

Automated Verification of Concurrent Data Structures!

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Abstract

This doctoral thesis considers the automatic verification of *parameterized systems*, i.e. systems with an arbitrary number of communicating components, such as mutual exclusion protocols, cache coherence protocols or heap manipulating programs. The components may be organized in various topologies such as words, multisets, rings, or trees. The task is to show correctness regardless of the size of the system and we consider two methods to prove safety: (i) a backward reachability analysis, using the well-quasi ordered framework and monotonic abstraction, and (ii) a forward analysis which only needs to inspect a small number of components in order to show correctness of the whole system. The latter relies on an abstraction function that views the system from the perspective of a fixed number of components. The abstraction is used during the verification procedure in order to dynamically detect cut-off points beyond which the search of the state-space need not continue. Our experimentation on a variety of benchmarks demonstrate that the method is highly efficient and that it works well even for classes of systems with undecidable property. It has been, for example, successfully applied to verify a fine-grained model of Szymanski's mutual exclusion protocol. Finally, we applied the methods to solve the complex problem of verifying highly concurrent data-structures, in a challenging setting: We do not a priori bound the number of threads, the size of the data-structure, the domain of the data to store nor do we require the presence of a garbage collector. We successfully verified the concurrent Treiber's stack and Michael&Scott's queue, in the aforementioned setting. To the best of our knowledge, these verification problems have been considered challenging in the parameterized verification community and could not be carried out automatically by other existing methods.

List of papers

This thesis is based on the following papers, which are referred in the text by their roman numerals.

- I

Verification of Heap Manipulating Programs with Ordered Data by Extended Forest Automata

Parosh Aziz Abdulla, Lukáš Holík, Bengt Jonsson, Ondřej Lengál, Cong Quy Trinh, Tomáš Vojnar.

In *Automated Technology for Verification and Analysis*, 2013.

ATVA'13
- II

Automated Verification of Linearization Policies

Parosh Aziz Abdulla, Bengt Jonsson, and Cong Quy Trinh.

In *Tools and Algorithms for the Construction and Analysis of Systems*, 2016.

SAS'16
- III

Fragment Abstraction for Concurrent Shape Analysis

Parosh Aziz Abdulla, Bengt Jonsson, and Cong Quy Trinh.

In *European Symposium on Programming*, 2018.

ESOP'18

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Sammanfattning på Svenska



Du har säkerligen använt en dator och varit väldigt nöjd, och detta för att den förenklade många långa och svåra uppgifter. Ibland gör den dock inte som du vill: du klickar på knappen “Gör det”, men den följer inte ditt kommando. Eller ännu värre, datorn hänger upp sig och svarar inte på något kommando alls. Viktiga arbetstider har kanske gått till spillo. Du tar ett djupt andetag, startar om datorn, och allting fungerar igen som det ska. Felet verkar komma från *programvaran* som styr datorn, inte från själva maskinen. Du undrar varför detta fel inte redan korrigerats, och dessutom, varför det från början inte *försäkrades* om att felet inte kunde inträffa.

Termen *buggar* används generellt för att beskriva alla typer av fel, vare sig felet finns i dem elektriska delarna av maskinen eller i programmeringen. Med tanke på hur komplexa dagens datorer är, är det inte förvånande att de innehåller många buggar. De behöver ju hantera en stor mängd parametrar och egenskaper som kan generera många olika möjliga beteenden. Och vi nämner inte ens den mänskliga faktorn som inför fel under programmeringsfasen! Därför finns det en trend att bygga datorer med flera mindre enheter, som då gör datorn enklare att hantera. Men då ställs vi inför ett nytt problem: dessa enheter kan när som helst kommunicera med varandra. På grund av denna oförutsägbarhet blir det väldigt svårt att ta hänsyn till alla scenarier.

Företag som tillverkar programvara har inget intresse av att lämna buggar, eftersom ett fel någonstans kan orsaka en rad andra fel i efterhand. De inkluderar därför en viktig kvalitetskontroll under utvecklingsfasen: om det finns för många buggar blir det helt enkelt inte kostnadseffektivt eller ens möjligt att senare eliminera dem. En del buggar är dock enklare att lösa än andra. Det spelar faktiskt ingen större roll om man inte kan svara i telefonen när man får ett samtal, eller om texteditorn förlorar dem senaste dokumentsuppdateringarna. Detta är väldigt tråkigt, men ingen fara på taket! Vi klarar oss och väntar bara på programuppdateringen som löser buggen. Å andra sidan får inga fel inträffa för dem så-kallade *kritiska systemen*, vars säkerhet är viktig i allra högsta grad. Alla fel måste elimineras, antingen i programvaran eller i maskinvaran. Det är oacceptabelt, till exempel, att en pacemaker slutar fungera när ägaren passerar säkerhetskontrollen på flygplatsen, att krockkudden inte utlöses snabbt nog vid en trafikolycka om radion är på, eller om två tåg kolliderar på grund av att signalsystemet inte fungerat. Därför är det nödvändigt att utveckla metoder som upptäcker ett felaktigt beteende i sådana system.

Den vanligaste metoden är *testing*: man kör programmet med olika värden för varje variabel och kollar om funktionaliteten överensstämmer med det som förväntas. Dessa testscenarier är genomtänkta för att beskriva ett maximalt antal beteenden. När komplexiteten däremot ökar, kan programmet befinna sig i väldigt många olika tillstånd. Det börjar då bli svårt att *garantera* att metoden täcker alla möjliga programkörningar. Det är även rent av omöjligt om systemet innehåller en parameter vars värde exempelvis är en siffra. Metoden är ett effektivt sätt att upptäcka enkla buggar snabbt, men subtila buggar, såsom dem som kommer från en oförutsägbar tajming, kvarstår obemärkta. För dem kritiska systemen där säkerheten är ytterst viktig, är det otänkbart att använda en metod som eventuellt tar sig igenom alla tester, men fortfarande innehåller buggar.

För att visa att ett program verkligen uppfyller vissa egenskaper, måste vi först komma överens om vad det ska överensstämma med —vilket formellt heter en *specifikation*. Genom att lista programkonfigurationerna som borde undvikas, eller som önskas, alternativt båda, så beskriver man dem två kategorierna av specifikationer: *säkerhetsegenskaper* och *livlighetsegenskaper* (Safety och Liveness på engelska). Till exempel, “pacemakern stannar aldrig”, “krockkudden måste öppnas inom mindre än x millisekonder”, eller “ingen enstaka process blockerar övriga processer” är säkerhetsegenskaper. Vi måste se till att programmet aldrig befinner sig i en felaktig konfiguration. Till motsatsen, “brevbäraren levererar till rätt mottagare”, “systemet gör framsteg”, eller “servern hanterar en HTTP-förfråga” är livlighetsegenskaper, där man observerar rätta konfigurationer, förknippade med vissa sannolikheter. Det är upp till oss att definiera vad specifikationen ska innehålla, för att verifiera det som önskas. Om ett låsningsfritt bromssystem (ABS) beräknar hur hårt bilens bromsning ska vara men gör detta för sent, bör systemet betraktas som felaktigt. Denna avhandling fokuserar på säkerhetsegenskaper.

Säkerhet: *Givet en specifikation, kan systemet hamna i en fel konfiguration?*

Snarare än att testa programmet i specifika scenarier eller att analysera källkoden, fokuserar vi i denna avhandling på *formell verifiering*, som är ett matematiskt ramverk, för att *bevisa* att programmet uppfyller sina egenskaper.

Dessutom vill vi bevisa det helt *automatiskt*, det vill säga utan tillsyn av användaren. Vi börjar genom att extrahera en modell som motsvarar det ursprungliga programmet. I samband med det borttar vi delarna som är irrelevanta för den egenskapen som verifieras. Men vad händer när programmet manipulerar obegränsade variabler? Vi talar då om system med ett oändligt antal tillstånd, och vi behöver framställa en approximation som är i linje med själva programmet. Målet är att utforma en metod som *garanterar* att inget fel kvarstår, och som även ger svar inom en rimlig tid.

Många system med ett oändligt antal tillstånd kan faktiskt karakteriseras av en familj som består av system med ett ändligt antal tillstånd, och en parameter

(eller flera) med värde från en obegränsad domän. För varje värde på parametern, innehåller systemet ett ändligt antal tillstånd. Parametern kan exempelvis vara antalet processer som är aktiva i en viss session av ett protokoll, antalet noder i ett nät, eller hur komponenterna av ett program kommunicerar med varandra. Hur som helst, system som innehåller en preliminärt okänd parameter bör ha ett korrekt beteende oavsett värdet på parametern. De betraktas därför som system med ett oändligt antal tillstånd och kallas för parametriserade system. I denna avhandling presenterar vi två metoder för att verifiera vissa säkerhetsegenskaper hos sådana parametriserade system.

Den första metoden körs baklänges. Den startar från dem felaktiga tillstånden, och beräknar vilka andra tillstånd som skulle kunna leda till ett fel. Med andra ord upptäcker metoden alla konfigurationer som direkt eller indirekt är felaktiga. Om initialtillstånden av programmet inte tillhör dem sistnämnda, anses programmet vara korrekt. Vi måste såklart först se till att approximationen av modellen motsvarar det ursprungliga programmet.

Den andra metoden startar från initialtillstånden. Den begränsar sig till små värden av parametern, och härleder ett tröskelvärde, efter vilket metoden inte behöver fortsätta: den har faktiskt all nödvändig information för att dra slutsatsen att det inte finns några felkonfigurationer för större värden på parametern över denna tröskel. För att förenkla, bryter metoden ner varje konfiguration i små bitar av en viss storlek, och rekombinerar bitarna på alla möjliga sätt, just för att skapa konfigurationer av större storlek (litegrann som Legobitar). Tanken är att samla alla dem små bitarna och se till att ingen rekombination matchar någon felaktig konfiguration. I så fall är tröskeln hittad. Annars börjar man om med lite större bitar.

Slutligen är det intressant att undra vilka av villkoren som behövs för att metodens beräkningar inte ska fortsätta för evigt. Problemet klassas som oavgörbart, det vill säga att det inte finns någon generell metod som kan lösa *alla* instanser av problemet. Däremot har dessa två metoder trots allt visat sig vara väldigt effektiva. Där andra metoder tog timmar, kunde dessa två metoder faktiskt verifiera vissa program inom sekunder, vilket var målet. I vissa fall kan man även ge garanti på att antalet beräkningar är begränsad.

