1 Q1

Let us take crossover probability as p.

Thus,
$$P_{error} = 1 - [(1-p)^6 + p(1-p)^5 + p(1-p)^5 + p(1-p)^5 + p(1-p)^5 + p(1-p)^5 + p(1-p)^5 + p^2(1-p)^4]$$

Which gives us $1 - [(1-p)^6 + 6p(1-p)^5 + p^2(1-p)^4]$

2 Q2

For this code, k=2, and n=8, giving us an information rate of 1/4. Let us assume there exists a more optimal code, i.e a [7,2,5] code. We can check this via the Plotkin bound.

$$M \leq 2 \lfloor \tfrac{d}{2d-n} \rfloor$$

$$\begin{array}{l} \lfloor \frac{d}{2d-n} \rfloor = 2 \lfloor \frac{5}{10-7} \rfloor = 2*1 \\ 2 \not > M \text{ where } M = 4 \end{array}$$

Thus the code is optimal, as no [7,2,5] code exists.

3 Q3

For this code, k = 3, and n = 6, giving us an information rate of 1/2. Let us see if a [5,3,3] code exists.

$$\begin{array}{l} 2\lfloor \frac{3}{6-5} \rfloor = 2*3 = 6 \\ 6 \not> M, \text{ where } M = 8 \end{array}$$

Thus the code is optimal, as no [5,3,3] code exists.