

# FAKULTÄT FÜR INFORMATIK

DER TECHNISCHEN UNIVERSITÄT MÜNCHEN

Master Thesis in Informatics

# **Adding C++ Support to MBEDDR**

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Adding C++ Support to MBEDDR

C++ Unterstützung für MBEDDR

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Date: September 16, 2013



Ich versichere, dass ich diese Masterarbeit se Quellen und Hilfsmittel verwendet habe.	lbständig verfasst und nur die angegebenen
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### Acknowledgments

If someone contributed to the thesis... might be good to thank them here.

### **Abstract**

An abstracts abstracts the thesis!

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### 1 Introduction

In embedded programming the C++ programming language is widely spread, [6]. Being a general purpose programming language, C++ does not provide, however, any special support for an embedded systems programmer.

By changing the language itself, together with a tool set for it, it is possible to get a better environment for a dedicated domain, for example, specifically for embedded programming.

The first possible approach is dropping some language features, to get the language, which is simpler. As an example, a subset of C++, called *Embedded C*++ can be brought, [2]. The approach taken in Embedded C++ is omitting very many core features of C++ off, allows for a higher degree of optimizations by compiler possible. Embedded C++ was intended to allow higher software quality through better understanding of the limited C++ by programmers, higher quality of compilers, better suitability for the embedded domain, [1]. This approach, however, has been criticized by the C++ community, specifically for the inability of the limited language to take advantage of the C++ standard library, which requires the C++ language features, absent in Embedded C++, [7].

The second approach to modify a language to get it more suitable for the embedded development, consists of extending the language with constructions specific to the domain. Such approach is taken, for example, in the *MBEDDR Project*, to improve on the C programming language, [11]. Extensions to C language developed in *MBEDDR Project* include state machines and decision tables.

A special language engineering environment is used to support modular and incremental language development in *MBEDDR Project*, *JetBrains MPS*. The language under development is split in a special class-like items, called concepts. As an example of a concept expression can be taken. Over the inheritance mechanisms, it is possible to extend languages, providing new concepts as children of the existing ones. For example, expression concept can be extended to support new sort of expressions, e.g. decision tables.

Building a general purpose programming language in a language engineering environment brings a basis to develop domain specific extensions to a well-known general purpose language.

Additionally to the language modification, the *Integrated Development Environment (IDE)* can be improved to support domain specific development. Various analyses can be built in into the code editor in order to detect inconsistencies, or, simply, "dangerous" constructs, and inform the programmer. Certain code formatting, or standard requirements could be enforced as well, and many more.

A mixture of the two approaches could be used in an attempt to achieve a "better" C++ for embedded development. A special *IDE* can be created together with a new C++ language flavor, which supports the embedded C++ programmer. This is the problem to be solved by this Master Thesis.

The new language together with the new *IDE* can later serve as a basis for extending the

C++ programming language with domain specific constructs for embedded programming. Creation of these extensions lies out of scope for this Master Thesis, and is left for further research.

The approach taken in this work goes further into exploring the language modularity on the basis of *JetBrains MPS*While building the C++ programming language itself with the goal of embedded domain specific extensions in mind, the C++ itself is being built itself as an extension to the C programming language, provided by the *MBEDDR Project*.

Although C++ is a separate from C language, the high degree of similarity allows to make use of the C programming language, implemented by the *MBEDDR Project* as a foundation. Not only reuse of the basic C is achieved, but also the embedded extensions from the *MBEDDR Project* are immediately supported by the newly built C++.

This work explores further the support, provided by *JetBrains MPS* for the modular language construction, [3], and reviews it from the architectural point of view.

The C++ programming language is provided with a number of automations and analyses for it. The automations include code generation and structuring. They are implemented as a programming on the *Abstract Syntax Tree* (*AST*) in a Java-like programming language. The analyses and automations intend to improve quality, security and better understanding of the code.

In the *JetBrains MPS Application Programming Interface (API)* it is not explicitly defined, when the analyses and checks take place, how much of the computational resource they can take advantage of. This may affect the overall *IDE* performance, as the analyses complexity may be high. The question of analyses run-time and complexity is raised and discussed in this work.

### 2 Foundations

Before describing the technologies on which the current work is based, as well as the work itself, it makes sense to describe more general foundations and principles, around which the technology is built.

