

Entang entropy

✓ \rightarrow Schmidt (see how Schmidt get S)

Renyi

✓ \rightarrow Entanglement entropy (estimate S_2 using random mes)

Other measure

\rightarrow concurrence - negativity

PPT crit

\rightarrow separability

Sharing of entanglement

\rightarrow Monogamy

Propagati: of entanglement

\rightarrow Area law

Testing existence of Q:

\rightarrow Bell tests.

Entanglement witness

\rightarrow Ent-witness

LOCC cannot create entanglement

\rightarrow LOCC.

Remarkable entangled state

\rightarrow GHz-W

Different pattern in entanglement:

\rightarrow entanglement structure

\rightarrow use: \rightarrow stabilizer

Regenerating entanglement:

\rightarrow purification.

Teleportati: w/ noise.

\rightarrow Bell

Choi channel

\rightarrow Choi

Noise model

\rightarrow Noise - entanglement

Competitor of QC

\rightarrow tensor-network

| Notion of entanglement theory |

0. Reminder

ρ is density matrix of state

x equivalent repres. of isolated sys. w/ state vector

x said to be the only valid one for open system

. Properties:

x $\rho^\dagger = \rho$

x $\text{Tr}(\rho) = 1$

x $\text{Tr}(\rho^2) \leq 1$; $= 1$ iff pure state

iff $\exists |\psi\rangle$ s.t. $\rho = |\psi\rangle\langle\psi|$

x definite positive

x diag are probabilities

. Remarkable ρ :

• pure

• mixed

• thermal mix

• classical mix

I Evaluating entanglement

1) Motivation: notion of locality

Superpositⁿ vs entanglement, a notion of locality

$$|\phi^+\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

$L_0 \equiv |0\rangle \quad L_0 \equiv |1\rangle$

$$= \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \rightarrow \text{only superposition?}$$

\Rightarrow no because of locality / separability:

$$\mathcal{H} = \mathcal{H}_1 \otimes \mathcal{H}_2$$

\uparrow
spatially
at 1 localⁿ

\uparrow spatially in
another

(superposition
within)



entanglement without

We say "qubit 1 entangled w/ qubit 2"

How to evaluate: Von Neumann entanglement
entropy.

2] Partial trace: an example

ex: $\rho_{AB} = |\phi\rangle\langle\phi|$

$$= \frac{1}{2} (|00\rangle\langle 00| + |00\rangle\langle 11| + |11\rangle\langle 00| + |11\rangle\langle 11|)$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{pmatrix} \begin{matrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{matrix}$$

$T_{A_2} \downarrow$

Basis of \mathcal{H}_2

$$\langle 0_2 | \rho_{AB} | 0_2 \rangle + \langle 1_2 | \rho_{AB} | 1_2 \rangle$$

$$= \frac{1}{2} \left(\begin{aligned} &1 \otimes \langle 0 | 00 \langle X 00 | 1 \otimes | 0 \rangle \\ &+ 1 \otimes \langle 0 | 00 \langle X 11 | 1 \otimes | 0 \rangle \\ &+ 1 \otimes \langle 0 | 11 \langle X 00 | 1 \otimes | 0 \rangle \\ &+ 1 \otimes \langle 0 | 11 \langle X 11 | 1 \otimes | 0 \rangle \\ &+ 1 \otimes \langle 1 | 00 \langle X 00 | 1 \otimes | 1 \rangle \\ &+ 1 \otimes \langle 1 | 00 \langle X 11 | 1 \otimes | 1 \rangle \\ &+ 1 \otimes \langle 1 | 11 \langle X 00 | 1 \otimes | 1 \rangle \\ &+ 1 \otimes \langle 1 | 11 \langle X 11 | 1 \otimes | 1 \rangle \end{aligned} \right)$$

$$= \frac{1}{2} \left(\begin{aligned} &|0\rangle\langle 0| \\ &+ 0 \\ &+ 0 \\ &+ 0 \\ &+ 0 \\ &+ 0 \\ &+ 0 \\ &+ 0 \\ &+ |1\rangle\langle 1| \end{aligned} \right)$$

$$= \frac{1}{2} (|0\rangle\langle 0| + |1\rangle\langle 1|)$$

$$= \frac{1}{2} \begin{pmatrix} 1 & \\ & 1 \end{pmatrix} = \rho_1$$

↑
partial trace

↑
reduced density mat

3] Partial trace : general

* In general : $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$

$\uparrow \qquad \qquad \uparrow$
 $|A_i\rangle \qquad |B_j\rangle$

$$\rho_A = \text{Tr}_B(\rho) = \sum_j \mathbb{1} \otimes \langle B_j | \rho (\mathbb{1} \otimes |B_j\rangle)$$

$$\rho_B = \text{Tr}_A(\rho) = \sum_i \langle A_i | \otimes \mathbb{1} \rho | A_i \rangle \otimes \mathbb{1}$$

