

SQL: Part II

Introduction to Database Management

CS348 Spring 2021

Basic SQL features

- Query
 - SELECT-FROM-WHERE statements
 - Set/bag (DISTINCT, UNION/EXCEPT/INTERSECT(ALL))
 - Subqueries (table, scalar, IN, EXISTS, ALL, ANY)
 - Aggregation and grouping (GROUP BY, HAVING)
 - Ordering (ORDER)
 - Outerjoins (and Nulls)
- Modification
 - INSERT/DELETE/UPDATE
- Constraints

Lecture 4

Incomplete information

- Example: *User* (*uid*, *name*, *age*, *pop*)
- Value **unknown**
 - We do not know Nelson's age
- Value **not applicable**
 - Suppose *pop* is based on interactions with others on our social networking site
 - Nelson is new to our site; what is his *pop*?

Solution 1

- Dedicate a value from each domain (type)
 - pop cannot be -1 , so use -1 as a special value to indicate a missing or invalid pop

```
SELECT AVG(pop) FROM User;
```

Incorrect answers

```
SELECT AVG(pop) FROM User WHERE pop <> -1;
```

Complicated

- Perhaps the value is not as special as you think!
 - the Y2K bug



Solution 2

- A valid-bit for every column
 - *User (uid,
name, name_is_valid,
age, age_is_valid,
pop, pop_is_valid)*

```
SELECT AVG(pop) FROM User WHERE pop_is_valid;
```

- Complicates schema and queries

Solution 3

- Decompose the table; missing row = missing value
 - *UserName* (uid, name)
 - *UserAge* (uid, age)
 - *UserPop* (uid, pop)
 - *UserID* (uid)
- Conceptually the cleanest solution
- Still complicates schema and queries
 - How to get all information about users in a table?
 - Natural join doesn't work!

SQL's solution

- A special value **NULL**
 - For every domain
 - Special rules for dealing with NULL's
- Example: *User* (*uid*, *name*, *age*, *pop*)
 - $\langle 789, \text{"Nelson"}, \text{NULL}, \text{NULL} \rangle$

Three-valued logic

TRUE = 1, FALSE = 0, UNKNOWN = 0.5

$x \text{ AND } y = \min(x, y)$

$x \text{ OR } y = \max(x, y)$

$\text{NOT } x = 1 - x$

- Comparing a NULL with another value (including another NULL) using =, >, etc., the result is NULL
- WHERE and HAVING clauses only select rows for output if the condition evaluates to TRUE
 - NULL is not enough
- Aggregate functions ignore NULL, except COUNT(*)

Unfortunate consequences

- Q1a = Q1b?

```
Q1a. SELECT AVG(pop) FROM User;
```

```
Q1b. SELECT SUM(pop)/COUNT(*) FROM User;
```

- Q2a = Q2b?

```
Q2a. SELECT * FROM User;
```

```
Q2b SELECT * FROM User WHERE pop=pop;
```

- Be careful: NULL breaks many equivalences

Another problem

- Example: Who has NULL pop values?

```
SELECT * FROM User WHERE pop = NULL;
```

Does not work!

```
(SELEC * FROM User)  
EXCEPT ALL  
(SELECT * FROM USER WHERE pop=pop);
```

Works, but ugly

- SQL introduced special, built-in predicates
IS NULL and **IS NOT NULL**

```
SELECT * FROM User WHERE pop IS NULL;
```

Outerjoin motivation

- Example: a master group membership list

```
SELECT g.gid, g.name AS gname,  
       u.uid, u.name AS uname  
FROM Group g, Member m, User u  
WHERE g.gid = m.gid AND m.uid = u.uid;
```

- What if a group is empty?
- It may be reasonable for the master list to **include empty groups** as well
 - For these groups, *uid* and *uname* columns would be NULL

Outerjoin examples

Group

gid	name
abc	Book Club
gov	Student Government
dps	Dead Putting Society
nuk	United Nuclear Workers

Group ⋈ Member

gid	name	uid
abc	Book Club	857
gov	Student Government	123
gov	Student Government	857
dps	Dead Putting Society	142
nuk	United Nuclear Workers	NULL
foo	NULL	789

Member

uid	gid
142	dps
123	gov
857	abc
857	gov
789	foo

A **full outerjoin** between R and S :

- All rows in the result of $R \bowtie S$, plus
- “Dangling” R rows (those that do not join with any S rows) padded with NULL’s for S ’s columns
- “Dangling” S rows (those that do not join with any R rows) padded with NULL’s for R ’s columns

Outerjoin examples

Group \bowtie Member

Group

gid	name
abc	Book Club
gov	Student Government
dps	Dead Putting Society
nuk	United Nuclear Workers

