

Quantum Noise Theory



Outline

- Noise Definition
- Noisy Intermediate-Scale Quantum Era
- Quantum Background
- Types of Noise
- Noise Characterization
- Noise Mitigation



What is noise in quantum computing?



Any deviation from the ideal intended quantum state of a quantum circuit.



Noisy Intermediate Scale Quantum Era

- NISQ termed by John Preskill in 2018.
- Noisy refers to limited uncontrollable noise.
- Intermediate-Scale refers to limited available number of qubits (ranging from 50 to a few hundred qubits)



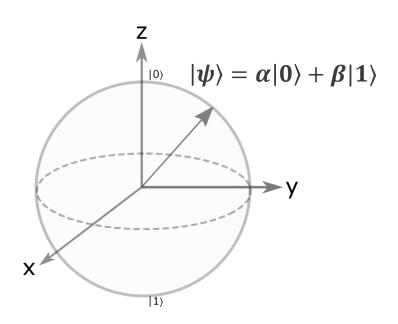
Quantum Background



Quantum Bits

- A classical bit can be either 0 or 1.
- A quantum bit (qubit) can be any linear combination of both 0 and 1.







Qubit Representation

- "\langle " is the Dirac notation also known as the bra-ket notation.
- "<|" is a row vector (also known as the bra).
- "|>" is the column vector (also known as the ket).
- Qubit at state $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$ or state $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- In the bra notation $\langle 0| = \begin{bmatrix} 1 & 0 \end{bmatrix}$ or state $\langle 1| = \begin{bmatrix} 0 & 1 \end{bmatrix}$
- Qubit can be in a linear combination of both states (Superposition !!)
 - $|\psi\rangle=\alpha|0\rangle+\beta|1\rangle$, where α and β are complex numbers
 - $|\psi\rangle = \begin{bmatrix} \alpha \\ 0 \end{bmatrix} + \begin{bmatrix} 0 \\ \beta \end{bmatrix}$
 - $|\psi\rangle = \begin{bmatrix} \alpha \\ \beta \end{bmatrix}$ How much the qubit is in state $|0\rangle$ How much the qubit is in state $|1\rangle$



Pure and Mixed States

- Pure state:
 - can be represented in a vector notation.

$$|\psi_{pure}\rangle = \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_{2^{n}-1} \end{bmatrix}$$

- · Lies on the surface of the Bloch sphere
- can also be represented as a density matrix

$$\rho = |\psi_{pure}\rangle\langle\psi_{pure}|$$



Mixed State:

- probability distribution of several pure states (not superposition!!!)
- Lies inside the Bloch sphere
- can only be represented using a density matrix

 $\rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|$, where p_i represents the probability to be in pure state $|\psi_i\rangle$

Quantum Gates

- Any unitary operator that can be applied to qubits to transform its state.
- Can be performed on single or multiple qubits.
- These gates are reversible.
- Example

$$\bullet \ \ H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$



Applying a Unitary Evolution

- For a state vector
 - $U|\psi\rangle$
- For a density matrix
 - $\rho = \sum_{i} p_{i} U |\psi_{i}\rangle\langle\psi_{i}|U^{\dagger} = U \rho U^{\dagger}$



State Fidelity

- Is the measure of closeness between two quantum states
- $F(\rho, \sigma) = \left\| \sqrt{\rho} \sqrt{\sigma} \right\|_{1}^{2}, \quad 0 \le F \le 1$
 - Where F represents fidelity and ρ and σ represent either pure or mixed state.
- To measure the closeness between the noisy state and the ideal one, the fidelity will be reduced to

$$F(\rho,\sigma) = \langle \psi_{\rho} | \sigma | \psi_{\rho} \rangle$$
, where $\rho = |\psi_{\rho}\rangle \langle \psi_{\rho}|$

• If both states are pure, the fidelity will be $|\langle \psi_{
ho} | \psi_{\sigma} \rangle|^2$



Example

- Compute the fidelity between 2 pure orthogonal states $|0\rangle$ and $|1\rangle$
 - Solution:

$$F = |\langle 0|1\rangle|^2 = \left| \begin{bmatrix} 1 & 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \right|^2 = 0$$



