Secure File Sharing on Cloud

Attribute-Based Access Control and Keyword Search over Encrypted Data

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Preliminaries

System Architecture and Security Model

System Workflow

Results and Evaluation

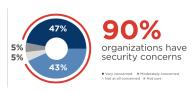
Future Work

Cloud Security

- Separation of Data Ownership and Management
- Threats
 - Malicious Outsiders
 - Malicious Insiders
- Data Leakage: e.g., iCloud 2014, Dropbox 2016



(a) Data Stored on Cloud Servers



(b) Security Concerns

Encrypted Cloud Storage

Data Leakage Prevention → Encrypted Cloud Storage





(c) Tech for Data Leakage Prevention

- (d) Encrypted Cloud
- Conventional Encryption Schemes: Confidentiality & Utility & Flexibility
- New Encryption Schemes: Confidentiality & Utility & Flexibility

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Attribute-Based Encryption





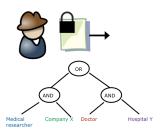
Authority

Functionality: output message if attributes

satisfy the access formulas

CT: associated with access formulas

Key Key: associated with attributes





{Nurse, Hospital Y} 🗶 {MR, Company X} ✓

Attribute-Based Keyword Search





Functionality: could search over the encrypted file if attributes satisfy the access formulas

CT: associated with access formulas Trapdoor: associated with attributes





Preliminaries

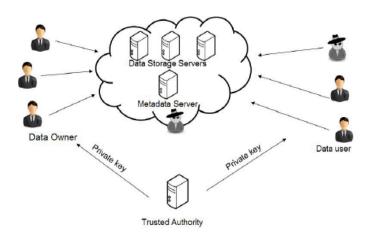
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System Architecture



Security Model

- Trusted Authority: fully trusted
- Cloud Server (Data Storage Server and Metadata Server):
 honest but curious
- Users (Data Owners and Data Users): want to obtain access privilege beyond their private keys

Preliminaries

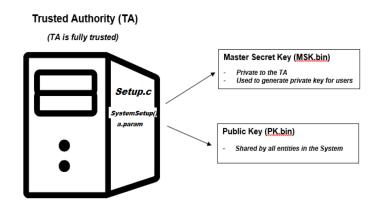
System Architecture and Security Mode

System Workflow

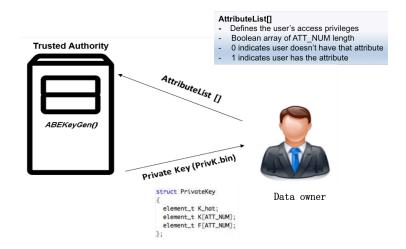
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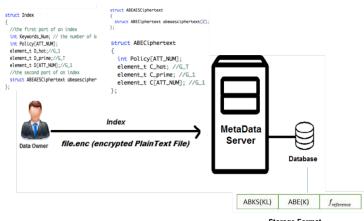
System Setup



Private Key Generation

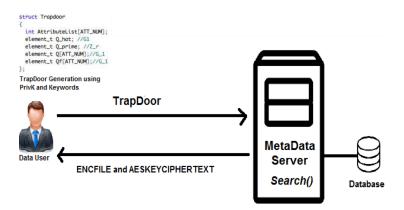


File Upload



Storage Format

Search Request and Response



Preliminaries

System Architecture and Security Mode

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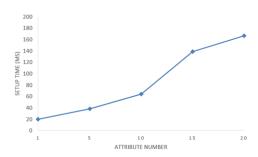
Results and Evaluation

Future Work

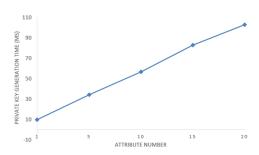
Algorithm Description

System Demonstration

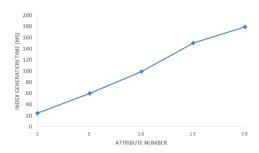
System Setup



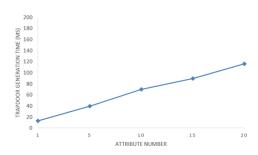
Private Key Generation



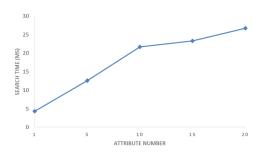
Secure Index Generation



Trapdoor Generation



Search



Preliminaries

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Future Work

- User Revocation
- ► Search over numerical values, e.g., range search
- efficiency enhancement, e.g., resource-constrained devices

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System Setup

SystemSetup() \rightarrow (**PK**, **MK**): It defines a bilinear group \mathbb{G}_1 of prime order p with a generator g. Thus, a bilinear map is defined as $e: \mathbb{G}_1 \times \mathbb{G}_1 \to \mathbb{G}_T$, which has the properties of bilinearity, computability and non-degeneracy. It selects random elements t_1, \dots, t_{3n} . Define a collision-resistant hash function $H: \{0,1\}^* \to \mathbb{Z}_p$. Let $T_k = g^{t_k}$ for each $k \in \{1,\cdots,3n\}$ such that for $1 \le i \le n$, T_i are referred to as positive attributes, T_{n+i} are for negative ones, and T_{2n+i} are thought of as don't care. Let $Y = e(g, g)^y$. The public key is $PK = (e, g, Y, T_1, \dots, T_{3n})$ and the master key is $MK = (y, t_1, \dots, t_{3n})$.