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How to read this thesis

This document is a comprehensive summary.

We first describe the domain of formal verification, where this thesis belongs, and approach the problem in a top down fashion.

We present two techniques in Chapter ?? and ?? to prove safety properties for a wide range of programs (listed in Chapter ??). Finally, we dedicate Chapter ?? to the problem of shape analysis. It deals with a class of programs that manipulate memory heaps concurrently. To finish, we give some conclusions and potential directions for future research topics.

1. Introduction

Computers have been used for a variety of applications in business, science, education, engineering and so on. They help to solve real-world problems that would otherwise be slow, impossible or extremely difficult to address without computers and software. However, sometimes they do not behave exactly as we expect them. In many cases, the consequences could be very serious, for example when errors in banking or flight control software result in unexpected behaviours. Errors in computer systems are mostly not caused by the machine itself, but typically originate from the software that controls the computer systems, so-called bugs. Bugs are quite common in complex software systems since they typically have complicated input and involve many features, which makes them difficult to design and make them perfect by human effort. Detecting and fixing software bugs are important tasks in software development process. Remaining undetected bugs in any software project may lead huge problems. They can be very hard to detect and correct, especially if they are discovered after the software has been delivered. Therefore, it is very important to allocate sufficient resources, both in terms of time and manpower, to ensure that developed software is as free of bugs as possible.

Some bugs are less serious than others. Some types of software, e.g., in user interfaces or entertainment software, can be useable even if it contains a small number of bugs. However, in the case of critical systems and system components such as software in libraries of programming languages, bugs can have far-reaching consequences, and must be avoided as much as possible. Some of libraries provide standard data structures such as stacks, queues, containers. Such data structures provide ways of storing and retrieving data in a way that suits the application at hand. For example, a stack allows inserting and removing elements in a particular order. Every time an element is inserted, that element is removed in reverse order of insertion. The simplest application of a stack is to reverse a word. You insert a given word to stack - letter by letter - and then remove letters from the stack. By using data structures data can be easily, and efficiently exchanged; it allows portability, comprehensibility, and adaptability of information.

A data structure can be both sequential or concurrent which is tricky and difficult to get correct. Concurrent data structures can be accessed and manipulated concurrently by many parallel threads are a central component of many parallel software applications. A data structure should ideally provide a simple interface to the software that uses it. An interface provides the set of operations with specifications about their types of arguments and returned

values. Data structures typically use heap-allocated memory to store their data. For example, the concurrent linked queue in `java.util.concurrent` uses a singly linked list to organize their data. The data structure may be quite complex like skiplists and binary trees which are used to implement sets.

The predominant method to ensure software quality is *testing*. It is a dynamic analysis where a program is executed under specific conditions, in so-called test cases, while checking whether the result for a given input matches the expected output. The test cases should be carefully designed to cover as many as possible cases of program executions. However, it is infeasible to cover all possible executions. Therefore, it is said by Edsger W. Dijkstra that testing can be used to show the presence of bugs, but never to show their absence. It would be nice to have techniques for checking that all executions conform to the interface of a data structure. A possibility is to use formal techniques which is the approach used in this thesis.

1.1 Formal Verification

Formal verification uses mathematical methods to check whether a program, or piece of software satisfies its specification. There are several approaches to formal verification, including equivalence checking, theorem proving, and model checking. Equivalence checking method decides whether a system is equivalent to its specification with respect to some notion of equivalence. In industry, this is mostly used for hardware designs. Theorem proving is a technique where both the behavior of the system and its desired properties are expressed in mathematical logic. Then, theorem proving, typically assisted by an interactive theorem prover, will try to prove that the system satisfies these properties.

Model checking takes as input a model of the program under consideration and a formal specification of the a property to be verified as inputs. The specification of a software component may consist of a number of such properties, each of which can be verified using model checking. The approach exhaustively explores all possible executions of the model. This is typically done exploring the set of reachable states of the model which can be finite or infinite. This works well if the set of reachable states is finite which typically happens for embedded controllers and hardware design. However, most softwares are infinite-state, e.g., a data structure may contain an unbounded amount of data. A common technique for handling this is to devise a symbolic representation of sets of states, such that a single symbolic representation represents an infinite set of states. However, it is difficult to find a suitable symbolic representation for data structures. In particular, complicated data structures such as trees and skip-lists where relationship between heap cells are complicated in both reachability and data aspects.

1.2 Research Challenges

Our challenge of this thesis is to develop techniques for automated verification of both sequential and concurrent data structures using dynamically heap-allocated memory. This requires to address several challenges in model checking:

- **Dynamically heap-allocated memory:** Data structures typically use dynamically heap allocated memory. In each cell of a heap, the domain of data values can be unbounded. In the area of formal verification, exist several approaches for heaps and for data, but not for combining them in suitable ways. In this thesis, we provide automate verification approach to sequential data structures where correctness depends on relationships between data values that are stored in the dynamically allocated structures. Such ordering relations on data are central for the operation of many data structures such as search trees, priority queues (based, e.g., on skip lists), key-value stores, or for the correctness of programs that perform sorting and searching. There exist many automated verification techniques dealing with these data structures, but only few of them can automatically reason about data properties. However, they are often limited to specific classes of structures mostly singly-linked lists (SLLs). Our approach is based on the notion of forest automata which has previously been developed for representing sets of reachable configurations of programs with complex dynamic linked data structures.
- **Unbounded number of threads:** For the case of concurrent data structures. We have to verify that the data structures are correct with any number of threads that access and manipulate the structures. We handle this challenge by extending the successful thread-modular approach which verifies a concurrent program by generating an invariant that correlates the global state with the local state of an arbitrary thread. By doing thread modular, we only verify each thread separately using an automatically inferred environment assumption that abstracts the possible steps of other threads.
- **Specification of correctness:** To ensure that a concurrent data structure is correct, we have to specify a correctness criterion that relates the concurrent interface to the interface of a corresponding sequential data structure. One way to capture such a correctness criterion is linearizability. Linearizability is generally accepted as the standard correctness criterion for such concurrent data structure implementations. Intuitively, it states that each operation on the concurrent data structure can be viewed as being performed atomically at some point (called linearization point (LP)) between its invocation and return. Existing approaches lack generality as they are limited to specific classes of concurrent data structures based on simple heaps such as singly linked lists, so far no technique (manual or automatic) for proving linearizability has been proposed that

is both sound and generic. In this thesis we provide a technique to specify linearizability of concurrent data structures.

- **Unbounded number of pointers:** In some data structures, the each cell can have unbound number of pointer fields, such as cells in skip-lists and arrays of lists. It is difficult to provide symbolic representation for these data structures. There are no techniques that have been applied to automatically verify concurrent algorithms that operate on such data structures. We propose a technique called *fragment abstraction* in which a heap is divided into small pieces called fragments. A fragment is an abstraction of a pair of heap cells that are connected by a pointer field. Our approach is general and precise enough to verify these complicated data structures.

The following sections are organized as follow: in Section 2, we present the general background about model checking, then in Section 3, we describe data structures. Thereafter, in Section 4, we describe how to specify linearizability. Our heap abstraction techniques are described in Section 5. Finally, in Section 6, we summarize and give future plans for our work.

2. Model Checking

The approach that we focus on this thesis is *model-checking*. This approach was introduced by Emerson and Clarke [39] and by Queille and Sifakis [101]. Model checking aims to check whether a model of a program satisfies a given specification. The method then computes and returns either "correct" when the specification is satisfied by the program, or "incorrect" when the program does not satisfy its specification. In the case of incorrect answer, the method can explain the reason by giving a counter-example. Models are typically transition systems consisting of states and transitions between states. A state in the model contains relevant information about the program. Alongside all the states of the system, the model depicts the transitions, i.e. how to move from one state to another state. Every behaviour of the system is represented as a succession of transitions, starting from some initial state. The number of states and transitions can be finite or infinite. A specification consists of properties of behaviors (i.e., of sequences of states), that one usually distinguishes safety properties ("something bad must never happen") from liveness properties ("something good must eventually happen"), and that most properties are safety properties. Then say that a model checking algorithm explores the set of states in a clever way, and explain the word "state-space". Model-checking aims to explore the state-space entirely from some initial states. One of the main problems with model checking is that the state-space is typically very large. However, when the state-space is of large size. It grows in-fact exponentially with the number of parameters or the size of their domain. Therefore, there have been several methods to address the state-space explosion problem.

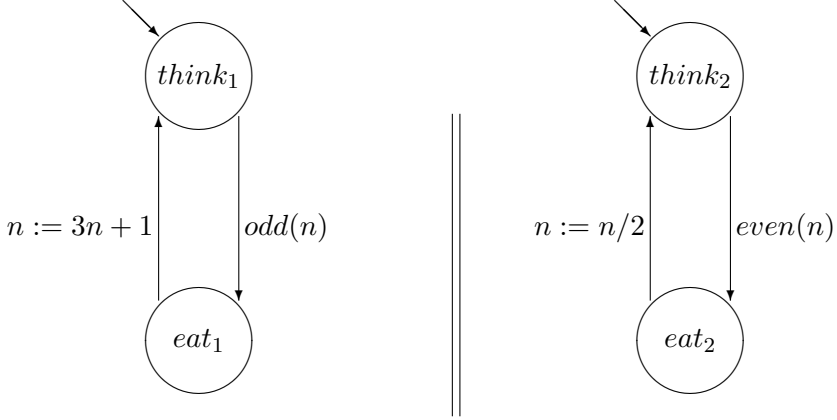
There are several techniques addressing the state-space explosion problem. One important approach is *Partial order* techniques, which aim at detecting and avoiding redundant situations, while retaining important dependencies among actions. This approach is based on the fact that, in some cases, exploring all orderings of events is not necessary because some states can be re-visited. The main approach to solve the state-space explosion problem is to use a *symbolic representation*. It avoids representing concretely all states of the system. Its main idea is to design a symbolic representation of sets of states. This designing process is done by *dropping* irrelevant details based on properties that we want to verify. Symbolic representations are of crucial help to combat the state-space explosion, accelerate the algorithms and get them to terminate in a reasonable amount of time.