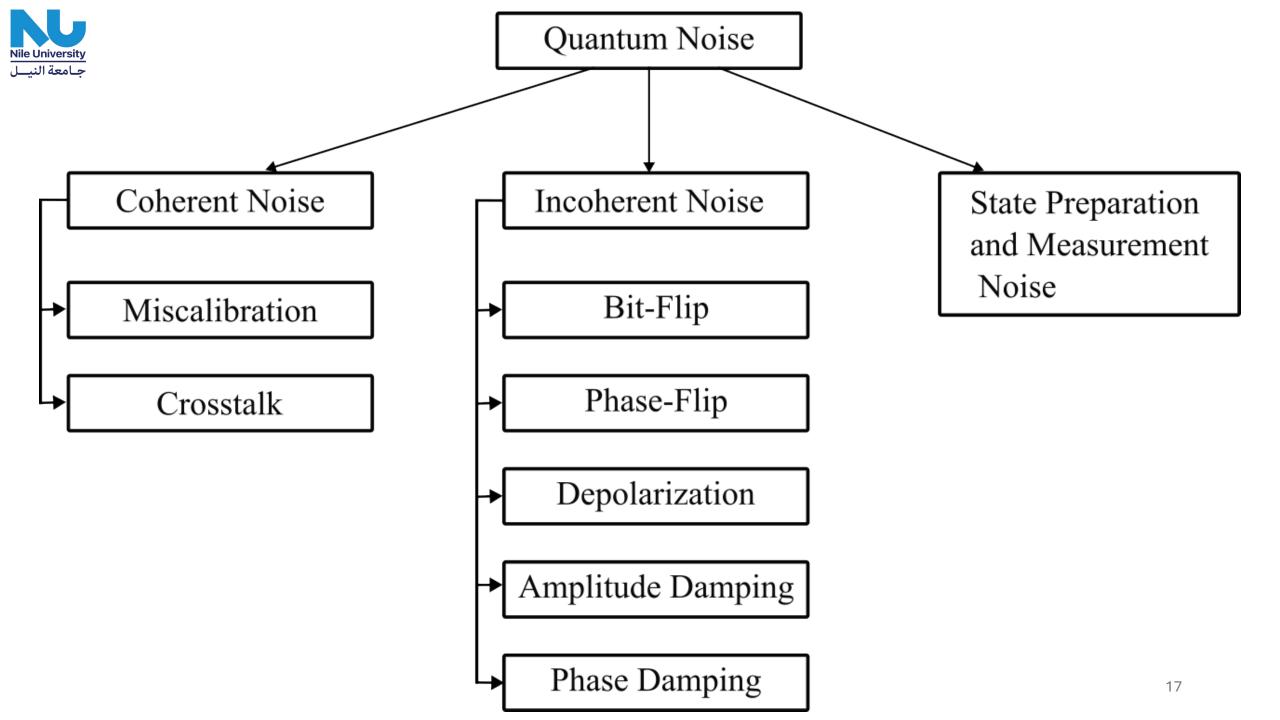
Example

• Compute the fidelity between pure state $\psi_{\rho} = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$ and

Solution:



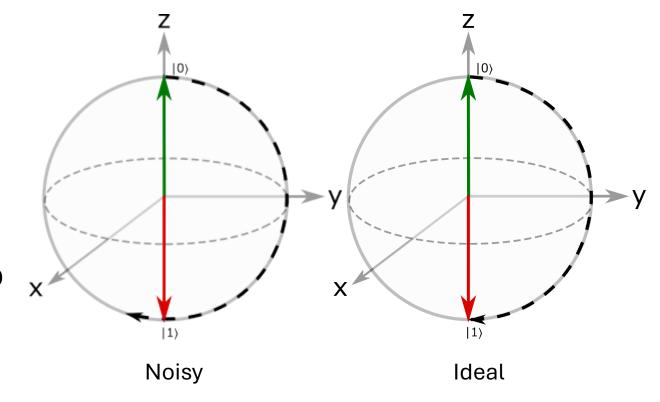
Types of Noise

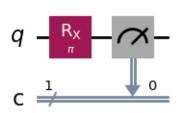


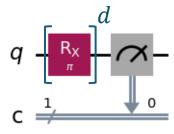


Coherent Errors

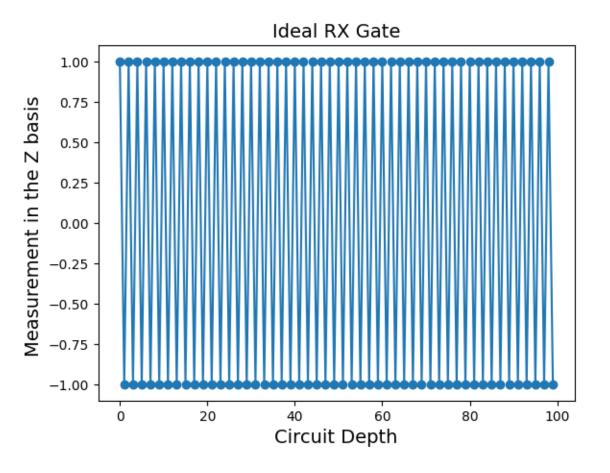
- Miscalibrated gates resulting in over or under rotations.
- Suppose, you have an X gate $(X = R_x(\pi))$,
- Due to noise, the state will end up $\tilde{X} = R_x(\pi + \epsilon)$, where ϵ is the additional error.

