2.1: Simple loops in MIPS assembly language

Topics:

loops in MIPS stepping through a MIPS program tracing a MIPS program with breakpoints adding output statements for debugging

Introduction:

In your CSc 256 lectures, you should have seen by now how we can rewrite loops with if-else statements, and implement them in MIPS. (If you haven't, you should probably wait until this material has been covered in class.) In this exercise, we'll look at a program with a simple for loop, and trace through it using spim.

Steps:

Earlier, we copied some files from ~whsu/csc256/LABS/PROGS/. One of these files is pretty much Example 2.4 from your lecture slides, illustrating a for loop. An excerpt:

```
# sum
        $s0
# i
        $s1
# limit $s2
         .data
endl:
        .asciiz "\n"
         .text
         .globl
                 loop
         .globl
                 cont
                                   #
main:
        li
                 $v0,5
                                       cin >> limit;
        syscall
        move
                 $s2,$v0
                 $s0,0
        li
        li
                 $s1,1
                                   #
                                       for (i=1; i<=limit; i++)
        bgt
                 $s1,$s2,cont
                 $s0,$s0,$s1
loop:
        add
                                   #
                                         sum = sum + i;
                 $s1,$s1,1
        add
                 $s1,$s2,loop
        ble
```

The main difference is we added these two lines:

```
.globl loop
.globl cont
```

As we mentioned before, **.text** is an assembler directive that says "the following is the start of the program code and should be read-only". Then comes the assembler directive **.globl loop**, which means "make loop a global label".

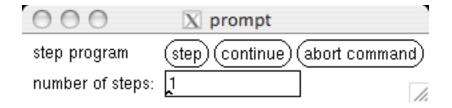
If we do not make **loop** a global label (i.e., if we delete the line .globl loop), the program will still run correctly. However, spim will not be able to "see" the label **loop** when we trace the program. Making **loop** a global label does not change how the programs runs; however, it helps us to trace and debug the program, if necessary.

The simplest way to walk through a program, instruction by instruction, is just to step through it using spim's **step** command. Let's invoke spim, load Example 2.4, and just hit **step** many times (note that hitting **[RETURN]** repeats the last command. Note that we start off in some code that doesn't appear in the source file for Example 2.4! These are instructions from spim's *kernel*; they set up the execution environment for your main program. Don't worry about them for now:

```
(spim) lo "2.4"
(spim) step
[0x00400000]
                0x8fa40000
                             lw $4, 0($29)
                                                               ; 183:
lw $a0 0($sp)
                              # argc
(spim) step
[0x00400004]
                0x27a50004
                             addiu $5, $29, 4
                                                               ; 184:
addiu $a1 $sp 4
                              # argv
(spim)
[0x00400008]
                0x24a60004
                             addiu $6, $5, 4
                                                               ; 185:
addiu $a2 $a1 4
                              # envp
(spim)
[0x0040000c]
                0x00041080
                             sll $2, $4, 2
                                                               ; 186:
sll $v0 $a0 2
(spim)
[0x00400010]
                                                               ; 187:
                0x00c23021
                             addu $6, $6, $2
addu $a2 $a2 $v0
(spim)
[0x00400014]
                0x0c100009
                             jal 0x00400024 [main]
                                                               ; 188:
jal main
(spim)
               finally we reach our main program here ...
[0x00400024]
                0x34020005 ori $2, $0, 5
                                                               ; 32: li
```

```
$v0,5
                #
                    cin >> limit;
(spim)
[0x00400028]
                0x0000000c syscall
                                                              ; 33:
syscall
3
          user enters limit...
(spim) step
[0x0040002c]
                             addu $18, $0, $2
                0x00029021
                                                              ; 34:
move $s2,$v0
(spim) step
[0x00400030]
                0x34100000
                             ori $16, $0, 0
                                                              ; 36: li
$s0,0
(spim) step
[0x00400034]
                0x34110001 ori $17, $0, 1
                                                              ; 37: li
                    for (i=1; i<=limit; i++)
$s1,1
(spim)
[0x00400038]
                0x0251082a slt $1, $18, $17
                                                              ; 38:
bgt
      $s1,$s2,cont
(spim)
[0x0040003c]
                0x14200005
                             bne $1, $0, 20 [cont-0x0040003c]
(spim)
[0x00400040]
                             add $16, $16, $17
                                                              ; 39:
                0x02118020
add
      $s0,$s0,$s1
                      #
                             sum = sum + i;
(spim)
[0x00400044]
                0x22310001
                             addi $17, $17, 1
                                                              ; 40:
      $s1,$s1,1
add
(spim)
                             slt $1, $18, $17
[0x00400048]
                0x0251082a
                                                              ; 41:
ble
      $s1,$s2,loop
(spim)
[0x0040004c]
                0x1020fffd beg $1, $0, -12 [loop-0x0040004c]
(spim)
```

This will work pretty much the same way in xspim as well. To step through one single instruction, just click on the **step** button. You'll see this pop-up window:



Click on the **step** button in the window. The line of code being executed will be highlighted in the Text Segment window; watch how it advances one instruction at a time.