Example:

Action System *Dining Mathematicians*

declare $n : \text{integer}$

initially $n \geq 1$



end

The above example consists of two processes, which share an integer variable n . The system has an infinite number of initial states, and the number of states is therefore infinite. Suppose we want to check that the two mathematicians are not eating at the same time, i.e., we want to check whether

$$\neg(@\text{eat}_1 \wedge @ \text{eat}_2)$$

is invariant. Simple inspection shows that this invariant indeed holds for the system, for the simple reason that when the left process is in eat_1 , then n is odd, and when the right process is in eat_2 , then n is even. However, this simple fact cannot be detected by naive reachability analysis, since the system has infinitely many states.

As an example, consider the dining mathematicians which is a protocol for implementing the classical mutual exclusion problem: There are two mathematicians living at the same place, whose life is focused on two activities, namely thinking (think_1 and think_2 , respectively) and eating (eat_1 and eat_2 , respectively). They do not want to eat at the same time. To ensure this, they agreed to have access to a common integer variable n . If value of n is even, the first mathematician is allowed to start eating. After finishing, the first mathematician sets $(n/2)$ to n . The second mathematician then check if n is odd, and the value he set set back after eating is $(3n+1)$. We want to verify that two mathematicians do not eat at the same time, this specification is stated by a safety property which says that the two mathematicians are not simultaneously in states eat_1 and eat_2 . if we keep all the information of the system,

then, a state represents the control state of each mathematician, and the value of the n . Therefore the number of states depends on the value of n . However, it is obvious that n is not fully needed to verify the property, so we could present it symbolically where we use two boolean predicates $\text{even}(n)$ and $\text{odd}(n)$ to represent whether n is even or odd. The dining mathematicians example is of course simple and does not reflect the complexity of today's software.

Finding a right symbolic representation is challenging, it might introduce behaviours that could turn out to be bad. Indeed, the method would return that the property is not satisfied and we would not know whether it comes from the approximation introduced by the symbolic representation or from the concrete system itself. For example in the dining mathematicians problem. If we would start with an over-approximation that ignores the variable n , we get a false positive, and we must then refine the symbolic representation to avoid it.

To deal with the imprecision caused by a too coarse over-approximation, it is possible to analyse the returned counter-example and find the origin of the problem. If it turns out to be a real concrete example, the method has in fact found a bug, and the property is surely not satisfied. Otherwise, the counter-example comes from the approximation, that is, there is a step in the sequence of events leading to that counter-example which is not performed by the original system but only by the abstract model. The approximation is refined by discarding this step and the method should be run anew.

Nevertheless, finding suitable over-approximations is a challenge on its own. This thesis now revolves around the problems of unboundedness including dynamically heap-allocated memory, unbounded number of threads and unbounded number of pointers which are described in previous research challenges section.

3. Data Structures

In general, a data structure is any data representation and its associated operations. Even an integer or floating point number stored on the computer can be viewed as a simple data structure. Typically, a data structure is meant to be an organization or structuring for a collection of data items. Each data structure has an interface which defines a set of possible values and a set of operations. More precisely, the interface consists of a set of operations or methods, each having a number of input and output parameters, and a specification of the effect of each operation.

For example, a sequential set is an data structure for storing a collection of elements, with the three operations as following:

- `add(e)` adds element `e` into the set, returning `true` if, and only if `e` was not already there.
- `remove(e)` removes element `e` from the set, returning `true` if, and only if `e` was there.
- `contains(e)` checks the existence of element `e` in the set, returns `true` if, and only if the set contains `e`.

For each method, we say that a call is successful if it returns `true`, and unsuccessful otherwise. It is typical that in applications using sets, there are significantly more `contains()` calls than `add()` or `remove()` calls. A set is implemented as a singly linked list of cells. Each cell has two fields. The `val` field present value of a cell. Cells are sorted according to `val` order, providing an efficient way to detect when an item is absent. The `next` field is a reference to the next cell in the list. The list has two sentinel cells, called `head` and `tail`, which are first and last list elements. Sentinel nodes are never added, removed, or searched for, and their values are the minimum and maximum integer values.

The implementation of a data structure should provide an efficient way to store data in computer memory and perform its operations in an efficient way. Data structures typically used heap-allocated memory to store their data. Various schemes can be used to organize the heap-allocated memory, such as singly-linked lists, doubly-linked lists, skip-lists, trees. The implementation can be sequential or concurrent.

A concurrent data structure is a way of storing and organizing data for access and manipulation by multiple computing threads (or processes) on a shared-memory computer. Each operation is implemented as a sequential method that is executed by a thread. Several features of shared-memory multiprocessors make concurrent data structures significantly more difficult to design and to verify as correct than their sequential counterparts. The primary source of

this additional difficulty is concurrency: because threads are executed concurrently possibly on different processors, and are subject to operating system scheduling decisions, interrupts, etc., we must think of the interaction between threads as completely asynchronous, so that the steps of different threads can be interleaved arbitrarily.

There are several techniques to construct concurrent data structures including coarse-grained locking, fine-grained locking, and lock-free programming. The simplest technique is coarse-grained locking, where a single lock is used to synchronize every access to an object. Coarse-grained locking is easy to reason about, however it works well only when the level of concurrency is low. However, if too many threads try to access an object at the same time, then the object becomes a sequential bottleneck, forcing threads to wait in line for access. Therefore, Fine-grained synchronization techniques address this problem by splitting the object into independently synchronized components, ensuring that method calls interfere only when trying to access the same component at the same time. Fine-grained locking requires very careful design of the data structure and its methods, since one must foresee what can happen when several threads access the same component in parallel. Fine-grained synchronization is often performed without locks, replacing them by less costly synchronization operations such as `compareAndSet()`. Each of these techniques can be applied (with appropriate customization) to a variety of common data structures (queues, stacks, sets) implemented by different linked data structures such as singly linked lists, skiplists, trees, or lists of lists.

As an example, Fig. 3.1 depicts a program `Lazy Set` [65] that implements a concurrent set containing integer elements with three operations `add`, `remove` and `contains`. It is just as the sequential version, but that each cell now has two additional fields `mark`, `lock`. The field `mark` is `true` if the node has been logically removed from the set. The `lock` field is a lock and the field `val` presents the data value which is integer in this case. The mechanism behind logically and physical removing is explained as following: it is impossible to atomically remove a cell from the list if other threads may concurrently access the adjacent cells. One reason is that one must both move a `next` pointer which reference to the cell and physically remove the cell. This cannot be done, e.g., if another thread currently is visiting the cell that is to be removed. Therefore, the task of removing a cell from the list can be split into two phases: the cell is logically removed simply by setting a `mark` field to be `true`, and later, the cell can be physically deleted by unlinking it from the rest of the data structure. The removal “actually happens” when an entry is marked, and the physical removal is just a way to clean up. The algorithm uses two global pointers, `head` that points to the first cell of the heap, and `tail` that points to the last cell. These two cells contain two values that are smaller and larger respectively than data

<pre> struct Node { bool lock; int val; Node* next; bool mark; } </pre>	<pre> locate(e): local p, c 1 while (true) 2 p := Head; 3 c := p.next; 4 while (c.val < e) 5 p := c; 6 c := c.next 7 lock(p); lock(c); 8 if (! p.mark && ! c.mark && p.next=c) 9 return (p,c); 10 else 11 unlock(p); 12 unlock(c); </pre>	<pre> add(e): local p, c, n, r 1 (p,c) := locate(e); 2 if (c.val <= e) • 3 n := new Node(0,e,c,false); 4 p.next := n; • 5 r := true; 6 else r := false; 7 unlock(p); unlock(c); 9 return r; </pre>
<pre> ctn(e): local c 1 c := Head; 2 while (c.val < e) 3 c := c.next 4 b := c.mark • 5 if (!b && c.val = e) 6 return true; 7 else 8 return false; </pre>	<pre> rmv(e): local p, c, n, r 1 (p,c) := locate(e); 2 if (c.val = e) • 3 c.mark := true; • 4 n := c.next; 5 p.next := n; 6 r := true; 7 else r := false; 8 unlock(p); unlock(c); 9 return r; </pre>	

Figure 3.1. Lazy Set Algorithm

values of all cells that may be inserted in the set. The algorithm also contains the subroutine `locate` that returns a structure containing the cells on either side of `e`. In more detail, the `locate` method traverses the list using two local variables `p` and `c`, starting at the head of the list and moves forward the list (line 2), comparing `c.val` to `e`. When `c` is set to the first cell whose the value of `val` is greater than or equal to `e`, the traversal stops, and the method locks cells pointed to by `p` and `c` (line 7) so that no other thread can update fields of `p` and `c`. Thereafter, if both `p.mark` and `c.mark` are `false` and `p.next = c` meaning that there is no added cell from other thread between `p` and `c` then the method returns the pair `(p, c)` (line 9). Otherwise, it unlocks cells pointed to by `p` and `c` and tries traversing again from the head of the list.

The `add(e)` method calls `locate(e)` at line 1 to locate the position in the list where `e` is to be inserted. Its local variables `p` and `c` are assigned the first and second values of the pair return by `locate(e)` respectively. If `c.val = e` meaning that a cell whose data value of `val` is equal to `e` is already in the list, the method unlocks `p` and `c` and returns `false` (line 7-8). Otherwise, a new cell `n` is created (line 3), and inserted into the list by linking it into the list between the `p` and `c` pointers returned by `locate` (line 3-4). Then, the method unlocks cells pointed to by `p` and `c` and returns `true`.

The `rmv(e)` method also calls `locate` at line 1 locate the position in the list where `e` is to be inserted. If `c.val \neq e` meaning that a cell with `val e` is not already in the list, the method unlocks `p` and `c` and returns `false`(line

7-9). Otherwise, cell *c* is logically removed (line 3) where the *mark* field of *c* gets assigned *true*, and unlinked from the list (line 4-5). Then, the method unlocks cells pointed to by *p* and *c* and returns *true*.

The *ctn(e)* method traverses the list by using local variable *c*, ignoring whether nodes are marked or not, until *c* is set to the first cell with a value of *val* greater than or equal to *e*. It simply returns *true* if and only if the cell pointed to by *c* is unmarked with the desired value of *val* equal to *e*.

