Team notebook

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Contents

1	Algo	prithms	1		
	1.1	Mo's algorithm on trees	1		
	1.2	Mo's algorithm	2		
2	2 Basics				
	2.1	default code	2		
3	Data	a structures	3		
	3.1	hash table	3		
	3.2	heavy light decomposition	3		
	3.3	persistent array	3		
	3.4	persistent seg tree	3		
	3.5	persistent trie	4		
	3.6	segment tree	4		
	3.7	sparse table	5		
	3.8	splay tree	5		
	3.9	STL order statistics tree II	6		
	3.10	STL order statistics tree	6		
	3.11	STL Treap	7		
	3.12	trie	7		
	3.13	wavelet tree	7		
4	Deletion 8				
	4.1	board	8		
	4.2	center 2 points + radious	8		
	4.3	counting	8		
	4.4	crt	8		
	4.5	cumulative sum of divisors	8		
	4.6	${\rm dates} \ . \ . \ . \ . \ . \ . \ . \ . \ . \ $	8		

	4.7	dijkstra
	4.8	eulerian path
	4.9	ext euclidean
	4.10	highest exponent factorial
	4.11	io
	4.12	Lucas theorem
	4.13	matrix
	4.14	mod integer
	4.15	mod inv
	4.16	mod mul
		mod pow
	4.18	planar graph (euler) 10
	4.19	polynomials
	4.20	primes
	4.21	query with lca
	4.22	sliding window
		squares
	4.24	totient function
		totient sieve
	4.26	totient
5	DP	Optimizations 13
	5.1	convex hull trick
	5.2	
		•
6	Geo	metry 13
7	Gra	phs 13
	7.1	bridges
	7.2	directed mst
	7.3	karp min mean cycle
	7.4	konig's theorem
	7.5	minimum path cover in DAG
	7.6	SCC kosaraju
	7.7	tarjan scc
	7.8	two sat (with kosaraju)

3	Math	16			
	8.1 fft	16			
	8.2 fibonacci properties	17			
	8.3 sigma function	17			
	8.4 triangles	17			
)	Matrix 18				
LO	Misc	18			
L1	Number theory	18			
	11.1 convolution	18			
	11.2 diophantine equations	18			
	11.3 discrete logarithm	18			
	11.4 miller rabin \dots	18			
	11.5 number theoretic transform $\dots \dots \dots$	19			
	11.6 pollard rho factorize	19			
L 2	Strings	19			
	12.1 Incremental Aho Corasick	19			
	12.2 minimal string rotation	20			
	12.3 suffix array	20			
	12.4 suffix automaton	21			
	12.5 z algorithm	21			
L3	Tmp	21			
	13.1 convexHull	21			
	13.2 seg intersection	21			
	13.3 windingNum	21			
	$1 {\bf Algorithms}$				

1.1 Mo's algorithm on trees

/**		
problem	s:	
- h	ttps://codeforces.com/gym/101161 problem E	
*/		
void fl	at(vector <vector<edge>> &g, vector<int> &a,</int></vector<edge>	
vec	tor <int> ≤, vector<int> &ri, vector<int> &cos</int></int></int>	٤t

```
int node, int pi, int &ts, int w) {
 cost[node] = w:
 le[node] = ts;
 a[ts] = node;
 ts++:
 for (auto e : g[node]) {
  if (e.to == pi) continue:
   flat(g, a, le, ri, cost, e.to, node, ts, e.w);
 ri[node] = ts:
 a[ts] = node;
 ts++:
* Case when the cost is in the edges.
void compute queries(vector<vector<edge>> &g) {
 // g is undirected
 int n = g.size();
 lca_tree.init(g, 0);
 vector<int> a(2 * n), le(n), ri(n), cost(n);
 // a: nodes in the flatten array
 // le: left id of the given node
 // ri: right id of the given node
 // cost: cost of the edge from the node to the parent
 int ts = 0; // timestamp
 flat(g, a, le, ri, cost, 0, -1, ts, 0);
 int q; cin >> q;
 vector<query> queries(q);
 for (int i = 0; i < q; i++) {
  int u, v;
   cin >> u >> v:
   u--; v--;
   int lca = lca_tree.query(u, v);
   if (le[u] > le[v])
    swap(u, v);
   queries[i].id = i;
   queries[i].lca = lca;
   queries[i].u = u;
   queries[i].v = v;
   if (lca == u) {
    queries[i].a = le[u] + 1;
     queries[i].b = le[v];
   } else {
     queries[i].a = ri[u];
     queries[i].b = le[v];
 solve_mo(queries, a, le, cost); // this is the usal algorithm
```

1.2 Mo's algorithm

```
const int MN = 5 * 100000 + 1:
const int SN = 708;
struct Query {
  int a, b, id;
  Querv() {}
  Query(int x, int y, int i) : a(x), b(y), id(i) {}
  bool operator<(const Query &o) const {</pre>
   if (a / SN != o.a / SN) return a < o.a;
   return a / SN & 1 ? b < o.b : b > o.b;
};
struct DS {
 DS() : {}
  void Insert(int x) {}
  void Erase(int x) {}
  long long Query() {}
}:
Query s[MN];
int ans[MN]:
DS active;
int main() {
  int n;
  cin >> n:
  vector<int> a(n);
  for (auto &i : a) cin >> i:
  int q;
  cin >> q;
  for (int i = 0: i < a: ++i) {
   int b, e;
    cin >> b >> e;
   b--:
   s[i] = Querv(b, e, i):
  sort(s, s + q);
  int i = 0:
  int j = -1;
  for (int k = 0; k < (int)q; ++k) {
   int L = s[k].a;
    int R = s[k].b;
    while (i < R) active.Insert(a[++i]):</pre>
    while (j > R) active.Erase(a[j--]);
    while (i < L) active.Erase(a[i++]);</pre>
    while (i > L) active.Insert(a[--i]):
```

```
ans[s[k].id] = active.Query();
}
for (int i = 0; i < q; ++i) {
   cout << ans[i] << endl;
}
return 0;
};</pre>
```

2 Basics

2.1 default code

```
#include<bits/stdc++.h>
using namespace std;
#define endl '\n'
#define pb emplace back
#define X first
#define Y second
#define SZ(a) ((int)a.size())
#define ALL(x) x.begin(), x.end()
#define CLR(x, y) memset(x, y, sizeof(x))
#define IOS ios::sync_with_stdio(false); cin.tie(nullptr)
#define rep(i, begin, end) for (__typeof(end) i = (begin) -
     ((begin) > (end)); i != (end) - ((begin) > (end)); i +=
     (begin > end ? -1 : 1))
#define debug(args...) { string _s = #args; replace(_s.begin(),
     _s.end(), ',', '); stringstream _ss(_s);
     istream_iterator<string> _it(_ss); err(_it, args); }
void err(istream_iterator<string> it) {}
template<typename T, typename... Args>
void err(istream_iterator<string> it, T a, Args... args) {
       cerr << *it << " = " << a << endl:
       err(++it, args...);
}
using ll = long long;
using vi = vector <int>;
using vii = vector <vi>;
using pii = pair <int, int>;
using pll = pair <11 , 11 >;
const int MOD = 1000000007;
const int INF = INT_MAX;
signed main () {
       IOS:
       return 0;
```

B Data structures

3.1 hash table

```
/*

* Micro hash table, can be used as a set. Very efficient vs std::set

*/

const int MN = 1001;

struct ht {
   int _s[(MN + 10) >> 5];
   int len;
   void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
   }

bool is_set(int id) {
   return _s[id >> 5] & (1LL << (id & 31));
   }
};
```

3.2 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
 vector<int> g[MAXN], c[MAXN];
 int s[MAXN]; // subtree size
 int p[MAXN]; // parent id
 int r[MAXN]; // chain root id
 int t[MAXN]; // index used in segtree/bit/...
 int d[MAXN]; // depth
 int ts:
 void dfs(int v, int f) {
   p[v] = f;
   s[v] = 1:
   if (f != -1) d[v] = d[f] + 1;
   else d[v] = 0;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
    int w = g[v][i]:
     if (w != f) {
      dfs(w. v):
      s[v] += s[w]:
 void hld(int v, int f, int k) {
   t[v] = ts++;
   c[k].push_back(v);
   r[v] = k;
   int x = 0, y = -1;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
    int w = g[v][i];
```

