Multipath extension for QUIC Draft-ietf-quic-multipath-13

QUIC meeting @ IETF-122 Bangkok

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Agenda

- ❖ Brief Summary: Changes from draft -11 to draft -13
- Hackathon Interop Reports: draft -13
- Issue
- Next Step

Diff from -11 to -12

- 1. Add PATH_CIDS_BLOCKED frame as an explicit signal for situations when endpoints don't have unused CIDs for the corresponding path
- 2. It is recommended to open a new path for active client migration.
- 3. Add guidance for handling PTO on multiple paths.

Clarifications:

- Packets which only contains PATH_ACK frames are not congestion controlled, but senders should carefully consider the load induced by these packets.
- 5. When endpoint sends PATH_ABANDON frame but don't receive a corresponding PATH_ABANDON, it's left to the implementation when to release the path resources, but endpoints SHOULD wait at least 3 PTOs before remove the issued CIDs of the path.
- 6. It is recommended (SHOULD) to always send the PATH_ABANDON frame on another path than the to-be-abandoned path.

Updated transport parameter for the new version because of new frame.

Diff from -12 to -13

- 1. Add Next Sequence Number (i) field to the PATH_CIDS_BLOCKED frame
- Recommend sending at least one unused CID for all paths up to the max path ID
- Changed the error code for MP frames in 0-RTT to align correctly with RFC9000

Clarifications:

- 4. Around the use, acking, and retransmissions of the PATHS_BLOCKED and PATH_CIDS_BLOCKED, MAX_PATH as well as PATH_AVAILABLE, and
- 5. On the use of IV and nonce.

Update transport parameter for the new version because of change in frame.

Hackathon Interop Reports: draft-13

server						
client ↓	xquic	picoquic	Rask	quiche	mp-aioquic	
xquic	HVDCSUABL	HVDCSUABL	HDVCSUABL	HVDCSUABL	HVDCSUAB	
picoquic	HVDCUA	HVDCU	HVDCUA	HVDC		
Rask	HVDCSURMB	HVDCSURM	HVDCSURMB	HVDCSUR		
quiche	HVDCISRA	HVDCISRA	HVDCISRA	HVDCISRA	HVDCIRA	
mp-aioquic						

