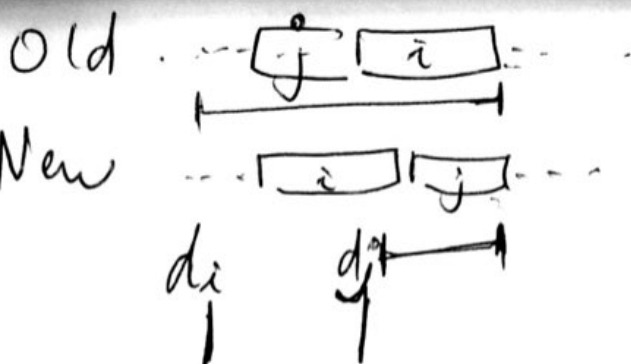
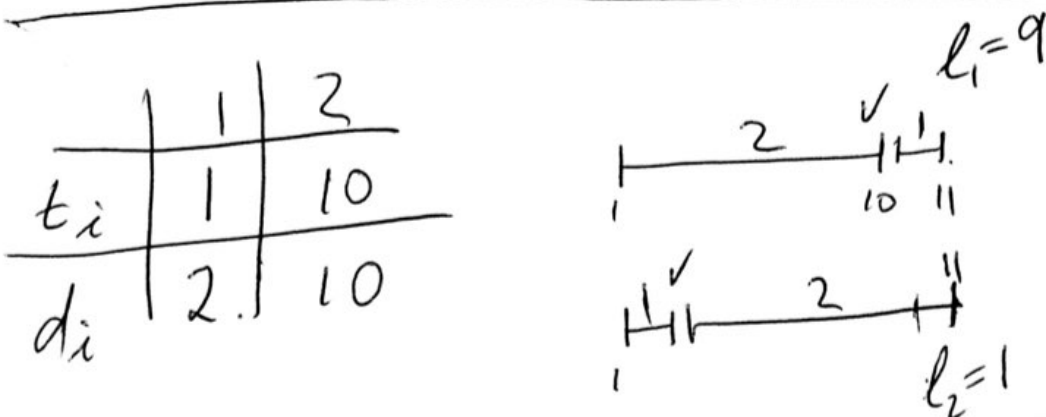
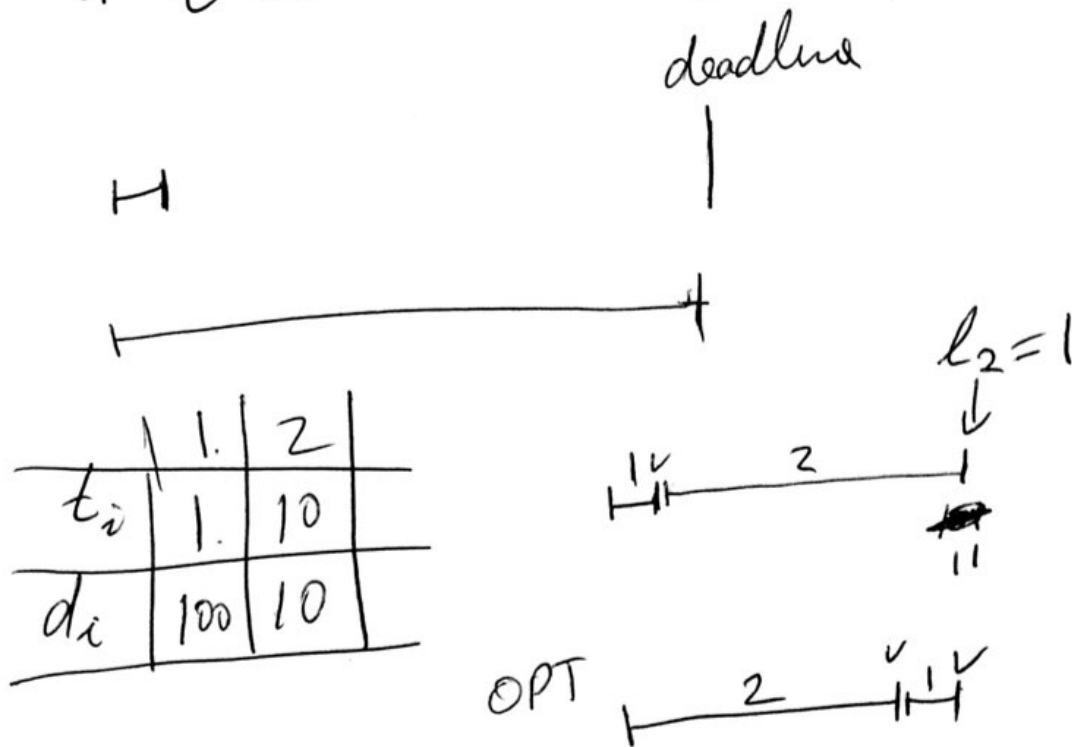


Homework Due Friday before class
4:40

Lecture on Friday cancelled.



Theorem: Greedy is optimal.

- Let S^* be an optimal solution that has the fewest number of inversions
- Let S be the greedy solution
- If S^* has 0 inversions, then $S = S^*$ and therefore Greedy is optimal.
- Otherwise.
 - Let i and j be an adjacent / consecutive inversion
 - Swapping i and j does not increase the max lateness but decreases the number of inversions

- This "swapped" schedule is also opt but has less inversions than S^* .

- Contradicts the fact that S^* was the optimal solution w/ the fewest inversions