# 嵌入式系統軟體設計 Embedded System Software Design

PA2

指導教授: 陳雅淑 教授

課程學生: M10907305 陳俊億

## • Part1:

1. Execution result of using mutex and barrier. 20%

```
chen@chen-VirtualBox:~/桌面/pa2$ ./part1.out
                   =====System Info======
Protect Shared Resource: Mutex
Synchronize: Barrier
Program ID : 0 Thread ID : 0 PID : 6785
Single-thread spend time : 121.299
                                              Core : 6
 =======Start Multi-Thread Matrix Multiplication========
Program ID : 0 Thread ID : 1
Program ID : 0 Thread ID : 0
                                PID: 6795
                                                Core: 1
                                PID : 6794
                                                 Core: 0
Program ID : 0 Thread ID : 3
Program ID : 0 Thread ID : 2
                                PID: 6797
                                                 Core: 3
                              PID : 6796
                                                Core: 2
Multi-thread spend time : 210.228
                  :=======Result=:
Program-0 obtain correct matrix multiplication result.
```

Fig1. Result of using mutex and barrier.

2. Describe how to synchronize thread. 10%

```
sharedSum = new int [PROGRAM_NUM];

/*-----Your code(PARTIGPART3)-----*/

pthread_barrier_init(&System::barr,NULL,PROGRAM_NUM*THREAD_NUM);

pthread_spin_init(&Thread::lock,pshared);

sem_init(&System::sem,1,PROGRAM_NUM*THREAD_NUM);

sem_init(&System::sem1,1,0);
```

Fig2. Code of synchronize thread. (1)

Fig3. Code of synchronize thread. (2)

```
210     obj->synchronize();
211     for (int i = obj->startCalculatePoint; i < obj->endCalculatePoint; i++){
212          memcpy (obj->matrix [i], obj->multiResult [i], obj->matrixSize * sizeof (int));
213     }
214     obj->synchronize();
```

Fig4. Code of synchronize thread. (3)

首先先建立 barrier 的初始設定,如圖 2 所示,透過測試需要在 memcpy 上下要進行同步,否則數值會有錯誤,可透過圖 3 的功能來進行同步。 3. Describe how to protect a shared resource. 5%

```
public:
   pthread_t pthreadThread;
   static pthread_spinlock_t lock;
   static pthread_mutex_t count_mutex;
```

Fig5. Code of protect a shared resource. (1)

Fig6. Code of protect a shared resource. (2)

先在 thread.h 宣告 mutex 變數,如圖 6 所示,而主要保護的共用資源為 shareSum,因此它的上下進行 mutex 的共用資源保護。

#### • Part2:

1. Execution result of using reentrant function. 15%

Fig7. Result of using reentrant function.

2. Describe how to modify non-reentrant function into reentrant function. 10%

Fig8. reentrant function.

要實現 reentrant function, 捨棄了共用資源(全局參數), 改用一般參數 (局部參數), 如圖 8 所示。

3. Describe the reason why using a non-reentrant function or a reentrant function could obtain better performance. 5%

從圖 1 與圖 7 的結果上看,使用 reentrant function 效果較 non-reentrant function 好,首先 non-reentrant 會使用到 global variables or static variables 因此在不同的 thread 中需要使用到類似於 mutex 去保護共用資源,以免被干擾,但是 mutex 會鎖死當下的 thread 等到做完之後才會讓其他 thread 動作,導致時間會比 reentrant 慢。

## • Part3:

1. Execution result of using spinlock. 10%

Fig9. Result of using spinlock.

2. Describe which method (mutex and spinlock) could obtain better performance under the benchmark we provided (5%) and why (5%).

Fig10. Code of using spinlock.

以圖1以圖9的結果來比較,使用 spinlock 較優於 mutex,其原因 spinlock 會一直 busy waiting 等待解鎖,而 mutex 只限定於當下的 thread 可以動作,而從特性中得知,spinlock 有利於保護單一變數(global variables),而 mutex 比較適合保護一段程式,在圖 10 中,基本上都算是保護單一參數,因此 spinlock 較 mutex 佳。

3. Show the benchmark your used (5%), explain the properties of such benchmark (5%) and the execution results (5%).

#### 表 1. 參數測試表:

測試組/參數名稱	PROGRAM_NUM	MATRIX_SIZE	MULTI_TIME
1	3	500	2
2	2	500	2

Fig12. execution results of part3.out.(測試組 1)

Fig13. execution results of part1.out. (測試組 1)

Fig14. execution results of part3.out.(測試組 2) Fig15. execution results of part1.out. (測試組 2)

PROGRAM\_NUM 在我理解為開設每顆 core 各開幾個 thread,假設program\_num 等於 2 就相當於在每個 core 上開兩個 thread,如圖 12 與圖 13 所示,program\_num=3 時每顆 core 各開 3 個 thread,而 Matrix\_size 為運算的矩陣大小,最後一個 Multi\_time 重複計算次數,但因為每次一個迴圈算完之後會更新matrix 的數值,所以每次迴圈完的結果會不一樣。從上述的圖 12 與圖 13 得知mutex 反而比 spinlock 效果好,主要原因是 mutex 多了 2 個 thread 如果有一個被擋住還可以切到其他同一顆 core 的 thread 繼續做,但是 spinlock 還是一直在 busywaiting,而圖 14 與圖 15 也是相同例子只是修改 PROGRAM\_NUM,結果與剛剛一樣。