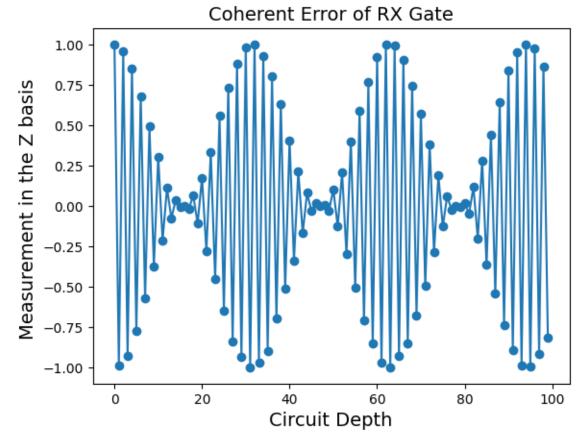














Crosstalk Noise

- Unwanted noise due to execution of multiple gates in parallel.
- Noise added to qubits other than those originally operated by the gate



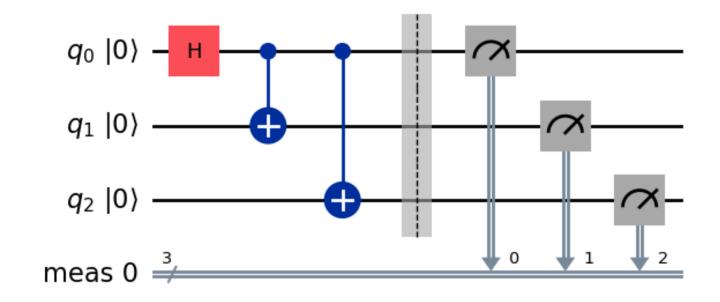
Incoherent Errors

- Unwanted interactions with the environment resulting in a probability distribution of several pure states
- The state can be represented as a mixed state

 $\rho = \sum_j p_j |\psi_j\rangle\langle\psi_j|$, where p_j represents the probability of being in the pure state $|\psi_i\rangle$.



Consider Simulating 3-qubit GHZ Circuit



Ideally, the state before measurement is

$$|\psi\rangle = \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$$



Bit Flip Error

- Flip the qubit state with probability p
- Can be applied after a certain unitary *U*

•
$$\mathcal{E}(\rho) = (1 - p)U\rho U^{\dagger} + p (XU\rho U^{\dagger}X)$$

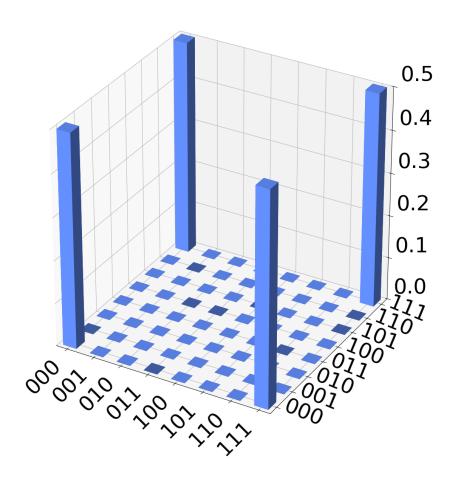


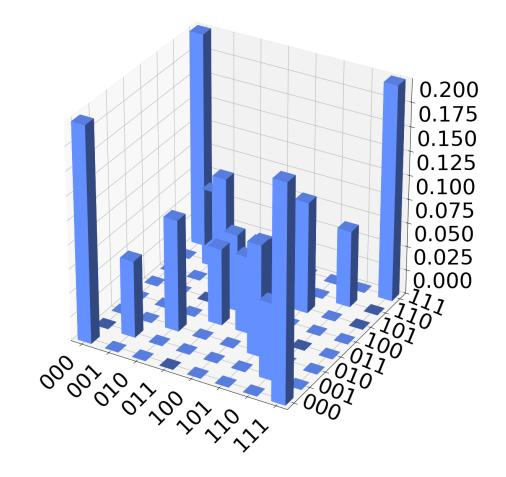
Ideal

Bit Flip (p=0.2)

Real Amplitude (ρ)









Phase Flip Error

- Flip the phase of a qubit state with probability p
- Can be applied after a certain unitary *U*

•
$$\mathcal{E}(\rho) = (1 - p)U\rho U^{\dagger} + p (ZU\rho U^{\dagger}Z^{\dagger})$$

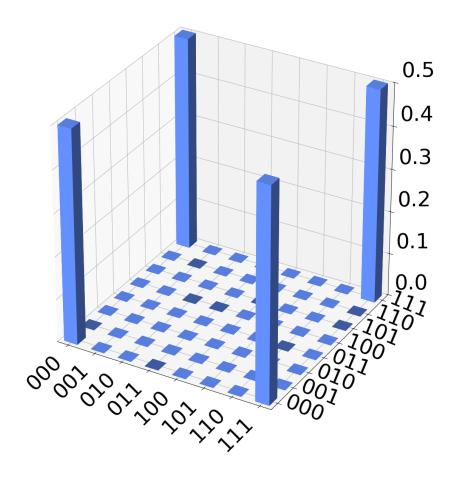


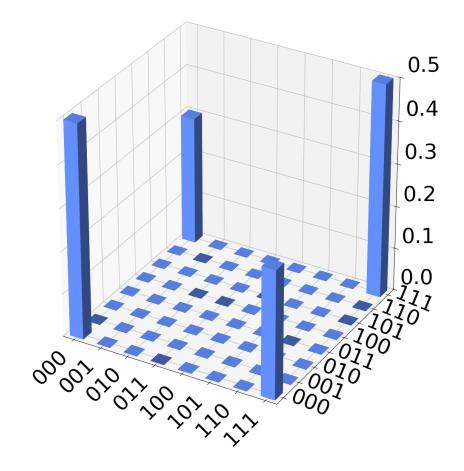
Ideal

Phase Flip (p=0.2)