Member

uid	gid
142	dps
123	gov
857	abc
857	gov
789	foo

gid	name	uid
abc	Book Club	857
gov	Student Government	123
gov	Student Government	857
dps	Dead Putting Society	142
nuk	United Nuclear Workers	NULL

- A **left outerjoin** ($R \bowtie S$) includes rows in $R \bowtie S$ plus dangling R rows padded with NULL's

Group \bowtie Member

gid	name	uid
abc	Book Club	857
gov	Student Government	123
gov	Student Government	857
dps	Dead Putting Society	142
foo	NULL	789

- A **right outerjoin** ($R \bowtie S$) includes rows in $R \bowtie S$ plus dangling S rows padded with NULL's

Outerjoin syntax

```
SELECT * FROM Group LEFT OUTER JOIN Member  
ON Group.gid = Member.gid;
```

$\approx \text{Group} \bowtie_{\text{Group.gid}=\text{Member.gid}} \text{Member}$

```
SELECT * FROM Group RIGHT OUTER JOIN Member  
ON Group.gid = Member.gid;
```

$\approx \text{Group} \bowtie_{\text{Group.gid}=\text{Member.gid}} \text{Member}$

```
SELECT * FROM Group FULL OUTER JOIN Member  
ON Group.gid = Member.gid;
```

$\approx \text{Group} \bowtie_{\text{Group.gid}=\text{Member.gid}} \text{Member}$

☞ A similar construct exists for regular (“inner”) joins:

```
SELECT * FROM Group JOIN Member ON Group.gid = Member.gid;
```

☞ For natural joins, add keyword NATURAL; don’t use ON

```
SELECT * FROM Group NATURAL JOIN Member;
```

SQL features covered so far

- SELECT-FROM-WHERE statements
- Set and bag operations
- Table expressions, subqueries
- Aggregation and grouping
- Ordering
- NULL's and outerjoins

☞ Next: data modification statements, constraints

INSERT

- Insert one row
 - User 789 joins Dead Putting Society

```
INSERT INTO Member VALUES (789, 'dps');
```

- Insert the result of a query
 - Everybody joins Dead Putting Society!

```
INSERT INTO Member  
  (SELECT uid, 'dps' FROM User  
   WHERE uid NOT IN (SELECT uid  
                     FROM Member  
                     WHERE gid = 'dps'));
```


DELETE

- Delete **everything** from a table

```
DELETE FROM Member;
```

- Delete according to a **WHERE** condition

- Example: User 789 leaves Dead Putting Society

```
DELETE FROM Member WHERE uid=789 AND gid='dps';
```

- Example: Users under age 18 must be removed from United Nuclear Workers

```
DELETE FROM Member  
WHERE uid IN (SELECT uid FROM User WHERE age < 18)  
AND gid = 'nuk';
```

UPDATE

- Example: User 142 changes name to “Barney”

```
UPDATE User  
SET name = 'Barney'  
WHERE uid = 142;
```

- Example: We are all popular!

```
UPDATE User  
SET pop = (SELECT AVG(pop) FROM User);
```

- But won't update of every row causes average *pop* to change?
 - ☞ Subquery is always computed over the old table

Constraints

- Restrictions on allowable data in a database
 - In addition to the simple structure and type restrictions imposed by the table definitions
- Why use constraints?
 - Protect data integrity (catch errors)
 - Tell the DBMS about the data (so it can optimize better)
- Declared as **part of the schema** and enforced by the DBMS

Types of SQL constraints

- NOT NULL
- Key
- Referential integrity (foreign key)
- General assertion
- Tuple- and attribute-based CHECK's

NOT NULL constraint examples

```
CREATE TABLE User  
(uid DECIMAL(3,0) NOT NULL,  
name VARCHAR(30) NOT NULL,  
twitterid VARCHAR(15) NOT NULL,  
age DECIMAL (2,0),  
pop DECIMAL(3,2));
```

```
CREATE TABLE Group  
(gid CHAR(10) NOT NULL,  
name VARCHAR(100) NOT NULL);
```

```
CREATE TABLE Member  
(uid DECIMAL(3,0) NOT NULL,  
gid CHAR(10) NOT NULL);
```

Key declaration examples

```
CREATE TABLE User
(uid DECIMAL(3,0) NOT NULL PRIMARY KEY,
name VARCHAR(30) NOT NULL,
twitterid VARCHAR(15) NOT NULL UNIQUE,
age DECIMAL(2,0),
pop DECIMAL(3,2));
```

At most one
primary key per
table

Any number of
UNIQUE keys per
table

```
CREATE TABLE Group
(gid CHAR(10) NOT NULL PRIMARY KEY,
name VARCHAR(100) NOT NULL);
```

