Chapter 11 Quasi-Newton Methods

ullet Motivation It takes $\mathcal{O}\left(d^3
ight)$ to calculate $abla^2 f(x)^{-1}$ or solve for $abla^2 f(\mathbf{x}_t) \Delta \mathbf{x} = -
abla f(\mathbf{x}_t)$.

The secant method

- ullet Motivation $rac{f(x_t)-f(x_{t-1})}{x_t-x_{t-1}}pprox f'\left(x_t
 ight)$, so $x_{t+1}:=x_t-rac{f(x_t)}{f'(x_t)}pprox x_{t+1}:=x_t-f'\left(x_t
 ight)rac{x_t-x_{t-1}}{f'(x_t)-f'(x_{t-1})}$
- In optimization regime, we want to find a similar matrix s.t. $\nabla f(\mathbf{x}_t) \nabla f(\mathbf{x}_{t-1}) = H_t(\mathbf{x}_t \mathbf{x}_{t-1})$ and so $\mathbf{x}_{t+1} = \mathbf{x}_t H_t^{-1} \nabla f(\mathbf{x}_t)$.
 - This is called secant condition

Quasi-Newton methods

- ullet Definition If H_t symmetric, and follows secant condition, the update method is *quasi-newton*.
- Lemma Exercise 71 $f \in C^2$ and $\nabla^2 f \neq 0$, then Newton's method is a Quasi-Newton method iff $f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^\top M\mathbf{x} \mathbf{q}^\top \mathbf{x} + c$ with M invertable symmetric.
 - Proof
 - Newtwon is quasi $\Leftrightarrow \nabla f(y) \nabla f(x) = \nabla^2 f(y)(y-x), \forall x, y$, take derivative w.r.t x, we get $\nabla^2 f(x) = \nabla^2 f(y), \forall x, y$ and this means f is a quadratic funciton, its invertable since every secant condition there is a solution. The other direction is straightforward.

Greenstadt's Approach

- We already have $H_{t-1}^{-1}, x_{t-1}, x_t$, we need H_t^{-1} , idea is $H_t^{-1} = H_{t-1}^{-1} + E_t$, and we want to minimized the general change $\|AEA^\top\|_F^2$.
- ullet Denote $H:=H_{t-1}^{-1}$, $H':=H_t^{-1}$, $E:=E_t$, $oldsymbol{\sigma}:=\mathbf{x}_t-\mathbf{x}_{t-1}$, $\mathbf{y}=
 abla f(\mathbf{x}_t)abla f(\mathbf{x}_{t-1})$, and $\mathbf{r}=oldsymbol{\sigma}-H\mathbf{y}$,
 - \circ then the update formula is H'=H+E, such that $H'\mathbf{y}=oldsymbol{\sigma}$ or equivalently $E\mathbf{y}=\mathbf{r}$,
 - \circ so the overall minimization is $ext{minimize}$ $extstyle \frac{1}{2} \|AEA^ op\|_F^2$, $ext{subject to } E\mathbf{y} = \mathbf{r}$ and $E^ op E = \mathbf{0}$.

Solving with Lagrange multiplier

- ullet Fact 11.2 If $f(E):=rac{1}{2}\|AEB\|_{F'}^2$ then $abla f(E)=A^ op AEBB^ op$ if define $abla f(E)=\left(rac{\partial f(E)}{\partial E_{ij}}
 ight)$.
 - Proof
 - By this def, we have $\nabla_E \mathrm{Tr}(AE) \nabla_E = \mathrm{Tr}(E^\top A^\top) = A^\top$, so $f(E) := \frac{1}{2} \|AEB\|_F^2 = \frac{1}{2} \mathrm{Tr}(B^\top E^\top A^\top AEB)$, then $\nabla f(E) = \nabla_E \frac{1}{2} \mathrm{Tr}(E^\top A^\top AE_0 BB^\top) + \nabla_E \frac{1}{2} \mathrm{Tr}(BB^\top E_0^\top A^\top AE) = A^\top AEBB^\top/2 + (BB^\top E^\top A^\top A)^\top/2$ $= A^\top AEBB^\top$
- Denote $\pmb{\lambda} \in \mathbb{R}^d$ as the multiplier for d constrains of $E\mathbf{y} = \mathbf{r}$, and $\Gamma \in \mathbb{R}^{d \times d}$ as the multiplier for $d \times d$ constraints of $E^ op E = 0$.
- ullet For each equation of $\partial_{E_{ij}}f=m{\lambda}^ op f_1+{
 m Tr}(\Gamma f_2)$, $m{\lambda}$ part yields a term of $\lambda_i y_i$ and Γ yields a term of $\Gamma_{ji}-\Gamma_{ij}$
- Lemma 11.3 The above equation gives the optimial conditional of $WE^*W = \lambda y^\top + \Gamma^\top \Gamma$, where $W := A^\top A$ is symmetric and posi definite.

