1. The system flow and the explanation of the process(functions).

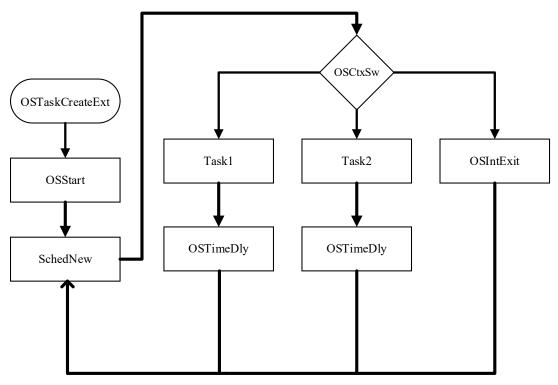


Fig 1. System flow

一開始,先從OSTaskCreateExt建立兩個task,進入到OSStart裡面的SchedNwe去設定參數,並將優先權更新,透過OSCtxSw去切換目前最高優先權,它會透過OSTimeDly來決定優先度,當有空窗期時,像tick(3)時,它會切至Taskidle再到OSIntExit再回到task(1)繼續動作,當如果遇到像 $6 \cdot 9 \cdot 12$ 等數字會先以,task(1)->task(2)->task(63),其中只會進OSIntExit都會是task(63),接下來就重複循環。

2. The screenshot of the result.

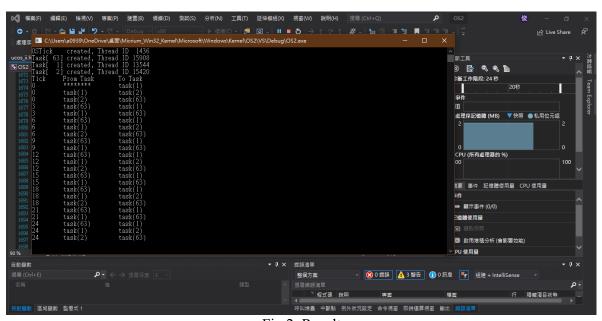


Fig 2. Result

3. A report that describes your implementation (please attach the screenshot of the code and **MARK** the modified part).

Fig 3. Code1

在圖4.中,OS\_Sched前面會先執行最高優先權的運算,因此我將printf放入其中,可以得知 Tick、From Task 、TO Task目前的狀態。

```
Eif OS_TASK_CREATE_EXT_EN > Ou

Eif defined(OS_TLS_TBL_SIZE) && (OS_TLS_TBL_SIZE > Ou)

| OS_TLS_TaskSw();
| endif |
| endif |
| printf("%d \t task(%d) \t task(%d)\n",OSTimeGet(), OSPrioCur, OSPrioHighRdy); //1019 add code |
| OS_TASK_SW(); /* Perform a context switch */
| }
| OS_EXIT_CRITICAL();
```

Fig 4. Code2

在圖5.中,透過講義中的Embedded OS p6得知,OSIntExit()也會計算最高優先權,因此我也將 printf放在OSIntCtxSw之前,已得知目前的狀態。

```
cif OS_TASK_CREATE_EXT_EN > Ou
cif defined(OS_TLS_TBL_SIZE) && (OS_TLS_TBL_SIZE > Ou)

OS_TLS_TaskSw();

endif

oS_TRACE_ISR_EXIT_TO_SCHEDULER();

printf("%d \t task(%d) \t task(%d) \n", OSTimeGet(), OSPrioCur, OSPrioHighRdy); //1019 add code

OSIntCtxSw(); /* Perform interrupt level ctx switch */

} else {

OS_TRACE_ISR_EXIT();

}

else {

OS_TRACE_ISR_EXIT();
```

Fig 5. Code3