In the Section 2.1 I describe two approaches to create an *IDE* for a certain language, and mainly the projectional approach, which originates from the area of building new *Domain Specific Languages (DSLs)*.

In the Section 2.2 I describe the modular approach toward alnguage engineering and extending, intensively used with the projectional approach to construct languages.

### 2.1 Building DSLs and IDEs

This section compares the traditional approach to build textual editors for the program code with the projectional approach, bringing up motivation for the least.

#### 2.1.1 Traditional Approach

Traditionally programming languages are used in a textual form in text files, forming programs. However the textual nature is not typical for the structure of programs themselves, being rather a low-level code representation, especially when talking abut syntax, which is only necessary for parsers to produce correct results, and not for the program intended semantics.

Parsers are used to construct so-called *AST*s from the textual program representation. *AST*s are structures in memory, usually graph-alike, reminding a control flow graph, where nodes are different statements and edges are the ways control passes from one statement to the next one.

For the developer, using an editor, the degree to which the editor can support the development process is important. For this, the editor has to recognize the programming language constructions and provide possible assistance. Among such assistance can be code formatting, syntax validation, source code transformations (including refactoring support), code analyses and verification, source code generation and others. Many of these operation rely indeed on the higher than text level notions related to program such as a method, a variable, a statement. A good editor has to be aware of these higher level program structures, to provide meaningful automations for the operations mentioned above.

Nowadays, most of the editors work with text, and, to provide assistance to a programmer, integrate with a parser/compiler front-end for the programming language. Such way to extract the program structure during editing is not perfect for several reasons.

First of all, the program being edited as text is not syntactically correct at every moment, being incomplete, for example. Under such circumstances the parsing front-end can not

be successfully invoked and returns error messages which are either not related to the program, when the code is completed, or false-positive warning and errors.

Secondly, after a minor editing of the code, usually the whole text file has to be processed again. Such compiler calls are usually computationally expensive, they slow down, sometimes significantly, the performance of the developer machine. Various techniques exist to speed it up, including partial and pre-compilation, but the problem is still relevant to a large extent.

Moreover, the textual nature of the code complicates certain operations additionally. As an example, we can take a refactoring to rename a method. Every usage of the method, being renamed, has to be found and changed. To implement it correctly an editor must take into account various possible name collisions, as well as presume a compilable state of the program prior to the start of the refactoring.

Not to mention the parsing problem itself. Parsing a program in a complex language like C++ is a difficult problem, it involves the need to resolve correctly scoping and typing, templates and related issues, work with pre-processor directives incorporated in the code. In this regard different compilers treat C++ in a different way, creating dialects, which may represent obstacles for the code to be purely cross-platform.

Listing 2.1: Closing several blocks

```
class MyClass {
  void doSomething() {
    while(true) {
        try {
            // ...
        }
        catch(Exception e) {
        }
    }
};
```

The textual representation of program code, involves the need in formatting and preserving syntax. These both tasks, indeed, have nothing to do with the functionality program, and additionally load the developer, reducing productivity. As an example, here I can mention the need to close several blocks ending at the same point correctly, indenting the closing brace symmetrically to opening one. The Listing 2.1 demonstrates it in the last few lines.

#### 2.1.2 Projectional Approach

Another approach which can be taken in the organization of an editor for a programming language is called *projectional approach*. Projectional editors do not work with a low-level textual representation of a program, but rather with a higher level concept, *AST*s. This approach is especially useful and used when constructing new *DSL*s.

Working with *AST*s directly has several advantages over the conventional textual code editing.

Firstly, all syntax errors are no longer possible, as there is no syntax.

Secondly, there is no need to format the code on the level of indentation and look, since it is only needed for textual code.

Thirdly, all features, which in textual approach require parsing, can be implemented without a parser involved, because *AST* is always known to the editor.

Additionally, as the compilers still expect a code in a textual form, code generation is used to convert the *AST* into the text code for the further use. The code generation step can be customized to provide support for a variety of compilers, when the compilers differ.

Projectional editors have to display the *AST* to the developer, in order for him/her to work with it. Such visualization of an *AST* is called "projection", giving a name to the editor class.

The model of code is stored as an *AST* in the projectional editor. As in the Model-View-Controller pattern the view for the model can be implemented separately, [4]. Thus the code may be presented in a number of different ways to the user. For example, the *AST* can be visualized as a graph, similar to control flow graph. This visualization, however, is not always advantageous being sometimes not compact and complicated to overview.