```
if (w != f) {
       if (s[w] > x) {
         x = s[w]:
        y = w;
   if (y != -1) {
     hld(y, v, k);
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f && w != y) {
       hld(w, v, w);
   }
  void init(int n) {
   for (int i = 0; i < n; ++i) {</pre>
     g[i].clear();
  void add(int a, int b) {
   g[a].push_back(b);
   g[b].push_back(a);
  void build() {
   ts = 0;
   dfs(0, -1);
   hld(0, 0, 0);
};
```

3.3 persistent array

```
struct node {
  node *1, *r;
  int val;

  node (int x) : l(NULL), r(NULL), val(x) {}
  node () : l(NULL), r(NULL), val(-1) {}
};

typedef node* pnode;

pnode update(pnode cur, int l, int r, int at, int what) {
  pnode ans = new node();

  if (cur != NULL) {
    *ans = *cur;
  }
  if (1 == r) {
    ans-> val = what;
}
```

3.4 persistent seg tree

```
/**
 * Important:
 * When using lazy propagation remember to create new
 * versions for each push_down operation!!!
struct node {
 node *1, *r;
 long long acc;
 node (int x) : 1(NULL), r(NULL), acc(x), flip(0) {}
 node () : 1(NULL), r(NULL), acc(0), flip(0) {}
typedef node* pnode;
pnode create(int 1, int r) {
 if (1 == r) return new node();
 pnode cur = new node();
 int m = (1 + r) >> 1:
 cur-> 1 = create(1, m);
 cur-> r = create(m + 1, r);
 return cur;
pnode copy_node(pnode cur) {
 pnode ans = new node():
 *ans = *cur:
 return ans;
void push_down(pnode cur, int 1, int r) {
 assert(cur):
 if (cur-> flip) {
   int len = r - 1 + 1:
   cur-> acc = len - cur-> acc:
   if (cur-> 1) {
     cur-> 1 = copy_node(cur-> 1);
     cur-> 1 -> flip ^= 1;
```

```
if (cur-> r) {
     cur-> r = copy_node(cur-> r);
     cur-> r -> flip ^= 1:
   cur-> flip = 0;
int get_val(pnode cur) {
 assert(cur):
 assert((cur-> flip) == 0);
 if (cur) return cur-> acc;
 return 0:
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = copy_node(cur);
 if (1 == r) {
   assert(1 == at):
   ans-> acc = what;
   ans-> flip = 0;
   return ans:
 int m = (1 + r) >> 1:
 push_down(ans, 1, r);
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
pnode flip(pnode cur, int 1, int r, int a, int b) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur:
 if (1 > b || r < a)
   return ans;
 if (1 >= a && r <= b) {
   ans-> flip ^= 1;
   push_down(ans, 1, r);
   return ans;
 int m = (1 + r) >> 1;
 ans-> 1 = flip(ans-> 1, 1, m, a, b);
 ans-> r = flip(ans-> r, m + 1, r, a, b);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
long long get_all(pnode cur, int 1, int r) {
 assert(cur):
```

3.5 persistent trie

```
// Persistent binary trie (BST for integers)
const int MD = 31;
struct node_bin {
 node_bin *child[2];
 int val;
  node_bin() : val(0) {
   child[0] = child[1] = NULL;
};
typedef node_bin* pnode_bin;
pnode_bin copy_node(pnode_bin cur) {
  pnode_bin ans = new node_bin();
  if (cur) *ans = *cur:
 return ans;
pnode_bin modify(pnode_bin cur, int key, int inc, int id = MD) {
  pnode_bin ans = copy_node(cur);
  ans->val += inc:
 if (id >= 0) {
   int to = (key >> id) & 1;
   ans->child[to] = modify(ans->child[to], key, inc, id - 1);
 return ans:
int sum_smaller(pnode_bin cur, int key, int id = MD) {
 if (cur == NULL) return 0;
  if (id < 0) return 0: // strictly smaller
 // if (id == - 1) return cur->val; // smaller or equal
  int ans = 0:
  int to = (key >> id) & 1;
  if (to) {
   if (cur->child[0]) ans += cur->child[0]->val:
   ans += sum_smaller(cur->child[1], key, id - 1);
 } else {
   ans = sum_smaller(cur->child[0], key, id - 1);
```

```
return ans;
// Persistent trie for strings.
const int MAX_CHILD = 26;
struct node {
 node *child[MAX CHILD]:
 int val;
 node() : val(-1) {
   for (int i = 0; i < MAX_CHILD; i++) {</pre>
     child[i] = NULL;
};
typedef node* pnode;
pnode copy_node(pnode cur) {
 pnode ans = new node();
 if (cur) *ans = *cur;
 return ans;
pnode set_val(pnode cur, string &key, int val, int id = 0) {
 pnode ans = copy_node(cur);
  if (id >= int(key.size())) {
   ans->val = val;
 } else {
   int t = key[id] - 'a';
   ans->child[t] = set_val(ans->child[t], key, val, id + 1);
 return ans;
pnode get(pnode cur, string &key, int id = 0) {
 if (id >= int(key.size()) || !cur)
   return cur;
 int t = key[id] - 'a';
 return get(cur->child[t], key, id + 1);
```

3.6 segment tree

```
const int MN = 1e5; // limit for array size
struct seg_tree {
  int n; // array size
  int t[2 * MN];
  seg_tree(int _n) : n(_n) {}

  void clear() {
    memset(t, 0, sizeof t);
  }

  void build() { // build the tree
```

```
for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
 // Single modification, range query.
 void modify(int p, int value) { // set value at position p
   for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] +
         t[p^1];
 int query(int 1, int r) { // sum on interval [1, r)
   int res = 0:
   for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
     if (l&1) res += t[l++]:
     if (r&1) res += t[--r]:
   return res:
};
// Range modification, single query.
void modify(int 1, int r, int value) {
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) t[1++] += value:
   if (r&1) t[--r] += value;
 }
}
int query(int p) {
 int res = 0:
 for (p += n; p > 0; p >>= 1) res += t[p];
 return res;
}
 * If at some point after modifications we need to inspect all
 * elements in the array, we can push all the modifications to
 * leaves using the following code. After that we can just
      traverse
 * elements starting with index n. This way we reduce the
      complexity
 * from O(n \log(n)) to O(n) similarly to using build instead of
      n modifications.
 * */
void push() {
 for (int i = 1; i < n; ++i) {</pre>
   t[i<<1] += t[i];
   t[i<<1|1] += t[i]:
   t[i] = 0;
}
// Non commutative combiner functions.
void modify(int p, const S& value) {
 for (t[p += n] = value; p >>= 1; ) t[p] = combine(t[p << 1],
       t[p<<1|1]):
```

```
S querv(int 1, int r) {
 S resl, resr;
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) resl = combine(resl, t[1++]);
   if (r&1) resr = combine(t[--r], resr);
 return combine(resl, resr);
 * segment tree for intervals
const int MN = 100000 + 100;
struct seg tree {
 int val[MN * 4 + 4];
  int pending[MN * 4 + 4];
  seg_tree() {
   memset(val. -1, sizeof val):
   memset(pending, -1, sizeof pending);
  void propagate(int node, int b, int e) {
   if (pending[node] != -1) {
     val[node] = pending[node];
     if (b < e) {
       pending[node << 1] = pending[node];</pre>
       pending[node << 1 | 1] = pending[node];</pre>
     pending[node] = -1;
  void set(int node, int b, int e, int from, int to, int v) {
   if (b > to || e < from) return:</pre>
   if (b >= from && e <= to) {</pre>
     pending[node] = v;
     propagate(node, b, e);
     return;
   int mid = (b + e) >> 1;
   set(node << 1, b, mid, from, to, v);
   set(node << 1 | 1, mid + 1, e, from, to, v):
  int query(int node, int b, int e, int pos) {
   propagate(node, b, e);
   if (b == e && b == pos) {
     return val[node]:
```

```
int mid = (b + e) >> 1;
if (pos <= mid)
    return query(node << 1, b, mid, pos);
return query(node << 1 | 1, mid + 1, e, pos);
}

void set(int from, int to, int v) {
    return set(1, 0, MN - 1, from, to, v);
}

int query(int pos) {
    return query(1, 0, MN - 1, pos);
}
};</pre>
```

3.7 sparse table

