



Faculty of Engineering & Technology
Electrical & Computer Engineering Department

ADVANCED DIGITAL SYSTEMS DESIGN
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ADVANCED DIGITAL SYSTEMS DESIGN

Project

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Abstract:

This project endeavors to construct a microprocessor characterized by a register file functioning as a compact and expeditious Random-Access Memory (RAM), complemented by an Arithmetic and Logic Unit (ALU) serving as the computational nucleus for processing operation codes (opcodes) provided to it. The proposed methodology involves the composition of Verilog code for each constituent component or module, subsequently interlinking them in a cohesive manner. Additionally, a comprehensive testbench will be developed to assess the functionality of the microprocessor. This testbench will encompass a series of instructions specifying opcodes and the corresponding addresses for the retrieval and storage of numerical data within the register file. The successful implementation of these procedures will serve as a validation of the constructed microprocessor's robustness and reliability.

Part One:

1.1 Arithmetic And Logic Unit (ALU):

This Unit do various operations based on 6-bit operation code (opcode) given on two 32-bit numbers symbolized by (a) and (b) and output the 32-bit result, the numbers given to it are from register file by addresses of them extracted from 32-bit instruction.

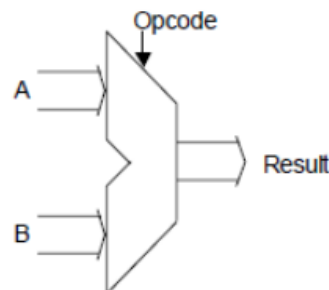


Fig.1.1: ALU

The operations that ALU can do following opcode based on last digit of my ID number 1212508 in decimal are listed below in Table:

<u>Opcode</u>	<u>Operation</u>	<u>Description</u>
1	$a + b$	Addition
6	$a - b$	Subtraction
13	$ a $	Absolute value of 'a'
8	$-a$	Negative value of 'a'
7	$\max(a, b)$	Maximum between 'a' and 'b'
4	$\min(a, b)$	Minimum between 'a' and 'b'
11	$\text{avg}(a, b)$	Average value of 'a' and 'b'
15	not a	reverse all bits of 'a'
3	$a \text{ OR } b$	OR between 'a' and 'b'
5	$a \text{ AND } b$	AND between 'a' and 'b'
2	$a \text{ XOR } b$	exclusive OR between 'a' and 'b'

How it works: We defined all opcodes as parameters with names associated with operation work, make a case statement to cover all opcodes to check if given

opcode equals one of them, and on positive edge result of operation will assigned the value immediately “blocking”, so on positive edge of current clock the result will appear.

How we tested it: We defined two numbers and change opcode to fit all opcodes with some delay between each case, to wait for operating the right result, since we define ‘a’ and ‘b’ as signed numbers then the result would be signed, from results since ‘a’ is positive then the absolute value also positive. Therefore, all operations successfully worked as expected in each operation, so our design is right.

The testbench results that verify the unit implemented successfully set in

Appendix 1

1.2 Register File:

This File as a small fast RAM but it’s implemented inside microprocessor, it stores 32 locations of wide 32-bit for each, these locations addresses from 0 to 31 store numbers that be passed to ALU at “out1” and “out2” to do various operations, also it can store the result of the ALU “in” in location addressed by address entered to register file “Address 3” extracted from 32-bit instruction.

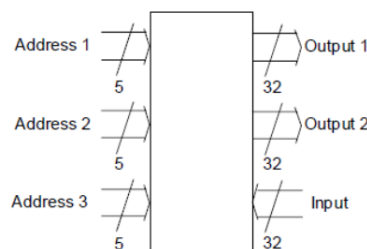


Fig.1.2: Register File

The register file contents based on second last digit of my ID number 1212508 which is zero got from project file from index 0 to index 31 below in Fig 1.2.1:

0	0	16	1238
1	7942	17	16008
2	13224	18	2426
3	15462	19	11930
4	8026	20	6724
5	3692	21	12790
6	9882	22	4842
7	8248	23	7108
8	3432	24	6296
9	178	25	3322
10	2378	26	10848
11	8302	27	14698
12	592	28	16378
13	7430	29	15456
14	10572	30	1260
15	7676	31	0

Fig.1.2.1: Table of stored values in Register file

Part Two:

2.1 Clocking the Register File:

To make sure that addresses and input to register file in right way, we may clock the register file on the clocks' edge, such that addresses of both numbers should be first passed to get the numbers from the file, then at next edge we got the address of result that should stored in the file at the end of successfully execution, in this case we used non-blocking assignment to put addresses and input value after positive edge of clock.

2.2 Enabling the Register File:

To use register file firstly we should enable it by checking the opcode if valid, if invalid then no matter to do anything; because it isn't enabled.

How it works: when enabled by valid opcode, get 3 addresses from input , first 2 to load 2 output values respectively based on addresses numbers, the third address is the address on which location we will store the value entered in input, load and store will not be on same moment to avoid any error or any unintended value, that is by “non-blocking” such that on positive edge all inputs come and after positive edge in order load before store and loaded value should be real value stored after this the value will be updated by store (just if any address or both first 2 equal last address of store), otherwise store can first normally, but to avoid any error as we mentioned we used non-blocking assignment.

How we tested it: we defined enable register as opcode valid , addresses , input value , input addresses and 2 output values, then initialize 3 addresses such that address two and address three equals each other, to test load and store are work successfully, and make clock, so at positive edge and valid opcode 2 outputs are ready then disabled to show the result with some delay, then again enable to test if new value stored in location of old value.

The testbench result shows that our register file design works right, as out1 loaded value shows the content of location addressed by addr1 and out2 also, but addr1 which equals addr2 store the value “0xF1234567” in same location, so non-blocking is right solution as load from location[addr2] must be before store new value in it. Therefore, from simulation results the load done before store such that old value was “0x00000D68” and new value is “0xF1234567” loaded and shown in last output line, so our design is right. The results are set in **Appendix 1**.

2.3 Creating the core of the microprocessor:

In order to manage ALU and Register file we should implement a simple microprocessor that connect them, and extract parameters of each one from the instruction provided. Therefore, the result of this microprocessor represents the result out from ALU after do operation on numbers.

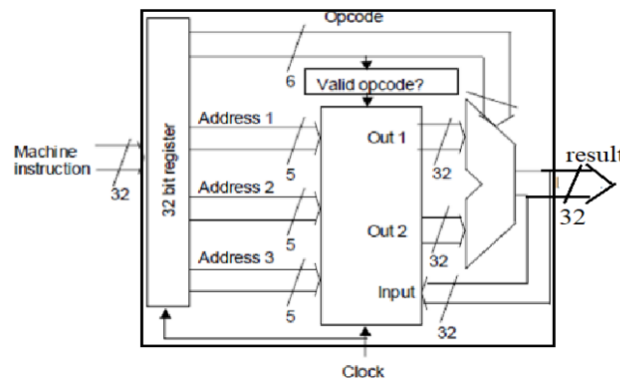


Fig.2.3: Microprocessor

Machine instruction 32-bit format:

|31 (11-bit) 21 | 20 (5-bit) 16 | 15 (5-bit) 11 | 10 (5-bit) 6 | 5 (6-bit) 0|

Unused	Address 3	Address 2	Address 1	Opcode
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How it works:

It connects ALU with register file, at positive edge get the 32-bit instruction, then extract all addresses of loaded values and value that should store. Therefore, first test if opcode valid to enable register file before give values from it, if not then not extract and again check till opcode valid or next instruction opcode valid. To add to that, in result of validity opcode giving addresses of input numbers to ALU and address of storing the result, then make some delaying on passing opcode to ALU till numbers come from the register file. After values come from register file result of them will be stored in the location of file addressed by address provided in the instruction, but showing the result of microprocessor after storing to guarantee that result stored correctly.