Real Amplitude (ρ)



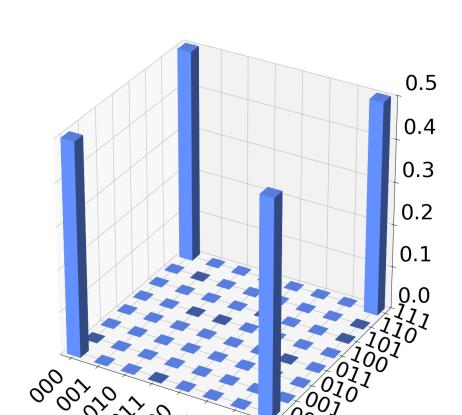






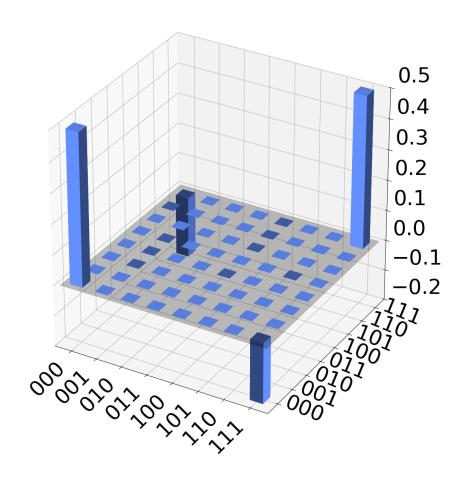
Ideal

Real Amplitude (ρ)



Phase Flip (p=0.7)

Real Amplitude (ρ)





Depolarizing Error

- Due to imperfect gates
- A unitary U is correctly applied with probability 1-p and followed by equally probable Pauli error after the gate with probability p

•
$$\mathcal{E}(\rho) = (1-p)U\rho U^{\dagger} + \frac{p}{P_n} \sum_{P \in P_n} PU\rho U^{\dagger} P^{\dagger}$$

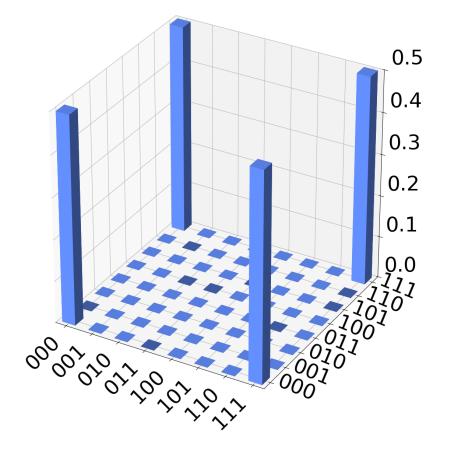
Consider the depolarization error on a single qubit

•
$$\mathcal{E}(\rho) = (1 - p)U\rho U^{\dagger} + \frac{p}{3}(XU\rho U^{\dagger}X^{\dagger} + YU\rho U^{\dagger}Y^{\dagger} + ZU\rho U^{\dagger}Z^{\dagger})$$



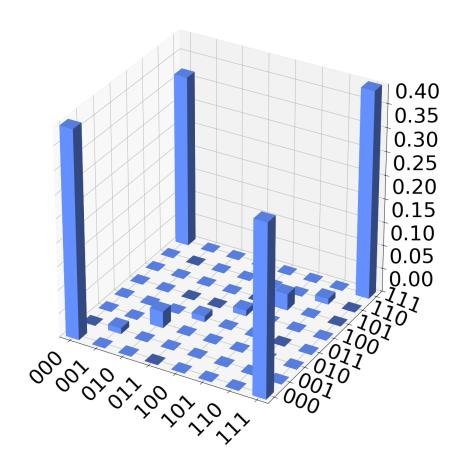
Ideal

Real Amplitude (ρ)



Depolarization p = 0.1

Real Amplitude (ρ)





Amplitude Damping

• Energy loss of the excited state of a qubit in $|1\rangle$ to the ground state energy $|0\rangle$.