```
const int MN = 100000 + 10: // Max number of elements
const int ML = 18; // ceil(log2(MN));
struct st {
  int data[MN];
  int M[MN][ML]:
  void init(const vector<int> &d) {
   n = d.size();
   for (int i = 0: i < n: ++i)
     data[i] = d[i]:
   build():
  void build() {
   for (int i = 0; i < n; ++i)
     M[i][0] = data[i];
   for (int j = 1, p = 2, q = 1; p \le n; ++j, p \le 1, q \le 1)
     for (int i = 0; i + p - 1 < n; ++i)
       M[i][j] = max(M[i][j-1], M[i+q][j-1]);
  int query(int b, int e) {
    int k = log2(e - b + 1);
    return max(M[b][k], M[e + 1 - (1<<k)][k]);
 }
};
```

3.8 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<endl;</pre>
```

```
typedef int T;
struct node{
 node *left, *right, *parent;
 node (T k) : key(k), left(0), right(0), parent(0) {}
};
struct splay_tree{
 node *root:
 void right_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   if (p->left = x->right) p->left->parent = p;
   x->right = p;
   p->parent = x;
 void left_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   if (p->right = x->left) p->right->parent = p;
   x->left = p;
   p->parent = x;
 void splay(node *x, node *fa = 0) {
   while( x->parent != fa and x->parent != 0) {
     node *p = x->parent;
     if (p->parent == fa)
       if (p->right == x)
        left_rot(x);
       else
         right_rot(x);
     else {
       node *gp = p->parent; //grand parent
       if (gp - > left == p)
        if (p->left == x)
          right_rot(x), right_rot(x);
          left_rot(x),right_rot(x);
       else
         if (p->left == x)
          right_rot(x), left_rot(x);
          left_rot(x), left_rot(x);
   if (fa == 0) root = x:
```

```
void insert(T kev) {
   node *cur = root;
   node *pcur = 0;
   while (cur) {
     pcur = cur;
     if (key > cur->key) cur = cur->right;
     else cur = cur->left;
   cur = new node(key);
   cur->parent = pcur;
   if (!pcur) root = cur;
   else if (key > pcur->key ) pcur->right = cur;
   else pcur->left = cur;
   splay(cur);
  node *find(T kev) {
   node *cur = root;
   while (cur) {
     if (key > cur->key) cur = cur->right;
     else if(key < cur->key) cur = cur->left;
     else return cur;
   return 0;
  splay_tree(){ root = 0;};
};
```

3.9 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;
order set X:
int get(int y) {
 int 1=0.r=1e9+1:
 while(l<r) {</pre>
   int m=l+((r-l)>>1);
   if(m-X.order_of_key(m+1)<y)</pre>
     l=m+1;
   else
     r=m:
 return 1;
```

```
main(){
 ios::sync_with_stdio(0);
 cin.tie(0):
 int n,m;
 cin>>n>>m;
  for(int i=0;i<m;i++) {</pre>
    char a:
    int b;
    cin>>a>>b;
    if(a=='L')
     cout<<get(b)<<endl;
     X.insert(get(b));
/***
Input
20 7
L 5
D 5
L 4
L 5
D 5
L 4
I. 5
Output
5
4
6
4
***/
```

3.10 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>

using namespace __gnu_pbds;
using namespace std;

typedef
tree<
   pair<int,int>,
   null_type,
   less<pair<int,int>>,
   rb_tree_tag,
   tree_order_statistics_node_update>
ordered_set;

main()
{
   ios::sync_with_stdio(0);
```

```
cin.tie(0);
   int n;
   int sz=0:
    cin>>n;
   vector<int> ans(n,0);
   ordered_set t;
   int x,y;
   for(int i=0;i<n;i++)</pre>
       cin>>x>>v:
       ans[t.order_of_key(\{x,++sz\})]++;
       t.insert({x.sz}):
   for(int i=0:i<n:i++)</pre>
        cout<<ans[i]<<'\n';</pre>
/***
Input
5
1 1
5 1
7 1
3 3
5.5
Output
1
2
1
1
0
***/
```

3.11 STL Treap

```
#include <ext/rope> //header with rope
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some
     additional stuff
int main()
{
   ios_base::sync_with_stdio(false);
   rope <int> v; //use as usual STL container
   int n. m:
   cin >> n >> m;
   for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
    int 1, r;
   for(int i = 0; i < m; ++i)</pre>
       cin >> 1 >> r;
       --1. --r:
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
```

```
v.insert(v.mutable_begin(), cur);
}
for(rope <int>::iterator it = v.mutable_begin(); it !=
          v.mutable_end(); ++it)
          cout << *it << " ";
    return 0;
}</pre>
```

3.12 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010: // Number of states.
struct trie{
 struct node{
   int c;
   int a[MN];
 };
  node tree[MS];
  int nodes:
  void clear(){
   tree[nodes].c = 0:
   memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++:
  void init(){
   nodes = 0;
   clear():
  int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]):
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
 }
};
```

3.13 wavelet tree

```
// this can be tested in the problem:
    http://www.spoj.com/problems/ILKQUERY/
```

```
struct wavelet {
 vector<int> values. ori:
  vector<int> map_left, map_right;
 int 1, r, m;
 wavelet *left, *right;
  wavelet() : left(NULL), right(NULL) {}
  wavelet(int a, int b, int c) : l(a), r(b), m(c), left(NULL),
       right(NULL) {}
};
wavelet *init(vector<int> &data, vector<int> &ind, int lo, int
  if (lo > hi || (data.size() == 0)) return NULL:
 int mid = ((long long)(lo) + hi) / 2;
 if (lo + 1 == hi) mid = lo; // handle negative values
  wavelet *node = new wavelet(lo, hi, mid);
  vector<int> data_1, data_r, ind_1, ind_r;
  int ls = 0, rs = 0;
  for (int i = 0; i < int(data.size()); i++) {</pre>
   int value = data[i];
   if (value <= mid) {</pre>
     data_l.emplace_back(value);
     ind_l.emplace_back(ind[i]);
     ls++:
   } else {
     data_r.emplace_back(value);
     ind_r.emplace_back(ind[i]);
     rs++;
   node->map_left.emplace_back(ls);
   node->map_right.emplace_back(rs);
   node->values.emplace_back(value);
   node->ori.emplace_back(ind[i]);
 if (lo < hi) {</pre>
   node->left = init(data_1, ind_1, lo, mid);
   node->right = init(data_r, ind_r, mid + 1, hi);
 return node;
int kth(wavelet *node, int to, int k) {
 // returns the kth element in the sorted version of (a[0],
       ..., a[to])
 if (node->1 == node->r) return node->m:
 int c = node->map_left[to];
 if (k < c)
   return kth(node->left, c - 1, k);
 return kth(node->right, node->map_right[to] - 1, k - c);
int pos_kth_ocurrence(wavelet *node, int val, int k) {
 // returns the position on the original array of the kth
       ocurrence of the value "val"
  if (!node) return -1:
```

```
if (node->l == node->r) {
   if (int(node->ori.size()) <= k)
      return -1;
   return node->ori[k];
}

if (val <= node->m)
   return pos_kth_ocurrence(node->left, val, k);
return pos_kth_ocurrence(node->right, val, k);
```

4 Deletion

4.1 board

```
struct board {
  int n, m, r;
  board(int a, int b, int c = 1) : n(a), m(b), r(c) {}

long long frec(int x, int y) {
    // returns how many squares of r x r contain the cell (x, y)
  long long a = min(x, n - r) - max(x - r + 1, 0) + 1;
  long long b = min(y, m - r) - max(y - r + 1, 0) + 1;
  return a * b;
}

bool valid(int x, int y) {
  return x >= 0 && x < n && y >= 0 && y < m;
}
};</pre>
```

4.2 center 2 points + radious

```
vector<point> find_center(point a, point b, long double r) {
  point d = (a - b) * 0.5;
  if (d.dot(d) > r * r) {
    return vector<point> ();
  }
  point e = b + d;
  long double fac = sqrt(r * r - d.dot(d));
  vector<point> ans;
  point x = point(-d.y, d.x);
  long double l = sqrt(x.dot(x));
  x = x * (fac / l);
  ans.push_back(e + x);
  x = point(d.y, -d.x);
  x = x * (fac / l);
  ans.push_back(e + x);
  return ans;
}
```

4.3 counting

```
const int MN = 1e5 + 100;
long long fact[MN];

void fill_fact() {
  fact[0] = 1;
  for (int i = 1; i < MN; i++) {
    fact[i] = mult(fact[i - 1], i);
  }
}

long long perm_rep(vector<int> &frec) {
  int total = 0;
  long long den = 1;
  for (int i = 0; i < (int)frec.size(); i++) {
    den = mult(den, mod_inv(fact[frec[i]]);
    total += frec[i];
  }
  return mult(fact[total], den);
}</pre>
```

4.4 crt