ted								
code	details							
Н	The handshake completes with successful negotiation of enable_multipath transport parameter (both ends indicate 0x01)							
V	Client sends PATH_CHALLENGE frame to open a new path and server replies with PATH_RESPONSE							
D	Stream data (of one of more streams) is send on all paths; ACK_MP frames are sent and processed							
С	Client closes a path with PATH_ABANDON frame							
Tested								
code	details							
1	A server offers new CIDs for a path using PATH_NEW_CONNECTION_ID to a client in advance. Upon some events, the client starts using a new server CID on one pat							
S	Client sends PATH_BACKUP and PATH_AVAILABLE frames							
U	One endpoint updates keys and sends at least one packet with the new key on all active paths							
Α	One endpoint sends data and the other endpoints sends PATH_ACK (randomly) on all path independent of where data is received							
R	One endpoint send an PATH_RETIRE_CONNECTION_ID for an active path							
М	Change CID and 4-tuple	e.g. port on an existing path						
В	Server configures a path limitation of two paths, client is configured to open three or more paths. When unable to open requested number of paths it sends PATHS_BLOCKED and server logs or increases number of paths by announcing more using PATH_NEW_CONNECTION_IDS.							
	wishes to open a path w established path but is u	th a valid Path ID or change the	e CID on an ire no unused					
	H V D C Tested code I S U A R M	code details H The handshake complete V Client sends PATH_CHA D Stream data (of one of m C Client closes a path with Tested code details I A server offers new CIDs S Client sends PATH_BAC U One endpoint updates ke A One endpoint sends data R One endpoint send an PA M Change CID and 4-tuple Server configures a path When unable to open re or increases number of p A sender can send a PA wishes to open a path with established path but is u	code details H The handshake completes with successful negotiation of V Client sends PATH_CHALLENGE frame to open a new D Stream data (of one of more streams) is send on all pat C Client closes a path with PATH_ABANDON frame Tested code details I A server offers new CIDs for a path using PATH_NEW_S Client sends PATH_BACKUP and PATH_AVAILABLE fr U One endpoint updates keys and sends at least one pac A One endpoint sends data and the other endpoints send R One endpoint send an PATH_RETIRE_CONNECTION_M Change CID and 4-tuple e.g. port on an existing path Server configures a path limitation of two paths, client is When unable to open requested number of paths it sen or increases number of paths by announcing more usin A sender can send a PATH_CIDS_BLOCKED frame (ty wishes to open a path with a valid Path ID or change the established path but is unable to do so because there a	code details H The handshake completes with successful negotiation of enable_multipath trans V Client sends PATH_CHALLENGE frame to open a new path and server replies v D Stream data (of one of more streams) is send on all paths; ACK_MP frames are C Client closes a path with PATH_ABANDON frame Tested code details I A server offers new CIDs for a path using PATH_NEW_CONNECTION_ID to a client sends PATH_BACKUP and PATH_AVAILABLE frames U One endpoint updates keys and sends at least one packet with the new key on a client sends path and the other endpoints sends PATH_ACK (randomly), R One endpoint sends data and the other endpoints sends PATH_ACK (randomly), R One endpoint send an PATH_RETIRE_CONNECTION_ID for an active path M Change CID and 4-tuple e.g. port on an existing path Server configures a path limitation of two paths, client is configured to open thre when unable to open requested number of paths it sends PATH_S_BLOCKED at or increases number of paths by announcing more using PATH_NEW_CONNECTION_ED and sends a path I wishes to open a path with a valid Path ID or change the CID on an established path but is unable to do so because there are no unused	code details The handshake completes with successful negotiation of enable_multipath transport parameter (both of V Client sends PATH_CHALLENGE frame to open a new path and server replies with PATH_RESPONSI D Stream data (of one of more streams) is send on all paths; ACK_MP frames are sent and processed C Client closes a path with PATH_ABANDON frame Tested code details I A server offers new CIDs for a path using PATH_NEW_CONNECTION_ID to a client in advance. Upor S Client sends PATH_BACKUP and PATH_AVAILABLE frames U One endpoint updates keys and sends at least one packet with the new key on all active paths A One endpoint sends data and the other endpoints sends PATH_ACK (randomly) on all path independe R One endpoint send an PATH_RETIRE_CONNECTION_ID for an active path M Change CID and 4-tuple e.g. port on an existing path Server configures a path limitation of two paths, client is configured to open three or more paths. When unable to open requested number of paths it sends PATH_S_BLOCKED and server logs or increases number of paths by announcing more using PATH_NEW_CONNECTION_IDS. A sender can send a PATH_CIDS_BLOCKED frame (type=0x15228c0e) when it wishes to open a path with a valid Path ID or change the CID on an established path but is unable to do so because there are no unused	code details The handshake completes with successful negotiation of enable_multipath transport parameter (both ends indicate 0x01) V Client sends PATH_CHALLENGE frame to open a new path and server replies with PATH_RESPONSE D Stream data (of one of more streams) is send on all paths; ACK_MP frames are sent and processed C Client closes a path with PATH_ABANDON frame Tested code details I A server offers new CIDs for a path using PATH_NEW_CONNECTION_ID to a client in advance. 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Upon some events, the client starts using a new: S Client sends PATH_BACKUP and PATH_AVAILABLE frames U One endpoint updates keys and sends at least one packet with the new key on all active paths A One endpoint sends data and the other endpoints sends PATH_ACK ((randomly) on all path independent of where data is received R One endpoint send an PATH_RETIRE_CONNECTION_ID for an active path M Change CID and 4-tuple e.g. port on an existing path Server configures a path limitation of two paths, client is configured to open three or more paths. When unable to open requested number of paths it sends PATHS_BLOCKED and server logs or increases number of paths by announcing more using PATH_NEW_CONNECTION_IDS. A sender can send a PATH_CIDS_BLOCKED frame (type=0x15228c0e) when it wishes to open a path with a valid Path ID or change the CID on an established path but is unable to open a path with a valid Path ID or change the CID on an established path but is unable to do so because there are no unused	

Issue 500: Is PATH_CIDS_BLOCKED frame really needed?

- Peer should provide at least one unused CID for all active paths, otherwise this seems to be rather an implementation error.
- We now recommend to provide CIDs for all unused Path IDs below the max limit of both side (see merged PR #501)
 - If you want to limit the number of paths, just set a lower max Path ID Limit.
- Still potential relevant use case: You want to open a new path but didn't get any CIDs for any unused Path ID (below max value)?

Proposed solution: Keep it for now as we already have it...

A client can choose to not implement and a server can choose to ignore the frame.

Next Steps

- Full editorial review on-going (including some restructuring of sections)
- Next version ready for WGLC!