3.1 Linearizability

In a concurrent program, the methods of the different executing threads can overlap in time. Thus, a *rmv* method that executes in parallel with an *add* method for the same key may or may not find the element in the set, depending on how the individual method statements overlap in time. For a user of the data structure, it is important to know precisely what can happen when several methods access a data structure concurrently without inspecting the code of each method. Such a user would want to have a criterion for how operations take effect, which considers only the points in time of method calls and returns. The most widely accepted such condition is linearizability. Linearizability defines consistency for the history of call and response events generated by an execution of the program at hand [69]. Intuitively, linearizability requires every method to take effect at some point (*linearization point*) between its call and return events. A linearization point is intuitively a moment where the effect of the method becomes visible to other threads. An execution of a (concurrent) system is modeled by a (concurrent) history, which is a finite sequence of method invocation and response events. A (concurrent) history is linearizable if and only if there is some order for the effects of the actions that corresponds to a valid sequential history. The valid sequence history can be generated by an execution of the sequential specification object. A concurrent object is linearizable iff each of its histories is linearizable.

Figure 3.2 provides a examples of trace of methods of concurrent program implementing sets. In the trace, each method takes effect instantaneously at its (called the *linearization point*) between call and return events [69]. When we order methods according to its linearization point, we get a total ordered sequence that respect the sequential specification of the set.

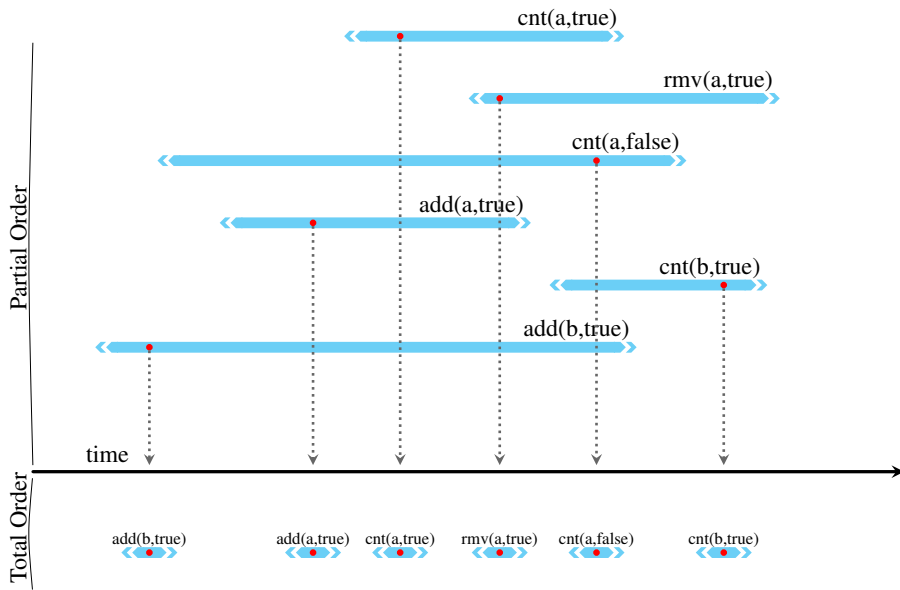


Figure 3.2. Linearizability, where the commit points are marked with `.`

4. Specifying Linearizability

In the previous sections, we describe the correctness criterion of linearizability for concurrent data structures. The goal is to specify linearizability in a way that is suitable for automated verification. We separate the problem of specifying linearizability into several ones.

- To specify the sequential semantics of a data structure in a way that is suitable for automated verification.
- To specify placement of linearization points in executions of a concurrent data structure.

For the first, we use the techniques of observers [3]. For the second, we present a new technique, in which methods are equipped with controllers. Controllers specify so-called “linearization policies”, which prescribe how LPs are placed in executions.

4.1 Observers

Data structures are, by nature, infinite-state objects, since they are intended to carry an unbounded number of data elements. For automated verification, it is desirable with specifications that are constructed without explicitly mentioning such infinite objects. This problem is addressed by observers [3]. Observers specify allowed sequences of operations by constraining their projection on a small number of data elements. Observers are finite automata extended with a finite set of observer *registers* that assume values in integer domain. At initialization, the registers are nondeterministically assigned arbitrary values, which never change during a run of the observer. Transitions are labeled by opera-

tions that may be parameterized on registers. Observers are used as acceptors of sequences of operations. The observer processes such sequences one operation at a time. If there is a transition, whose label, after replacing registers by their values, matches the operation, such a transition is performed. If there is no such transition, the observer remains in its current state. The observer accepts a sequence if it can be processed in such a way that an accepting state is reached. We use observers to give exact specifications of the behaviors of

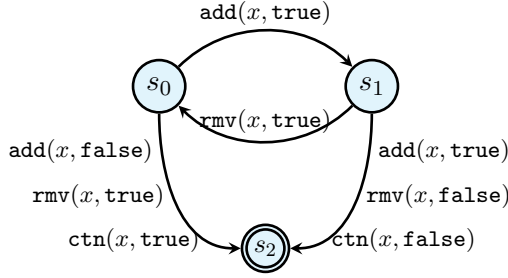


Figure 4.1. A stack observer

data structures such as sets, queues, and stacks. The observer is defined in such a way that it accepts precisely those sequences of abstract operations that are not allowed by the semantics of the data structure. This is best illustrated by an example. Fig. 3 depicts an observer that accepts the set of method invocations that are not allowed by sequential specification of a set. The observer have three states s_0 , s_1 and s_2 . The initial state s_0 corresponds to positions in the runs where the non-deterministically tracked value stored in the observer variable z is not present in the set (i.e. each time it has been inserted it got deleted afterwards). The state s_1 corresponds to positions in the runs where the tracked value is present in the set (i.e. it has not been deleted since it was last inserted). The accepting state s_2 corresponds to positions in the runs where the bad specification captured by the observer has been observed. The captured bad specification are those where a data value is deleted or found although it is not present in the set, or a data value is not found or cannot be deleted although it is already present in the set

4.2 Linearization Policies

In order to prove linearizability, the most intuitive approach is to find a linearization point (LP) in the code of the implementation, and show that it is the single point where the effect of the operation takes place [3, 22, 120]. However, for a large class of linearizable implementations, it is not possible to assign fixed LPs in the code of their methods, but depend on actions of other threads in each particular execution. For example, in the `Lazy Set` algorithm, a successful `rmv(e)` method has its LP at line 3, and an unsuccessful `rmv(e)` has its LP at line 2 when the test `c.val = e` evaluates to `false`. A successful `add(e)` method has its LP at line 4 and an unsuccessful `add` has its LP at line 2 when the test `c.val <> e` evaluates to `false`. The successful `ctn` is linearized at line 4 then the value of `b` is `true`. However it is not possible to assign a fixed LP in the code of the `ctn` method.

To see why unsuccessful `ctn` method invocations may not have fixed LPs, note that the naive attempt of defining the LP at line 4 provided that the test at

add

ρ_1 : **when** • **provided** $pc=4$ **emit** $add(e, true)$ **broadcast** $add(e)$
 ρ_2 : **when** • **provided** $pc=2 \ \&\& \ (c.val = e)$ **emit** $add(e, false)$

rmv

ρ_3 : **when** • **provided** $pc=3$ **emit** $rmv(e, true)$
 ρ_4 : **when** • **provided** $pc=2 \ \&\& \ c.val <> e$ **emit** $rmv(e, false)$

ctn

ρ_5 : **when** • **provided** $pc=4 \ \&\& \ !b \ \&\& \ c.val = e$ **emit** $ctn(e, true)$
 ρ_6 : **from** q_0 **when** • **provided** $pc=4 \ \&\& \ (!b \ \&\& \ c.val = e)$ **emit** $ctn(e, false)$ **goto** q_0
 ρ_7 : **from** q_0 **when** $(add(e), b)$ **provided** $1 \leq pc \leq 4$ **emit** $ctn(e, false)$ **goto** q_1

Figure 4.2. Reaction Rules for Controllers of Lazy Set.

line 5 fails will not work. Namely, the `ctn` method may traverse the list and arrive at line 4 in a situation where the element e is not in the list (either e is not in any cell, or the cell containing e is marked). However, before executing the command at line 4, another thread performs an `add` operation, inserting a new cell containing the element e into the list. The problem is now that the `ctn` method cannot “see” the new cell, since it is unreachable from the cell currently pointed to by the variable b of the `ctn` method. If the `ctn` method would now try to linearize an unsuccessful `ctn`, this would violate the semantics of a set, since the `add` method just linearized a successful insertion of e .

There have been several previous works dealing with the problems of non-fixed linearization points [132, 42, 41, 108, 46, 107, 121, 35]. However, they are either manual approaches without tool implementation or not strong general enough to cover various types of concurrent programs. In this thesis we handle non-fixed linearization points by providing semantic for specifying linearization policies by a mechanism for assigning LPs to executions, which we call *linearization policies*.

A linearization policy is expressed by defining for each method an associated controller, which is responsible for generating operations announcing the occurrence of LPs during each method invocation. The controller is occasionally activated, either by its thread or by another controller, and mediates the interaction of the thread with the observer as well as with other threads.

To add controllers, we first declare some statements in each method to be *triggering*: these are marked by the symbol • as in Figure 3.1. We specify the behavior of the controller, belonging to a method m , by a set of *reaction rules*. To define these rules, we first define different types of events that are used to specify their behaviors. Recall that an *operation* is of the form $m(d^{in}, d^{out})$ where m is a method name and d^{in}, d^{out} are data values. Operations are emitted by the controller to the observer to notify that the thread executing the method performs a linearization of the corresponding method with the given input and output values. Next, we fix a set Σ of *broadcast messages*, each with a fixed arity, which are used for synchronization between controllers. A message is formed by supplying data values as parameters. In reaction rules, these data

values are denoted by expressions over the variables of the method, which are evaluated in the current state when the rule is invoked. In an operation, the first parameter, denoting input, must be either a constant or the parameter of the method call.

- A *triggered* rule, of form **when** • **provided** *cond* **emit** *op* **broadcast** *se*, specifies that whenever the method executes a triggering statement and the condition *cond* evaluates to `true`, then the controller performs a *reaction* in which it emits the operation obtained by evaluating *op* to the observer, and broadcasts the message obtained by evaluating *se* to the controllers of other threads. The broadcast message *se* is optional.
- A *receiving* rule, of form **when** $\langle re, ord \rangle$ **provided** *cond* **emit** *op*, specifies that whenever the observer of some other thread broadcasts the message obtained by evaluating *re*, and *cond* evaluates to `true`, then the controller performs a reaction where it emits the operation obtained by evaluating *op* to the observer. Note that no further broadcasting is performed. The interaction of the thread with the observer may occur either *before* or *after* the sender thread, according to the flag *ord*.