```
/**
  * Chinese remainder theorem.
  * Find z such that z % x[i] = a[i] for all i.
  * */
long long crt(vector<long long> &a, vector<long long> &x) {
  long long z = 0;
  long long n = 1;
  for (int i = 0; i < x.size(); ++i)
    n *= x[i];

  for (int i = 0; i < a.size(); ++i) {
    long long tmp = (a[i] * (n / x[i])) % n;
    tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
    z = (z + tmp) % n;
}

  return (z + n) % n;
}</pre>
```

4.5 cumulative sum of divisors

```
/**
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,
SOD(24) = 2+3+4+6+8+12 = 35.
```

```
The function CSOD(n) (cumulative SOD) of an integer n, is
    defined as below:

csod(n) = \sum_{i = 1}^{n} sod(i)

It can be computed in O(sqrt(n)):
*/

long long csod(long long n) {
  long long ans = 0;
  for (long long i = 2; i * i <= n; ++i) {
    long long j = n / i;
    ans += (i + j) * (j - i + 1) / 2;
    ans += i * (j - i);
  }
  return ans;
}</pre>
```

4.6 dates

```
// Time - Leap years
// A[i] has the accumulated number of days from months previous
273, 304, 334 }:
// same as A, but for a leap year
const int B[13] = { 0, 0, 31, 60, 91, 121, 152, 182, 213, 244,
     274, 305, 335 }:
// returns number of leap years up to, and including, y
int leap_years(int y) { return y / 4 - y / 100 + y / 400; }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 == 0 && y %
     100 != 0); }
// number of days in blocks of years
const int p400 = 400*365 + leap vears(400):
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B[m]
      : A[m]) + d:
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400 block?
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1; // are we in the top year of a 4 block?
 y = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (days-1) \% p400 + 1;
```

```
if (d > p100*3) top100 = true, d -= 3*p100, y += 300;
else y += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;

if (d > p4*24) top4 = true, d -= 24*p4, y += 24*4;
else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;

if (d > p1*3) top1 = true, d -= p1*3, y += 3;
else y += (d-1) / p1, d = (d-1) % p1 + 1;

const int *ac = top1 && (!top4 || top100) ? B : A;
for (m = 1; m < 12; ++m) if (d <= ac[m + 1]) break;
d -= ac[m];</pre>
```

4.7 dijkstra

```
struct edge {
 int to;
 long long w;
 edge () {}
 edge (int a, long long b) : to(a), w(b) {}
 bool operator < (const edge &o) const {</pre>
   return w > o.w:
};
typedef vector<vector<edge>> graph;
const long long inf = 1000000LL * 10000000LL;
pair<vector<int>, vector<long long>> dijkstra(graph &g, int
     start) {
 int n = g.size();
 vector<long long> d(n, inf);
 vector<int> p(n, -1);
 d[start] = 0;
 priority_queue<edge> q;
 q.push(edge(start, 0));
 while (!q.empty()) {
   int node = q.top().to;
   long long dist = q.top().w;
   q.pop();
   if (dist > d[node]) continue;
   for (int i = 0; i < (int)g[node].size(); i++) {</pre>
     int to = g[node][i].to;
     long long w_extra = g[node][i].w;
     if (dist + w_extra < d[to]) {</pre>
       p[to] = node:
       d[to] = dist + w extra:
       q.push(edge(to, d[to]));
```

```
return {p, d};
```

4.8 eulerian path

```
// Taken from
     https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail
struct Euler {
 ELV adi: IV t:
  Euler(ELV Adj) : adj(Adj) {}
  void build(int u) {
   while(! adj[u].empty()) {
     int v = adj[u].front().v;
     adi[u].erase(adj[u].begin());
     build(v);
   t.push_back(u);
};
bool eulerian trail(IV &trail) {
 Euler e(adj);
  int odd = 0, s = 0;
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false;
    if (diff == 1) {
    if (++odd > 2) return false;
    if (out[v] > in[v]) start = v;
    */
  e.build(s);
  reverse(e.t.begin(), e.t.end());
  trail = e.t;
  return true:
```

4.9 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long
    long &y, long long &g) {
    x = 0, y = 1, g = b;
    long long m, n, q, r;
    for (long long u = 1, v = 0; a != 0; g = a, a = r) {
        q = g / a, r = g % a;
        m = x - u * q, n = y - v * q;
        x = u, y = v, u = m, v = n;
    }
}
```

4.10 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

4.11 io

```
// taken from :
     https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo
// this is very fast as well :
     https://github.com/lbv/pc-code/blob/master/code/input.cpp
typedef unsigned int u32:
#define BUF 524288
struct Reader {
 char buf[BUF]; char b; int bi, bz;
 Reader() { bi=bz=0; read(); }
 void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
 void skip() { while (b > 0 && b <= 32) read(); }</pre>
 u32 next u32() {
   u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48;
        return v: }
 int next_int() {
   int v = 0; bool s = false;
   skip(); if (b == '-') { s = true; read(); }
   for (; 48<=b&&b<=57; read()) v = v*10 + b-48; return s ? -v
 char next_char() { skip(); char c = b; read(); return c; }
};
```

4.12 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$
 where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$
 and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n}=0$ if $m\leq n$.

4.13 matrix

```
const int MN = 111:
const int mod = 10000;
struct matrix {
 int r, c;
 int m[MN][MN];
 matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
 void print() {
   for (int i = 0; i < r; ++i) {
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl;</pre>
   }
 }
 int x[MN][MN];
 matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)</pre>
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {</pre>
          x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod))
   memcpy(m, x, sizeof(m));
   return *this;
};
void matrix_pow(matrix b, long long e, matrix &res) {
 memset(res.m. 0, sizeof res.m);
 for (int i = 0; i < b.r; ++i)</pre>
   res.m[i][i] = 1;
 if (e == 0) return;
 while (true) {
   if (e & 1) res *= b;
   if ((e >>= 1) == 0) break;
   b *= b:
 }
```

4.14 mod integer

```
template<class T, T mod>
```

```
struct mint_t {
   T val;
   mint_t() : val(0) {}
   mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
   return (val + o.val) % mod;
   }

mint_t operator - (const mint_t& o) const {
   return (val - o.val) % mod;
   }

mint_t operator * (const mint_t& o) const {
   return (val * o.val) % mod;
   }

stypedef mint_t<long long, 998244353> mint;
```

4.15 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
    return 0;
  return (x + m) % m;
}
```

4.16 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
        x = (x + y) % mod;
    y = (y * 2) % mod;
    b /= 2;
  }
  return x % mod;
}
```

4.17 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
    if (exp & 1)
        ans = mod_mul(ans, a, mod);
    a = mod_mul(a, a, mod);
}
```

```
exp >>= 1;
}
return ans;
```

4.18 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

```
f + v = e + 2
```

It can be extended to non connected planar graphs with c connected components:

```
f + v = e + c + 1
```

4.19 polynomials

```
const double pi = acos(-1):
struct poly {
 deque<double> coef;
 double x_lo, x_hi;
  double evaluate(double x) {
   double ans = 0;
   for (auto it : coef)
     ans = (ans * x + it):
   return ans;
  double volume(double x, double dx=1e-6) {
   dx = (x_hi - x_lo) / 1000000.0;
   double ans = 0;
   for (double ix = x_lo; ix <= x; ix += dx) {
     double rad = evaluate(ix):
     ans += pi * rad * rad * dx;
   return ans;
};
```

4.20 primes

```
namespace primes {
  const int MP = 100001;
```

```
bool sieve[MP];
long long primes[MP];
int num_p;
void fill_sieve() {
 num_p = 0;
 sieve[0] = sieve[1] = true;
 for (long long i = 2; i < MP; ++i) {</pre>
   if (!sieve[i]) {
     primes[num_p++] = i;
     for (long long j = i * i; j < MP; j += i)
       sieve[j] = true;
 }
// Finds prime numbers between a and b, using basic primes up
     to sqrt(b)
// a must be greater than 1.
vector<long long> seg_sieve(long long a, long long b) {
 long long ant = a;
 a = max(a, 3LL);
 vector<bool> pmap(b - a + 1);
 long long sqrt_b = sqrt(b);
 for (int i = 0; i < num_p; ++i) {</pre>
   long long p = primes[i];
   if (p > sqrt_b) break;
   long long j = (a + p - 1) / p;
   for (long long v = (j == 1) ? p + p : j * p; v <= b; v +=
         p) {
     pmap[v - a] = true;
 vector<long long> ans;
 if (ant == 2) ans.push_back(2);
 int start = a % 2 ? 0 : 1;
 for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
   if (pmap[i] == false)
     ans.push_back(a + i);
 return ans;
vector<pair<int, int>> factor(int n) {
 vector<pair<int, int>> ans;
 if (n == 0) return ans;
 for (int i = 0; primes[i] * primes[i] <= n; ++i) {</pre>
   if ((n % primes[i]) == 0) {
     int expo = 0;
     while ((n % primes[i]) == 0) {
      expo++:
       n /= primes[i];
     ans.emplace_back(primes[i], expo);
 if (n > 1) {
   ans.emplace_back(n, 1);
 return ans;
```

+

4.21 query with lca