... we can partial trace more.

eg: $\rho_A = \rho_A^+$; $\text{Tr} \rho_A = 1$; etc...

* Remarks :

> If initially pure, ρ_{reduced} can be mixed.
x if in " , is mixed A & B are entangled.

x if initially pure, ρ_{reduced} pure $\Rightarrow A$ & B not entangled \Rightarrow the state is separable and

$$\rho = \rho_A \otimes \rho_B \quad (\text{prove it})$$

x if initially mixed, ρ_A pure $\Rightarrow A$ separable & only B is mixed (not interesting)

x if mixed mixed $\Rightarrow A ? B$
 \Rightarrow need a measure \Rightarrow VN ent. ent.

4) Bipartite Von Neumann entanglement entropy

A) Formula

Von Neumann entang. entropy (bipartite)

$$\begin{aligned} S_{A/B} &= -\text{Tr}_A (\rho_A \log \rho_A) \\ &= -\text{Tr}_B (\rho_B \log \rho_B) \end{aligned}$$

Meaning of $\log \rho_A$: $\rho_A = \rho_A^\dagger \Rightarrow \text{diag}$

$$\begin{aligned} \text{i.e. } \log(\rho_A) &= \log(U D U^\dagger) \\ &= U \log(D) U^\dagger \\ &= U \log \begin{pmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{pmatrix} U^\dagger \\ &= U \begin{pmatrix} \log \lambda_1 & & \\ & \ddots & \\ & & \log \lambda_n \end{pmatrix} U^\dagger \end{aligned}$$

What if $\lambda_i = 0$?

$$\Rightarrow \text{Tr}(\rho_A \log \rho_A) = \text{Tr} \left(U \begin{pmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{pmatrix} U^\dagger U \begin{pmatrix} \log \lambda_1 & & \\ & \ddots & \\ & & \log \lambda_n \end{pmatrix} U^\dagger \right)$$

abuse of notation for explanation

$$= \text{Tr} \left(\begin{pmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{pmatrix} \begin{pmatrix} \log \lambda_1 & & \\ & \ddots & \\ & & \log \lambda_n \end{pmatrix} U^\dagger U \right)$$

cyclic rotation in the trace

$$= \text{Tr} \begin{pmatrix} \lambda_1 \log \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \log \lambda_n \end{pmatrix}$$

$$= \sum_i \lambda_i \log \lambda_i$$

if $\lambda_i = 0$, could be infinite

$$= \sum_{i | \lambda_i \neq 0} \lambda_i \log \lambda_i$$

$1 \geq \lambda_i \geq 0$ by def.

B) A measure of entanglement

$S_{A/B}$ is a measure of entanglement i.e.

x non-negativity: $S_{A/B} \geq 0$

x $S_{A/B} = 0$ iff $\rho = \rho_A \otimes \rho_B$ (ρ_A, ρ_B pure)

x Monotonically decrease under LOCC (local operation and classical communication)

↳ i.e. operation that, well, do not increase entanglement.

i.e. operation that you can do on A or

B ~~or~~ the operation where you need both A & B to do.

i.e. separable operation

i.e. operation to ~~cannot~~ turn a separable state into a non separable one.