A controller may also use a finite set of states, which restrict the possible sequences of reactions by a controller in the standard way. Whenever such states are used, the rule includes source and target states using keywords **from** and **goto**. In Figure 4.2, the rule ρ_7 changes the state from q_0 to q_1 , meaning that no further applications of rules ρ_6 or ρ_7 are possible, since they both start from state q_0 . Rules that do not mention states can be applied regardless of the controller state and leave it unchanged.

Let us illustrate how the reaction rules for controllers in Figure 4.2 specify LPs for the algorithm in Figure 3.1. Here, a successful `rmv` method has its LP at line 3, and an unsuccessful `rmv` has its LP at line 2 when the test `c.val = e` evaluates to `false`. Therefore, both these statements are marked as triggering. The controller has a reaction rule for each of these cases: in Figure 4.2: rule ρ_3 corresponds to a successful `rmv`, whereas rule ρ_4 corresponds to an unsuccessful `rmv`. Rule ρ_4 states that whenever the `rmv` method executes a triggering statement, from a state where `pc=2` and `c.val <> e`, then the operation `rmv(e, false)` will be emitted to the observer.

A successful `add` method has its LP at line 4. Therefore, the controller for `add` has the triggered rule ρ_1 which emits the operation `add(e, true)` to the observer. In addition, the controller also broadcasts the message `add(e)`, which is received by any controller for a `ctn` method which has not yet passed line 4, thereby linearizing an unsuccessful `ctn(e)` method by emitting `ctn(e, false)` to the observer. The keyword **b** denotes that the operation `ctn(e, false)` will be presented before `add(e, true)` to the observer. Since the reception of `add(e)` is performed in the same atomic step as the triggering statement at line 4 of the `add` method, this describes a linearization pattern, where a `ctn` method, which has not yet reached line 4, linearizes an

unsuccessful `ctn`-invocation just before some other thread linearizes a successful `add` of the same element.

To see why unsuccessful `ctn` method invocations may not have fixed LPs, note that the naive attempt of defining the LP at line 4 provided that the test at line 5 fails will not work. Namely, the `ctn` method may traverse the list and arrive at line 4 in a situation where the element e is not in the list (either e is not in any cell, or the cell containing e is marked). However, before executing the command at line 4, another thread performs an `add` operation, inserting a new cell containing the element e into the list. The problem is now that the `ctn` method cannot “see” the new cell, since it is unreachable from the cell currently pointed to by the variable `b` of the `ctn` method. If the `ctn` method would now try to linearize an unsuccessful `ctn(e)`, this would violate the semantics of a set, since the `add` method just linearized a successful insertion of e .

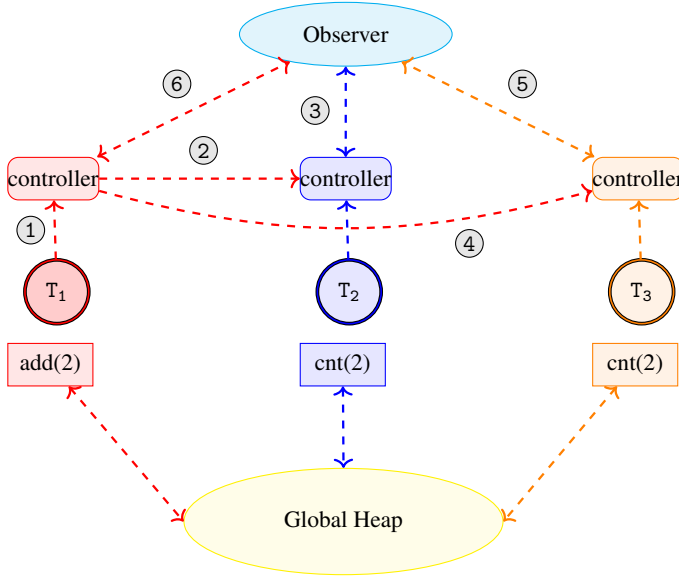


Figure 4.3. An example of linearization policies of Lazy Set algorithm

Let us describe an example of how the `controller` handles non-fixed linearization point of unsuccessful `ctn(e)` method. Figure 4.3 give an example of Lazy Set with three threads 1, 2, and 3. The thread 1 is executing the `add(e)` method to insert the cell whose value of `val` is equal to e into the set, while both 2 and 3 are executing `ctn(e)` to search for the cell whose value of `val` is equal to e in from set. When thread 1 reaches the triggering statement at line 4 of the `add(e)` method at the step 1. The controller rule ρ_1 is active. to informs the observer that an `add` operation with argument e has been performed, and that the outcome of the operation is `true` (the operation was successful) in the step 6. However, before that the controller will help

other threads by to linearize. This is done by *broadcasting* a message $\text{add}(e)$ to the threads 2, 3 which is executing $\text{ctn}(e)$ in steps 2, 3. These threads get message then the rule ρ_7 is active to inform the observer that the element e is not in the list in 3, 4 respectively.

Verifying Linearization Policies

By using an observer to specify the sequential semantics of the data structure, and defining controllers that specify the linearization policy, the verification of linearizability is reduced to establishing four conditions: (i) each method invocation generates a non-empty sequence of operations, (ii) the last operation of a method conforms to its parameters and return value, (iii) only the last operation of a method may change the state of the observer, and (iv) the sequence of all operations cannot drive the observer to an accepting state. Our verification framework automatically reduces the establishment of these conditions to a problem of checking control state reachability. This is done by augmenting the observer by a *monitor*. The monitor is automatically generated. It keeps track of the state of the observer, and records the sequence of operations and call and return actions generated by the threads. For each thread, it keeps track of whether it has linearized, whether it has caused a state change in the observer, and the parameters used in its last linearization. Using this information, it goes to an error state whenever any of the above four conditions is violated.

5. Shape Analysis

Pointers and heap-allocated storage are features of all modern imperative programming languages. They are among the most complicated features of imperative programming language: updating pointer variables (or pointer-fields of records) may have large side effects. For example, dereferencing a pointer that has been freed will lead to segmentation fault in a C or C++ program. Such side effects also make program analysis harder, because they make it difficult to compute aliasing relationships among different pointers in a program. Aliasing arises when the same memory location can be accessed using different names. For instance, consider the instruction statement $x.f := y$ in an imperative language, where x and y are pointer variables. Its effect is to assign the value of the pointer y to the cell pointed to by $x.f$. In order to update aliasing information of y . We have to require information about all the cell pointed by $x.f$ which is not an easy task.

In verification and program analysis, it is a problem to deduce and describe how the heap-allocated memory is organized. E.g., program invariants must often describe how the heap-allocated memory is structured in order to infer the effects of statements that dereference pointer fields. This is the topic of “shape analysis”.

5.1 Previous Approaches

Shape analysis is a generic term denoting static program-analysis techniques that attempt to discover and verify properties of the heap contents in (usually imperative) programs. The shape analysis problem becomes more challenging in concurrent programs that manipulate pointers and dynamically allocated objects, which are usually complicated. In concurrent programs, different threads interact in complex ways, which are difficult to foresee in the analysis. Several approaches for representing the possible structures of the heap have been proposed. TVLA (Three-Valued Logic Analysis) [104] is one of the first and one of the most popular shape analysis methods. It is based on a three-valued first-order predicate logic with transitive closure. Intuitively, concrete heap structure is represented by a finite set of abstract summary cells, each of them representing a set of concrete cells. Summary cells are obtained by merging several heap cells that agree on the values of a chosen set of unary abstraction predicates. A unique important aspect of TVLA is that it automatically generates the abstract transformers from the concrete semantics; these transformers

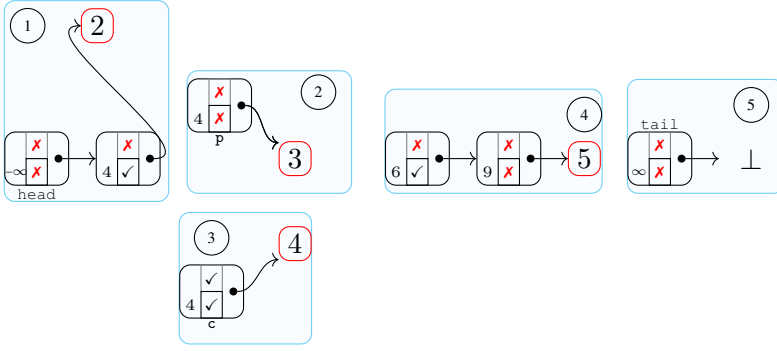


Figure 5.2. The forest representation of the heap configuration in Fig. 5.1.

or having several incoming edges. The tree components can refer to the roots of each other, and hence they are “separated” much like heaps described by formulae joined by the separating conjunction in separation logic [102]. Using this decomposition, sets of heaps with a bounded number of cut-points can then be represented by so called forest automata (FA). Each of the tree automata within the tuple accepts trees whose leaves may refer back to the roots of any of these trees. A forest automaton then represents exactly the set of heaps that may be obtained by taking a single tree from the language of each of the component tree automata and by gluing the roots of the trees with the leaves referring to them. Moreover, they allow alphabets of FA to contain nested FA, leading to a hierarchical encoding of heaps, allowing us to represent even sets of heaps with an unbounded number of cut-points (e.g., sets of DLL, skip-list).

Let us take an example of how to split a heap into small tree components. Figure 5.2 shows five tree components obtained by splitting the heap in figure 5.1. These components are named as 1,2,3,4 and 5 from left to right. Each root of a tree component is a cut-point in the heap in figure 5.1. These cut-points are cells pointed to by variables *head*, *tail*, *p*, *c* and the cell which has two incoming pointers. In each tree component, the red node show which tree component it refers. For instance, the tree component 1 refers to tree component 2, and both tree components 2 and 3 refer to tree component 4.

This forest automata approach is fully automatic and able to verify various classes of data structures, including complicated structures such as trees and skip-lists with bounded number of levels. However, the approach can not verify properties related to data values of heap cells such as sortedness in the lazy set algorithm. Therefore, in this thesis, we extend their work to verify data properties.

The last approach that we will mention is based on graph grammars describing heap graphs [64, 63]. The approach is to model heap states hypergraphs, and represent both pointer operations and abstraction mappings by hypergraph transformations. The presented approaches differ in their degree of specialisation for a particular class of data structures, their efficiency, and their level of

dependence on user assistance (such as definition of loop invariants or inductive predicates for the considered data structures).

