

Homework 3 has been posted

The date of Midterm ~~has~~
is: 9/23

In Class. During regular time.

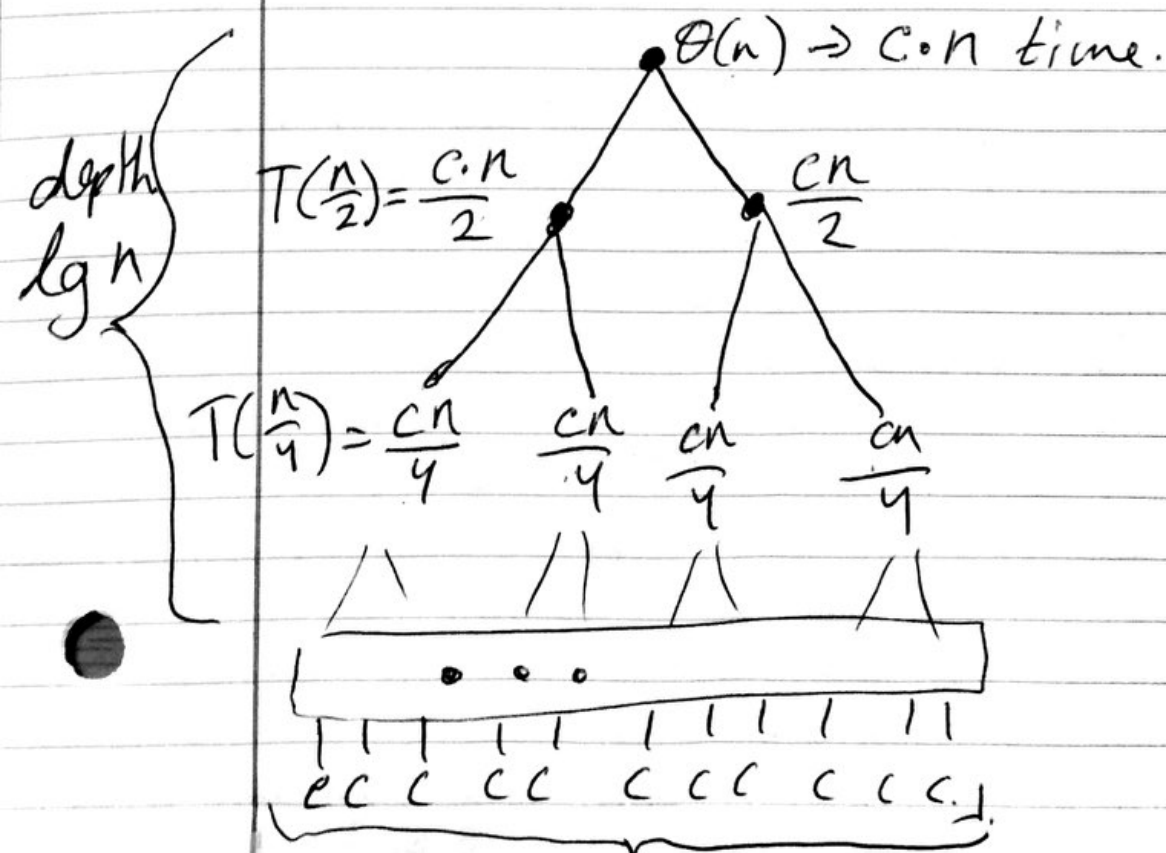
~~Here~~ Coming Recitation: Solutions for HW2.

Solving Recurrence relations

$$T(n) = 2T\left(\frac{n}{2}\right) + \theta(n) = 2T\left(\frac{n}{2}\right) + cn$$

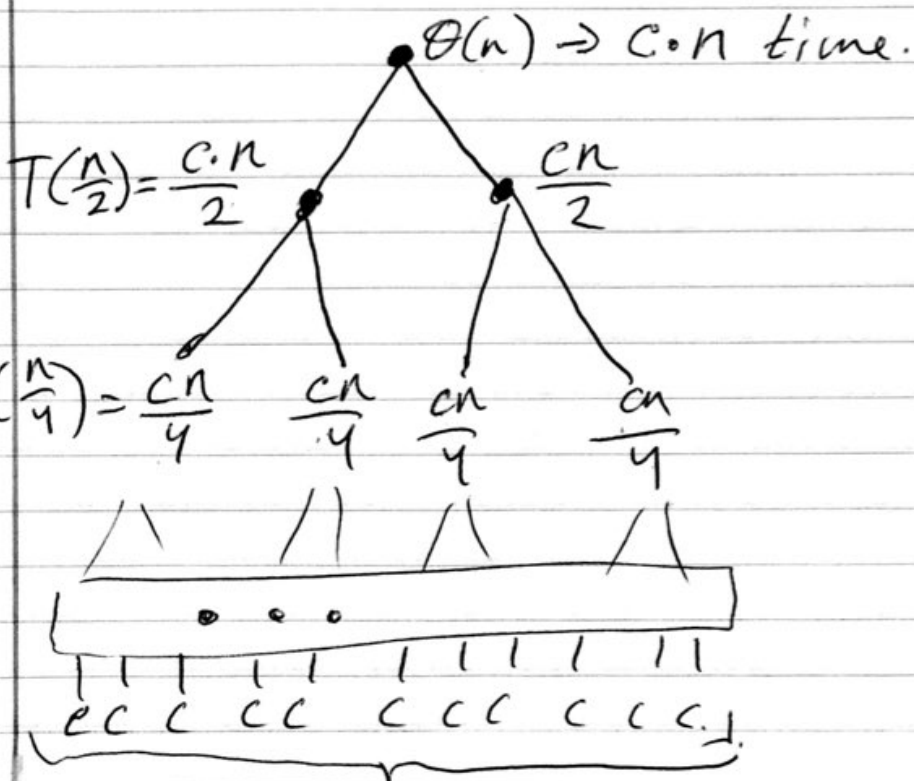
↑ runtime for inputs of size n for some c
How to solve this