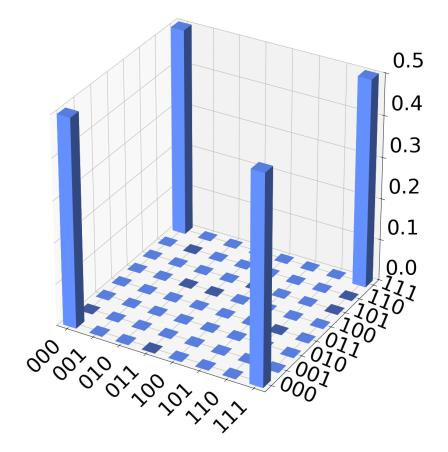
•
$$E_0 = \begin{bmatrix} 1 & 0 \\ 0 & \sqrt{1-\gamma} \end{bmatrix}$$
, $E_1 = \begin{bmatrix} 0 & \sqrt{\gamma} \\ 0 & 0 \end{bmatrix}$

•
$$\mathcal{E}_{AD}(\rho) = E_0 \rho E_0^{\dagger} + E_1 \rho E_1^{\dagger}$$



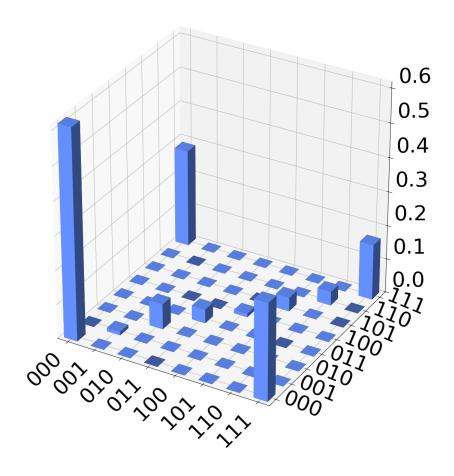
Ideal

Real Amplitude (ρ)



Amplitude Damping p = 0.2

Real Amplitude (ρ)





Phase Damping

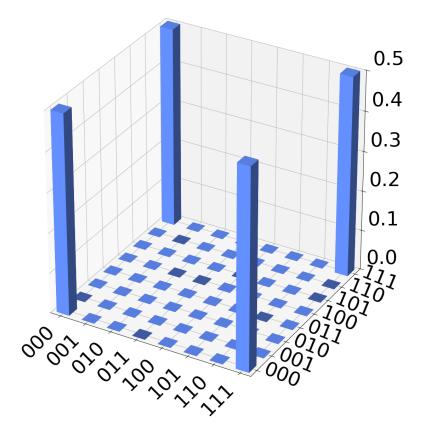
- Decay of phase due to interactions with the environment.
- Pure state is transformed to mixed state.

•
$$E_0 = \begin{bmatrix} 1 & 0 \\ 0 & \sqrt{1-\gamma} \end{bmatrix}$$
, $E_1 = \begin{bmatrix} 0 & 0 \\ 0 & \sqrt{\gamma} \end{bmatrix}$

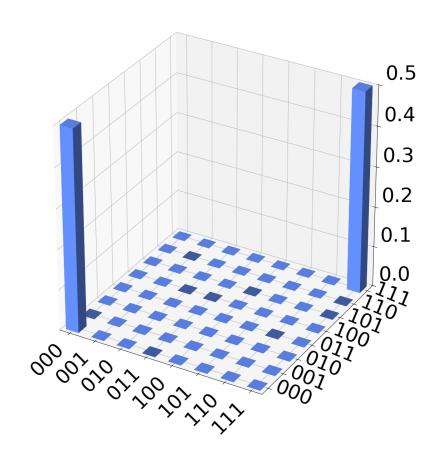
•
$$\mathcal{E}_{PD}(\rho) = E_0 \rho E_0^{\dagger} + E_1 \rho E_1^{\dagger}$$



Ideal Real Amplitude (ρ)



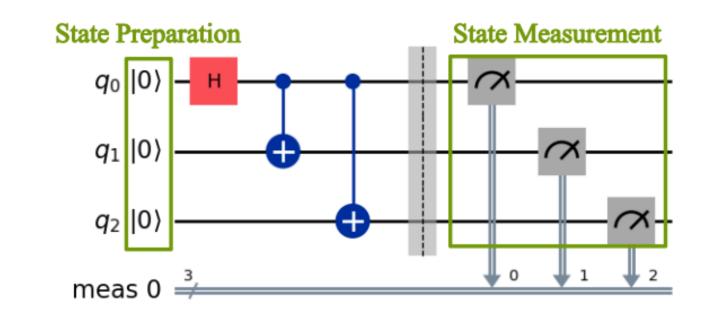
Phase Damping ($\gamma = 1$) Real Amplitude (ρ)





State Preparation and Measurement Errors (SPAM Errors)

- Errors in state initialization or preparation
 - Initialized $|0\rangle$ but ended up in $|1\rangle$
- Readout Errors in measurement
 - Should measure |1>
 but measured |0>
 instead.





Noise Characterization



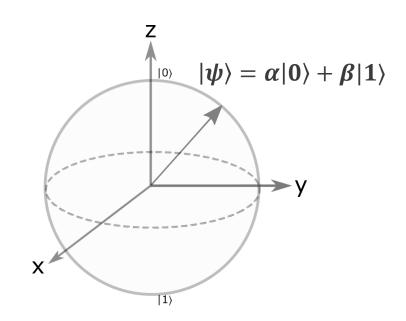
Simple Approach

- Multiply all success rates of gates in the circuit along with qubits' readout errors.
 - $P_{success} = \prod_{g \in G} (1 \epsilon_g) \cdot \prod_{q \in Q} (1 \epsilon_q)$
 - ϵ represents error, g represents gates and q represents qubits.