This form is
required for multi-
attribute keys

```
CREATE TABLE Member
(uid DECIMAL(3,0) NOT NULL,
gid CHAR(10) NOT NULL,
PRIMARY KEY(uid,gid));
```

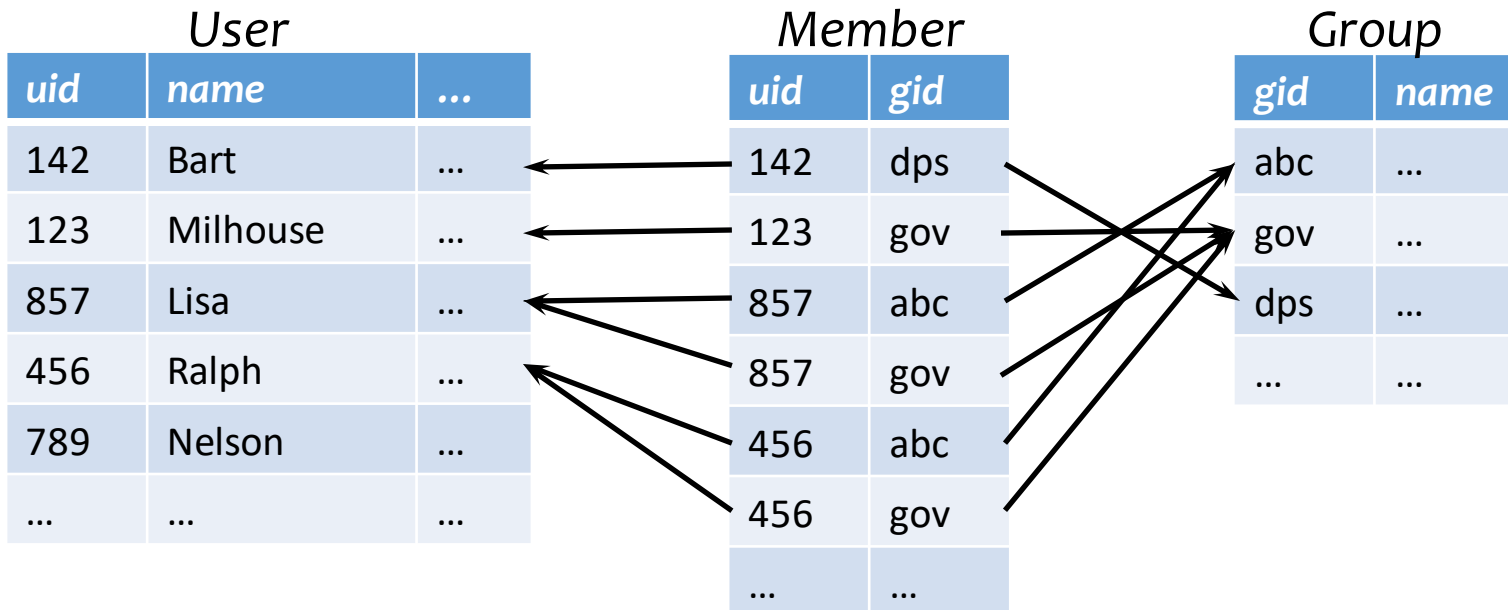
```
CREATE TABLE Member
(uid DECIMAL(3,0) NOT NULL PRIMARY KEY,
gid CHAR(10) NOT NULL PRIMARY KEY,
```

Incorrect!

Referential integrity example

- If an *uid* appears in *Member*, it must appear in *User*
 - *Member.uid* references *User.uid*
- If a *gid* appears in *Member*, it must appear in *Group*
 - *Member.gid* references *Group.gid*

☞ That is, no “dangling pointers”



Referential integrity in SQL

- Referenced column(s) must be **PRIMARY KEY**
- Referencing column(s) form a **FOREIGN KEY**
- Example

```
CREATE TABLE Member  
(uid DECIMAL(3,0) NOT NULL REFERENCES User(uid),  
gid CHAR(10) NOT NULL,  
PRIMARY KEY(uid,gid),  
FOREIGN KEY (gid) REFERENCES Group(gid));
```



This form is required for multi-attribute foreign keys

```
CREATE TABLE MemberBenefits  
(....  
FOREIGN KEY (uid,gid) REFERENCES Member(uid,gid));
```


Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Insert or update a *Member* row so it **refers to a non-existent uid**
 - **Reject**

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	456	abc
...	456	gov
			000	gov

Reject

Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Delete or update a *User* row whose *uid* is referenced by some *Member* row
 - Multiple Options (in SQL)

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	456	abc
...	456	gov
...