1st Step! Guess a solution using a recursion tree.



1st Step! Guess a solution using a recursion tree.

depth
 $\lg n$



$$\frac{\text{Sum}}{c \cdot n}$$

$$\frac{c \cdot n}{2} + \frac{c \cdot n}{2} = c \cdot n$$

$$\frac{c \cdot n}{4} \cdot 4 = c \cdot n$$

\vdots

$c \cdot n$

Each level of tree takes $c \cdot n$ time.
there are $\lg n$ levels.
Total time = $c \cdot n \cdot \lg n$. ← This is our guess.

Let's verify that $T(n) = \Theta(n \lg n)$.

Need to show two things.

$$\underline{T(n) = O(n \lg n)} \quad \text{and} \quad \underline{T(n) = \Omega(n \lg n)}.$$

$T(n) = O(n \lg n) \leftarrow$ let's prove this!

$$T(n) \leq d \cdot n \lg n \quad \text{for some constant } d$$

we use induction on input size with IH
equal $\rightarrow \ominus$

$$T(n) \leq 2T\left(\frac{n}{2}\right) + cn$$

$$T(n) \leq 2\left[\frac{dn}{2} \lg \frac{n}{2}\right] + cn. \quad \text{IH implies } T\left(\frac{n}{2}\right) \leq d \cdot \frac{n}{2} \lg \frac{n}{2}$$

$$\uparrow \leq dn \lg n$$

need to show $2\left[\frac{dn}{2} \lg \frac{n}{2}\right] + cn = \underline{\frac{dn}{2} \lg \frac{n}{2}} + cn$

$$= \underline{dn \lg n} - \underbrace{dn + cn}_{\text{residual should be } \leq 0}$$

$$\begin{aligned} & \left[\lg n - \lg 2 \right] \\ &= \lg n - 1 \end{aligned}$$

$$\begin{aligned} -dn + cn &\leq 0 \\ cn &\leq dn \\ c &\leq d \end{aligned}$$

d cannot depend on n
but can depend on c .

Lets show $T(n) = \Omega(n \lg n)$.

Equivalently, need to give a constant $d > 0$

$$\text{s.t } \forall n, T(n) \geq d n \lg n.$$

$$T(n) = 2T\left(\frac{n}{2}\right) + cn$$

$$\text{by IH.} \rightarrow \geq 2\left[d \frac{n}{2} \lg \frac{n}{2}\right] + cn$$

$$= d n \lg \frac{n}{2} + cn = d n \lg n - \underbrace{dn + cn}$$

need to show $\rightarrow \geq d n \lg n$

residual should be ≥ 0 .

$$-dn + cn \geq 0$$

$$cn \geq dn$$

$$c \geq d. \quad \checkmark$$

Lets try to prove $T(n) = \Omega(n^3)$ and lets see what goes wrong.

$$\text{IH: } T(n) \geq d n^3$$

$$T(n) = 2T\left(\frac{n}{2}\right) + cn$$

$$= 2\left(d \left(\frac{n}{2}\right)^3\right) + cn = \frac{2dn^3}{8} + cn = \frac{dn^3}{4} + cn$$

need to show $\rightarrow \geq dn^3$

$$\frac{dn^3}{4} + cn = dn^3 - \frac{3n^3d}{4} + cn \quad \left[\frac{dn^3}{4} = \frac{dn^3 - 3dn^3}{4} \right]$$

residual should be ≥ 0

$$-\frac{3}{4}n^3d + cn \geq 0$$

$$cn \geq \frac{3}{4}n^3d$$

$$c \geq \frac{3n^2}{4}d$$

$$d \leq \frac{4c}{3n^2} \quad (\forall n)$$

~~*~~ Remember - d cannot depend on n

Such a value of d cannot exist,
therefore, we have not proven.

$$T(n) = \Omega(n^3)$$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + T\left(\frac{n}{8}\right) + n$$

Recursion Tree for a guess