```
struct lowest_ca {
 int T[MN], L[MN], W[MN];
 int P[MN][ML], MI[MN][ML], MA[MN][ML];
 void dfs(vector<vector<edge> > &g, int root, int pi = -1) {
   if (pi == -1) {
     L[root] = W[root] = 0;
     T[root] = -1:
   for (int i = 0; i < (int)g[root].size(); ++i) {</pre>
     int to = g[root][i].v;
     if (to != pi) {
       T[to] = root;
       W[to] = g[root][i].w;
       L[to] = L[root] + 1;
       dfs(g, to, root);
    }
   }
  void init(vector<vector<edge> > &g, int root) {
   // g is undirected
   dfs(g, root);
   int N = g.size(), i, j;
   for (i = 0; i < N; i++) {</pre>
     for (j = 0; 1 << j < N; j++) {
       P[i][i] = -1;
       MI[i][j] = inf;
   for (i = 0: i < N: i++) {</pre>
    P[i][0] = T[i];
     MI[i][0] = W[i];
   for (i = 1: 1 << i < N: i++)
     for (i = 0; i < N; i++)</pre>
       if (P[i][j - 1] != -1) {
        P[i][j] = P[P[i][j-1]][j-1];
        MI[i][j] = min(MI[i][j-1], MI[P[i][j-1]][j-1]);
 }
 int query(int p, int q) {
   int tmp, log, i;
   int mmin = inf:
   if (L[p] < L[q])</pre>
     tmp = p, p = q, q = tmp;
   for (log = 1; 1 << log <= L[p]; log++);</pre>
```

```
log--;
    for (i = log: i >= 0: i--)
     if (L[p] - (1 << i) >= L[q]) {
       mmin = min(mmin, MI[p][i]);
       p = P[p][i];
    if (p == q)
     // return p;
     return mmin;
    for (i = log; i >= 0; i--)
     if (P[p][i] != -1 && P[p][i] != P[q][i]) {
       mmin = min(mmin, min(MI[p][i], MI[q][i]));
       p = P[p][i], q = P[q][i];
    // return T[p]:
   return min(mmin, min(MI[p][0], MI[q][0]));
  int get_child(int p, int q) { // p is ancestor of q
   if (p == q) return -1;
   int i, log;
    for (log = 1; 1 << log <= L[q]; log++) {}
   log--;
    for (i = log; i >= 0; i--)
     if (L[q] - (1 << i) > L[p]) {
       q = P[q][i];
    assert(P[q][0] == p);
   return q;
  int is_ancestor(int p, int q) {
   if (L[p] >= L[q])
     return false;
    int dist = L[q] - L[p];
    int cur = q;
    int step = 0;
    while (dist) {
     if (dist & 1)
       cur = P[cur][step];
     step++;
     dist >>= 1;
    return cur == p;
};
```

4.22 sliding window

```
* Given an array ARR and an integer K, the problem boils down
      to computing for each index i: min(ARR[i], ARR[i-1],
      .... ARR[i-K+1]).
 * if mx == true, returns the maximun.
      http://people.cs.uct.ac.za/~ksmith/articles/sliding_window_minimum.html
vector<int> sliding_window_minmax(vector<int> & ARR, int K,
     bool mx) {
 deque< pair<int, int> > window;
 vector<int> ans;
 for (int i = 0; i < ARR.size(); i++) {</pre>
   if (mx) {
     while (!window.empty() && window.back().first <= ARR[i])</pre>
       window.pop_back();
   } else {
     while (!window.empty() && window.back().first >= ARR[i])
       window.pop_back();
   window.push_back(make_pair(ARR[i], i));
   while(window.front().second <= i - K)</pre>
    window.pop_front();
   ans.push_back(window.front().first);
 return ans;
}
```

4.23 squares

```
typedef long double ld;
const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
    return ( x <= y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}
struct point{
    ld x, y;
    point(ld a, ld b) : x(a), y(b) {}
    point() {}
};
struct square{
    ld x1, x2, y1, y2,
        a, b, c;
    point edges[4];
    square(ld _a, ld _b, ld _c) {
        a = _a, b = _b, c = _c;
    }
</pre>
```

```
x1 = a - c * 0.5;
    x2 = a + c * 0.5:
    v1 = b - c * 0.5:
    v2 = b + c * 0.5;
    edges[0] = point(x1, y1);
    edges[1] = point(x2, y1);
    edges[2] = point(x2, y2);
    edges[3] = point(x1, y2);
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x.
    y = a.y - b.y;
  return sqrt(x * x + y * y);
bool point_in_box(square s1, point p) {
  if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
     cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
    return true:
  return false:
bool inside(square &s1, square &s2) {
  for (int i = 0; i < 4; ++i)
   if (point_in_box(s2, s1.edges[i]))
     return true;
  return false;
bool inside_vert(square &s1, square &s2) {
  if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) != 1))
    return true;
 return false;
bool inside_hori(square &s1, square &s2) {
  if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) != 1))
    return true:
 return false;
ld min_dist(square &s1, square &s2) {
  if (inside(s1, s2) || inside(s2, s1))
   return 0:
  ld ans = 1e100:
  for (int i = 0; i < 4; ++i)</pre>
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));
  if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.v1, s2.v2) != -1)
     ans = min(ans, s1.y1 - s2.y2);
```

```
if (cmp(s2.y1, s1.y2) != -1)
    ans = min(ans, s2.y1 - s1.y2);
}

if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
    if (cmp(s1.x1, s2.x2) != -1)
        ans = min(ans, s1.x1 - s2.x2);
    else
    if (cmp(s2.x1, s1.x2) != -1)
        ans = min(ans, s2.x1 - s1.x2);
}

return ans;
}
```

4.24 totient function

```
const int MN = 1e6 + 5;
int E[MN];
void Init(){
  for(int i = 1; i < MN; i++){
    E[i] = i;
  }
  for(int i = 1; i < MN; i++){
    for(int x = (i << 1); x < MN; x += i){
       E[x] -= E[i];
    }
}</pre>
```

4.25 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

4.26 totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
      while ((n % primes[i]) == 0) n /= primes[i];
      ans -= ans / primes[i];
   }
}</pre>
```

```
if (n > 1) {
   ans -= ans / n;
}
return ans;
}
```

5 DP Optimizations

5.1 convex hull trick

```
/**
 * Problems:
     http://codeforces.com/problemset/problem/319/C
     http://codeforces.com/contest/311/problem/B
     https://csacademy.com/contest/archive/task/squared-ends
     http://codeforces.com/contest/932/problem/F
 * */
struct line {
 long long m, b;
 line (long long a, long long c) : m(a), b(c) {}
 long long eval(long long x) {
   return m * x + b;
 }
};
long double inter(line a, line b) {
 long double den = a.m - b.m;
 long double num = b.b - a.b;
 return num / den:
 * min m_i * x_j + b_i, for all i.
      x_j \le x_{j+1}
      m_i >= m_{j+1}
struct ordered_cht {
 vector<line> ch;
 int idx; // id of last "best" in query
 ordered cht() {
   idx = 0;
 void insert_line(long long m, long long b) {
   line cur(m, b):
   // new line's slope is less than all the previous
   while (ch.size() > 1 &&
      (inter(cur, ch[ch.size() - 2]) >= inter(cur,
           ch[ch.size() - 1]))) {
       // f(x) is better in interval [inter(ch.back(), cur).
            inf)
       ch.pop_back();
   ch.push_back(cur);
```

```
long long eval(long long x) { // minimum
   // current x is greater than all the previous x,
    // if that is not the case we can make binary search.
    idx = min<int>(idx, ch.size() - 1);
    while (idx + 1 < (int)ch.size() && ch[idx + 1].eval(x) <=
         ch[idx].eval(x))
     idx++;
    return ch[idx].eval(x);
};
 * Dynammic convex hull trick
typedef long long int64;
typedef long double float128;
const int64 is_query = -(1LL<<62), inf = 1e18;</pre>
struct Line {
 int64 m, b;
  mutable function<const Line*()> succ:
  bool operator<(const Line& rhs) const {</pre>
   if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    if (!s) return 0:
    int64 x = rhs.m:
   return b - s->b < (s->m - m) * x;
};
struct HullDynamic : public multiset<Line> { // will maintain
     upper hull for maximum
  bool bad(iterator y) {
   auto z = next(y);
    if (y == begin()) {
     if (z == end()) return 0;
     return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
    return (float128)(x->b - y->b)*(z->m - y->m) >=
         (float128)(y->b - z->b)*(y->m - x->m);
  void insert_line(int64 m, int64 b) {
    auto y = insert({ m, b });
    y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
    if (bad(y)) { erase(y); return; }
    while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
  int64 eval(int64 x) {
    auto 1 = *lower_bound((Line) { x, is_query });
    return 1.m * x + 1.b:
```

```
}
};
```