Figure 2.1: Example projection of an AST, "source code" view

One of the well-accepted way to visualize *AST* is by visualizing its textual representation, as if it would be written as a text code in the programming language, see Figure 2.1. There can be in principle many such textual visualizations, supporting different ways the code looks. Normally in the traditional approach this has to be achieved by reformatting, and thus changing, the source code. This is performed for the code to look similar across the developed software, and standards or coding guidelines are written to enforce the way to format the text code. Compare to the projectional approach, where such formatting guidelines are not needed, when arguing about the low-level code formatting, like indentation.

The textual projection of the *AST* looks similar to the text code. However the projectional nature of it has certain outcomes, which may be unusual for a programmer, who is used to editing the code as text.

The statements in the projectional editor are only selected as whole. There is now way to just select the "while" word for cut or copy, without selecting the condition and the block belonging to the statement. This behavior represents the position of the condition

and while-body in the *AST* as children of the while statement. The statement can be selected all together only, including all of its children. Alternatively, one could select just an expression in the condition part.

Every block delimiters are just a part of the block visualization. They are organized in a proper way automatically, and there is no way to delete or confuse them, as well as to type them initially. Each closing brace can marked with the parent statement name (through implementing such behavior in the *AST* visualization), enhancing navigation through the displayed code.

As one can see, the textual projection of the *AST* looks almost the same, as a text code in a conventional textual editor. This can cause some confusion for the developer at first, as attempts to edit this textual visualization as a real text will sometimes fail.

Eventually, however, advantages of such visualization overwhelm the disadvantages. Among the benefits of the textual projection over text code are quicker code construction after short learning, better way to select code fragments, since not individual characters or lines, but rather *AST* nodes or groups of nodes are selected, plus, all the advantages, the projectional editing brings by itself, as discussed above.

I discuss additionally the projectional approach and some of its basic principles in the Part ??.

### 2.2 Modular Language Engineering

#### 2.2.1 Describing a Language in Projection

When building a projectional editor for a language, the language must be given as a certain description of the AST. As AST represent a graph, the nodes and edges types, as well as their possible relationships must be described<sup>1</sup>.

The nodes of the *AST* are described through giving their types. The node type in projectional editing is called a *concept*. *Concepts* are very similar to classes in object oriented programming languages. They feature inheritance, they can implement interfaces, they can have internal data, similar to member fields, and they can feature behavior, similar to member functions. The difference with classes, however, is that the member fields are not usually encapsulated.

The edges of the *AST* are not described on their own, but instead as a properties of nodes. A node can have children, or can reference other nodes. An example of child relationship, can be a condition expression of the if statement. An example of reference relationship can be a local variable usage, referencing the declaration of the local variable. Child and reference relationships can have different cardinality, with minimal border from 0 or 1, to the maximal border of 1 or N, where N stands for just "many", or several.

The cardinality itself, is not usually enough, to restrict as desirable node relationships. Special constraints can be added and checked for each relationship, which describe precisely, or provide procedure to check, the validity of a relationship being established. The

<sup>&</sup>lt;sup>1</sup>Compare this with the textual approach, where a grammar for the language must be built, which is generally speaking complex, and some times even not a possible task, which leads to the increasing parser complexity, known problem in particular in the C++ area

projectional editor must inform the user, every time, the constraint was not satisfied, so that the user has a chance to correct the code, to match a valid *AST* description.

For the user to be able to manipulate the *AST* for each node *concept* an editor has to be created. The editor defines, how a node of a given *concept* should be represented to the user, which editing operations, and how, the user may perform on the node.

The minimal set of data was described above, which has to be defined for a language, to enable the projectional editing for it.

Additionally, constraints may be refined, involving some usual for typed languages type restrictions and checks. Generators can be added to transform *ASTs* given in a language. Text generators can be defined to generate a text code from an *AST*.

Behavior can be defined for a concept, to provide some method-like functionality to it. Additionally, some user-invokable functions can be described, to perform manipulations with the *AST*.

The process of defining a modular language in the *JetBrains MPS* environment is described additionally with practical details in the Section 3.1.

#### 2.2.2 Language Modularity

As *concepts* feature inheritance, it is possible<sup>2</sup> to use a child concept at a place where a parent concept could be used.