Data Structure

The Public Key data structure in C:

```
struct PublicKey
  element_t g;//G_1
  element_t Y;//G_T
  element_t T[PARAM_NUM];//G_1
};
struct MasterKey
  element_t y;//Z_r
  element_t t[PARAM_NUM];//Z_r
};
```

Key Generation

ABEKeyGen(intAttributeList[], structPrivateKey * PrivK): For every attribute, TA selects random number r_i from \mathbb{Z}_p hence $r = \sum_{i=1}^n r_i$. K_hat is set as g^{y-r} . For AttributeList[i] = 1, set $K_i = g^{\frac{r_i}{l_i}}$ and $K_i = g^{\frac{r_i}{l_{n+1}}}$ otherwise. Finally, let F_i be $g^{\frac{r_i}{l_{2n+1}}}$. The secret key $PrivK = (K_hat, \{K_i, F_i\})_{i \in [1, \cdots, ATT_NUM]}$

Data Structure in C:

```
struct PrivateKey
{
   element_t K_hat;
   element_t K[ATT_NUM];
   element_t F[ATT_NUM];
};
```

Data Upload

SecureIndexGeneration(int Policy[], char * keywords[]): Generate a random $s \in \mathbb{Z}_p$, set $D_hat = g^s$, $D_prime = Y^s$. Given the policy $\mathbf{Policy} = \{\mathbf{Policy}[\mathbf{i}]_{\mathbf{i} \in [1, \cdots, \mathbf{ATT_NUM}]}\}$, for each $i \in [1, \cdots, ATT_NUM]$. If Policy[i] = 1, $D_i = T^s_i$, if Policy[i] = 2 $D_i = T^s_{n+i}$, else $D_i = T^s_{2n+i}$. The generated keyword ciphertext is (D_hat, D_prime, D) .

Data Structure in C:

```
struct Index
{
    //the first part of an index
    int Keywords_Num; // the number of keywords in the file
    int Policy[ATT_NUM];
    element_t D_ntat;//G_1
    element_t D_prime;//G_T
    element_t D_Prime;//G_T
    //the second part of an index
    struct ABEAESCiphertext abeaesciphertext;
};
```

Data Upload

ABEEncrypt(int Policy[], element_t plaintext, struct ABECiphertext* abeciphertext): Generate a random $s \in \mathbb{Z}_p$, set $C_prime = g^s$, $C_hat = Y^s$. Given the policy Policy = $\{Policy[i]_{i \in [1, \dots, ATT_NUM]}\}$, for each $i \in [1, \dots, ATT_NUM]$. If Policy[i] = 1, $C_i = T_i^s$, if Policy[i] = 2 $C_i = T_{n+i}^s$, else $C_i = T_{2n+i}^s$. The generated keyword ciphertext is (C_hat, C_prime, C) .

Data Structure in C:

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Search Request

TrapdoorGeneration(struct PrivateKey PrivK, char* keyword):

```
Generate a random u \in \mathbb{Z}_p. Let Q\_hat = K\_hat^u and Q_i = K_i^u, Qf_i = F_i^u. For the same i', Q_{i'} = K_{i'}^{H(w') \cdot u}, where w' is the keyword of interest and Qf_{i'} = F_{i'}^{H(w') \cdot u}. The trapdoor is (Q\_hat, Q\_prime, Q, Qf).
```

Data Structure in C:

```
struct Trapdoor
{
  int AttributeList[ATT_NUM];
  element_t Q_hat; //G1
  element_t Q_prime; //Z_r
  element_t Q[ATT_NUM];//G_1
  element_t Qf[ATT_NUM];//G_1
};
```

Search Response

SearchIndex(struct Index ind, struct Trapdoor trapdoor): First compare the Policy[] in the index and the AttributeList[] in trapdoor. If data user's AttributeList[] does not satisfy data owner's Policy[], then skip this index, since the data user does not have the search and access privilege of this file. Or else, the following algorithm is run: check whether $D_prime^{Q_prime} = e(D_hat, Q_hat) \cdot \Pi_{i=1}^{ATT_NUM} e(D_i, Q_i^*) \text{ holds,}$ where $Q_i^* = Q_i$ if data user has that attribute, $Q_i^* = Qf_i$ otherwise.

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System Demonstration

```
https://www.youtube.com/watch?v=t7rhjdX030Y
http:
//gitlab.thothlab.org/rgupta36/SecureFileStorageCloud
```

Thanks!