5.2 divide and conquer

```
/*
  * recurrence:
  * dp[k][i] = min dp[k-1][j] + c[i][j - 1], for all j > i;
  *
  * "comp" computes dp[k][i] for all i in O(n log n) (k is fixed)
  */

void comp(int 1, int r, int le, int re) {
  if (l > r) return;
  int mid = (l + r) >> 1;
  int best = max(mid + 1, le);
  dp[cur][mid] = dp[cur ^ 1][best] + cost(mid, best - 1);
  for (int i = best; i <= re; i++) {
    if (dp[cur][mid] > dp[cur ^ 1][i] + cost(mid, i - 1)) {
      best = i;
      dp[cur][mid] = dp[cur ^ 1][i] + cost(mid, i - 1);
  }
}

comp(l, mid - 1, le, best);
  comp(mid + 1, r, best, re);
}
```

6 Geometry

7 Graphs

7.1 bridges

```
struct Graph {
  vector<vector<Edge>> g;
  vector<int> vi, low, d, pi, is_b;
  int bridges_computed;

int ticks, edges;

Graph(int n, int m) {
    g.assign(n, vector<Edge>());
    is_b.assign(m, 0);
    vi.resize(n);
    low.resize(n);
    d.resize(n);
    edges = 0;
    bridges_computed = 0;
}
```

```
void AddEdge(int u, int v) {
   g[u].push back(Edge(v. edges)):
   g[v].push_back(Edge(u, edges));
   edges++;
 void Dfs(int u) {
   vi[u] = true;
   d[u] = low[u] = ticks++;
   for (int i = 0; i < (int)g[u].size(); ++i) {</pre>
     int v = g[u][i].to;
     if (v == pi[u]) continue;
     if (!vi[v]) {
       pi[v] = u;
       Dfs(v):
       if (d[u] < low[v]) is_b[g[u][i].id] = true;</pre>
       low[u] = min(low[u], low[v]);
     } else {
       low[u] = min(low[u], d[v]);
 // Multiple edges from a to b are not allowed.
 // (they could be detected as a bridge).
 // If you need to handle this, just count
 // how many edges there are from a to b.
 void CompBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0);
   fill(low.begin(), low.end(), 0);
   fill(d.begin(), d.end(), 0);
   ticks = 0;
   for (int i = 0; i < (int)g.size(); ++i)</pre>
    if (!vi[i]) Dfs(i);
   bridges_computed = true;
 map<int, vector<Edge>> BridgesTree() {
   if (!bridges_computed) CompBridges();
   int n = g.size();
   Dsu dsu(g.size());
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (!is_b[e.id]) dsu.Join(i, e.to);
   map<int, vector<Edge>> tree;
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (is b[e.id])
         tree[dsu.Find(i)].emplace_back(dsu.Find(e.to), e.id);
   return tree;
};
```

7.2 directed mst

```
const int inf = 1000000 + 10:
struct edge {
 int u, v, w;
 edge() {}
 edge(int a, int b, int c) : u(a), v(b), w(c) {}
* Computes the minimum spanning tree for a directed graph
 * - edges : Graph description in the form of list of edges.
 * each edge is: From node u to node v with cost w
 * - root : Id of the node to start the DMST.
 * - n : Number of nodes in the graph.
 * */
int dmst(vector<edge> &edges, int root, int n) {
 int ans = 0;
 int cur_nodes = n;
 while (true) {
   vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
   for (int i = 0; i < edges.size(); ++i) {</pre>
     int u = edges[i].u, v = edges[i].v, w = edges[i].w;
     if (w < lo[v] and u != v) {
      lo[v] = w;
       pi[v] = u;
   lo[root] = 0;
   for (int i = 0; i < lo.size(); ++i) {</pre>
     if (i == root) continue;
     if (lo[i] == inf) return -1;
   int cur_id = 0;
   vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
   for (int i = 0: i < cur nodes: ++i) {</pre>
     ans += lo[i];
     while (u != root and id[u] < 0 and mark[u] != i) {</pre>
       mark[u] = i;
       u = pi[u]:
     if (u != root and id[u] < 0) { // Cycle
        for (int v = pi[u]; v != u; v = pi[v])
         id[v] = cur_id;
        id[u] = cur id++:
   if (cur_id == 0)
     break:
   for (int i = 0; i < cur_nodes; ++i)</pre>
     if (id[i] < 0) id[i] = cur_id++;</pre>
   for (int i = 0; i < edges.size(); ++i) {</pre>
```

```
int u = edges[i].u, v = edges[i].v, w = edges[i].w;
  edges[i].u = id[u];
  edges[i].v = id[v];
  if (id[u] != id[v])
    edges[i].w -= lo[v];
}
cur_nodes = cur_id;
root = id[root];
}
return ans;
}
```

7.3 karp min mean cycle

```
* Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 * test: uva, 11090 - Going in Cycle!!
const int MN = 1000;
struct edge{
 int v;
 long long w;
 edge(){} edge(int v, int w) : v(v), w(w) {}
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
 g.resize(n + 1): // this is important
 for (int i = 0; i < n; ++i)</pre>
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
 for(int i = 0:i < n:++i)
   fill(d[i],d[i]+(n+1),INT_MAX);
 d[n - 1][0] = 0:
 for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
   if (d[u][k - 1] == INT_MAX) continue;
   for (int i = g[u].size() - 1; i >= 0; --i)
     d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] +
          g[u][i].w):
 bool flag = true;
```

```
for (int i = 0; i < n && flag; ++i)</pre>
 if (d[i][n] != INT_MAX)
   flag = false;
if (flag) {
 return true; // return true if there is no a cycle.
double ans = 1e15;
for (int u = 0; u + 1 < n; ++u) {</pre>
 if (d[u][n] == INT_MAX) continue;
 double W = -1e15:
 for (int k = 0; k < n; ++k)
   if (d[u][k] != INT MAX)
     W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));
 ans = min(ans, W):
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;</pre>
return false;
```

7.4 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

7.5 minimum path cover in DAG

Given a directed acyclic graph G=(V,E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

```
Vout = \{v \in V : v \text{ has positive out} - degree\} Vin = \{v \in V : v \text{ has positive in} - degree\} E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}
```

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necesarily disjoints, find the transitive closure and solve the problem for disjoint paths.

7.6 SCC kosaraju

```
struct SCC {
 vector<vector<int> > g, gr;
 vector<bool> used:
 vector<int> order, component;
 int total_components;
 SCC(vector<vector<int> > &adj) {
   g = adj;
   int n = g.size();
   gr.resize(n);
   for (int i = 0; i < n; i++)
    for (auto to : g[i])
       gr[to].push_back(i);
   used.assign(n, false);
   for (int i = 0; i < n; i++)
    if (!used[i])
       GenTime(i);
   used.assign(n, false);
   component.assign(n, -1);
   total components = 0:
   for (int i = n - 1; i \ge 0; i--) {
     int v = order[i]:
     if (!used[v]) {
       vector<int> cur_component;
       Dfs(cur_component, v);
       for (auto node : cur_component)
         component[node] = total_components;
       total_components++;
  void GenTime(int node) {
   used[node] = true;
   for (auto to : g[node])
     if (!used[to])
       GenTime(to);
   order.push back(node):
  void Dfs(vector<int> &cur. int node) {
   used[node] = true;
   cur.push_back(node);
   for (auto to : gr[node])
     if (!used[to])
       Dfs(cur, to);
  vector<vector<int>> CondensedGraph() {
   vector<vector<int>> ans(total_components);
   for (int i = 0; i < int(g.size()); i++) {</pre>
```

```
for (int to : g[i]) {
    int u = component[i], v = component[to];
    if (u != v)
        ans[u].push_back(v);
    }
    return ans;
}
```