5.2 Our Approaches

In this thesis, we propose three approaches for heap abstractions. In paper I, we propose a novel approach of extending the forest automata approach [110] by expressing relationships between data elements associated with cells of the heap by two classes of constraints.

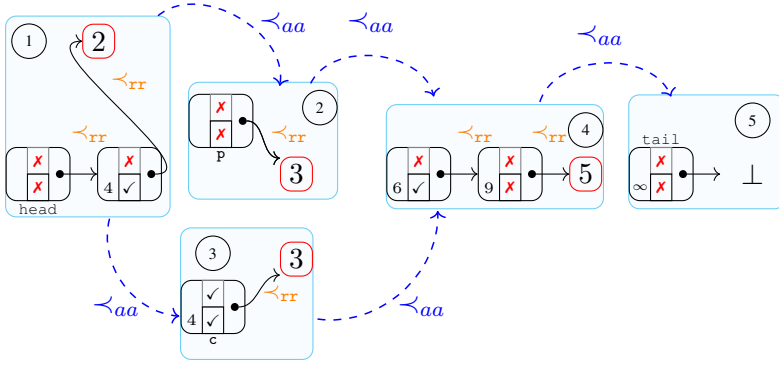


Figure 5.3. A graph and its forest representation

- Local data constraints are associated with transitions of each TA and capture relationships between data of neighboring cells in a heap; they can be used, e.g., to represent ordering internal to data structures such as sorted linked lists and binary search trees.
- Global data constraints are associated with states of TA and capture relationships between data in distant parts of the heap. Intuitively, a global data constraint between two TAs captures the data relationship between cells of two heaps accepted by these two TAs.

Figure 5.3 shows an example of how to represent a heap in 5.1 by a set of tree automata with added data constraints. In the figure, the local data constraints are located along the solid arrows between cells, whereas global constraints are located along the blue arrows. The global constraint \prec_{aa} means that data values of all cells in the left hand side are smaller than data values of all cells in the right hand side. The local constraint \prec_{rr} means that left hand side cell is smaller than data value of the cell in the right hand side. For instance, in tree component 1, the data value of the cell pointed to by head is smaller than its successor whose data value is 4, and data values of all cells in tree component 1 are smaller than data values of all cells in tree component 2. We just show

here small examples of data constraints, the detail about different types of constraints can be found in paper I.

This approach was applied to verification of sequential heap manipulation programs. This approach is general and able to verify properties of many types of sequential programs without any manual step. However, due to the complexity of tree automata operations, this approach is not optimal to handle concurrent programs where a large number of states and computation are needed. In order to verify concurrent data structures with unbounded number of threads, thread-modular is a promising approach for this challenge. Its high efficiency is achieved by abstracting the interaction between threads. The approach verifies a concurrent data structure by generating an invariant that correlates the global state with the local state of an arbitrary thread. In other words, it only keep track of the shape viewed by one thread, while abstracting away all the other threads. The thread-modular approach includes a step where it takes the information about one thread and intersect it with the information from another thread, in order to take into account the interference of all the other threads on the first thread. Forest automata approach is not suitable for this thread-modular. The reason is that computing intersection between FAs is not efficient in concurrent systems where the number of FAs is huge.

Therefore, in paper II, we adapt FA to the new setting by providing a symbolic encoding of the heap structure, that is less precise than the forest automata approach in paper I. However it is precise enough to allow the verification of the concurrent data structures, and efficient enough to make the verification procedure feasible in practice. The main idea of the abstraction is to have a more precise description of the parts of the heap that are visible (reachable) from global variables, and to make a succinct representation of the parts that are local to the threads. Intuitively, a heap segment can be characterized by a TA with data constraints. More concretely, we will extract a set of heap segments between two cut-points which are same as cut-points in the forest automata approach. For each segment, we will store a summary of the content of the heap along the segment. This summary consists of three parts, each part contains different pieces of information, including

- the values of data fields of cells along the segment which have finite values. Note that, these values do not include values of cut-points, and
- the ordering among data values of fields of cells along the segment which have integer values.
- The sequence of observer registers that appear in the segment.

Figure 5.4 gives our symbolic abstraction of the heap in figure 5.1 where the observer register z is equal to 3. In this figure, there are four segments obtained from five cut-points. In each segment, the red box contains ordering information between data values of cells along the segment, the green box contains information about the values of fields `mark`, and `lock` of cells along the segment. Finally, the blue box contains the information about the sequence of observer registers. In this example, in the first segment between the two

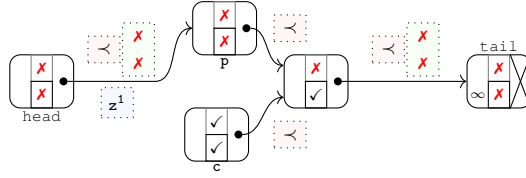


Figure 5.4. A symbolic representation of the configuration of `lazy set` in Figure 5.1

cells pointed to by `head` and `p`. The sequence of observer registers is z^1 , it means that between two cut-points there is exactly one cell whose data value of `val` field is equal to the observer register `z`.

This approach is very efficient but it is not optimal for complicated concurrent data structures like trees, lists of lists or skiplists. The approach is now specialized for SLLs.

In paper III, we present an approach which can handle concurrent programs implemented from simple to complex data structures. More precise, we can handle well-known data structures like singly-linked lists, skiplists, and lists of lists.

Heap abstraction for singly-linked list

The main idea of the approach is to represent a set of heap states by a set of fragments. A fragment represents two heap cells that are connected by a pointer field. For each of its cells, the fragment represents the contents of its non-pointer fields, together with information about how the cell can be reached from the program's pointer variables, under the applied data abstraction. The fragment contains both

- *local* information about the cell's fields and variables that point to it, as well as
- *global* information, representing how each cell in the pair can reach to and be reached from (by following a chain of pointers) a small set of globally significant heap cells.

A set of fragments represents the set of heap structures in which each pair of pointer-connected nodes is represented by some fragment in the set. Put differently, a set of fragments describes the set of heaps that can be formed by “piecing together” pairs of pointer-connected nodes that are represented by some fragment in the set. This “piecing together” must be both locally consistent (appending only fragments that agree on their common node), and globally consistent (respecting the global reachability information). Figure ?? shows a set of fragments that is sufficient to represents the configuration in

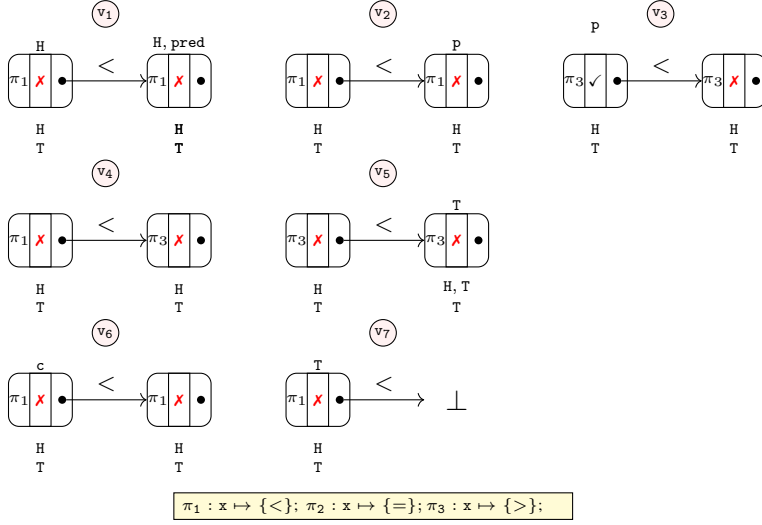


Figure 5.5. Fragment abstraction of the configuration in 5.1

Figure 5.1. There are 8 fragments, named v_1, \dots, v_8 . Fragment of v_8) consists of a tag that points to \perp . All other fragments consist of a pair of pointer-connected tags. The two left-most cells in Figure ?? are represent by the level 1-fragment v_1 in Figure 5.7. Here, the variable `preds[3]` is represented by `preds[higher]`. The mapping π_1 represents the data abstraction of the key field, here saying that it is smaller than the value 9 of the observer register. Note that, in our approach, the data abstraction is a mapping function from data value to a set of observer registers. The two left-most cells are also represented by a higher-level fragment, viz. v_6 . The pair consisting of the two sentinel cells (with keys $-\infty$ and $+\infty$) are represented by the higher-level fragment v_7 . In each fragment, the abstraction values of non-pointer fields are shown represented inside each tag of the fragment. The data relation is shown as a label on the arrow between two tags. Above each tag is reach from information, whereas the second row is reach to information.

Heap abstraction for skiplist

Let us illustrate how pairs of heap cells can be represented by fragments. Before going detail to the example, let us describe in short the definition of skiplist. A skiplist consists of a collection of sorted linked lists, each of which is located at a *level*, ranging from 1 up to a maximum value. Each skiplist node has a key value and participates in the lists at levels 1 up to its *height*. The skiplist has sentinel head and tail nodes with maximum heights and key values $-\infty$ and $+\infty$, respectively. The lowest-level list (at level 1) constitutes an ordered list of all nodes in the skiplist. Higher-level lists are increasingly sparse sublists of the lowest-level list, and serve as shortcuts into lower-level lists.

