

**UNIVERSITY BORDEAUX 1**  
**MASTER INFORMATIC - SOFTWARE ENGINEERING**

**A PROJECT REPORT**  
**ON**

**FORMAL DESIGN**



**BY**

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# ABSTRACT

Formal Methods in System Design allows the designing, implementing, and validating the correctness of the system. In this project, I will implement and prove the correctness of the machines that check a given array if it's sorted in ascending order, sorted (ascending or descending), check two given arrays if they are identicals, if one array included the other, and sort a given array in ascending order

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# 1 INTRODUCTION

In the last lesson of Formal Design, we were given a project to finish. Our job is to design several machines that :

- check if a given array is sorted in ascending order
- check if a given array is sorted (ascending or descending order)
- check if two given arrays are identical
- given two arrays, check if one array included the other one
- given an array with pairwise distinct values, sort this array in ascending order
- given an array, sort this array in ascending order

I used a software, named **Rodin**, to do this project. Some theories will be proved in this paper, if cannot be proved in the program.

## 2 ARRAY'S CONTEXT

In this project, we defined an array as a morphism

$$f : X \rightarrow \mathbb{Z}, X = \{i | i \in \mathbb{N} \wedge 1 \leq i \leq n\} \quad (n \text{ is the size of our array})$$

## 3 MACHINE'S SPECIFICATION

Our machine's job is to evaluate an array if it satisfies the condition, means that it is sorted in ascending order. So we have to give the definition of the condition first

If an array satisfies the condition, we must have :

$$\forall i, i \in 1..n - 1 \implies array(i) \leq array(i + 1) \quad (n \text{ is the size of array})$$

Means that if an array doesn't satisfy the condition, we will have :

$$\exists i, i \in 1..n - 1 \implies array(i) > array(i + 1) \quad (n \text{ is the size of array})$$

## 4 MACHINE'S IMPLEMENTATION

For examining an array, we need to run through all of the members of the array. So I introduced a variable *indice* which hold the current position of the last array's element that was visited.

$$indice \in 1..size$$

The machine will stop when it find an element that doesn't satisfy the condition in section 3. That means all the member before the *indice* must satisfy the condition.

$$\forall i \cdot i \in 1..size \wedge i < indice \implies array(i) \leq array(i + 1)$$

At first *indice* start with value of 1. Then, as long as it satisfies the condition below, we increase *indice* by 1

$$indice < size \wedge array(indice) \leq array(indice + 1)$$

Now, our array is sorted in ascending order if  $indice \geq size$ , and it isn't if  $indice < size \wedge array(indice) > array(indice + 1)$

To test that our machine will not stop half way, we need to justify that the expression below is true

$$(indice \geq size) \vee (indice < size \wedge array(indice) > array(indice + 1))$$

$$\vee (indice < size \wedge array(indice) \leq array(indice + 1))$$

To test that our machine will not fall into an infinite loop, we need to point out a value that would decrease each time we go through a loop, and is greater than zero. I choose this value :  $size - indice$ . This value is indeed decreased each time we go through the loop, as  $indice$  would increase and  $size$  is fixed, and because of the definition of  $indice$ ,  $size - indice \geq 0$

## 5 CONCLUSION

After finishing this project, I have obtained many knowledge, first of all, is the usage of *Rodin* and the way to define a system in a logical way. I have to think about the problems that I have never thought about before when programming. And most importantly, is that I could find out that all my knowledges about boolean algebra are still useful, and it could contribute to build a system in a clear and logical way.