How we tested it: We initialize all inputs and outputs of other components that connect them, input instruction, clock and output result that obtained from ALU when opcode valid such that all valid opcodes in range determined in “if statements” compare each given opcode with them. To add to that, instantiate components and pass parameters then at positive edge extract opcode from the instruction and check it, if valid then enable register file, make some delay on passing opcode to ALU till values come then calculate the result, store it then output it from micro.

In testbench, we initialized the register file values and valid opcode again to calculate expected value of any given instruction, then we initialize array of 11 instructions to pass them in micro., now at positive edge instruction passed to micro. Calculate expected result by function calculate() then after 2 cycles (to guarantee that result obtained in right way) check if result identical with expected or not, by task called check() that take result and the expected, then print “Pass” if both identical and “Fail” if not.

***Note:** Making the array of instructions done by a website that take the format of instruction and make instructions by typing operation as assembly (i.e.: Add r31, r1, r2 that $r31 = r1 + r2$) make instruction of ADD opcode and addresses depending on numbers after ‘r’.

The testbench results shown in **Appendix 1** approved that our design is correct and works as expected.

Testbench code and results:

Testbench code and results:

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Microprocessor (Top):

Testbench code and results:

```

/*
Testbench of simple microprocessor
*/
module tp_mp;
    reg clk;
    reg signed [31:0] location [0:31]; // 32 location each one 32-bit to store the values in it start from index 0
    reg [5:0] opcodes [0:15]; // opcodes that ALU can do them
    reg [31:0] instructions [0:25]; // array of instructions
    reg [31:0] instruction; // instruction given from the array
    reg signed [31:0] result; // result of microprocessor
    reg signed [31:0] expected_result; // expected result of microprocessor
    integer i = 0; // index of stored instructions
    integer size; // size: last index of last instruction

    // instantiate microprocessor
    mp_top mp( .clk(clk), .instruction(instruction), .result(result) );

    // Clock generation
    initial begin
        clk = 0;
        forever #5 clk = ~clk;
    end

    initial begin
        // initialize the values stored in register file
        location[0] = 32'h0;
        location[1] = 32'h1F06;
        location[2] = 32'h33A8;
        location[3] = 32'h3C66;
        location[4] = 32'h1F5A;
        location[5] = 32'hE6C;
        location[6] = 32'h269A;
        location[7] = 32'h293B;
        location[8] = 32'hD68;
        location[9] = 32'h82;
        location[10] = 32'h94A;
        location[11] = 32'h209E;
        location[12] = 32'h250;
        location[13] = 32'h1D08;
        location[14] = 32'h294C;
        location[15] = 32'h1DFC;
        location[16] = 32'h4D6;
        location[17] = 32'h3E0B;
        location[18] = 32'h97A;
        location[19] = 32'h2E9A;
        location[20] = 32'h1A44;
        location[21] = 32'h31F6;
        location[22] = 32'h12EA;
        location[23] = 32'h1BC4;
        location[24] = 32'h1898;
        location[25] = 32'hCFA;
        location[26] = 32'h2A60;
        location[27] = 32'h396A;

        // initialize valid opcodes
        opcodes[0] = 6'h1;
        opcodes[1] = 6'h6;
        opcodes[2] = 6'hD;
        opcodes[3] = 6'h0;
        opcodes[4] = 6'h7;
        opcodes[5] = 6'h4;
        opcodes[6] = 6'hB;
        opcodes[7] = 6'hF;
        opcodes[8] = 6'h3;
        opcodes[9] = 6'h5;
        opcodes[10] = 6'h2;

        // initialize array of instructions
        instructions[0] = 32'h01F1041;
        instructions[1] = 32'h001FF886;
        instructions[2] = 32'h008E83C;
        instructions[3] = 32'h001300C8;
        instructions[4] = 32'h001910C7;
        instructions[5] = 32'h001C1884;
        instructions[6] = 32'h0000F28B;
        instructions[7] = 32'h0016008F;
        instructions[8] = 32'h001F1F43;
        instructions[9] = 32'h001A0C45;
        instructions[10] = 32'h00120642;

        size = 10;
        /*
        operation addr3, addr1, addr2
        add r31, r1, r2
        sub r31, r2, r31
        abs r14, r15
        neg r19, r3
        max r25, r3, r2
        min r28, r2, r3
        avg r6, r18, r30
        not r22, r2
        or r31, r29, r3
        and r26, r17, r19
        xor r18, r25, r6
        */
    end

    always @(posedge clk) begin
        // on positive edge
        if(i == size+1)
            $finish;
        instruction = instructions[i];
        expected_result = calculate(location[instruction[10:6]], location[instruction[15:11]], instruction[5:0], opcodes);
        $display("===== %0d =====", i+1);
        $display("Instruction: 0x%h \nopcode = %0d, a = location[%0d] = %0d, b = location[%0d] = %0d", instruction, instruction[5:0], instruction[10:6], location[instruction[10:6]], instruction[15:11], location[instruction[15:11]]);
        #5 check(expected_result, result); // after two clock cycles the result obtained from micro. then check if it is correct or not
        location[instruction[20:16]] = expected_result;
        i = i + 1;
    end

    function signed [31:0] calculate; // as functionality of ALU just for verification
        input signed [31:0] a, b;
        input [5:0] opcode;
        input [5:0] opcodes[0:10];

        case(opcode)
            /* Addition operation */
            opcodes[0]:
                return a + b;

            /* Subtraction operation */
            opcodes[1]:
                return a - b;

            /* Absolute value of input 'a' operation */
            opcodes[2]:
                if ( a < 0 ) // if 'a' is negative then change it positive by multiplying it with "-1"
                    return (-1) * a;
                else // else if not negative then the result equals 'a' whatever positive or zero
                    return a;

            /* Negate the value of 'a' operation */
            opcodes[3]:
                return (-1) * a; // just multiplying it with minus one whatever positive, negative or zero

            /* Maximum of 'a' and 'b' operation */
            opcodes[4]:
                if ( a > b ) // if 'a' greater than 'b', then result is 'a'
                    return a;
                else if ( b > a ) // if the opposite then result is 'b'
                    return b;
                else // if 'a' equals 'b' then result anyone of them, I chose 'a'
                    return a;
        endcase
    end