→ see <<2_locc.ipynb>>, but not the 1st one to check

C| Examples

$$\begin{aligned} \text{ex } S_1(\phi^+) &= -\hat{h} \left(\frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \log \left(\frac{1}{2} \begin{pmatrix} 1 & \\ & 1 \end{pmatrix} \right) \right) \\ &= \hat{h} \begin{pmatrix} \frac{1}{2} \log 2 & \\ & \frac{1}{2} \log 2 \end{pmatrix} \\ &= \log 2 : \text{ the max that can be} \\ &\quad \text{obtained w/ 2 qubit} \end{aligned}$$

\Rightarrow Bell pairs are "maximally entangled states"
 $S \rightarrow$ not measurable

D] Schmidt decomposition & SVD

* Schmidt decomp & max entang. (cf Schmidt-decomp)

If $\dim \mathcal{H}_A = N_A$; $\dim \mathcal{H}_B = N_B$, then

if $|\Psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \Rightarrow \exists \alpha_i, i \leq \min\{N_A, N_B\}$

s.t. $|\Psi\rangle = \sum_i \alpha_i |\phi_i\rangle \otimes |\varphi_i\rangle$ where

$\{|\phi_i\rangle\}_{i \leq N_A}$ basis of \mathcal{H}_A ; $\{|\varphi_i\rangle\}_{i \leq N_B}$ basis of \mathcal{H}_B

$$\Rightarrow S_{A/B} \leq \log(\min\{N_A, N_B\})$$

i.e. a qubit entangled w/ the universe...

only shares a max of $\log 2$ entanglement w/ it.

* there are several algo to get the α_i , but the singular value decomposition is a classic:

i.e. if Π matrix (not necessarily square):

$\exists U, V$ unitary s.t.

$$\Pi = U D V^{-1}$$

D diagonal not square.

why? $|\phi_i\rangle \otimes |\varphi_i\rangle \rightarrow |\phi_i\rangle \langle \varphi_i|$ in the Schmidt decomposition.

\rightarrow see << Schmidt decomposition >>

5] Renyi-entropies

$$\text{Renyi entropy : } S_{\alpha}(p) = \frac{1}{1-\alpha} \log(p_A^{\alpha}) \quad \alpha \geq 0$$

estimable \nearrow
 $\alpha \rightarrow 0 \rightarrow$ capture all non-zero terms & put it to 1.

$$\begin{aligned} p_A^{\alpha} &= e^{\alpha \log p_A} = U e^{\alpha \log(D)} U^{-1} \\ &= U \prod_i e^{\alpha \log \lambda_i} U^{-1} \\ &= U \left(\begin{matrix} \lambda_i^{\alpha} & & \\ & \ddots & \\ & & 0 \end{matrix} \right) U^{-1} \\ &= U \frac{1}{N} \left(\begin{matrix} 1 & & \\ & \ddots & \\ & & 0 \end{matrix} \right) U^{-1} \end{aligned} \quad \left. \begin{array}{l} \text{quick \& } \\ \text{dirty} \end{array} \right\}$$

$\lambda_i \neq 0$

$\alpha \rightarrow 0$

$\alpha \rightarrow 0$: "rich get richer"

$\alpha \rightarrow 1$: coincide w/ VNEE.

\rightarrow see << 3- Renyi-entropy >> for estimation

6) Other measures

A) Concurrence (bipartite)

$$C_A(\rho) = \sqrt{2(1 - T_2(\rho_A^2))}$$

For 2 qubits: $C(\rho) \equiv \max(0, \lambda_1 - \lambda_2 - \lambda_3 - \lambda_4)$

where $\lambda_1 > \lambda_2 > \lambda_3 > \lambda_4$ v.p. of $R = \sqrt{\rho_A \tilde{\rho} \rho_A}$; $\tilde{\rho} = (\sigma_y \otimes \sigma_y) \rho^* (\sigma_y \otimes \sigma_y)$
in \mathbb{Z} basis

B) Negativity

$$\mathcal{N}(\rho) = \frac{\|e^{\tau_A}\|_1 - 1}{2}$$
$$= \sum_i \frac{|\lambda_i| - \lambda_i}{2}$$

e^{τ_A} : partial transpose
Not a density mat; can have $\lambda < 0$

$$\|X\|_1 = T_2(|X|) = T_2 \sqrt{X^\dagger X}$$

→ see <4 - Concurrence negativity>

C) Log negativity

$$E_N(\rho) = \log_2 \|e^{\tau_A}\|_1 = \log_2 (2\mathcal{N} + 1)$$

D) Mutual information (tripartite)

* Not an entanglement monotone

$$I(A:B) = S(\rho^A) + S(\rho^B) - S(\rho^{AB})$$
$$= S(\rho^{AB} \| \rho^A \otimes \rho^B)$$

III Estimating & evaluating entanglement

1) Confrontation with hidden variables

ex: teleportation; is there really no Δ^{ab} that we have no access to that says in advance that Alice's envelope was not meant to be red all along, & Bob's envelope blue.

