

**BRAC UNIVERSITY**  
Merul Badda, Dhaka, Bangladesh  
**CSE331 : Automata and Computability**  
**Assignment 3**

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**1. Use the pumping lemma to show that the following languages are not regular:**

- A.  $L(M) \rightarrow \{0^{2n}1^n \mid n \geq 0\}$ , where  $\Sigma = \{0, 1\}$ .

See example 4 of the provided video resource on the Pumping Lemma for RL.

- B.  $L(M) \rightarrow \{0^n1^m0^{f(n,m)} \mid n, m \geq 0\}$ , where  $f(n, m) = n * m$  and  $\Sigma = \{0, 1\}$ .

See example 7 of the provided video resource on the Pumping Lemma for RL.

- C.  $L(M) \rightarrow \{0^n1^m \mid n > m\}$ , where  $\Sigma = \{0, 1\}$ .

See example 3 of the provided video resource on the Pumping Lemma for RL.

- D.  $L(M) \rightarrow \{0^{n^2} \mid n \geq 0\}$ , where  $\Sigma = \{0, 1\}$ .

See example 10 of the provided video resource on the Pumping Lemma for RL.

- E.  $L(M) \rightarrow \{w \text{ is not a palindrome}\}$ , where  $\Sigma = \{0, 1\}$ .

See example 18 of the provided video resource on the Pumping Lemma for RL.

**2. Write a CFG for the following CFL:**

- A.  $L(M) \rightarrow \{0^n1^m \mid n, m \geq 0 \text{ and } 2n = 3m\}$ , where  $\Sigma = \{0, 1\}$

$S \rightarrow 000S11 \mid \epsilon$

- B.  $L(M) \rightarrow \{0^n1^m2^m3^n \mid n, m > 0\}$ , where  $\Sigma = \{0, 1, 2, 3\}$

$S \rightarrow 0S3 \mid A \mid \epsilon$

$A \rightarrow 1A2 \mid \epsilon$

- C.  $L(M) \rightarrow \{w = 0^i1^j2^k \mid i, j, k \geq 0 \text{ and } j < i + k\}$ , where  $\Sigma = \{0, 1, 2\}$

$S \rightarrow XABY \mid XA \mid BY$

$A \rightarrow 0A1 \mid \epsilon$

$B \rightarrow 1B2 \mid \epsilon$

$X \rightarrow 0X \mid 0$

$Y \rightarrow 2Y \mid 2$

- D.  $L(M) \rightarrow \{w_1 \# w_2 \mid \text{the number of } 00 \text{ in } w_1 \text{ is equal to the number of } 11 \text{ in } w_2\}$ , where  $\Sigma = \{0, 1\}$

$S \rightarrow 0B1 \mid XS \mid SY \mid Z$

$X \rightarrow 1X \mid 01X \mid \epsilon$

$Y \rightarrow 0Y \mid 01Y \mid \epsilon$

$Z \rightarrow 0\# \mid \#1 \mid 0\#1 \mid \#$

$A \rightarrow 0B \mid 1A$

B → 0D | 1A

C → D1 | C0

D → E1 | 0C

E → B | #

E.  $L(M) \rightarrow \{w\#x \mid w^R \text{ is a substring of } x\}$ , where  $\Sigma = \{0, 1\}$

S → AB

A → 0A0 | 1A1 | #B

B → 0B | 1B | ε