The inheritance works over the language borders, allowing to create the child concept in a language L2, separate from the language L1, where the original parent concept has been described. Thus the language L2 can be seen as a modular extension to the language L1.

The modular *DSL* creation is discussed in [5], [8]. The language modularity, in a context of *IetBrains MPS* is described in [3], [10].

The main focus of this work is a construction of the C++ programming language on top of the *MBEDDR Project* C implementation. Thus the modularity in the language engineering plays a key role in the work.

While extending the C language of the *MBEDDR Project* with C++ specific *concepts*, all the aspects of the language description (see previous subsection) have to be extended. The newly introduced *concepts* for nodes of the *AST*, typical for C++, must inherit from some *concepts* of the original *MBEDDR Project* C language. Not only new nodes and edges are introduced, but also constraints and other language description aspects have to be made incorporating the newly introduced concepts. The practical side of the language modularity and extensibility is discussed throughout this work.

<sup>&</sup>lt;sup>2</sup>to some extent, the extensibility is described separately in one of the following chapters

### 3 Technologies in Use

The C++ programming language developed through out this Master Thesis is based on two technologies, which are introduced in this chapter. The first technology is the *JetBrains MPS* language engineering environment, which provides the core foundations and means for incremental language construction. The second technology is the *MBEDDR Project*.

### 3.1 Jetbains MPS

*JetBrains MPS* stands for Jetbrains Meta Programming System. In this *IDE*-like software it is possible to develop *DSLs*. The approach taken in the *JetBrains MPS* is rather unique, and it is considered to be advantageous in many ways, [9].

In general, the way the language is describe in *JetBrains MPS* coresponds to the way, described in the Section 2.2. Here I will describe the process in practical details, as it is crucial for understanding the practical part of this Master Thesis.

In this section I will go through a definition of one *concept*, describing the facilities, *Jet-Brains MPS* provides to support it. Each *concept* is described by several views on it.

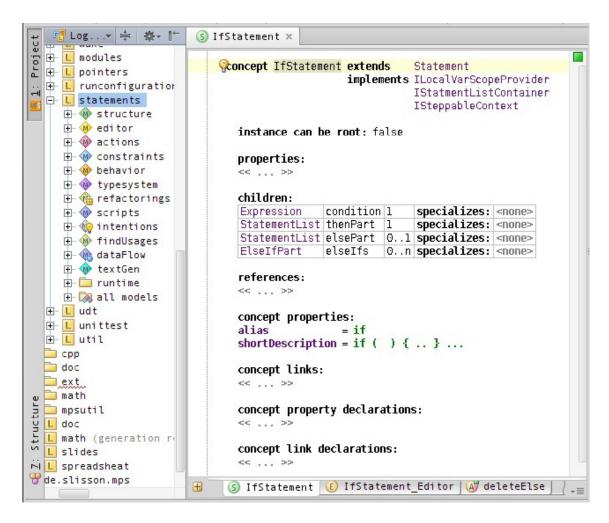


Figure 3.1: *JetBrains MPS* User Interface, if Statement Concept

Among such views there are *concept* declaration (or language structure) view, editor view, behavior view, constraints view, type system view, generator view and few more views. As a language consists mostly of the included in it concepts, the whole language is presented by the mentioned views as well, where each view on the language contains all views of the kind on each language *concept*. Left part of the Figure 3.1 demonstates the views on the *statements* language.

The representation of a new language created in concepts and views to them, present a seamless approach to creating a new language with a projectional editor for it.

The MBEDDR Project is a software, separate from JetBrains MPS but based on it, representing an extensible C language implementation with extensions. I will use concepts from the MBEDDR Project' throughout this section as examples to demonstrate MPS. Different C language parts are going to be decomposed into concepts and these concepts are going to be defined using JetBrains MPS. The reader should not confuse though the MBEDDR Project and JetBrains MPS itself: the former is a software, developed in the latter and is used to demonstrate the latter.

#### 3.1.1 Concept Declaration

*Concepts*, as it is described in the Section 2.2 represent a class-like types for nodes of an *AST*. This terminology is kept in *JetBrains MPS* and MPS concept has the same meaning as *concept* term used in the Section 2.2. I use the term "concept" both in general, to describe an *AST* node type, and in particular referring to an MPS concept.