→ can build a mathematical theory out of it

→ Has \neq predictⁿ than QM

→ QM was tested right (Nobel)

2) CHSH & Bell inequalities (a stat. constraint)

Clauser - Horne - Shimony - Holt

$$|S| \leq 2 \quad \text{if hidden } \Delta^{ab} \quad (\leq 2\sqrt{2} \text{ if Q.})$$

$$S = E(a,b) - E(a,b') + E(a',b) + E(a',b')$$

In QM of course

$$\left\{ \begin{array}{l} E(a,b) = \langle \psi_{a,b} | Z \otimes Z | \psi_{a,b} \rangle \\ |\psi_{a,b}\rangle = \underbrace{A(a) \otimes B(b)}_{\downarrow \text{define a precise state of the device}} | \psi_{\text{ref}} \rangle \end{array} \right. \quad \hookrightarrow \text{e.g. } |00\rangle$$

$$\Rightarrow E(a,b) = \frac{N_{++} - N_{+-} - N_{-+} + N_{--}}{N_{++} + N_{+-} + N_{-+} + N_{--}}$$

$\hookrightarrow a, b$ are the "hidden Δ^{ab} "

See << 5 - Bell_test >>

3] Entanglement witnesses

An entanglement witness W is an operator / $E_{\mathcal{H}}^{\text{al}}$ that is \oplus on all separable (biseparable) states.

If strictly negative, it is an entangled state.

↑ not iff.

→ see << 6 - entanglement witness >>

4] PPT criteria (an entang. witness)

ex PPT criteria

requirements: partial trace:

$$\rho = \sum \alpha_{ijkl} |iXj\rangle \otimes |kXl\rangle$$

↘ T_B : partial trace

$$\rho^{T_B} = \sum \alpha_{ijkl} |iXj\rangle \otimes |lXk\rangle$$

$$= \sum \alpha_{ijlk} |iXj\rangle \otimes |kXl\rangle$$

PPT: if ρ separable, n.p. of $\rho^{T_B} \in \mathbb{R}^+$

hence if $\exists \lambda_i < 0$ n.p. of ρ^{T_B} , then ρ entangle

note: n.p. of $\rho^{T_B} \in \mathbb{R}^+ \nRightarrow \rho$ separable

ρ entangles $\nRightarrow \exists \lambda_i < 0$ n.p. of ρ^{T_B}

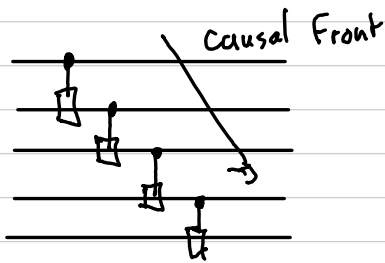
→ see << 7 - separability >>

IV Entanglement pattern

1) Propagation of entanglement

Typically, through 2-by-2 local interaction.

2 patterns:



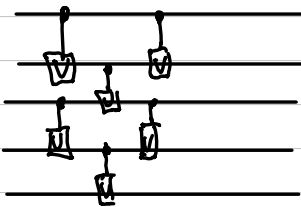
After a single front

1 entangled w/ N , but

1 is not modified by N

past state (while N is)

→ need a return



Local constant interact:

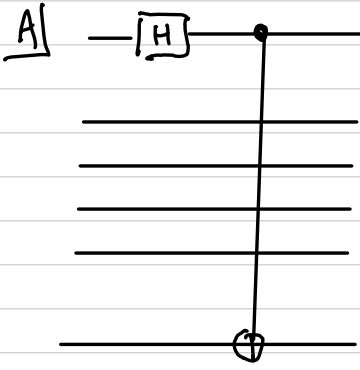
From the second layer onwards,

Vanishing entanglement b/w

1 & N . But no causal link.