State Tomography

- a method to perform a series of projections to a circuit to reconstruct its state.
- For a single qubit state reconstruction, project on x, y and z axes.
- For a two-qubit state, project on xx, xy, xz, ..., zz axes.
- For a n-qubit state, 3^n projections will be required.





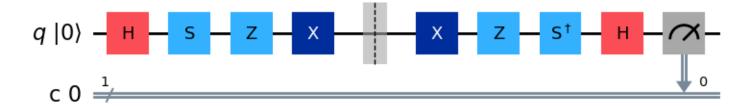
State Reconstruction steps

- 1.Generate a circuit
- 2. Calculate the ideal output state vector of this circuit
- 3. Choose the type of projectors (e.g., Pauli)
- 4. For a given circuit, the tomographic measurements are appended at the end of the replicas of the same circuit resulting in 3^n circuits.
- 5. Repeat the previous step multiple shots.
- 6.Reconstruct the density matrix ρ from these projective measurements using maximum likelihood technique.
- 7.To quantify the noise in the device, measure the fidelity between the ideal state vector and the reconstructed density matrix, $F(\rho, \sigma) = \langle \psi_{\rho} | \sigma | \psi_{\rho} \rangle$.



Randomized Benchmarking

- A scalable method to measure the Clifford gates fidelity
- Steps:
 - Initialize qubits to $|0\rangle$
 - Apply random sequences of Clifford gates of different lengths
 - For each random sequence, apply its inverse at the end of the circuit
 - Measure the circuit
- Ideally, the initial state $|0\rangle$ should be measured.
- Noise will cause the fidelity to decay as the sequence length increases.





Noise Mitigation



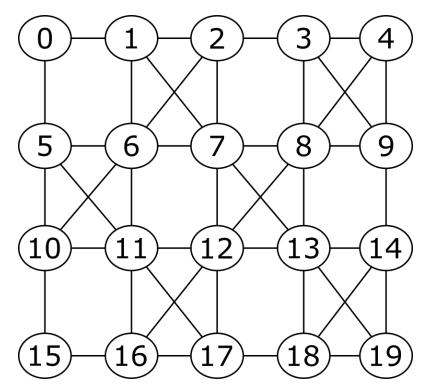
Approaches

- Mapping Optimization
- Circuit Optimization
- Quantum Error Correction Codes
- Fault Tolerant Quantum Computing



Mapping Optimization

- Assign logical qubits to physical ones based on
 - Error rates
 - Hardware-Connectivity (reduce extra swap gates!!!)



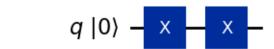
Hardware-Connectivity



Circuit Optimization

Gates Cancellation

$$UU^{\dagger} = U^{\dagger}U = I$$



Gates Replacement

$$q_0$$
 H $=$ q_0 q_1 H $=$ q_1

• Gates Combination (ex. Combine 2 consecutive RX gates)

$$RX(a)RX(b) = RX(a+b)$$
 $q - \frac{R_X}{\pi/2} - \frac{R_X}{\pi} - \equiv q - \frac{R_X}{3\pi/2} - \frac{R_X}{\pi}$



Quantum Error Correction



Outline

- Repetition Codes
 - Bit-Flip code
 - Phase-Flip code
 - Or Both (Shor !!)
- Steane Code



Classical Repetition codes

- Assume we want to send classical 0 over a channel
 - Send three consecutive 0s
 - Decode based on "Majority Vote"
 - Success: $(000,001,010,100) \rightarrow 0$
 - Failure if 2 or more bits flipped $(011,101,110,111) \rightarrow 1$
- Similarly, if we want to send classical 1

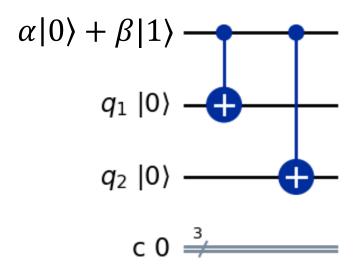


Three-qubit bit-flip code

- Encode each qubit using three qubits:
 - $|0\rangle \rightarrow |000\rangle$
 - $|1\rangle \rightarrow |111\rangle$
- Can correct at most 1 error
- Error can be detected using parity checks.



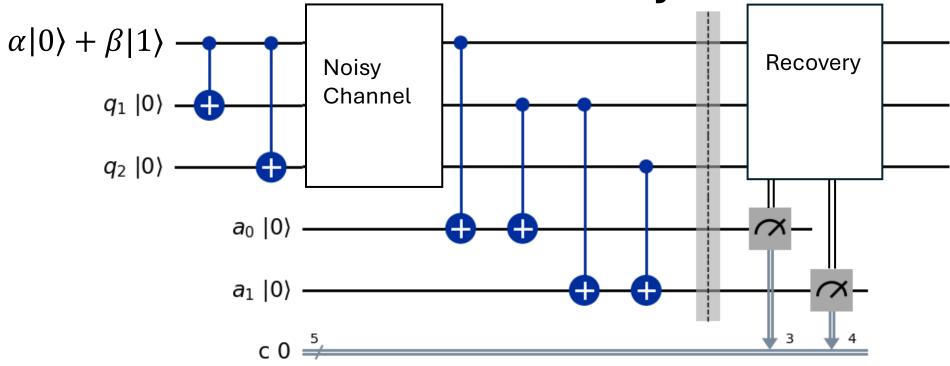
Encoding Circuit



$$(\alpha|0\rangle + \beta|1\rangle)|00\rangle \rightarrow \alpha|000\rangle + \beta|111\rangle$$



