

Algebra - MATH310

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Preface

Helo! I'm Jack :3.

I'm a student that need sto type out courses in order to make sure they properly understand them. So I put them out into the world! They might help you more than they help me :D. They are given as they are, with no guarantee of quality but guarantee of goodwill, bla bla bla. You know the gist of it.

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Chapter 1

Introduction

Week 1

Algebra rests on 3 basic principles, which are equivalent in nature.

1. **Induction:** Let $S \subset \mathbb{N}$ such that $0 \in S$ and $n \in S \Rightarrow n + 1 \in S$. Then, $S = \mathbb{N}$.
2. **Well-ordering principle:** For any non-empty $A \subset \mathbb{N}$, there exists an element $a : \forall b \in A, a \leq b$.
3. **Strong induction:** Let $S \subset \mathbb{N}$ such that $0 \in S$ and $\{0, \dots, n\} \in S \Rightarrow n + 1 \in S$. Then, $S = \mathbb{N}$.

It is well-established that these three principles are equivalent. Let us prove it.

Theorem 1.1. $I \Rightarrow WOP \Rightarrow SI \Rightarrow I$.

Proof. We will prove each induction separately.

1. $1. \Rightarrow 3.$ Let S be the construction from the strong induction definition, and let us consider $P(n) = \{0, 1, \dots, n\} \subset S$. We can prove it by induction:

Base: $0 \in S$ by construction $\Rightarrow \{0\} \subset S$.

Induction: Let us prove that $P(k) \Rightarrow P(k + 1)$ for some k .

$$\begin{aligned}
 \{0, 1, \dots, k\} \subset S & \text{ [by IH]} \Rightarrow k \in S & \text{ [by construction]} \\
 & \Rightarrow k + 1 \in S & \text{ [by definition]} \\
 & \Rightarrow \{0, 1, \dots, k, k + 1\} \in S
 \end{aligned}$$

Since it is hereditary and true for 0, it is true $\forall n \in \mathbb{N}$ by the induction principle.

Since $\{0, 1, \dots, n\} \subset \mathbb{N} \forall n$, then $S = \mathbb{N}$.

2. $2 \Rightarrow 1.$ Suppose $S \subset \mathbb{N}$ such that $0 \in S$ and $n \in S \Rightarrow n + 1 \in S$. Consider $S' = \mathbb{N} \setminus S$, which we assume to be nonempty by absurd. By the well-ordering principle, we can pick a least element in $k \in S'$, which is by definition not in S . k cannot be zero, since $0 \in S$ by definition, but it can also not be non-zero, since $k \neq 0 \Rightarrow k = m + 1$ for some $m < k$ (therefore not in S'). $m \in S$, so by construction, $m + 1 = k \in S$ as well, which is a contradiction. S' has to be empty, so $S = \mathbb{N}$.
3. $3 \Rightarrow 2.$ Done in a Problem Set, found in appendix A.

□

Chapter 2

Primes

2.1 Divisors and primes

Definition 2.1. Let $a, b \in \mathbb{Z}$. We say that a *divides* b (notate: $a|b$) if there exists $k \in \mathbb{Z}$ such that $b = ka$.

Definition 2.2. A number $p \in \mathbb{Z}$ is prime if $p > 1$ and the only numbers that divide it are itself and 1.

Theorem 2.3. Any $n > 1$ has a prime divisor.

Proof. Let $S = \{n \in \mathbb{N} : n > 1 \wedge n \text{ has no prime divisors}\}$. We suppose S to be nonempty, meaning it contains a least element $k \in S$. k cannot be prime, since $k|k \forall k$. Therefore, it has to be true that $k = ab$ for $a, b < k \in \mathbb{N}$. Since k was the least element, then, $a \notin S$, meaning that there exists a prime p such that $a = pt$ for $t \in \mathbb{N}$. Therefore, $k = ab = ptb \Rightarrow p|k$, contradicting our construction of S . Therefore, S must be empty. \square

Theorem 2.4. Any $n > 1$ can be expressed by the product of primes.

This proof was done in an exercise set, and can be found in the appendix.

Theorem 2.5. The prime number factorization of a number is unique.