→ see << 8-Area-law >>

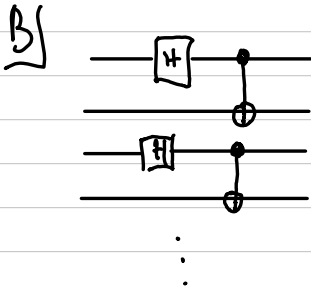
2) Range of entanglements



a long range

any S convex locally $\neq 0$

but actually not a lot of entq.
(separable, albeit not completely)



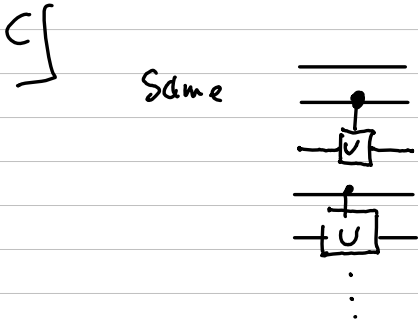
} 2-by-2 entang.



Some $S = 0$, some only as much
as above.

↪ short-range

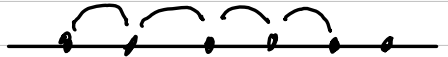
(separable, albeit not completely)



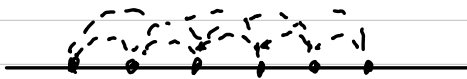
→ not separable, $S \neq 0$ ever.

but not all entangled state
are possible

e.g. the 1st qubit is
vanishingly entangled w/ the last



D) Possible to entangle at longer range, but at the
cost of shorter range:



↪ see 9-entanglement-structure

10-GHz-W

3] Entanglement distillation & purification

A & B bought $|\phi^+\rangle$ at Charlie, not Charlie!

They have ^(copies of)

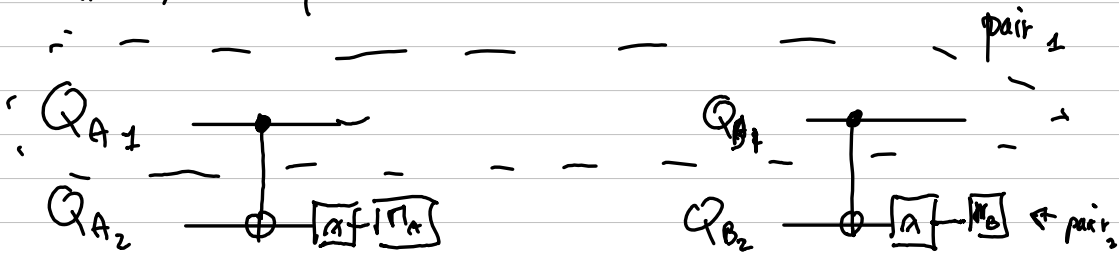
$$\rho = F |\phi^+\rangle\langle\phi^+| + \frac{1-F}{3} (|\phi^-\rangle\langle\phi^-| + |\psi^+\rangle\langle\psi^+| + |\psi^-\rangle\langle\psi^-|)$$
w/ $F > 0,5$

How can they do better, knowing that they cannot meet? (Ideally, they want $|\phi^+\rangle\langle\phi^+|$ only)

→ BDPSSW protocols.

Idea: sacrificing copies (a. distillation) to increase the purity of the other (purification)

Ex w/ 2 copie



Alice calls Bob. If $M_A = M_B \rightarrow$ they keep pair 1. Otherwise, they throw it.

The pairs left look like:

$$\rho' = F' |\phi^+\rangle\langle\phi^+| + \frac{1-F'}{3} (\dots)$$

w/ $|F'| > F$.

→ see < 12 - Purification >

II Conclusions

- x Entanglement distinguishes itself from superposition in terms of locality
- x When two part of a system are not entangled, we say that they are separable / biseparable
- x Entanglement, however, is an abstract q'ty difficult to capture. It can be theoretically measured using an entanglement measure ... but only few of them can be estimated w/ actual system measurements.
- x More often, we use entanglement witnesses or entang. criteria, that are + accessible, but less informative.
- x Using measure or witnesses, it is possible to have some understanding of patterns of entanglement existing in a state. In Qc, these patterns are directly related to the shape of the circuit when the depth is small.
- x Entanglement leads to non-classical correlations. Not causality!
- x «Recuperating» info from entanglement can be called entanglement distillatⁿ. It is costly in terms of states.
- x Entanglement is exclusively initiated locally
- x ... but can be enhanced at a distance w/ purificatⁿ.