7.7 tarjan scc

```
const int MN = 20002;
struct tarjan_scc {
 int scc[MN], low[MN], d[MN], stacked[MN];
 int ticks, current_scc;
 deque<int> s; // used as stack.
 tarjan_scc() {}
 void init () {
   memset(scc, -1, sizeof scc);
   memset(d, -1, sizeof d);
   memset(stacked, 0, sizeof stacked);
   s.clear();
   ticks = current_scc = 0;
  void compute(vector<vector<int> > &g, int u) {
   d[u] = low[u] = ticks++;
   s.push_back(u);
   stacked[u] = true;
    for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i];
     if (d[v] == -1)
       compute(g, v);
     if (stacked[v]) {
       low[u] = min(low[u], low[v]);
    if (d[u] == low[u]) { // root
     do {
       v = s.back():s.pop back():
       stacked[v] = false;
       scc[v] = current_scc;
     } while (u != v):
     current_scc++;
 }
};
```

7.8 two sat (with kosaraju)

```
* Given a set of clauses (a1 v a2)^(a2 v a3)....
 * this algorithm find a solution to it set of clauses.
 * test: http://lightoj.com/volume_showproblem.php?problem=1251
#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'
vector<int> G[MAX]:
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC:
bool visited[MAX];
int n:
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i];
   if (visited[curr]) continue;
   dfs1(curr):
 Ftime.push_back(n);
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1:
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i]:
   if (visited[curr]) continue;
   dfs2(curr, scc);
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0: i < 2 * n : ++i) {
   if (!visited[i]) dfs1(i);
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue:
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
```

```
* After having the SCC, we must traverse each scc, if in one
      SCC are -b y b, there is not a solution.
 * Otherwise we build a solution, making the first "node" that
      we find truth and its complement false.
bool two sat(vector<int> &val) {
 kosaraju();
 for (int i = 0: i < SCC.size(): ++i) {</pre>
   vector<bool> tmpvisited(2 * n. false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
     else {
       val[SCC[i][i]] = 0:
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][j]] = 1;
 return 1;
// Example of use
int main() {
 int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number, m = clauses number
 while (t--) {
   cin >> m >> n;
   Ftime.clear();
   SCC.clear();
   for (int i = 0; i < 2 * n; ++i) {
    G[i].clear():
     GT[i].clear();
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m; ++i) {</pre>
    cin >> u >> v;
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
     G[p 1].push_back(q);
     G[q ^ 1].push_back(p);
     GT[p].push_back(q ^ 1);
     GT[q].push_back(p ^ 1);
   vector < int > val(2 * n, -1);
   cout << "Case " << ++nc <<": ";
   if (two sat(val)) {
     cout << "Yes" << endl:
```

```
vector<int> sol;
for (int i = 0; i < 2 * n; ++i)
    if (i % 2 == 0 and val[i] == 1)
        sol.push_back(i / 2 + 1);
    cout << sol.size();

for (int i = 0; i < sol.size(); ++i) {
        cout << " " << sol[i];
    }
    cout << endl;
} else {
    cout << "No" << endl;
}
return 0;
}</pre>
```

8 Math

8.1 fft

```
* Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f \operatorname{star} g)[n] = \operatorname{sum}_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
using namespace std:
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'
const int MN = 262144 << 1;</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0):
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
   real = real:
   image = _image;
 cpx(){}
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
cpx operator - (const cpx &c1, const cpx &c2) {
```

```
return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image,
       c1.real*c2.image + c1.image*c2.real);
int rev(int id, int len) {
 int ret = 0;
 for (int i = 0; (1 << i) < len; i++) {</pre>
   ret <<= 1;
   if (id & (1 << i)) ret |= 1;</pre>
 return ret;
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s):
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
   for(int k = 0; k < len; k += m) {</pre>
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
       w = w * wm:
   }
 if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /=
       len, A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return;
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d);
 for (int i = 0; i < n; ++i) {</pre>
   cin >> t:
   d[t] = true;
 }
 int m;
 cin >> m;
 vector<int> q(m);
 for (int i = 0; i < m; ++i)
   cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
     in[i] = cpx(1, 0):
```

```
else
     in[i] = cpx(0, 0);
  FFT(in, MN, 1);
  for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
  int ans = 0:
  for (int i = 0; i < q.size(); ++i) {</pre>
   if (in[q[i]].real > 0.5 || d[q[i]]) {
     ans++:
 cout << ans << endl;
int main() {
 ios_base::sync_with_stdio(false);cin.tie(NULL);
  int n;
  while (cin >> n)
   solve(n);
  return 0;
```

8.2 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

8.3 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d \mid n} d^x$$

when x=0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^{k} g(e_k + 1)$$

where q(x) is the sum of the first x positive numbers:

$$q(x) = (x * (x + 1))/2$$

8.4 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

9 Matrix

10 Misc

11 Number theory

11.1 convolution

```
//check x is 2^a
inline bool is_pow2(LL x);
inline int ceil_log2(LL x) {
 int ans = 0:
 --x;
 while (x != 0) {
  x >>= 1;
   ans++;
 }
 return ans:
/* Returns the convolution of the two given vectors in time
     proportional to n*log(n).
 * The number of roots of unity to use nroots_unity must be set
      so that the product of the first
 * nroots_unity primes of the vector nth_roots_unity is greater
      than the maximum value of the
 * convolution. Never use sizes of vectors bigger than 2^24, if
      you need to change the values of
 * the nth roots of unity to appropriate primes for those sizes.
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b,
     int nroots_unity = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1;
 for (int times = 0; times < nroots_unity; times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0; i < a.size(); i++) fA[i] = a[i];</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv modulo = mod inv(modulo % prime, prime):
   LL normalize = mod_inv(N, prime);
   ntfft(fA. 1. nth roots unitv[times]):
   ntfft(fB, 1, nth roots unitv[times]):
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;</pre>
   ntfft(fC, -1, nth roots unitv[times]):
   for (int i = 0; i < N; i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) % prime;
     k = (k * inv_modulo) % prime;
     ans[i] += modulo * k:
   modulo *= prime;
 return ans;
```

11.2 diophantine equations

```
long long gcd(long long a, long long b, long long &x, long long
     &v) {
  if (a == 0) {
   x = 0:
   y = 1;
   return b;
 long long x1, y1;
 long long d = gcd(b \% a, a, x1, y1);
 x = v1 - (b / a) * x1;
 y = x1;
 return d;
bool find_any_solution(long long a, long long b, long long c,
     long long &x0,
   long long &y0, long long &g) {
  g = gcd(abs(a), abs(b), x0, y0);
  if (c % g) {
   return false;
  x0 *= c / g;
  y0 *= c / g;
 if (a < 0) x0 = -x0;
  if (b < 0) v0 = -v0;
  return true;
void shift_solution(long long &x, long long &y, long long a,
     long long b,
   long long cnt) {
 x += cnt * b:
 y -= cnt * a;
long long find_all_solutions(long long a, long long b, long
     long c.
   long long minx, long long maxx, long long miny,
   long long maxv) {
  long long x, y, g;
  if (!find_any_solution(a, b, c, x, y, g)) return 0;
  a /= g;
  b /= g;
  long long sign_a = a > 0 ? +1 : -1;
 long long sign_b = b > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x) / b);
  if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
  if (x > maxx) return 0;
  long long lx1 = x;
```

```
shift_solution(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift_solution(x, y, a, b, -sign_b);
long long rx1 = x;

shift_solution(x, y, a, b, -(miny - y) / a);
if (y < miny) shift_solution(x, y, a, b, -sign_a);
if (y > maxy) return 0;
long long lx2 = x;

shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
long long rx2 = x;

if (lx2 > rx2) swap(lx2, rx2);
long long lx = max(lx1, lx2);
long long rx = min(rx1, rx2);

if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
}
```

11.3 discrete logarithm

```
// Computes x which a \hat{x} = b \mod n.
long long d_log(long long a, long long b, long long n) {
 long long m = ceil(sqrt(n));
 long long aj = 1;
 map<long long, long long> M;
 for (int i = 0; i < m; ++i) {</pre>
   if (!M.count(aj))
     M[ai] = i:
   aj = (aj * a) % n;
 long long coef = mod_pow(a, n - 2, n);
 coef = mod_pow(coef, m, n);
 // coef = a^{(-m)}
 long long gamma = b;
 for (int i = 0; i < m; ++i) {</pre>
   if (M.count(gamma)) {
     return i * m + M[gamma];
   } else {
     gamma = (gamma * coef) % n;
 return -1;
```

11.4 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not prime, 1 < a < n
bool witness(long long a, long long n) {</pre>
```