Proof. Let $k = \prod_i^n p_i = \prod_j^m q_j$ two distinct prime sets. Suppose without loss of generality that $q_1 > p_1$ and let $t = (q_1 - p_1)q_2 \dots q_m > 0$. Then:

$$\begin{aligned} t &= (q_1 - p_1)q_2 \dots q_m \\ &= q_1 q_2 \dots q_m - p_1 q_2 \dots q_m \\ &= k - p_1 q_2 \dots q_m > 0 \Rightarrow p_1 | t \end{aligned}$$

We know that $p_1 \neq q_j$ for all j , so we focus on the only “weird” term:

$$\begin{aligned} (q_1 - p_1) &= sp_1 \\ \Rightarrow q_1 &= (s + 1)p_1 \end{aligned}$$

Which is a contradiction because q_1 is supposed to be prime. Therefore, the prime factorization is unique. \square

2.2 Integer arithmetic

Definition 2.6 (Euclidian division). *Let $n \in \mathbb{Z}, d \in \mathbb{Z}^*$. There exists a unique pair $q, r \in \mathbb{Z}$ such that $n = qd + r$ with $0 < r < d$.*

Proof. **Existence.** Consider the set of all numbers $S = \{n - kd\}_{k \in \mathbb{Z}} \cap \mathbb{N} = \{n - kd, kd \leq n\}_{k \in \mathbb{Z}}$.

We know that S is not empty, because:

- ▷ if $n \geq 0$, then we set $k = 0$, meaning $n \in S$
- ▷ if $n < 0$, then we set $k = |n| + 1$, meaning $kd > |n|$ and $n + kd \in S$.

Since it's never empty, we can pick the least element of S by means of the well-ordering principle. Let's call it r . Therefore, we have $r = n - kd$ for some k . To prove $r < d$, we assume towards absurdity that $r \geq d$, meaning that

$$n - (k + 1)d = n - kd - d = r - d \geq 0$$

meaning r wasn't minimal, which is a contradiction.

Uniqueness. Suppose $n = q_1d + r_1 = q_2d + r_2$. Without loss of generality, assume $q_1 > q_2$. Then:

$$(q_1 - q_2)d + r_1 = r_2 \geq d$$

Since r_1 and $q_1 - q_2$ are positive. This contradicts the definition of r_2 , and is therefore absurd. \square

Definition 2.7. *Let $a, b \in \mathbb{Z}$. We define the greatest common divisor (gcd) of two numbers as*

$$\gcd(a, b) = \max\{x \in \mathbb{Z} : x|a \wedge x|b\}$$

Theorem 2.8. *For $n, q \in \mathbb{Z}, d \in \mathbb{Z}^*$, such that $n = qd + r$, it is always the case that:*

$$\gcd(n, d) = \gcd(d, r)$$

Proof. By inspection of the relationship $n = qd + r$, it's clear that if $x|n \wedge x|d$ then $x|r$, and if $x|d \wedge x|r$ then $x|n$. \square

Method This induces a special algorithm to compute the gcd of two numbers! Let $d_1, d_2 \in \mathbb{Z}$. Then:

$$\begin{aligned} d_1 &= q_1d_2 + d_3 \\ d_2 &= q_2d_3 + d_4 \\ &\dots \\ d_k &= q_kd_{k+1} + 0 \end{aligned}$$

The relationship $\gcd(d_{i-1}, d_i) = \gcd(d_i, d_{i+1})$ holds down the tree, meaning that by the end

$$\gcd(d_1, d_2) = d_{k+1}$$

Additionally, we have:

Corollary 2.9. *For any $a, b \in \mathbb{Z}^+$, there exist $x, y \in \mathbb{Z}$ such that*

$$\gcd(a, b) = xa + yb$$

This is obtained by running Euclid “up the tree”.

Example 1. *TODO*

Special consequence of corollary 2.9 is the following

Corollary 2.10. *If $a, b \in \mathbb{Z}^+$ are such that $d = \gcd(a, b)$, then the equation:*

$$c = ax + by$$

has solutions (x, y) if and only if $\exists k > 0 : c = kd$, and they can be found as the solutions in corollary 2.9 multiplied by k .

Final consequence of these facts is the well-known Bézout’s theorem.

Theorem 2.11. *Two numbers $a, b \in \mathbb{Z}^+$ are relatively prime if and only if the equation*

$$1 = ax + by$$

has integer solutions.

Definition 2.12. *For any $n \in \mathbb{Z}^+$, Euler’s totient function is defined as:*

$$\varphi(n) = |\{k \in \{1, \dots, n\} : \gcd(k, n) = 1\}|$$

meaning the number of positive integers less than n that are coprime to it.

