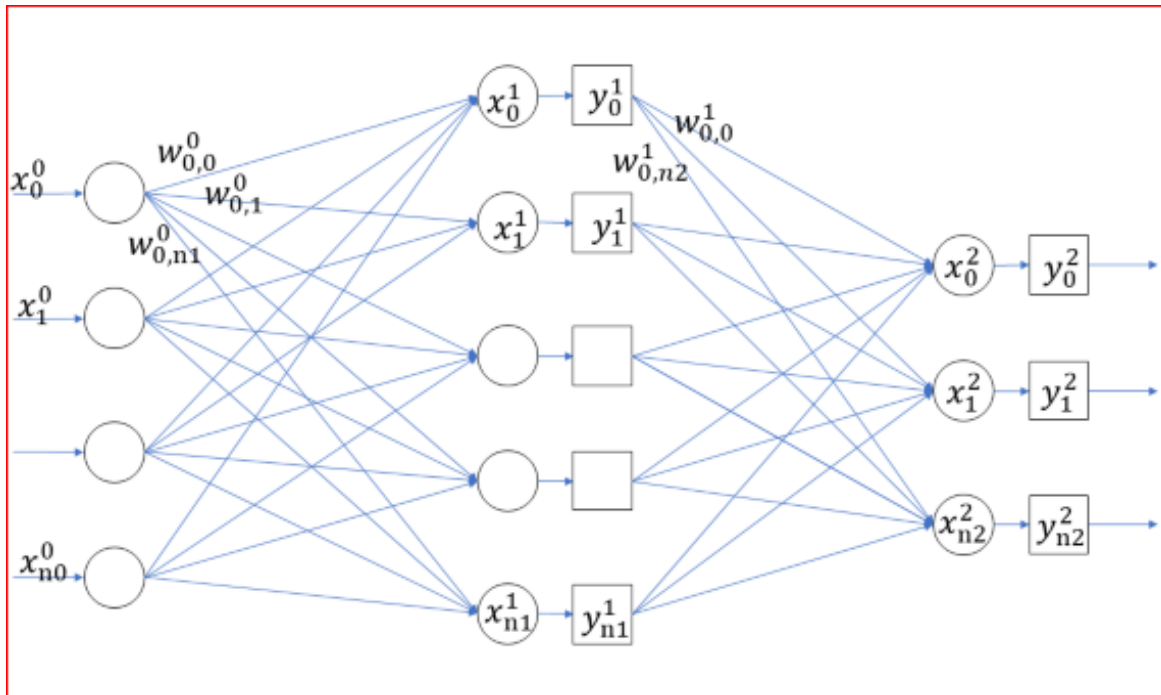
- There is no binding between RDMA and lossless network
- But implementing more sophisticated transport protocol in hardware is a challenge

RDMA virtualization for the container networking



- A router acts as a proxy for the containers
- Shared memory for improved performance
- Zero copy possible

RDMA for DNN



- TCP does not work for distributed DNN training
- For 16-GPU, 2-host speech training with CNTK, TCP communications dominant the training time (72%), RDMA is much faster (44%)

RDMA Programming

- How many LOC for a “hello world” communication using RDMA?
- For TCP, it is 60 LOC for client or server code
- For RDMA, it is complicated ...
 - IBVerbs: 600 LOC
 - RCMA CM: 300 LOC
 - Rsocket: 60 LOC

RDMA Programming

- Make RDMA programming more accessible
 - Easy-to-setup RDMA server and switch configurations
 - Can I run and debug my RDMA code on my desktop/laptop?
 - High quality code samples
- Loosely coupled vs tightly coupled (Send/Recv vs Write/Read)

Summary: RDMA for data centers!

- RDMA is experiencing a renaissance in data centers
 - RoCEv2 has been running safely in Microsoft data centers for two and half years
- Many opportunities and interesting problems for high-speed, low-latency RDMA networking
- Many opportunities in making RDMA accessible to more developers

Acknowledgement

- Yan Cai, Gang Cheng, Zhong Deng, Daniel Firestone, Juncheng Gu, Shuihai Hu, Hongqiang Liu, Marina Lipshteyn, Ali Monfared, Jitendra Padhye, Gaurav Soni, Haitao Wu, Jianxi Ye, Yibo Zhu
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- Arista Networks, Cisco, Dell, Mellanox partners

Questions?