Solving Greenstadt family

- The minimization has now turn into three linear equation (i) $E\mathbf{y}=\mathbf{r}$, (ii) $E^{\top}-E=0$ and (iii) $WEW=\boldsymbol{\lambda}\boldsymbol{y}^{\top}+\Gamma^{\top}-\Gamma$
- To eliminate Γ , by plug (iii) into (ii) we get $M\left(\mathbf{\lambda}\mathbf{y}^{\top}-\mathbf{y}\mathbf{\lambda}^{\top}+2\Gamma^{\top}-2\Gamma\right)M=0$, where $M=W^{-1}$, so $\Gamma^{\top}-\Gamma=\frac{1}{2}\left(\mathbf{y}\mathbf{\lambda}^{\top}-\mathbf{\lambda}\mathbf{y}^{\top}\right)$
 - \circ then $E = rac{1}{2} M \left(oldsymbol{\lambda} \mathbf{y}^ op + \mathbf{y} oldsymbol{\lambda}^ op
 ight) M$
- Then to eliminate $m{\lambda}$, we plug in the secant condition (i) we get $m{\lambda} = rac{1}{\mathbf{y}^{ op} M \mathbf{y}} ig(2 M^{-1} \mathbf{r} \mathbf{y} m{\lambda}^{ op} M \mathbf{y} ig)$
 - \circ multiply with $\mathbf{y}^ op M$, we get $z = oldsymbol{\lambda}^ op M \mathbf{y} = rac{\mathbf{y}^ op \mathbf{r}}{\mathbf{y}^ op M \mathbf{y}}$, so $oldsymbol{\lambda} = rac{1}{\mathbf{y}^ op M \mathbf{y}} \Big(2 M^{-1} \mathbf{r} rac{(\mathbf{y}^ op \mathbf{r})}{\mathbf{y}^ op M \mathbf{y}} \mathbf{y} \Big)$
- Plug this into E, we get $E = \frac{1}{2}M\left(\mathbf{\lambda}\mathbf{y}^{\top} + \mathbf{y}\mathbf{\lambda}^{\top}\right)M = \frac{1}{\mathbf{y}^{\top}M\mathbf{y}}\left(\mathbf{r}\mathbf{y}^{\top}M + M\mathbf{y}\mathbf{r}^{\top} \frac{(\mathbf{y}^{\top}\mathbf{r})}{\mathbf{y}^{\top}M\mathbf{y}}M\mathbf{y}\mathbf{y}^{\top}M\right)$, by definition $\mathbf{r} = \boldsymbol{\sigma} H\mathbf{y}$, we get
 - $egin{aligned} \circ \stackrel{\cdot}{E}^{\star} &= rac{1}{\mathbf{y}^{ op} M \mathbf{y}} \Big(oldsymbol{\sigma} \mathbf{y}^{ op} M + M \mathbf{y} oldsymbol{\sigma}^{ op} H \mathbf{y} \mathbf{y}^{ op} M M \mathbf{y} \mathbf{y}^{ op} H rac{1}{\mathbf{y}^{ op} M \mathbf{y}} ig(\mathbf{y}^{ op} oldsymbol{\sigma} \mathbf{y}^{ op} H \mathbf{y} ig) M \mathbf{y} \mathbf{y}^{ op} M ig) \end{aligned}$

BFGS (Broyden, Fletcher, Goldfarb and Shanno)

- Definition BFGS is when $M = H' = H_t^{-1}$, $M\mathbf{y} = H'\mathbf{y} = \boldsymbol{\sigma}$, even we don't know H', but it never appears in the solution, so $E^\star = \frac{1}{\mathbf{y}^\top \boldsymbol{\sigma}} \left(-H\mathbf{y} \boldsymbol{\sigma}^\top \boldsymbol{\sigma} \mathbf{y}^\top H + \left(1 + \frac{\mathbf{y}^\top H\mathbf{y}}{\mathbf{y}^\top \boldsymbol{\sigma}}\right) \boldsymbol{\sigma} \boldsymbol{\sigma}^\top \right)$, where $H = H_{t-1}^{-1}$, $\boldsymbol{\sigma} = \mathbf{x}_t \mathbf{x}_{t-1}$, $\mathbf{y} = \nabla f(\mathbf{x}_t) \nabla f(\mathbf{x}_{t-1})$.