Properties Properties of the totient function include:

- ▷ $\varphi(p) = p - 1$ for any prime p .
- ▷ $\varphi(pq) = (p - 1)(q - 1)$ for any pair of distinct primes p, q .
- ▷ More generally, $\varphi(mn) = \varphi(m)\varphi(n)$ for any m, n coprime.

Chapter 3

Groups

Week 2

3.1 Base definitions

Definition 3.1. A *group* is a set G with a binary operation $\cdot : G \times G \rightarrow G$, satisfying the following axioms:

- ▷ \cdot is *associative*: $(a \cdot b) \cdot c = a \cdot (b \cdot c)$
- ▷ There exists a *neutral element* e such that $a \cdot e = e \cdot a = a \ \forall a \in G$.
- ▷ For any $a \in G$ there exists an *inverse* a^{-1} such that $a^{-1} \cdot a = a \cdot a^{-1} = e$.

We say that G is a *finite* group if $|G| < \infty$. In that case, we say that G is of *order* $|G|$. We say that G is *abelian* (or commutative) if $a \cdot b = b \cdot a \ \forall a, b \in G$.

Definition 3.2. $H \subset G$ is a *subgroup* if it contains the neutral element e_G and if it is closed with respect to \cdot_G , meaning that for every $a, b \in H$, $a \cdot b \in H$, and to inverses.

We can note that any group has a subgroup generated by a single element:

$$\langle g \rangle = \{e, g^1, g^2, \dots, g^{-1}, g^{-2}, \dots\}$$

Since $g^i \cdot g^j = g^{i+j}$ by definition of the group operation, this set is closed under it, meaning it is a subgroup.

Definition 3.3. If it exists, the minimal $n \in \mathbb{N}^*$ such that $g^n = e$ is called the *order* of g . It is finite for every element in a finite group.

3.2 Cosets

Definition 3.4. Let $H \subset G$ be a subgroup of G . The *left coset* of g with respect to H , denoted gH , is the following set:

$$gH = \{gh, h \in H\}$$

Theorem 3.5. Let $H \subset G$ finite. Then:

1. Two left-cosets xH, yH are either disjoint ($xH \cap yH = \emptyset$) or equal.

2. For any element $g \in G$ there exists a left coset of H such that $g \in H$.

3. $|xH| = |H| \forall x \in G$

Proof. We will prove each part separately:

1. Suppose xH, yH are such that $xH \cap yH \neq \emptyset$. This means that there exist h_1, h_2 such that $xh_1 = yh_2$. Therefore,

$$x = yh_2h_1^{-1} = yh_3 \in yH \Rightarrow xh = yh_3h \forall h \in H$$

This means that if there exists an element of xH that is in yH , then every element in xH can be written as an element in yH , meaning they are equal.

2. For any $g \in G$, one can construct $gH = \{e, g, g^2, \dots\}$, which naturally contains g .

3. The mapping

$$\begin{aligned} f(h) : H &\rightarrow xH \\ h &\mapsto xh \end{aligned}$$

is surjective, by definition of $xH = \{xh, h \in H\}$, and it is also injective, since $xh_1 = yh_2 \Leftrightarrow h_1 = h_2$. This means it defines a bijection between H and xH , indicating they have the same cardinality.

Example Let $G = (\mathbb{Z}, +, 0), H = 3\mathbb{Z} \subset \mathbb{Z}$. The left coset of 0 with respect to H is :

$$\{0 + 3k\}_{k \in \mathbb{Z}} = H = \{3 + 3k\}_{k \in \mathbb{Z}}$$

The left coset of 1 is

$$\{1 + 3k\}_{k \in \mathbb{Z}} = \{1, 4, 7, -2, \dots\}$$

□

Theorem 3.6 (Lagrange). *Let G be a finite group, $H \subset G$ a subgroup. Then, $|H|$ divides $|G|$.*

Proof. Each $g \in G$ belongs to a left coset of H , which are either disjoint or equal. This means:

$$\begin{aligned} G &= \bigcup_{i=0}^r x_i H && \text{[disjoint union of finitely many sets]} \\ \Rightarrow |G| &= \sum_{i=0}^r |x_i H| \\ \Rightarrow |G| &= \sum_{i=0}^r |H| && \text{[since } |xH| = |H| \text{]} \\ \Rightarrow |G| &= r|H| \end{aligned}$$

with $r \in \mathbb{N}$, meaning that $|H|$ divides $|G|$.