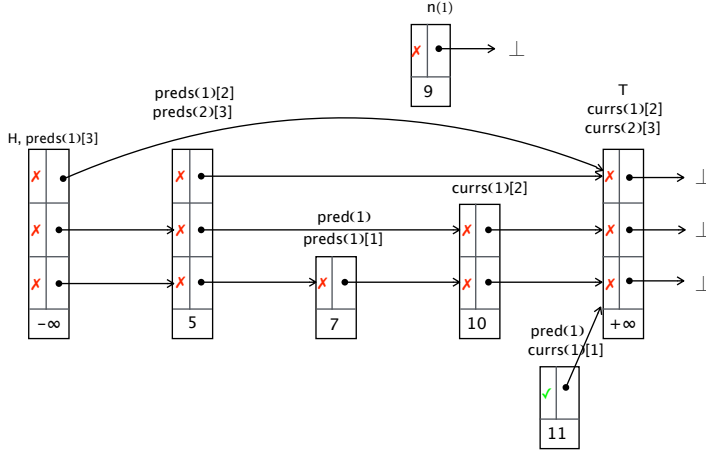


Figure 5.6. A concrete shape of 3-level skiplist

Figure 5.6 shows an example heap state of the skiplist algorithm with three levels. Each heap cell is shown with the values of its `mark` and `key` fields. Let us illustrate how pairs of heap nodes can be represented by fragments. Figure 5.7 shows a set of fragments that is sufficient to represent the configuration in Figure 5.6. There are 11 fragments, named v_1, \dots, v_{11} . Three of these (v_9 , v_{10} and v_{11}) consist of a tag that points to \perp . All other fragments consist of a pair of pointer-connected tags. The fragments $v_1, v_2, v_3, v_4, v_5, v_9$ and v_{10} are level-1-fragments, whereas the other fragments are higher level-fragments. The `private` field of the input tag of v_7 is `true`, whereas the `private` field of other tags of other fragments are `false`. The two left-most cells in Figure 5.6 are represented by the level 1-fragment v_1 in Figure 5.7. Here, the variable `preds[3]` is represented by `preds[higher]`. The mapping π_1 represents the data abstraction of the `key` field, here saying that it is smaller than the value 9 of the observer register. Note that, in our approach, the data abstraction is a mapping function from data value to a set of observer registers. The two left-most cells are also represented by a higher-level fragment, viz. v_6 . The pair consisting of the two sentinel cells (with keys $-\infty$ and $+\infty$) are represented by the higher-level fragment v_7 . In each fragment, the abstraction values of non-pointer fields are shown represented inside each tag of the fragment. The data relation is shown as a label on the arrow between two tags. Above each tag is pointer variables. The first row under each tag is reach from information, whereas the second row is reach to information.

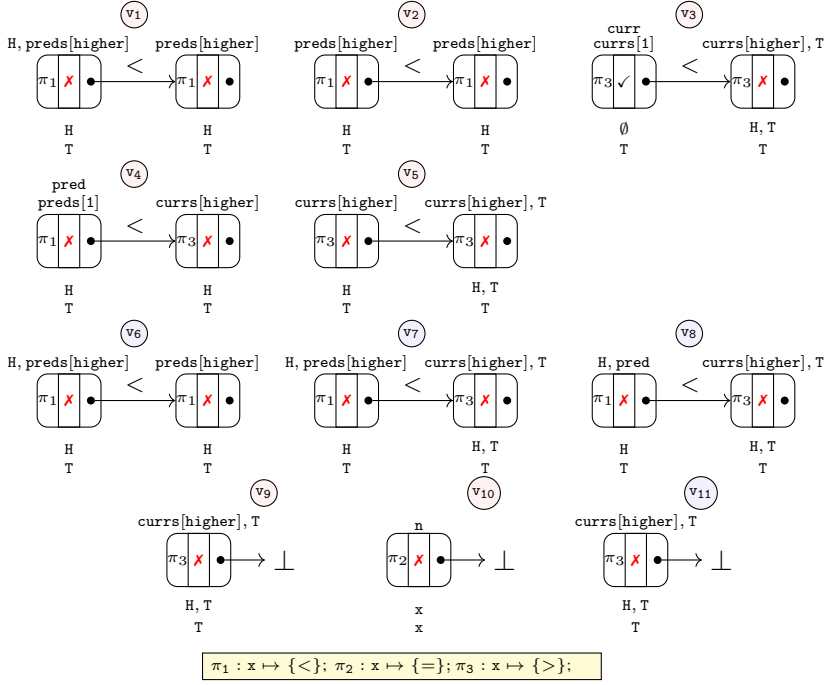


Figure 5.7. Fragment abstraction of skiplist algorithm

6. Conclusion and Future Work

We have presented, in this thesis, approaches to verify the complex problem of both sequential and concurrent heap manipulating programs. Such programs induce an infinite-state space in several dimensions: they (i) consist of an unbounded number of concurrent threads, (ii) use unbounded dynamically allocated memory, and (iii) the domain of data values is unbounded. (iv) consist of an unbounded number of pointers. In addition, the linearization points of some programs are not fixed. They are depended on the future executions of these programs. In this thesis, we focus on proving both safety properties, and linearization properties for the system, regardless of the value of this parameter. In order to prove safety properties, we define an abstract model of the program, and we employ approximation techniques to that ignore irrelevant information so that we can reduce the problem into a finite-state model. In fact, we use an over-approximation, such that the abstract model cover all the behaviors of the original system. However, it might cover other behaviors which are not in the original system. If the bad states are not reached during the computation of reachable states of the abstract model, then the abstract model is considered safe, and so is the original system is also safe. If the bad state is reachable, we have to refine the abstraction. In order to verify linearization properties of a program, we add a specification which expresses its data structure, using the technique of observers. In our approaches, the user have to provide linearization policies which specify how the program is linearized. We use a technique call `controller` to specify linearization policies. We then verify that in any concurrent execution of a collection of method calls, the sequence of announced operations satisfies the semantics of the data structure. This check is performed by an observer, which monitors the sequence of announced operations. This reduces the problem of checking linearizability to the problem of checking that in this cross-product, the observer cannot reach a state where the semantics of the set data structure has been violated. To verify that that the observer cannot reach a state where a violation is reported, we compute a symbolic representation of an invariant that is satisfied by all reachable configuration of the cross-product of a program and an observer.

There are two main possible lines of future work we would like to work on from this thesis. The first line is to extend the type programs we consider, by allowing more complicated data structures such as trees and design methods that allow the automatic synthesis of the controllers. Another possible line of work is to extend the view abstraction to multi-threaded programs running on machines with different memory models. Such hardware systems employ

store buffers and cache systems that could be modeled using views. This is an interesting challenge since it would help programmers to write their code under a given memory model that is simpler to reason around, and verify that the behaviour of the program is the same under another less-restricted memory model.

7. Summary of Contributions

In this chapter, we give a short overview of our three peer-reviewed papers. We will explain the problem addressed by each paper and its main contributions.

Automated verification of programs that manipulate complex dynamic linked data structures is a challenging problem in software verification. The problem becomes even more challenging when program correctness depends on relationships between data values that are stored in the dynamically allocated structures.

In this paper, we present a general framework for verifying these programs. The underlying formalism of our framework is that of forest automata (FA) which has previously been developed for representing sets of reachable configurations of programs with complex dynamic linked data structures without data stored in nodes. In the FA framework, a heap graph is represented as a composition of tree components. Sets of heap graphs can then be represented by tuples of tree automata (TA). We extend FA by adding constraints between data elements associated with nodes of the heaps represented by FA, and we present extended versions of all operations needed for using the extended FA in a fully-automated verification approach, based on abstract interpretation. Technically, we express relationships between data elements associated with nodes of the heap graph by two classes of constraints. Local data constraints are associated with transitions of TA and capture relationships between data of neighbouring nodes in a heap graph; they can be used, e.g., to represent ordering internal to some structure such as a binary search tree. Global data constraints are associated with states of TA and capture relationships between data in distant parts of the heap. In order to obtain a powerful analysis based on such extended FA, the entire analysis machinery must have been redesigned, including a need to develop mechanisms for propagating data constraints through FA, to develop a new inclusion check between extended FAs, and to define extended abstract transformers

The resulting method allows for verification of pointer programs where the needed inductive invariants combine complex shape properties with constraints over stored data, such as sortedness. The method is fully automatic, quite general, and its efficiency is comparable with other state-of-the-art analyses even though they handle less general classes of programs and/or are less automated.

We have implemented our approach as an extension of the Forester tool and successfully applied it to a number of programs dealing with data structures such as various forms of singly- and doubly-linked lists, binary search trees, as well as skip lists. We presented experimental results from verifying programs dealing with variants of (ordered) lists and trees. To the best of our

knowledge, our method is the first one to cope fully automatically with a full C implementation of a 3-level skip list.

7.1 Paper II: Automated Verification of Linearization Policies

Data structures that can be accessed concurrently by many parallel threads are a central component of many software applications, and are implemented in several widely used libraries (e.g., `java.util.concurrent`). Linearizability is the standard correctness criterion for such concurrent data structure implementations. It states that each operation on the data structure can be considered as being performed atom

ically. The controller is occasionally activated, either by its thread or by another controller, and mediates the interaction of the thread with the observer as well as with other threads. Secondly, we handle the challenge of an unbounded number of threads by extending the successful thread-modular approach which verifies a concurrent program by generating an invariant that correlates the global state with the local state of an arbitrary thread. Finally, we present a novel symbolic representation of singly-linked heap structures. We have implemented our technique in a tool, and applied it to specify and automatically verify linearizability of all the implementations of concurrent set, queue, and stack algorithms known to us in the literature, as well as some algorithms for implementing atomic memory read/write operations. To use the tool, the user needs to provide the code of the algorithm together with the controllers that specify linearization policies. To our knowledge, this is the first time all these examples are verified fully automatically in the same framework.

7.2 Paper III: Fragment Abstraction for Concurrent Shape Analysis

A major challenge in automated verification is to develop techniques that are able to reason about fine-grained concurrent algorithms that consist of an unbounded number of concurrent threads, which operate on an unbounded domain of data values, and use unbounded dynamically allocated memory. Existing automated techniques consider the case where shared data is organized into singly-linked lists.

In this paper, we present a technique for automatic verification of concurrent data structure implementations that operate on dynamically allocated heap structures which are more complex than just singly-linked lists. Our approach is the first framework that can automatically verify concurrent data structure

implementations that employ singly linked lists, skiplists [15, 23, 39], as well as arrays of singly linked lists [11], at the same time as handling an unbounded number of concurrent threads, an unbounded domain of data values (including timestamps), and an unbounded shared heap.