 Iteration cost is $O(d^2)$
- Lemma Exercise 74.1 If f convex, $\mathbf{y}^{\top} \sigma > 0$, unless $\mathbf{x}_t = \mathbf{x}_{t-1}$ or $f(\lambda \mathbf{x}_t + (1-\lambda)\mathbf{x}_{t-1}) = \lambda f(\mathbf{x}_t) + (1-\lambda)f(\mathbf{x}_{t-1})$ for all $\lambda \in (0,1)$.
 - Proof
 - ullet By property of convexity $\mathbf{y}^ op \sigma = (
 abla f(\mathbf{x}_t)
 abla f(\mathbf{x}_{t-1}))^ op (\mathbf{x}_t \mathbf{x}_{t-1}) \geq 0$
 - $\text{If } (\nabla f(\mathbf{x}_t) \nabla f(\mathbf{x}_{t-1}))^\top (\mathbf{x}_t \mathbf{x}_{t-1}) = 0 \text{ while } \exists \lambda \text{ s.t. } f(\lambda \mathbf{x}_t + (1-\lambda)\mathbf{x}_{t-1}) < \lambda f(\mathbf{x}_t) + (1-\lambda)f(\mathbf{x}_{t-1}),$
 - $\text{ then by convexity } f(\lambda \mathbf{x}_t + (1-\lambda)\mathbf{x}_{t-1}) \geq f(\mathbf{x}_{t-1}) + \lambda \nabla f(\mathbf{x}_{t-1})^\top (\mathbf{x}_t \mathbf{x}_{t-1}) \text{ this means } \\ f(\mathbf{x}_t) f(\mathbf{x}_{t-1}) > \nabla f(\mathbf{x}_{t-1})^\top (\mathbf{x}_t \mathbf{x}_{t-1}), \text{ similarly } f(\mathbf{x}_{t-1}) f(\mathbf{x}_t) > \nabla f(\mathbf{x}_t)^\top (\mathbf{x}_{t-1} \mathbf{x}_t)$
 - lacksquare add them together, we get $(\nabla f(\mathbf{x}_t) \nabla f(\mathbf{x}_{t-1}))^{\top}(\mathbf{x}_t \mathbf{x}_{t-1}) > 0$ contradiction.
- Observation 11.6 $H' = \left(I \frac{\sigma \mathbf{y}^{\top}}{\mathbf{y}^{\top} \sigma}\right) H \left(I \frac{\mathbf{y} \sigma^{\top}}{\mathbf{y}^{\top} \sigma}\right) + \frac{\sigma \sigma^{\top}}{\mathbf{y}^{\top} \sigma}$
 - o Proof
 - $E^{\star} + H = H + \frac{1}{\mathbf{y}^{\top} \boldsymbol{\sigma}} \left(-H \mathbf{y} \boldsymbol{\sigma}^{\top} \boldsymbol{\sigma} \mathbf{y}^{\top} H + \left(1 + \frac{\mathbf{y}^{\top} H \mathbf{y}}{\mathbf{y}^{\top} \boldsymbol{\sigma}} \right) \boldsymbol{\sigma} \boldsymbol{\sigma}^{\top} \right) = \frac{\boldsymbol{\sigma} \boldsymbol{\sigma}^{\top}}{\mathbf{y}^{\top} \boldsymbol{\sigma}} + H (I \frac{\mathbf{y} \boldsymbol{\sigma}^{\top}}{\mathbf{y}^{\top} \boldsymbol{\sigma}}) \boldsymbol{\sigma} \mathbf{y}^{\top} H + \frac{\boldsymbol{\sigma} \mathbf{y}^{\top} H \mathbf{y} \boldsymbol{\sigma}^{\top}}{\mathbf{y}^{\top} \boldsymbol{\sigma}} = \text{QED}$
- Lemma Exercise 74.2 If $H \succeq 0$ and $\mathbf{y}^{\top} \sigma > 0$, then also H' is positive definite.
 - Proof
 - $\blacksquare \text{ Since } C := \left(I \frac{\mathbf{y} \boldsymbol{\sigma}^\top}{\mathbf{y}^\top \boldsymbol{\sigma}}\right) = \left(I \frac{\boldsymbol{\sigma} \mathbf{y}^\top}{\mathbf{y}^\top \boldsymbol{\sigma}}\right)^\top \text{, so } H' = C^\top H C + \frac{\boldsymbol{\sigma} \boldsymbol{\sigma}^\top}{\mathbf{y}^\top \boldsymbol{\sigma}}, C^\top H C \text{ is semi positive definite and so is } \frac{\boldsymbol{\sigma} \boldsymbol{\sigma}^\top}{\mathbf{y}^\top \boldsymbol{\sigma}}.$
 - When $z \perp \sigma^{\top}$, we have $Cz = z \neq 0$, so the two quadratic form will not be zero at the same time, this means positive definiteness.
- Remark Usually Newton or Quasi-Newton are performed with scaled steps $\mathbf{x}_{t+1} = \mathbf{x}_t \alpha_t H_t^{-1} \nabla f(\mathbf{x}_t)$, either line search or backtracking line search (when $\alpha_t = 1$ is not good enough, do $\alpha_t/2$).