□

Definition 3.7. The number of left cosets of H of G is called the *index* of G :

$$[G : H] = |G|/|H| \in \mathbb{N}^*$$

This means that the order of any element $g \in G$ (notated $\text{ord}(g)$) divides the order of the group $|G|$, since every element generates a subgroup $\langle g \rangle$. Additionally, it implies

Corollary 3.8. $g^{|G|} = (g^{\text{ord}(g)})^k = e^k = e$ for some k .

3.3 RSA

Theorem 3.9 (Euler's theorem). Let $a, n \in \mathbb{Z}^+$ such that $\gcd(a, n) = 1$. Then,

$$a^{\varphi(n)} \equiv 1 \pmod{n}$$

Proof. Consider $G = (\mathbb{Z}/n\mathbb{Z}, \cdot, 1)$. Then,

$$a^{\varphi(n)} = a^{|G|} \stackrel{3.8}{=} 1$$

□

Theorem 3.10 (Fermat's little theorem). Let $a \in \mathbb{Z}^+$, p prime such that p does not divide a . Then, $a^{p-1} = 1$.

Proof. Consider $G = (\mathbb{Z}/p\mathbb{Z}, \cdot, 1)$. Then, $|G| = \varphi(p) = p - 1$. By Euler's theorem,

$$a^{\varphi(p)} = a^{(p-1)} = 1$$

□

RSA The RSA cryptosystem for message transmission works as follows:

1. Choose two distinct large primes p, q .
2. Compute $m = pq \Rightarrow \varphi(m) = (p - 1)(q - 1)$.
3. Choose $e \leq m$ an encryption key such that $\gcd(e, \varphi(m)) = 1$.
4. Use Euclid's algorithm to determine d such that $ed - k\varphi(m) = 1$ for some integer k .
5. The encoding key is the pair (m, e) , and it can be published. To decode, you use the decoding key (m, d) which is to be kept private.

To send a message x to someone, you need their public pair (m, e) . You first compute $c \equiv x^e \pmod{m}$, which can be sent publicly. To decode, the person will use their private pair (m, d) , computing $x \equiv c^d \pmod{m} \equiv x^{ed} \pmod{m}$.

Why is it the case that $x^{ed} \equiv x \pmod{m}$? Well...

Theorem 3.11. Let p, q be two distinct primes, and $m = pq$. Let $e : \gcd(e, \varphi(m)) = 1$, and let $d \in \mathbb{Z} : ed - k\varphi(m) = 1$ for some $k \in \mathbb{Z}$. Then,

$$x^{ed} \equiv x \pmod{m}$$

for all $x \in \{1, \dots, m\}$.

Appendix A

Proofs from exercise sets

This barely needs such pompous titles but oh well. It's fun.

A.1 Theorem 1.1

Theorem. *Strong Induction \Rightarrow Well-ordering principle.*

Proof. We can prove this by induction. Suppose there exists a subset $Y \subset \mathbb{N}$ such that it contains no least element. Consider $P(n) = "n \notin Y"$.

Base: If 0 was in Y , then it would be its least element, since there are no smaller elements of \mathbb{N} . As such, it cannot be that $0 \in Y$, meaning $P(0)$ is true.

Induction: Assume $P(k)$ is true for any $k \in \{0, 1, \dots, n\}$. Then, if it was in Y , $n + 1$ would be its smallest element, since every smaller element is not in Y . As such, $P(n + 1)$ holds as well. Since P is hereditary and true for 0, it is true for any $n \in \mathbb{N}$.

□

A.2 Theorem 2.4

Theorem A.1. *Any $n > 1$ can be expressed by the product of primes.*

Proof. Consider $S = \{n \in \mathbb{N} : n > 1 \wedge n \text{ does not have a prime factorization}\}$. Then, we take the smallest element in this set, k . k cannot be a prime, so there exists $a, b < k$ such that $k = ab$. However, a and b cannot be in S , since they are smaller than k . This means that $a = \prod p_{a,k}^{a_k}, b = \prod p_{b,k}^{b_k}$, meaning that their product is the product of primes. Therefore, k cannot be in S , so S has to be empty. □