Error Detection and Recovery



$ \psi angle$ After Noisy Channel	M_{a_0}	M_{a_1}	Recovery	$\ket{\psi}$ After Recovery
$\alpha 000\rangle + \beta 111\rangle$	0	0	$I \otimes I \otimes I$	$\alpha 000\rangle + \beta 111\rangle$
$\alpha 001\rangle + \beta 110\rangle$	0	1	$I \otimes I \otimes X$	$\alpha 000\rangle + \beta 111\rangle$
$\alpha 100\rangle + \beta 011\rangle$	1	0	$X \otimes I \otimes I$	$\alpha 000\rangle + \beta 111\rangle$
$\alpha 010\rangle + \beta 101\rangle$	1	1	$I \otimes X \otimes I$	$\alpha 000\rangle+\beta 111\rangle$

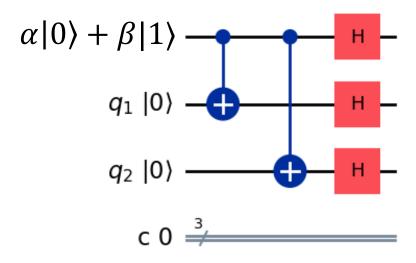


Three-qubit phase-flip code

- Encode each qubit using three qubits:
 - $|0\rangle \rightarrow |+++\rangle$
 - $|1\rangle \rightarrow |---\rangle$
- Can correct at most 1 error
- Error can be detected using parity checks.



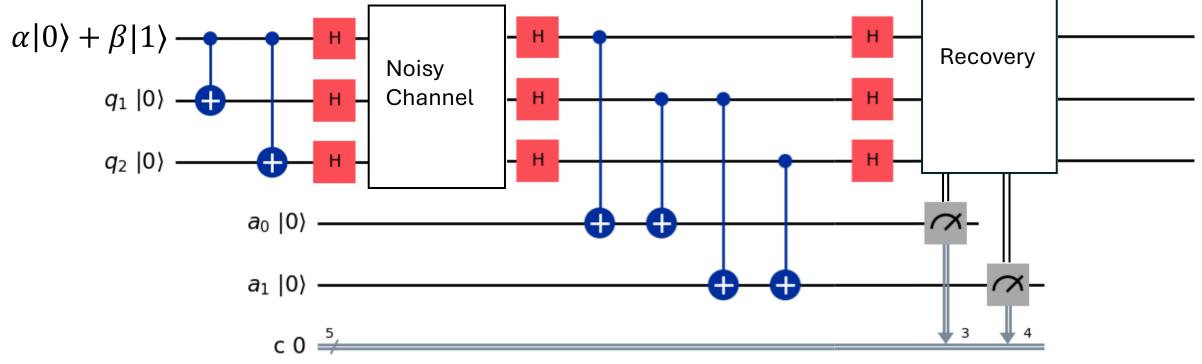
Encoding Circuit



$$(\alpha|0\rangle + \beta|1\rangle)|00\rangle \rightarrow \alpha|+++\rangle + \beta|---\rangle$$



Error Detection and Recovery

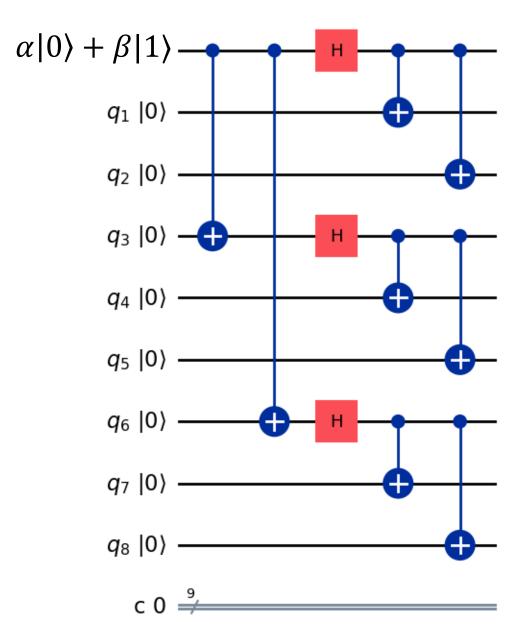


$ \psi angle$ After Noisy Channel	M_{a_0}	M_{a_1}	Recovery	$\ket{\psi}$ After Recovery
$\alpha +++\rangle+\beta \rangle$	0	0	$I \otimes I \otimes I$	$\alpha +++\rangle+\beta \rangle$
$\alpha ++-\rangle+\beta +\rangle$	0	1	$I \otimes I \otimes Z$	$\alpha +++\rangle+\beta \rangle$
$\alpha -++\rangle+\beta +\rangle$	1	0	$Z \otimes I \otimes I$	$\alpha +++\rangle+\beta \rangle$
$\alpha +-+\rangle+\beta -+-\rangle$	1	1	$I \otimes Z \otimes I$	$\alpha +++\rangle+\beta \rangle$