```

```

261 /* Minimum of 'a' and 'b' operation */
262 opcodes[1]:
263     if (a < b) // if 'a' less than 'b', then result is 'a'
264         return a;
265     else if (b < a) // if the opposite then result is 'b'
266         return b;
267     else // if 'a' equals 'b' then result anyone of them, I chose 'a'
268         return a;
269
270 /* Average value of 'a' and 'b' operation */
271 opcodes[2]:
272     return (a + b) / 2;
273
274 /* invert all bits of 'a' (bit-wise NOT) operation */
275 opcodes[3]:
276     return ~a; // Inverting all bits of 'a' then set the result of that to output 'result'
277
278 /* bit-wise OR between 'a' and 'b' operation */
279 opcodes[4]:
280     return a | b; // ORing all bits of 'a' and 'b' then set the result of that to output 'result'
281
282 /* bit-wise AND between 'a' and 'b' operation */
283 opcodes[5]:
284     return a & b; // ANDing all bits of 'a' and 'b' then set the result of that to output 'result'
285
286 /* bit-wise XOR between 'a' and 'b' operation */
287 opcodes[6]:
288     return a ^ b; // XORing all bits of 'a' and 'b' then set the result of that to output 'result'
289
290 endcase
291
292 endfunction
293
294 task check;
295     input signed [31:0] expected_result;
296     input signed [31:0] result;
297
298     if(expected_result == result) begin
299         $display("Result: %0d, Expected Result: %0d", result, expected_result);
300         $display("*****");
301         $display("*****");
302     end
303     else begin
304         $display("Result: %0d, Expected Result: %0d", result, expected_result);
305         $display("*****");
306         $display("*****");
307     end
308 endtask
309
310 endmodule

```