```
// check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0:
 while (u % 2 == 0) {
   t++:
   u >>= 1;
 long long next = mod pow(a, u, n):
 if (next == 1) return false;
 long long last;
 for (int i = 0: i < t: ++i) {
   last = next;
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1;
 return next != 1;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller rabin(999999999971LL) == 1):
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false:
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 for (int i = 0: i < it: ++i) {</pre>
   long long a = rand() % (n - 1) + 1;
   if (witness(a, n)) {
     return false:
   }
 }
 return true;
}
```

11.5 number theoretic transform

```
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
  int n = a.size():
  LL prime = root unity.first:
  LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
  if (dir < 0) basew = mod_inv(basew, prime);</pre>
  for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1;
   LL w = 1:
    for (int i = 0; i < mh; i++) {</pre>
     for (int j = i; j < n; j += m) {</pre>
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) % prime;
       a[j] = (a[j] + a[k]) % prime;
       a[k] = (w * x) \% prime:
     w = (w * basew) % prime;
   basew = (basew * basew) % prime;
  int i = 0:
  for (int j = 1; j < n - 1; j++) {</pre>
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
}
```

11.6 pollard rho factorize

```
long long pollard_rho(long long n) {
 long long x, y, i = 1, k = 2, d;
 x = y = rand() \% n;
 while (1) {
   ++i:
   x = mod_mul(x, x, n);
   x += 2;
   if (x \ge n) x -= n;
   if (x == v) return 1:
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d;
   if (i == k) {
    y = x;
     k *= 2:
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans:
 if (miller rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
```

```
d = pollard_rho(n);
vector<long long> dd = factorize(d);
ans = factorize(n / d);
for (int i = 0; i < dd.size(); ++i)
ans.push_back(dd[i]);
}
return ans;
}</pre>
```

12 Strings

12.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26;
 static const int AlphabetBase = 'a';
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum;
   Node(): fail(NULL), next{}, sum(0) { }
 struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0;
   strings.clear();
   roots.clear():
   sizes.clear();
   que.resize(totalLen);
 void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push back(nodes.data() + nNodes):
   sizes.push back(1):
   nNodes += (int)str.size() + 1;
   auto check = [&]() { return sizes.size() > 1 &&
         sizes.end()[-1] == sizes.end()[-2]; };
   if(!check())
     makePMA(strings.end() - 1, strings.end(), roots.back(),
          que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m:
     if(!check())
```

```
makePMA(strings.end() - m * 2, strings.end(),
          roots.back(), que);
 }
int match(const string &str) const {
 int res = 0;
 for(const Node *t : roots)
   res += matchPMA(t, str);
 return res;
static void makePMA(vector<String>::const iterator begin.
     vector<String>::const_iterator end, Node *nodes,
     vector<Node*> &que) {
 int nNodes = 0;
 Node *root = new(&nodes[nNodes ++]) Node();
 for(auto it = begin: it != end: ++ it) {
   Node *t = root;
   for(char c : it->str) {
     Node *&n = t->next[c - AlphabetBase];
     if(n == nullptr)
      n = new(&nodes[nNodes ++]) Node():
     t = n:
   t->sum += it->sign;
  int qt = 0;
  for(Node *&n : root->next) {
   if(n != nullptr) {
     n->fail = root;
     que[qt ++] = n;
   } else {
     n = root;
  for(int qh = 0; qh != qt; ++ qh) {
   Node *t = que[qh];
   int a = 0;
   for(Node *n : t->next) {
     if(n != nullptr) {
       que[qt ++] = n;
       Node *r = t->fail;
       while(r->next[a] == nullptr)
        r = r->fail;
       n->fail = r->next[a];
       n->sum += r->next[a]->sum;
static int matchPMA(const Node *t, const string &str) {
 int res = 0:
 for(char c : str) {
   int a = c - AlphabetBase;
   while(t->next[a] == nullptr)
     t = t->fail:
```

```
t = t-\text{next}[a];
     res += t->sum;
   return res;
 vector<Node> nodes:
 int nNodes;
 vector<String> strings;
 vector<Node*> roots:
 vector<int> sizes;
 vector<Node*> que:
int main() {
 int m;
 while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac:
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
     if(ty == 1) {
       iac.insert(s, +1);
     } else if(ty == 2) {
       iac.insert(s, -1);
     } else if(ty == 3) {
       int ans = iac.match(s);
       printf("%d\n", ans);
       fflush(stdout);
     } else {
       abort();
   }
 }
 return 0;
```

12.2 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
    string s;
    cin >> s;
    int n = s.size();
    s += s;
    vector<int> f(s.size(), -1);
    int k = 0;
    for (int j = 1; j < 2 * n; ++j) {
        int i = f[j - k - 1];
        while (i != -1 && s[j] != s[k + i + 1]) {
            if (s[j] < s[k + i + 1])
            k = j - i - 1;
            i = f[i];
        }
}</pre>
```

```
if (i == -1 && s[j] != s[k + i + 1]) {
    if (s[j] < s[k + i + 1]) {
        k = j;
    }
    f[j - k] = -1;
} else {
    f[j - k] = i + 1;
    }
}
return k;
}</pre>
```

12.3 suffix array

```
/**
 * 0 (n log^2 (n))
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf
      for reference
struct entry{
 int a, b, p;
 entrv(){}
 entry(int x, int y, int z): a(x), b(y), p(z){}
 bool operator < (const entry &o) const {</pre>
   return (a == o.a) ? (b == o.b) ? (p < o.p) : (b < o.b) :
          (a < o.a);
};
struct SuffixArray{
 const int N;
 string s;
 vector<vector<int> > P;
 vector<entry> M;
 SuffixArray(const string &s) : N(s.length()), s(s), P(1,
       vector<int> (N, 0)), M(N) {
    for (int i = 0; i < N; ++i)</pre>
     P[0][i] = (int) s[i];
    for (int skip = 1, level = 1; skip < N; skip *= 2, level++)</pre>
     P.push back(vector<int>(N, 0)):
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip] :</pre>
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)
       P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a \text{ and}
             M[i].b == M[i - 1].b) ? P[level][M[i - 1].p] : i;
  vector<int> getSuffixArray(){
```

```
vector<int> &rank = P.back();
   vector<pair<int, int> > inv(rank.size());
   for (int i = 0: i < rank.size(): ++i)</pre>
     inv[i] = make_pair(rank[i], i);
   sort(inv.begin(), inv.end());
   vector<int> sa(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     sa[i] = inv[i].second:
   return sa;
 // returns the length of the longest common prefix of
       s[i...L-1] and s[j...L-1]
 int lcp(int i, int i) {
   int len = 0:
   if (i == j) return N - i;
   for (int k = P.size() - 1; k >= 0 && i < N && j < N; --k) {</pre>
     if (P[k][i] == P[k][j]) {
       i += 1 << k:
       j += 1 << k;
       len += 1 << k;
   return len:
};
```

12.4 suffix automaton

```
* Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * The overall complexity is O(n)
      https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */
struct state {
 int len. link:
 long long num_paths;
 map<int, int> next;
};
const int MN = 200011:
state sa[MN << 1]:
int sz, last;
long long tot_paths;
void sa_init() {
 sz = 1:
```

```
last = 0;
 sa[0].len = 0;
 sa[0].link = -1:
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot_paths = 0;
void sa_extend(int c) {
 int cur = sz++:
 sa[cur].len = sa[last].len + 1;
 sa[cur].next.clear();
 sa[cur].num_paths = 0;
 int p;
  for (p = last; p != -1 && !sa[p].next.count(c); p =
       sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
 if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++:
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
       sa[p].next[c] = clone;
       sa[q].num_paths -= sa[p].num_paths;
       sa[clone].num_paths += sa[p].num_paths;
     sa[q].link = sa[cur].link = clone;
 last = cur;
```

12.5 z algorithm

```
vector<int> z(n,0);
int 1,r;
r = 1 = 0;
```

```
for(int i = 1; i < n; ++i){</pre>
   if(i > r) {
     while(r < n and s[r - 1] == s[r])r++;
     z[i] = r - 1;r--;
   }else{
     int k = i-1;
     if(z[k] < r - i +1) z[i] = z[k]:
     else {
      1 = i:
       while (r < n \text{ and } s[r - 1] == s[r])r++;
       z[i] = r - 1; r--;
 return z;
int main(){
 //string line;cin>>line;
 string line = "alfalfa";
 vector<int> z = compute_z(line);
 for(int i = 0; i < z.size(); ++i ){</pre>
   if(i)cout<<" ";
   cout<<z[i];
 cout<<endl:
 // must print "0 0 0 4 0 0 1"
 return 0;
```

13 Tmp

13.1 convexHull

13.2 seg intersection

13.3 windingNum