L-BFGS (limited memory version)

- Idea Only use information from the previous m iterations, for some small value of m.
- Lemma 11.7 If an oracle can compute $\mathbf{s} = H\mathbf{g}$ for any vertor \mathbf{g} , then $\mathbf{s'} = H'\mathbf{g'}$ can be computed with one oracle call of $\mathbf{s} = H\mathbf{g}$, and O(d) arithmetic operation, assuming σ, \mathbf{y} known.
 - o Proof

$$\blacksquare \ H'\mathbf{g}' = \Big(I - \frac{\sigma \mathbf{y}^\top}{\mathbf{y}^\top \sigma}\Big) H \underbrace{\Big(I - \frac{\mathbf{y}^\top}{\mathbf{y}^\top \sigma}\Big) \mathbf{g}'}_{\mathbf{g}} + \underbrace{\frac{\sigma \sigma^\top}{\mathbf{y}^\top \sigma} \mathbf{g}'}_{\mathbf{h}}$$

- $\mathbf{g}, \mathbf{h}, \mathbf{s}, \mathbf{w}, \mathbf{z}$ all are computed in O(d).
- The idea is that we need $H_t^{-1} \nabla f_t$, and we can borrow from $H_{t-1}^{-1} \nabla f_t$, etc, and recurse back to t=0, This gives the BFGS-step:
- Algorithm (BFGS-STEP)
 - \circ Input (k, \mathbf{g})
 - \circ If k=0 then return $H_0^{-1}\mathbf{g}'$
 - Else
 - lacksquare Set $\mathbf{h} = oldsymbol{\sigma} rac{oldsymbol{\sigma}_k^ op \mathbf{g}'}{\mathbf{y}_k^ op oldsymbol{\sigma}_k}$, and $\mathbf{g} = \mathbf{g}' \mathbf{y} rac{oldsymbol{\sigma}_k^ op \mathbf{g}'}{\mathbf{y}_k^ op oldsymbol{\sigma}_k}$
 - $\mathbf{s} = \mathrm{BFGS\text{-}STEP}\ (k-1,\mathbf{g})$ (recursive call)
 - $\mathbf{w} = \mathbf{s} oldsymbol{\sigma}_k rac{\mathbf{y}_k^{ op} \mathbf{s}}{\mathbf{y}_k^{ op} \sigma_k}$
 - $\mathbf{z} = \mathbf{w} + \mathbf{h}$
 - return z
- Remark If H_0 can be computed in O(d) the total runtime is O(td), this is acceptable when $t \leq d$. It's natual to think of a cut-off version
- Algorithm (L-BFGS-STEP)
 - Input (k, l, \mathbf{g})
 - \circ If l=0 then return $H_0^{-1}\mathbf{g}'$
 - Else

• Set
$$\mathbf{h} = oldsymbol{\sigma} rac{oldsymbol{\sigma}_k^ op \mathbf{g}'}{\mathbf{y}_k^ op oldsymbol{\sigma}_k}$$
, and $\mathbf{g} = \mathbf{g}' - \mathbf{y} rac{oldsymbol{\sigma}_k^ op \mathbf{g}'}{\mathbf{y}_k^ op oldsymbol{\sigma}_k}$

• $\mathbf{s} = \mathbf{L}\text{-BFGS-STEP}\left(k-1,l-1,\mathbf{g}\right)$ (recursive call)

$$lacksquare \mathbf{w} = \mathbf{s} - oldsymbol{\sigma}_k rac{\mathbf{y}_k^ op \mathbf{s}}{\mathbf{y}_k^ op oldsymbol{\sigma}_k}$$

$$\mathbf{z} = \mathbf{w} + \mathbf{h}$$

■ return z