Text Segments					
[0x00400000]	; 183: lw \$a0 0(\$sp)# arg ; 184: addiu \$a1 \$sp 4# a ; 185: addiu \$a2 \$a1 4# e ; 186: s11 \$v0 \$a0 2 ; 187: addu \$a2 \$a2 \$v0 ; 188: ja1 main ; 189: nop ; 191: li \$v0 10				

		Text Segments	
[0x00400000]	0x8fa40000	lw \$4, 0(\$29)	; 183: lw \$a0 0(\$sp)# arg
[0x00400004]	0x27a50004	addiu \$5, \$29, 4	; 184: addiu \$a1 \$sp 4# a
[0x00400008]	0x24a60004	addiu \$6, \$5, 4	; 185: addiu \$a2 \$a1 4# e
[0x0040000c]	0x00041080	sll \$2, \$4, 2	; 186: sll \$v0 \$a0 2
[0x00400010]	0x00c23021	addu \$6, \$6, \$2	; 187: addu \$a2 \$a2 \$v0
[0x00400014]	0x0c100009	jal 0x00400024 [main]	; 188: jal main
[0x00400018]	0x000000000	nop	; 189: nop
[0x0040001c]	0x3402000a	orî \$2, \$0, 10	; 191: li \$v0 10

		Text Segments	-
	0x8fa40000	lw \$4, 0(\$29)	; 183: lw \$a0 0(\$sp)# arg
[0x00400004] [0x00400008]	0x27a50004 0x24a60004	addiu \$5, \$29, 4 addiu \$6, \$5, 4	; 184: addiu \$a1 \$sp 4# a ; 185: addiu \$a2 \$a1 4# e
[0x0040000c]	0x00041080	sll \$2, \$4, 2	; 186: sll \$ v 0 \$a0 2
[[0x00400010] [[0x00400014]	0x00c23021 0x0c100009	addu \$6, \$6, \$2 jal 0x00400024 [main]	; 187: addu \$a2 \$a2 \$v0 ; 188: jal main
[0x00400018]	0x000000000	nop	; 189: nop
[0x0040001c]	0x3402000a	ori \$2, \$0, 10	; 191: li \$v0 10

And so on...

When stepping through Example 2.4 earlier, you may have noticed something weird. Example 2.4 has an instruction **bgt** \$s1, \$s2, cont. But when we step through it, we see two instructions instead of one:

What's going on here? Remember that MIPS assembly language instructions still have to be translated into low-level *machine language*, before they are executed by the MIPS hardware. Some of these assembly language instructions are actually *pseudo-instructions*, which may translate into more than one machine language instruction. The bgt instruction actually splits into the **slt** (set-less-than)

instruction, and the **bne** (branch-not-equals) instruction. Don't worry about the details for now; we'll see what is actually going on in Chapter 6 on machine language.

Obviously, stepping through a long program is very tedious! We want to execute many lines of a program, and only stop at a few important instructions. Breakpoints let us do this.

We saw earlier how we make certain labels global; this lets us mark them as **breakpoints**. Breakpoints are simply special marked places in the program code. When we run our program and come to an instruction that is marked as a breakpoint, the program stops. This is very useful for debugging.

Now let's try to run Example 2.4, and make it stop at one of our labels, before the program is actually finished! First we invoke spim (or xspim), and load 2.4. To make a global label a breakpoint, use the **bre** command:

To look at the list of breakpoints, use the **list** command:

```
(spim) list
Breakpoint at 0x00400040
(spim)
```

(Note that the actual numeric address may be different, if you run spim on a different system.) The message says "Breakpoint at 0x00400040" because the address of loop is 0x00400040. We can check by printing what's at address 0x400040, with the **print** command (abbreviated **pr**):

```
(spim) pr 0x400040
*[0x00400040] 0x02118020 add $16, $16, $17 ; 39: add $s0,$s0,$s1 # sum = sum + i;
```

It's the same line at the label **loop**. We've made **loop** a breakpoint, and the program will stop when we get to **loop**. Note the asterisk in the left margin; that also indicates that we are stopped at a breakpoint. Let's also set a breakpoint at **cont**, to stop the program when we leave the for loop. Then we start running the program:

The program waits for the user to enter limit; let's enter 10:

3 Breakpoint encountered at 0x00400040 (spim)

As expected, we've stopped at **loop**. Let's check our important registers, \$s0 (sum), \$s1 (i), \$s2 (limit):

They contain what we expect, of course! To continue running the program, and stop at the next breakpoint, we use the **continue** command (abbreviated **cont**), and check \$s0, \$s1, \$s2 on the next iteration:

Then one more iteration, stopping at **loop** again:

Then we go through the last iteration, and stop at the final breakpoint, which is **cont** (0x400050). Again we display \$s0, \$s1 and \$s2:

Setting breakpoints in xspim is very similar. Start xspim, load Example 2.4, then clock on the **breakpoints** button. You'll see this pop-up window:



In the textbox, you can type in the names of the two global labels we used as breakpoints earlier (**loop** and **cont**; one by one, please!), then click on the **add** button.