Our technique is based on a novel shape abstraction, called fragment abstraction, which in a simple and uniform way is able to represent several different classes of unbounded heap structures. Its main idea is to represent a set of heap states by a set of fragments. A fragment represents two heap cells that are connected by a pointer field. For each of its cells, the fragment represents the contents of its non-pointer fields, together with information about how the cell can be reached from the program’s pointer variables. The latter information consists of both: (i) local information, saying which pointer variables point directly to them, and (ii) global information, saying how the cell can reach to and be reached from (by following chains of pointers) other heap cells that are significant from a global perspective, typically since they are pointed to by global variables. A set of fragments represents the set of heap states in which any two pointerconnected nodes is represented by some fragment in the set. Thus, a set of fragments describes the set of heaps that can be formed by “pieced together” fragments in the set. The combination of local and global information in fragments supports precise reasoning about the sequence of cells that can be accessed by threads that traverse the heap by following pointer fields in cells and pointer variables: the local information captures properties of the cell fields that can be accessed as a thread dereferences a pointer variable or a pointer field; the global information also captures whether certain significant accesses will at all be possible by following a sequence of pointer fields. This support for reasoning about patterns of cell accesses enables automated verification of reachability and other functional properties. Fragment abstraction can (and should) be combined, in a natural way, with data abstractions for handling unbounded data domains and with thread abstractions for handling an unbounded number of threads. For the latter we adapt the successful thread-modular approach [5], which represents the local state of a single, but arbitrary thread, together with the part of the global state and heap that is accessible to that thread. Our combination of fragment abstraction, thread abstraction, and data abstraction results in a finite abstract domain, thereby guaranteeing termination of our analysis. We have implemented our approach and applied it to automatically verify correctness, in the sense of linearizability, of a large number of concurrent data structure algorithms, described in a C-like language. More specifically, we have automatically verified linearizability of most linearizable concurrent implementations of sets, stacks, and queues, and priority queues, which employ singly-linked lists, skiplists, or arrays of timestamped singly-linked lists, which are known to us in the literature on concurrent data structures. For this verification, we specify linearizability using the simple and powerful technique of observers [1, 7, 21, 3], which reduces the criterion of linearizability to a simple reachability property. To verify implementations of

stacks and queues, the application of observers can be done completely automatically without any manual steps, whereas for implementations of sets, the verification relies on light-weight user annotation of how linearization points are placed in each method.

7.3 Related Work

This chapter reviews related work, along the main topics of this thesis including: verification of linearizabilities of concurrent algorithms, sequential shape analysis and concurrent techniques.

7.3.1 Linearizabilities

Much previous work has been devoted to the *manual* verification of linearizability for concurrent programs such as [117]. In [95], O’Hearn *et al.* define a *hindsight lemma* that provides a non-constructive evidence for linearizability. The lemma is used to prove linearizability of an optimistic variant of the lazy set algorithm. Vafeiadis [119] uses forward and backward simulation relations together with history or prophecy variables to prove linearizability. These approaches are manual, and without tool implementations. *Mechanical* proofs of linearizability, using interactive theorem provers, have been reported in [42, 46, 108, 107]. For instance, Colvin *et al.* [42] verify the lazy set algorithm in PVS, using a combination of forward and backward simulations.

There are several works on *automatic* verification of linearizability. In [121], Vafeiadis develops an automatic tool for proving linearizability that employs instrumentation to verify logically pure executions. However, this work can handle non-fixed LPs only for read-only methods, i.e, methods that do not modify the heap. This means that the method cannot handle algorithms like the *Elimination* queue [91], *HSY* stack [66], *CCAS* [61], *RDCSS* [61] and *HM* set [68] that we consider in this paper. In addition, their shape abstraction is not powerful enough to handle algorithms like *Harris* set [60] and *Michael* set [89] that are also handled by our method. Chakraborty *et al.* [67] describe an “aspect-oriented” method for modular verification of concurrent queues that they use to prove linearizability of the Herlihy/Wing queue. Bouajjani *et al.* [27] extended this work to show that verifying linearizability for certain fixed abstract data types, including queues and stacks, is reducible to control-state reachability. We can incorporate this technique into our framework by a suitable construction of observers. The method can not be applied to sets. The most recent work of Zhu *et al.* [132] describe a tool that is applied for specific set, queue, and stack algorithms. For queue algorithms, their technique can handle queues with helping mechanism except for *HW* queue [69] which is handled by our paper. For set algorithms, the authors can only handle those that perform an optimistic contains (or lookup) operation by applying the *hindsight lemma*

from [95]. Hindsight-based proofs provide only *non-constructive* evidence of linearizability. Algorithms with non-optimistic contains (or lookup) operation like *HM* [68], *Harris* [60] and *Michael* [89] sets cannot be verified by their technique. Vechev *et al.* [124] check linearizability with user-specified non-fixed LPs, using a tool for finite-state verification. Their method assumes a bounded number of threads, and they report state space explosion when having more than two threads. Dragoi *et al.* [48] describe a method for proving linearizability that is applicable to algorithms with non-fixed LPs. However, their method needs to rewrite the implementation so that all operations have linearization points within the rewritten code. Černý *et al* [123] show decidability of a class of programs with a bounded number of threads operating on concurrent data structures. Finally, the works [2, 22, 118] all require fixed linearization points.

We have not found any report in the literature of a verification method that is sufficiently powerful to automatically verify the class of concurrent set implementations based on sorted and non-sorted singly-linked lists having non-optimistic contains (or lookup) operations we consider. For instance the lock-free sets of *HM* [68], *Harris* [60], or *Michael* [89], or unordered set of [131],

7.3.2 Shape Analysis

Sequential Shape Analysis

Our approach builds on the fully automated FA-based approach for shape analysis of programs with complex dynamic linked data structures [96, 17]. We significantly extend this approach by allowing it to track ordering relations between data values stored inside dynamic linked data structures.

For shape analysis, many other formalisms than FA have been used, including, e.g., separation logic and various related graph formalisms [129, 50, 36, 49], other logics [104, 73], automata [34], or graph grammars [64]. Compared with FA, these approaches typically handle less general heap structures (often restricted to various classes of lists) [129, 49], they are less automated (requiring the user to specify loop invariants or at least inductive definitions of the involved data structures) [50, 36, 49, 64], or less scalable [34].

Verification of properties depending on the ordering of data stored in SLLs was considered in [30], which translates programs with SLLs to counter automata. A subsequent analysis of these automata allows one to prove memory safety, sortedness, and termination for the original programs. The work is, however, strongly limited to SLLs. In our technique, we get inspired by the way that [30] uses for dealing with ordering relations on data, but we significantly redesign it to be able to track not only ordering between simple list segments but rather general heap shapes described by FA. In order to achieve this, we had to not only propose a suitable way of combining ordering relations with FA, but we also had to significantly modify many of the operations used over FA.

In [5], another approach for verifying data-dependent properties of programs with lists was proposed. However, even this approach is strongly limited to SLLs, and it is also much less efficient than our current approach.

Verification of properties of programs depending on the data stored in dynamic linked data structures was considered in the context of the TVLA tool [74] as well. Unlike our approach, [74] assumes a fixed set of shape predicates and uses inductive logic programming to learn predicates needed for tracking non-pointer data. The experiments presented in [74] involve verification of sorting and stability properties of several programs on SLLs (merging, reversal, bubble-sort, insert-sort) as well as insertion and deletion in BSTs. We do not handle stability, but for the other properties, our approach is much faster. Moreover, for BSTs, we verify that a node is greater/smaller than all the nodes in its left/right subtrees (not just than the immediate successors as in [74]).

An approach based on separation logic extended with constraints on the data stored inside dynamic linked data structures and capable of handling size, ordering, as well as bag properties was presented in [36]. Using the approach, various programs with SLLs, DLLs, and also AVL trees and red-black trees were verified. The approach, however, requires the user to manually provide inductive shape predicates as well as loop invariants. Later, the need to provide loop invariants was avoided in [100], but a need to manually provide inductive shape predicates remains.

Another work that targets verification of programs with dynamic linked data structures, including properties depending on the data stored in them, is [129]. It generates verification conditions in an undecidable fragment of higher-order logic and discharges them using decision procedures, first-order theorem proving, and interactive theorem proving. To generate the verification conditions, loop invariants are needed. These can either be provided manually or sometimes synthesized semi-automatically using the approach of [104]. The latter approach was successfully applied to several programs with SLLs, DLLs, trees, trees with parent pointers, and 2-level skip lists. However, for some of them, the user still had to provide some of the needed abstraction predicates. Several works, including [31], define frameworks for reasoning about pre- and post-conditions of programs with SLLs and data. Decidable fragments, which can express more complex properties on data than we consider, are identified, but the approach does not perform fully automated verification, only checking of pre-post condition pairs.

Concurrent Shape Analysis

There are several approaches for automated verification of concurrent algorithms are limited to the case of singly-linked lists [2, 62, 15, 109, 121]. Furthermore, many of these techniques impose additional restrictions on the considered verification problem, such as bounding the number of accessing threads [18, 124, 123].

In [2], concurrent programs operating on SLLs are analyzed using an adaptation of a transitive closure logic [24], combined with tracking of simple sorted-

ness properties between data elements; the approach does not allow to represent patterns observed by threads when following sequences of pointers inside the heap, and so has not been applied to concurrent set implementations. In our recent work [15], we extended this approach to handle SLL implementations of concurrent sets by adapting a well-known abstraction of singly-linked lists [?] for concurrent programs. The resulting technique is specifically tailored for singly-links. Our fragment abstraction is significantly simpler conceptually, and can therefore be adapted also for other classes of heap structures. The approach of [15] is the only one with a shape representation strong enough to verify concurrent set implementations based on sorted and non-sorted singly-linked lists having non-optimistic contains (or lookup) operations we consider, such as the lock-free sets of *HM* [68], *Harris* [60], or *Michael* [89], or unordered set of [131]. Our fragment abstraction can handle them as well as also algorithms employing skiplists and arrays of singly-linked lists.

There is no previous work on automated verification of skiplist-based concurrent algorithms. The work [105] generates verification conditions for statements in sequential skiplist implementations. All these works assume that skiplists have the well-formedness property that any higher-level lists is a sublist of any lower-level list, which is true for sequential skiplist algorithms, but false for several concurrent ones, such as [68, 81].

Concurrent algorithms based on arrays of SLLs, and including timestamps, e.g., for verifying the algorithms in [47] have shown to be rather challenging. Only recently has the TS stack been verified by non-automated techniques [32] using a non-trivial extension of forward simulation, and the TS queue been verified manually by a new technique based on partial orders [77, 111].. We have verified both these algorithms automatically using fragment abstraction,

Our fragment abstraction is related in spirit to other formalisms that abstract dynamic graph structures by defining some form of equivalence on its nodes (e.g., [126, 104]). These have been applied to verify functional correctness fine-grained concurrent algorithms for a limited number of SLL-based algorithms. Fragment abstraction's representation of both local and global information allows to extend the applicability of this class of techniques.

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