# • Bonus: using semaphore.

1. Describe how to use semaphore.

```
sem_init(&System::sem,1,PROGRAM_NUM*THREAD_NUM);
sem_init(&System::sem1,1,0);
```

Fig16. Define the Sme\_init.

```
then \
sed -i "/#define PROTECT_SHARED_RESOURCE/c\#define PROTECT_SHARED_RESOURCE SPINLOCK" ./src/config.h; \
sed -i "/#define SYNCHRONIZE/c\#define SYNCHRONIZE SEMAPHORE" ./src/config.h; \
else \
sed -i "/#define PROTECT_SHARED_RESOURCE/c\#define PROTECT_SHARED_RESOURCE MUTEX" ./src/config.h; \
sed -i "/#define PROTECT_SHARED_RESOURCE/c\#define PROTECT_SHARED_RESOURCE MUTEX" ./src/config.h; \
sed -i "/#define SYNCHRONIZE/c\#define SYNCHRONIZE SEMAPHORE" ./src/config.h; \
fi
```

Fig17. Modify the makefile.

```
#elif SYNCHRONIZE == SEMAPHORE
   sem wait(sem);
   pthread mutex lock (ioMutex);
    ++sem1 key;
   pthread_mutex_unlock (ioMutex);
    if( sem1 key == THREAD NUM*PROGRAM NUM){
        pthread mutex lock (ioMutex);
        for(int i = 0; i < THREAD NUM*PROGRAM NUM; <math>++i){
            sem post(sem1);
       pthread mutex unlock (ioMutex);
    sem_post(sem);
   sem wait(sem1);
   sem2_key = 0;
    -- sem1 key;
   while(1){
        if(sem1_key==0 or sem2_key ==1){
   sem2_key = 1;
```

Fig18. Semaphore function

透過 Semaphore 實現同步功能,首先跟 barrier 一樣先做初始化設定,為了實現同步功能,我使用到兩個 semaphore 去實現,如圖 14 所示,第一個 semaphore 也是要決定當下有幾個人可以拿 key,在此我設定為program\_num\*thread\_num 的數量,也是同時我有幾個 thread 的可以執行,但是每個 thread 的執行速度不一,不能確保說他們速度都一樣,有時候像是 thread 1 跑很快,有機會 thread 8 還沒拿到 key,thread1 就搶先拿到 key 執行第二次計算,會導致在 memcpy 那裡出現錯誤,所以透過另一個 smeaphore 去擋住其他已做好的 thread,現在期可做到同步功能,且在 makefile 需要設定為 semaphore,如圖 15 所示。

2. Modify benchmark to show execution result.

#### 表 2. 參數測試表:

測試組/參數名稱	PROGRAM_NUM	MATRIX_SIZE	MULTI_TIME
1	3	500	2

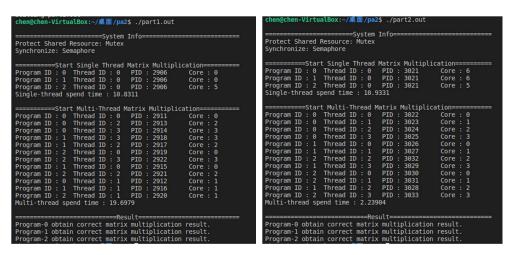


Fig19. Result of part1.out.

Fig20. Result of part2.out.

```
chen@chen-VirtualBox:~/桌面/pa2$ ./part3.out
 Protect Shared Resource: Spinlock
Synchronize: Semaphore
Program ID : 0 Thread ID : 0 PID : 2790 Core : 7
Program ID : 1 Thread ID : 0 PID : 2790 Core : 7
Program ID : 1 Thread ID : 0 PID : 2790 Core : 7
Program ID : 2 Thread ID : 0 PID : 2790 Core : 2
Single-thread spend time : 11.7135
                                =Start Multi-Thread Matrix Multiplication==
: 0 Thread ID : 0 PID : 2794 Core
: 1 Thread ID : 1 PID : 2799 Core
                                                                                                                                               Core : 0
Core : 1
Core : 3
Core : 2
Core : 3
Core : 1
 Program ID :
Program ID :
                                                                                                PID : 2799
PID : 2801
PID : 2796
PID : 2797
PID : 2795
PID : 2805
  Program ID
                                                  Thread ID
 Program ID :
Program ID :
Program ID :
Program ID :
                                                 Thread ID:
                                                 Thread
                                                 Thread
Thread
                                                                     ID
                                        0
2
2
1
Program ID : 2 Thread ID : 3 P
Program ID : 1 Thread ID : 0 P
Program ID : 1 Thread ID : 2 P
Program ID : 2 Thread ID : 2 P
Program ID : 2 Thread ID : 2 P
Multi-thread spend time : 25.5049
                                                                                                PID : 2803
PID : 2798
PID : 2800
PID : 2804
PID : 2802
                                                                                                                                               Core : 3
Core : 1
Core : 0
Core : 2
Core : 0
 Program-0 obtain correct matrix multiplication result.
Program-1 obtain correct matrix multiplication result.
Program-2 obtain correct matrix multiplication result.
```

Fig21. Result of part3.out