After entering the two breakpoints, click on the **list** button. If you look in the bottom-most window in xspim, you'll see the list of breakpoints (at addresses 0x400040 for **loop**, and 0x400050 for **cont**):

```
SPIM Version 8.0 of January 8, 2010
Copyright 1990-2010, James R. Larus.
All Rights Reserved.
See the file README for a full copyright notice.
Loaded: /afs/sfsu.edu/f1/whsu/spim-8.0/lib/exceptions.s
Breakpoint at 0x00400050
Breakpoint at 0x00400040
```

Then we can run the program by clicking on **run**; we stop when the program waits for us to enter **limit**. Then we'll stop the first time we reach the **loop** label, indicated (again) by the asterisk in the Text Segment window:

	Text Segments								
	[0x00400028] [0x0040002c]	0x0000000c 0x00029021	syscall addu \$18, \$0, \$2			sysca move	11 \$s2,\$ v 0		
20000	[0x00400030]	0x34100000	ori \$16, \$0, 0	;	36:	li	\$s0,0		_
	[[0x00400034] [[0x00400038]	0x34110001 0x0251082a	ori \$17, \$0, 1 slt \$1, \$18, \$17		37: 38:	li bat	\$s1,1 \$s1,\$s2,	# .con	t∎
	[0x0040003c] *[0x00400040]	0x14200005 0x02118020	bne \$1, \$0, 20 [cont-0x00400030 add \$16, \$16, \$17			_	\$30,\$30		
l	[0x00400040]	0x22110020 0x22310001	addi \$17, \$17, 1	;		add add	\$80,\$80,		*
\vdash			5 . 6 .						-

We can track the contents of registers by looking at the top window, where all the registers are displayed (\$\$0 = 0, \$\$1 = 1, \$\$2 = 3 at the moment):

```
000
                                             X xspim
                                       00400040
  PС
              00400040
                                                    Cause
                                                               00000024
                                                                            BadVAddr= 00000000
                           EPC
  Status
              3000ff10
                           _{
m HI}
                                       00000000
                                                   LO
                                                               00000000
                                        General Registers
 R0
      (r0)
              00000000
                          R8
                               (t0)
                                       00000000
                                                  RĨ6
                                                             = 00000000
                                                                           R24
                                                                                (t8)
                                                                                        00000000
                                                       (s0)
 R1
      (at)
              00000000
                          R9
                               (t1)
                                       00000000
                                                  R17
                                                        (s1)
                                                             =
                                                               00000001
                                                                           R25
                                                                                (t9)
                                                                                        00000000
 R2
              00000003
                          R10
                               (t2)
                                       00000000
                                                  R18
                                                        (s2)
                                                             =
                                                               00000003
                                                                           R26
                                                                                (k0)
      (v0)
 R3
      (v1)
                          R11
                               (t3)
                                       00000000
                                                  R19
                                                        (83)
                                                             =
                                                               00000000
                                                                           R27
                                                                                 (k1)
 R4
                          R12
                                       00000000
                                                  R20
                                                             =
                                                               00000000
                                                                           R28
      (a0)
                               (t4)
                                                        (84)
 R5
                          R13
                                       00000000
                                                  R21
                                                             =
                                                               00000000
      (a1)
                               (t5)
                                                        (s5)
                                                                           R29
                                                                                (sp)
 R6
              7ffffc3c
                          R14
                               (t6)
                                    =
                                       00000000
                                                  R22
                                                        (36)
                                                             =
                                                               00000000
                                                                           R30
      (a2)
                                                                                (s8)
 R7
              00000000
                          R15
                               (t7)
                                    =
                                      00000000
                                                  R23
                                                       (s7)
                                                             = 00000000
                                                                           R31
                                                                                        00400018
```

Then click on **continue** a few more times to run through the program, stopping at a breakpoint each time. Again, this is very similar to what happened in spim.

Here's a summary of what we learned about breakpoints:

- 1. Breakpoints are places in the code where program execution stops.
- 2. To make a label in a MIPS program a breakpoint, we first have to make the label a global label. This is usually done right after the .text directive:

- 3. Just because a label is globalized doesn't mean it's a breakpoint! To make it a breakpoint, we have to go into spim, and use the **bre** command. (**bre label** makes a breakpoint at "label".)
- 4. To list all breakpoints, use the command **list.**
- 5. When we're stopped at a breakpoint, spim will indicate to us the address of the breakpoint that it stopped at (since there might be several breakpoints).
- 6. To continue running a program until the next breakpoint (or until the end of the program), type **continue.**
- 7. To look at register contents, use the **print** command.