The Figure 3.1 on the right part demonstrates a declaration of IfStatement concept from the *MBEDDR Project statements* language. It corresponds to the if statement of the C language.

At first, the *concept* is named, and the inheritance is defined. The IfStatment *concept* inherits from the Statement *concept*. This allows the IfStatement to be used at any place at which the Statement could be used, and inherits like in object oriented programming all data and behavior of the Statement.

Next, it is defined, which interfaces the *concept* is going to implement. For example, by implementing ISteppableContext, the IfStatement supports the *MBEDDR Project* debugger when stepping in the body of the IfStatement.

The "instance can be root" property defines, if it is meaningful, to create a *concept* without a parent *concept* for it. In the case of the IfStatment it does not make sense, as the statement should belong to some block. The true value can be used, e.g. for modules, as they do not have any outer concepts, and can be seen as a document in *JetBrains MPS*.

The "properties" part defines if the described *concept* instances should have some primitive type data fields (string, boolean, int). An example of a property could be a name property of a variable declaration. The IfStatement *concept* does not specify any properties, neither does it inherit any from the Statement concept.

The "children" section describe which nodes can be children on the AST of the IfStatement. Each child is assigned with a role and cardinality. For example, the IfStatement should have exactly one child of concept Expression, it has a role "condition".

The "references" section describes in the similar way as in the "children" section, which nodes could be referenced by the node of a given *concept*. Referencing can be used, to bind a given node, to a node, located somewhere else on the *AST*. As an example, a variable

usage in expression shall reference the variable declaration, to express precisely, which variable is used.

Finally, some attributes of a *concept* follow, which do not have a primary importance for the discussion here. The "alias" is used to name a *concept* in a short way, to allow for quick instantiation in the editor. The "short description" is shown to hint a programmer on the alias meaning. A *concept* can be made abstract in the "concept properties" section. Abstract *concepts* are purely used in inheritance to create other non-abstract concepts with a common parent.

#### 3.1.2 Editor View

The editor view is designed to give a look for a concept, and a way to input it. This is where the projection of an AST node is defined. As editors are mostly defined to look like text, a program in the C++ programming languageimplemented in MPS looks almost like a regular C++ code.

#### 3.1.3 Behavior View

The behavior view, can be used to define certain behavioral methods for a concept. A concept is represented there similar to a java class, and it is possible to define the methods in a Java-like language.

#### 3.1.4 Constraints View

The constraints view can be used, to limit in a desirable way relationships the concept can have to other concepts.

#### 3.1.5 Type System View

The type system view is used mainly to define a type for each instance of a concept. The type is used later (for example, in the constraints), to determine suitability of the concept instance for a place, where it is used.

#### 3.1.6 Generator View

The generator view, as well as the textgen view, is used to define the way, the concept instance is going to be transformed when generation of the AST is invoked by the MPS user. In general, the generation is needed in the end of the programming, to obtain the source code from the AST in a textual representation. This textual representation is needed because all existing tools (like compilers) expect text code nowadays. It is worth mentioning that, in theory, one could implement a compiler, which works directly from the AST in the editor, eliminating thus the need in generators to text.

- 3.1.7 Textgen View
- 3.1.8 Action View
- 3.1.9 Intentions
- 3.1.10 Other MPS Instruments
- 3.2 MBEDDR Project

# 4 Projectional C++ Implementation

# 5 Evaluation

# 6 Conclusion

# Appendix

# Glossary

**API** Application Programming Interface. 2

**AST** Abstract Syntax Tree. 2–7, 11

**concept** is a class-like type, describing a node type in the AST when talking about projectional editing. 1, 6, 7, 9, 11

**DSL** Domain Specific Language. 3, 4, 7, 9

**Embedded C++** is a language subset of the C++ programming language, intended to support embedded software development. 1

**IDE** Integrated Development Environment. 1–3, 9

**JetBrains MPS** is a language engineering environment allowing to construct incrementally defined domain specific languages. 1, 2, 7, 9–11

**MBEDDR Project** is a JetBrains MPS based language workbench, representing C language and domain specific extensions for the embedded software development. 1, 2, 7, 9, 11

**projectional approach** is an approach to create an editor for a language, when the editor is aware of the AST for the code, and shows the code to the user, projecting the AST itself, and allowing to edit the AST directly. 4

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