Option 1: Reject

```
CREATE TABLE Member
(uid DECIMAL(3,0) NOT NULL
REFERENCES User(uid)
ON DELETE CASCADE,
....);
```

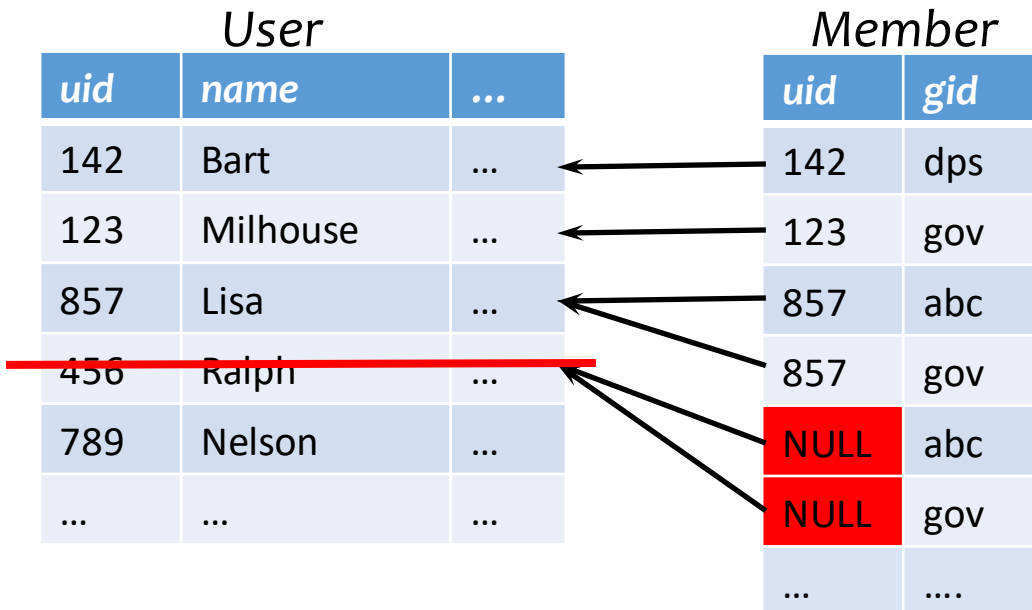
Option 2: Cascade
(ripple changes to all referring rows)

Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Delete or update a *User* row whose *uid* is referenced by some *Member* row
 - Multiple Options (in SQL)

```
CREATE TABLE Member
(uid DECIMAL(3,0) NOT NULL
REFERENCES User(uid)
ON DELETE SET NULL,
....);
```



Option 3: Set NULL
(set all references to NULL)

Deferred constraint checking

- Example:

```
CREATE TABLE Dept
(name CHAR(20) NOT NULL PRIMARY KEY,
chair CHAR(30) NOT NULL
REFERENCES Prof(name));
```

```
CREATE TABLE Prof
(name CHAR(30) NOT NULL PRIMARY KEY,
dept CHAR(20) NOT NULL
REFERENCES Dept(name));
```

- The first INSERT will always violate a constraint!
- **Deferred constraint checking** is necessary
 - Check only at the end of a transaction
 - Allowed in SQL as an option
- Curious how the schema was created in the first place?
 - **ALTER TABLE ADD CONSTRAINT** (read the manual!)

General assertion

- `CREATE ASSERTION assertion_name
CHECK assertion_condition;`
- *assertion_condition* is checked for each modification that could potentially violate it
- Example: *Member.uid* references *User.uid*

```
CREATE ASSERTION MemberUserRefIntegrity  
CHECK (NOT EXISTS  
      (SELECT * FROM Member  
        WHERE uid NOT IN  
          (SELECT uid FROM User)));
```

Tuple- and attribute-based CHECK's

- Associated with a single table
- Only checked when a tuple/attribute is inserted/updated
 - Reject if condition evaluates to FALSE
 - TRUE and UNKNOWN are fine
- Examples:

```
CREATE TABLE User(...  
  age INTEGER CHECK(age IS NULL OR age > 0),  
  ...);
```

```
CREATE TABLE Member  
(uid INTEGER NOT NULL,  
  CHECK(uid IN (SELECT uid FROM User)),  
  ...);
```

Exercise Question:
How does it differ
from a referential
integrity constraint
(slides 26-27)?

SQL features covered so far

- Query
 - SELECT-FROM-WHERE statements
 - Set and bag operations
 - Table expressions, subqueries
 - Aggregation and grouping
 - Ordering
 - Outerjoins (and NULL)
 - Modification
 - INSERT/DELETE/UPDATE
 - Constraints
- ☞ Next: triggers, views, indexes