Shor Code

- Introduced by Peter Shor in 1995.
- Combines both bit-flip and phase-flip repetition codes
- Each logical qubit is encoded to 9 physical qubits





Steane Code

- Encodes a logical qubit using 7 qubits (more compact than Shor!!)
- Based on the hamming code (7,4), where a 4-bit data word is encoded to 7 bits.
- $|0_L\rangle = \frac{1}{\sqrt{8}}(|0000000\rangle + |1010101\rangle + |0110011\rangle + |1100110\rangle + |0001111\rangle + |1011010\rangle + |0111100\rangle + |1101001\rangle)$
- $|1_L\rangle = \frac{1}{\sqrt{8}}(|11111111\rangle + |0101010\rangle + |1001100\rangle + |0011001\rangle + |1110000\rangle + |0100101\rangle + |1000011\rangle + |0010110\rangle)$



Other Correction Codes

- Surface codes
- Color codes
- LDPC codes



Fault Tolerant Quantum Computing



Definition

• The ability to apply transformations to the quantum state even in the presence of noise.



Fault Tolerance Properties

- The error propagates to at most one qubit in each block
- To achieve full fault tolerance, each element in the circuit should be fault tolerant (state preparation, gates, error correction, measurement)
- The physical error rate should be lower than the code threshold in order not to introduce additional noise

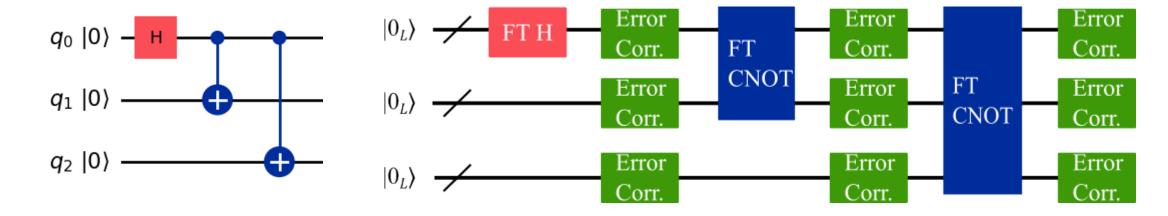


Steps to achieve fault tolerance

- Choose an error correction code
- Transform original circuit into its encoded version
- Apply gates in the original circuit transversally or use magic state injection
- Correct errors on regular intervals



Fault Tolerant Circuit Transformation



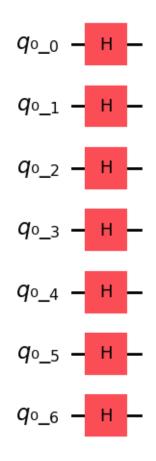
Original Circuit

Fault Tolerant Circuit



Implementing H Gate

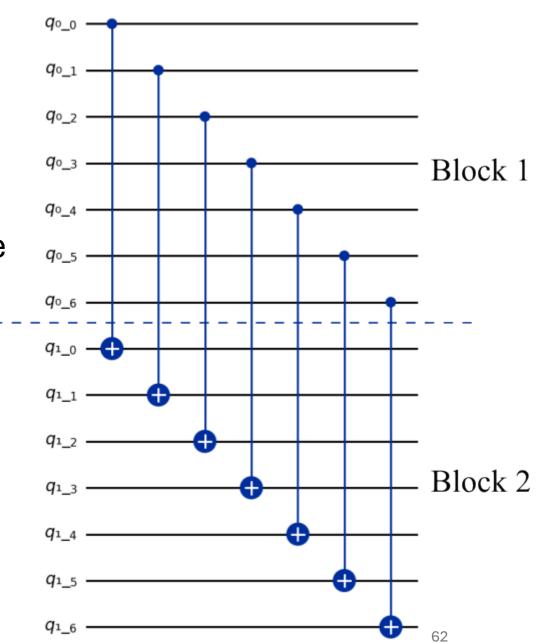
- Hadamard gate (H) can be implemented transversally on each qubit of the 7-qubit Steane encoding.
- If an error occurred, it won't propagate to the rest of qubits.





Implementing CNOT Gate

- CNOT gates can also be implemented transversally between a qubit in the first block with its corresponding in the second block
- Error will at most propagate to one qubit from each block





Other Gates

- Not all gates can be implemented transversally (example T Gate in Steane code) by any error correction code according to Eastin-Knill theorem.
- Use gate teleportation (magic state injection)



Future

• IBM Quantum Starling, a large-scale fault tolerant quantum computer, will be introduced by 2029 capable of executing 100 million gates on 200 logical qubits.



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Thank you